

Ques Range sum queries

1D array

q queries

0	1	2	3	4
1	2	3	4	5

1 query

s = 2

e = 4

e0 e1 e2 e3 e4

→ 12

for (i = s; i <= e; i++)

sum += arr[i];

TC → $O(n * q)$

psum

0	1	2	3	4
1	3	6	10	15

Sum of all the elements from 0 to 3

What

why

how

→ $psum[e] - psum[s] + arr[s]$

$(e0 + e1 + e2 + e3 + e4) - (e0 + e1 + e2)$

+ e2

~~e0 + e1 + e2 + e3 + e4~~ - ~~e0 - e1 - e2~~

+ e2

⇒ $e2 + e3 + e4$

TC ⇒ $O(n + q)$

for (i = 1; i < n; i++)

psom[i] = psom[i-1] + arr[i];

$\Rightarrow O(n)$

answer 1 query $\Rightarrow TC \Rightarrow O(1)$

q $\Rightarrow TC \Rightarrow O(n)$

$TC \Rightarrow O(n + q)$

Query Submatrix Sum Queries 2D array
q queries

	0	1	2	3	4
0	1	2	3	4	6
1	5	3	8	1	2
2	4	6	7	5	5
3	2	4	8	9	4

1 query {

tr = 1

trc = 2

br = 2

brc = 3

}

for (r = tlr; r <= br; r++) {

for (c = tlc; c <= brc; c++) {

sum += arr[r][c];

}

}

TC $\Rightarrow O(\underline{n \times m} \times q)$

↓

1 query

	0	1	2	3	4
0	1	2	3	4	6
1	5	3	8	1	2
2	4	6	7	5	5
3	2	4	8	9	4

arr

	tlr brc				
		c		a	
tlr			tlr tlc		
brr		b		s	

psum

what
why
how

tlr = 1

tlc = 2

brr = 2

brc = 3

a(tlr-1, brc)

b(brr, tlc-1)

c(tlr-1, tlc-1)

s(brr, brc)

ans = s - a - b + c

= s - a - b + c

$$TC \Rightarrow O(n \times m + q)$$

	0	1	2	3	4
0	1	2	3	4	6
1	5	3	8	1	2
2	4	6	7	5	5
3	2	4	8	9	4

arr

	0	1	2	3	4
0	1	3	6	10	16
1	6	11	22	27	35
2	10				
3	12				

sum

```

public void fillPrefixSum() {
    sum[0][0] = matrix[0][0];

    // filling first row
    for(int i = 1; i < sum[0].length; i++)
        sum[0][i] = sum[0][i - 1] + matrix[0][i];

    // filling first column
    for(int i = 1; i < sum.length; i++)
        sum[i][0] = sum[i-1][0] + matrix[i][0];

    for(int r = 1; r < sum.length; r++) {
        for(int c = 1; c < sum[0].length; c++) {
            sum[r][c] = sum[r - 1][c]
                + sum[r][c-1]
                - sum[r - 1][c - 1]
                + matrix[r][c];
        }
    }
}

```