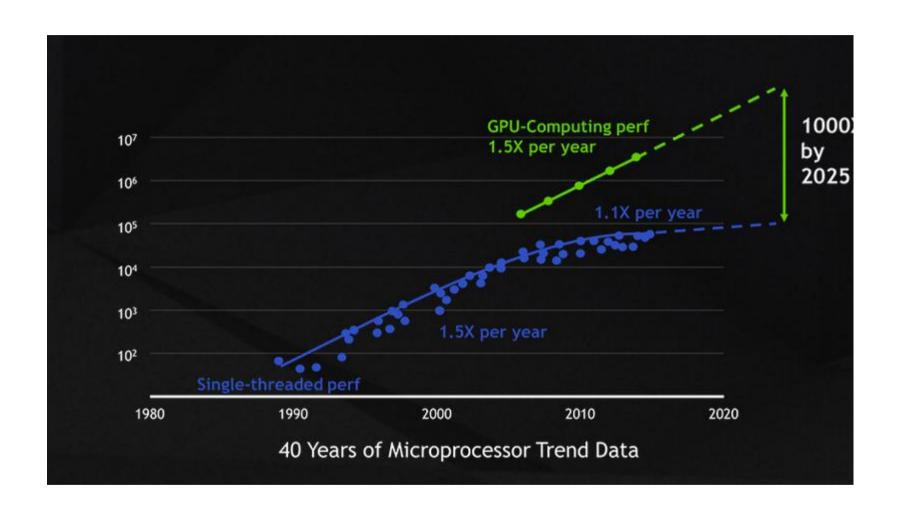
## High Performance Computing on GPUs

Mark Silberstein mark@ee.technion.ac.il

### Why GPUs?



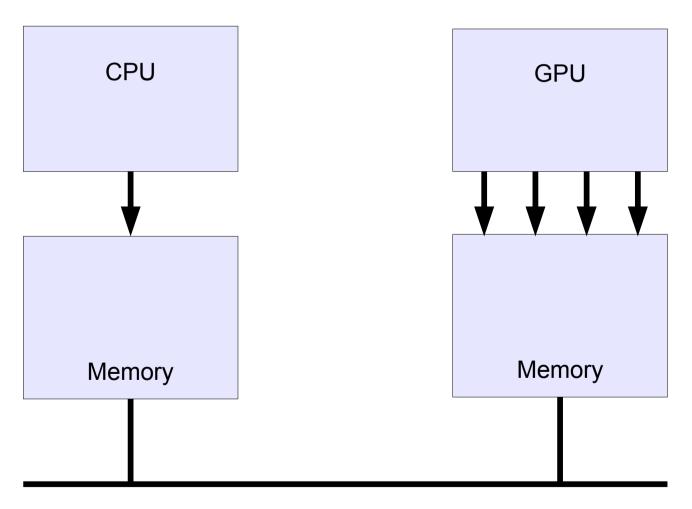
(from NVIDIA)

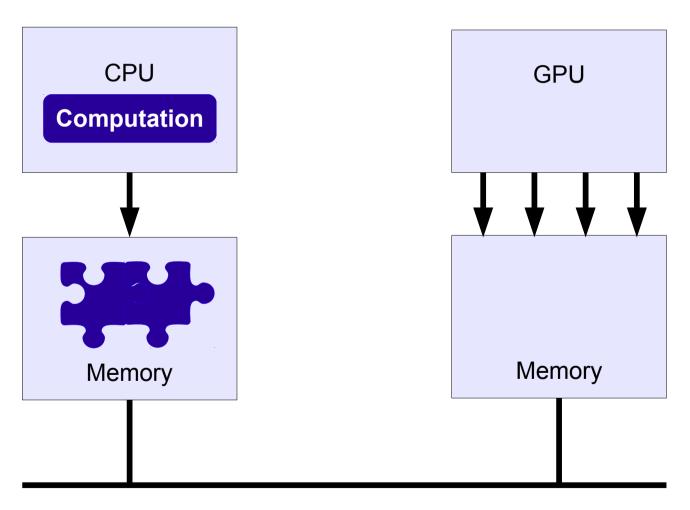
#### Is it a miracle? NO!

