EXP NO: 09 DATE:

DEVELOP THE BACK-END OF A COMPILER THAT TAKES THREE-ADDRESS CODE (TAC) AS INPUT AND GENERATES CORRESPONDING 8086 ASSEMBLY LANGUAGE CODE AS OUTPUT.

AIM:

To design and implement the back-end of a compiler that takes three-address code (TAC) as input and produces 8086 assembly language code as output. The three-address code is an intermediate representation used by compilers to break down expressions and operations, while the 8086 assembly code is a machine-level representation of the program that can be executed by a processor.

ALGORITHM:

1. Parse the Three-Address Code (TAC):

Input: Three-Address Code, which is an intermediate representation. For example:

t0 = b + c t1 = t0 * d a = t1

Output: 8086 assembly language code. For example:

MOV AX, [b]; Load b into AX ADD AX, [c]; Add c to AX MOV [t0], AX; Store result in t0

2. Process Each TAC Instruction:

- 1. Initialize Registers:
 - o Set up the registers in 8086 assembly (e.g., AX, BX, CX, etc.) for storing intermediate results and final outputs. o Maintain a temporary register counter for naming temporary variables in TAC (e.g., t0, t1).
- 2. For each TAC instruction, based on its operation:
 - o Identify the components: operands and operator. o Choose an appropriate register (AX, BX, etc.) for storing intermediate results. o If the operation involves multiple operands or temporary variables, map them to registers.

- 3. Translating TAC to 8086 Assembly:
 - Addition/Subtraction (e.g., t0 = b + c):
 - o Load operands into registers and perform the operation:

MOV AX, [b] ; Load b into AX ADD AX, [c] ; Add c to AX MOV [t0], AX ; Store result in t0

- Multiplication (e.g., t1 = t0 * d):
 - o Load operands into registers and perform the operation:

MOV AX, [t0] ; Load t0 into AX
MOV BX, [d] ; Load d into BX
MUL BX : Multiply AX by BX (result in

MUL BX; Multiply AX by BX (result in AX) MOV [t1],

AX; Store result in t1

- Assignment (e.g., a = t1):
 - o Move the value from a temporary variable to the target variable:

MOV [a], [t1]; Move value of t1 into a

- Division (e.g., t2 = b / c):
 - Division is a bit more complex due to the 8086's limitations with the DIV instruction. For example, the result might need to be stored in AX or DX:AX (if it's a 32-bit result):

MOV AX, [b]; Load b into AX

MOV DX, 0; Clear DX (important for division)

MOV BX, [c]; Load c into BX

DIV BX; AX = AX / BX (quotient in AX, remainder in DX) MOV [t2],

AX; Store quotient in t2

4. Manage Memory and Registers:

- Variables: Variables like a, b, c are stored in memory, so you will use memory addressing modes such as [variable name] to access them.
- Temporary Variables: Temporary variables like t0, t1, t2, etc., are stored in registers (AX, BX, etc.) or memory if there are more variables than registers available.

5. Handle Control Flow (Optional):

If the TAC contains control structures (such as loops, if-else statements, or function calls), you will need to generate labels and jump instructions in 8086 assembly.

• If Statements: For example, if $(x > 0) \{ y = 1; \}$ could generate:

```
MOVAX, [x]
           CMPAX, 0
           JG positive case; Jump if greater
           JMP end if
    PROGRAM:
#include <stdio.h>
#include <string.h>
#define MAX LINES 100
#define MAX LEN 100
int main() {
  char tac[MAX_LINES][MAX_LEN];
  int count = 0;
  printf("Enter TAC instructions (Ctrl+Z to stop):\n");
  // Read all input first
  while (fgets(tac[count], sizeof(tac[count]), stdin)) {
tac[count][strcspn(tac[count], "\n")] = '\0'; // Remove newline
count++;
  }
  printf("\n--- 8086 Assembly Output ---\n");
  // Process each line after input is complete
                                              for (int i = 0; i <
                  char lhs[20], op1[20], op2[20], op;
count; i++) {
(sscanf(tac[i], "%s = %s %c %s", lhs, op1, &op, op2) == 4) {
if (op == '+')
```

```
printf("MOV AX, \lceil \% s \rceil \ n", op1);
printf("ADD AX, [%s]\n", op2);
printf("MOV [%s], AX\n\n", lhs);
       } else if (op == '-') {
printf("MOV AX, [%s]\n", op1);
printf("SUB AX, [%s]\n", op2);
printf("MOV [%s], AX\n\n", lhs);
       } else if (op == '*') {
printf("MOV AX, \lceil \% s \rceil \ n", op1);
printf("MOV BX, [%s]\n", op2);
printf("MUL BX\n");
printf("MOV [%s], AX\n\n", lhs);
       } else if (op == '/') {
printf("MOV AX, [%s]\n", op1);
printf("MOV DX, 0\n");
printf("MOV BX, [%s]\n", op2);
printf("DIV BX\n");
printf("MOV [%s], AX\n\n", lhs);
    } else if (sscanf(tac[i], "%s = %s", lhs, op1) == 2) {
printf("MOV AX, [%s]\n", op1);
       printf("MOV [%s], AX\n\n", lhs);
     } else {
       printf("; Unsupported TAC format: %s\n\n", tac[i]);
  return 0;
```

OUTPUT:

MOV AX, [b]
ADD AX, [c]
MOV [t0], AX

MOV AX, [t0]
MOV BX, [d]
MUL BX
MOV [t1], AX

MOV AX, [t1]
MOV [a], AX

Implementation	
Output/Signature	

RESULT:

Thus the above example provides a foundational approach to converting TAC to 8086 assembly using C. For a complete compiler back-end, you would need to handle additional aspects such as register allocation memory management, and more complex control flow constructs.