CS492: Probabilistic Programming Markov Chain Monte Carlo

Hongseok Yang KAIST

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Really about: Metropolis-Hastings algorithm

(doquery:lmh induce-fn [ints2 outs2])

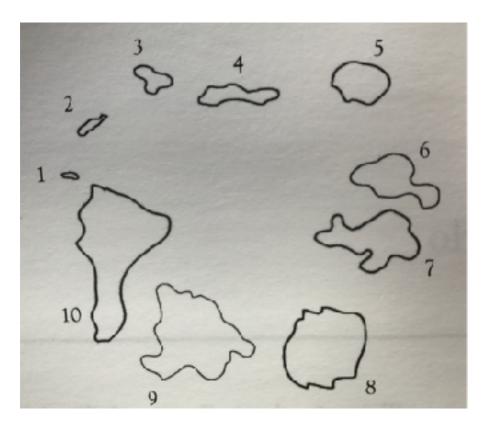
(doquery: lmh induce-fn [ints2 outs2])



Lightweight Metropolis Hastings algorithm* (LMH).

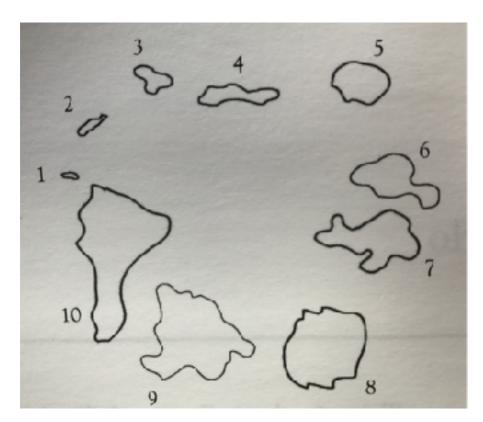
Learning outcome

- Can explain Metropolis-Hastings algorithm.
- Can say when and why this algo. is correct.
- Can develop an instance of the algorithm.



Markov rules 10 islands.

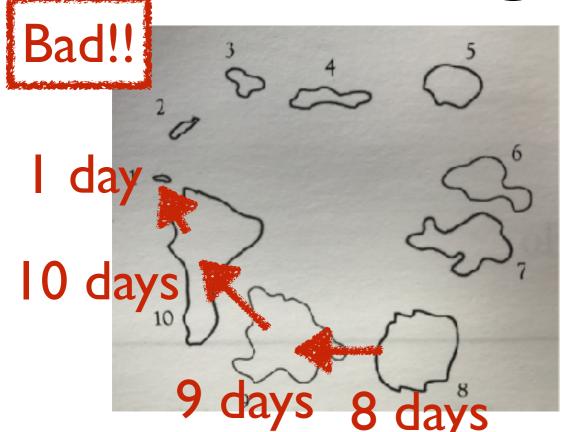
100i people live in island i.



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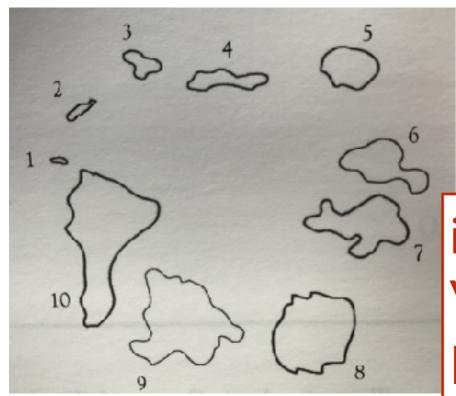
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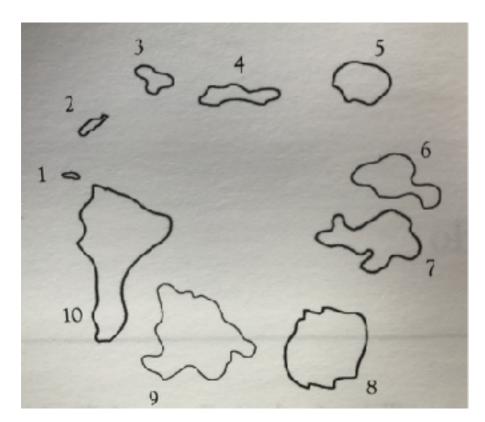
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i ~ discrete(1,2,...,10).Visit i.Repeat.



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Solution

k_n — island that the king visits at step n.

Repeat the following steps:

- I. Flip a coin with prob. 0.5. If head, pick next k' clockwise. If tail, use k' counterclockwise.
- 2. $\alpha := \min(1,k'/k_n)$.
- 3. Flip a coin with prob. α . If head, $k_{n+1} := k'$. Otherwise, $k_{n+1} := k_n$.