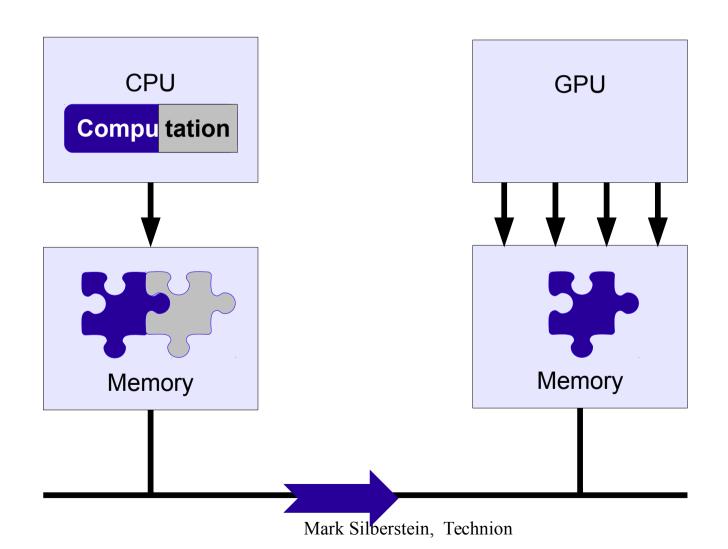
- Architectural solutions prefers parallelism!
- Example problem I have 100 apples to eat
  - 1) "high performance": finish one apple faster
  - 2) "high throughput": finish all apples faster
- The 1<sup>st</sup> option is unsustainable
- Performance = parallel hardware + scalable parallel program!