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[Q] Why correct? What does correctness even mean?

Sequence by the algo.: k_1 , k_2 , ..., k_n , ...

Corr. informally: Frequency represents probability.

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Corr. formally: For any $f:\{1,..,10\} \rightarrow \mathbb{R}$, $(\sum_{j\leq n}f(k_j))/n \longrightarrow \mathbb{E}_{p(i)}[f(i)]$ as $n\longrightarrow \infty$ with prob. I, where p(i)=i/55, target prob. for visiting island i.

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Holds because I) the random move of the algo. has p as invariant; 2) the algo. can move between any two islands in finitely many (>0) steps.

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[Q] Prove 1) and 2).

Metropolis algorithm

Goal: Generate samples from target r(x)/Z, where $Z=\int r(x)dx$, the normalising constant.

Parameter: Conditional distribution q(x'|x).

- Should be symmetric: q(x'|x) = q(x|x').
- Represents a random move.
- Called proposal kernel.

E.g.
$$r(i)=i$$
, $Z=55$, $q(j|i)=0.5\times[(j-i) \mod 10 \in \{1,0,-1\}]$

Metropolis algorithm

Target r(x)/Z. Symmetric proposal q(x'|x).

- I. initialise x_1 randomly; n:=1
- 2. repeat:
 - a) $x' \sim q(x'|x_n)$; $\alpha := \min(1, r(x')/r(x_n))$
 - b) $u \sim uniform(0, I)$
 - c) $x_{n+1} := (if (u \le \alpha) then x' else x_n); n:=n+1$

Metropolis Noisy greedy exploration.

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 - a) $x' \sim q(x'|x_n)$; $\alpha := \min(1, r(x')/r(x_n))$
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- \geq I for better x'
- < I for worse x'
- c) $x_{n+1} := (if (u \le \alpha) then x' else x_n); n:=n+1$

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May use flip(α) instead, as in our sol. for King Markov

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[QI] Does each step preserve r(x)/Z as invariant?

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```
r(x) = p(y|x)p(x), Noisy greedy exploration.

q(x'|x) = normal(x'_1|x_1,\epsilon_1) No need to know Z.
               \times normal(x'_2|x_2, \varepsilon_2)
```

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- (1) $x \in \mathbb{R}$, $y \in \mathbb{R}^2$
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- (3) $x \in \mathbb{R} \cup \mathbb{R}^2$, $y \in \mathbb{R}^2$

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Prob. distr. on result

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Execute all sample exprs. Prob. distr. on all samples.

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$$\overline{r(x_n) \times q(x'|x_n)}$$
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- c) $x_{n+1} := (if (u \le \alpha) then x' else x_n); n:=n+1$ [QI] q by re-execution for Anglican. What is q(x'|x)?

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 $r(x')\times q(x_n|x')$

[QI] q by re-execution for Anglican. What is q(x'|x)? [AI] q(x'|x) = p(x').

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Target r(x)/Z. Symmetric pro

When independent proposal is used.

- I. initialise x_1 randomly; n:=1
- 2. repeat:

$$\frac{r(x) \times q(x_n)}{r(x_n) \times q(x')}$$

- a) $x' \sim q(x'|x_n)$; $\alpha := \min(1, \frac{r(x_n) \times q(x')}{r(x_n)})$
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[Q3] Does each step have r(x)/Z as invariant?

Recap of the MH algo.

- Generate samples from unnormalised r(x). No need to know $Z=\int r(x)dx$.
- Noisy greedy exploration using q(x'|x).

[Thm I] Each step of MH has r/Z as inv. dist.

```
MH samples: x_1, x_2, x_3, ..., x_n, ... [Thm2] For all f:X \to \mathbb{R} with \mathbb{E}_{r(x)/Z}[f(x)] defined, \sum_{i \le n} f(x_i)/n \longrightarrow \mathbb{E}_{r(x)/Z}[f(x)] \text{ as } n \longrightarrow \infty \text{ with prob. I,} if the MH with q is r/Z-irreducible.
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The estimate converges to the right value.

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MH moves well. For any x,x' with r(x')>0, the MH can go from x to x' with non-zero prob.

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Consequence of a general result in ergodic theory.

Thm I plays a crucial role in the proof.

Reference

I looked at Chapters 5 and 6 of Robert & Casella's "Monte Carlo Statistical Methods".

Not recommended for general reading.

But details and pointers can be found there.