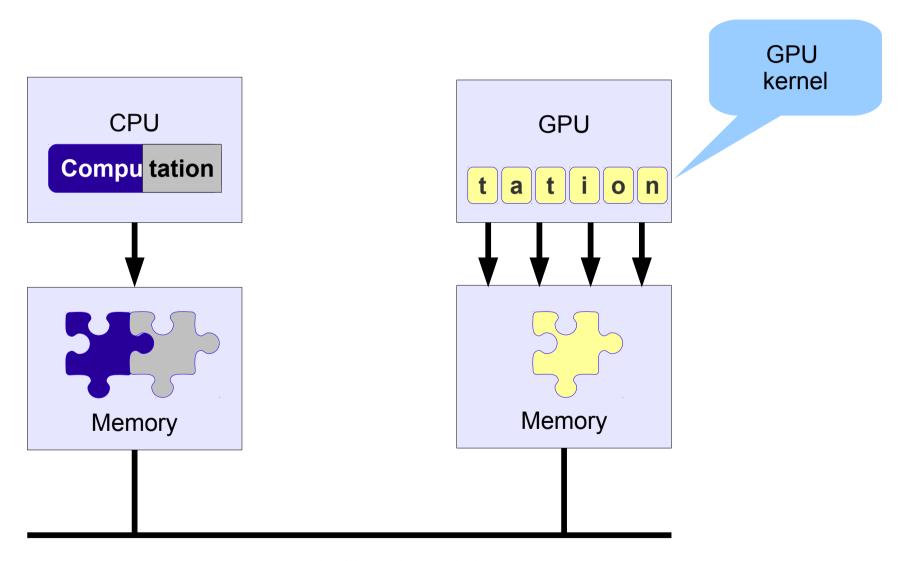
#### Simplified GPU model

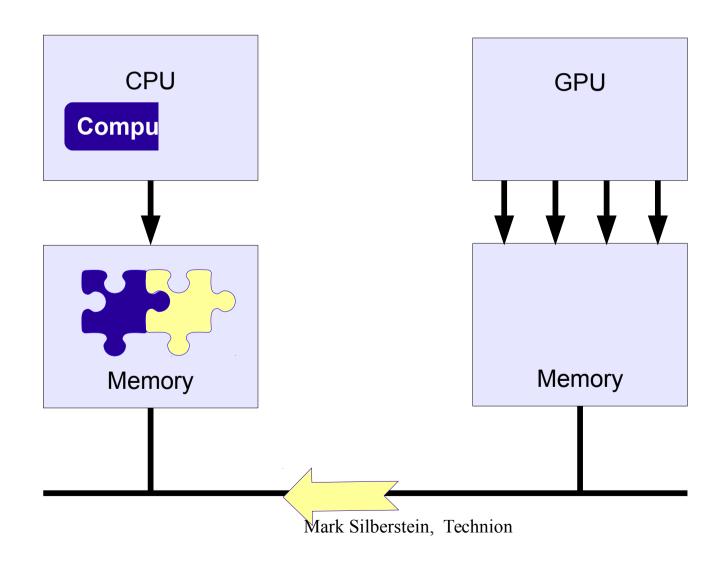
### **GPU 101**



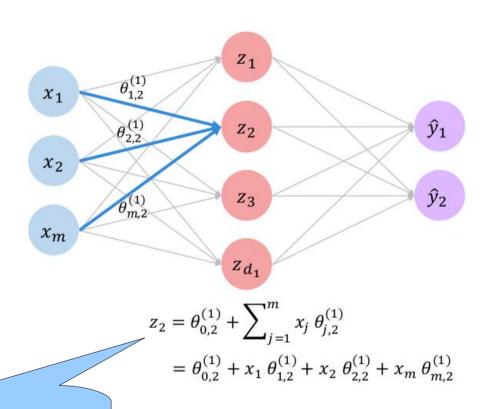






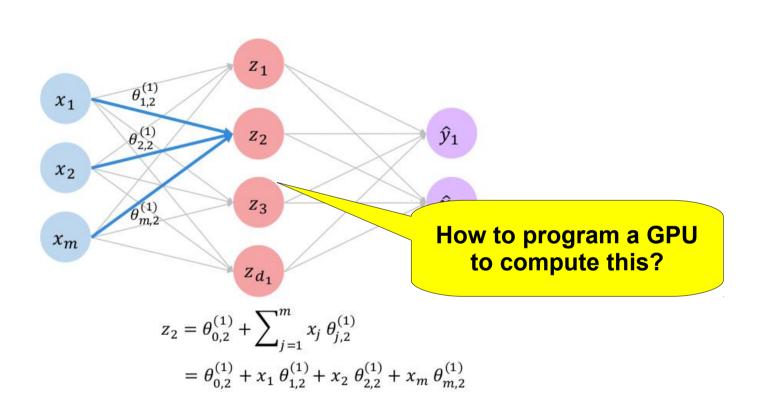


### GPUs in ML – Linear Algebra Accelerators



Compute Matrix Product
ON GPU

### GPUs in ML – Linear Algebra Accelerators



## Simple GPU program: exploiting data parallelism

- Idea: same set of operations is applied to different data chunks in parallel
- Algorithmic challenge identify data-parallel tasks
- Implementation
  - Every thread runs the same code on different data chunks.
  - GPU concurrently runs thousands of parallel threads

#### Vector sum C=A+B

Sequential algorithm

```
For every element i

C[i]=A[i]+B[i]
```

#### Vector sum C=A+B

Sequential algorithm

```
For every i

C[i]=A[i]+B[i]
```

Parallel algorithm

```
In parallel For every i
C[i]=A[i]+B[i]
```

## Implementation for a vector of length 1024

GPU kernel (this program runs in every thread)
 C[threadId]=A[threadId]+B[threadId]

Per-thread hardware-supplied ID

## Implementation for a vector of length 1024

GPU kernel

```
C[threadId]=A[threadId]+B[threadId]
```

- CPU
  - 1.Allocate three arrays (in GPU memory)
  - 2. Make data accessible to GPU (CPU->GPU copy)
  - 3. Invoke kernel with 1024 threads
  - 4. Wait until complete and make data accessible to CPU (GPU->CPU copy)

### Complete example

```
CPU:
void vector sum(float* A, float* B, float* C, int n)
{
   float* qA=GPU get reference(A);
   float* qB=GPU get reference(B);
   float* gC=GPU allocate mem(n);
  GPU set num threads(n);
     // GPU will invoke n threads
   GPU run(vector sum kernel(gA,gB,gC));
  GPU retrieve(C,gC);
GPU:
void vector sum kernel(float* qA, float* qB, float*qC)
{
     int my=HardwareThreadID;
     gC[my]=gA[my]+gC[my];
}
```

### Complete example

```
CPU:
void vector_sum(float* A, float* B, float* C, int n)
{
    float* gA=GPU_get_reference(A);
    float* gB=GPU_get_reference(B);
    float* gC=GPU_allocate_mem(n);

GPU_set_num_threads(n);
    // GPU will invoke n threads

GPU_run(vector_sum_kernel(gA,gB,gC));

GPU_retrieve(C,gC);
}
```

```
GPU:
void vector_sum_kernel(float* gA, float*
{
    int my=HardwareThreadID;
    gC[my]=gA[my]+gC[my];
}
GPU programming
```

### Complete example

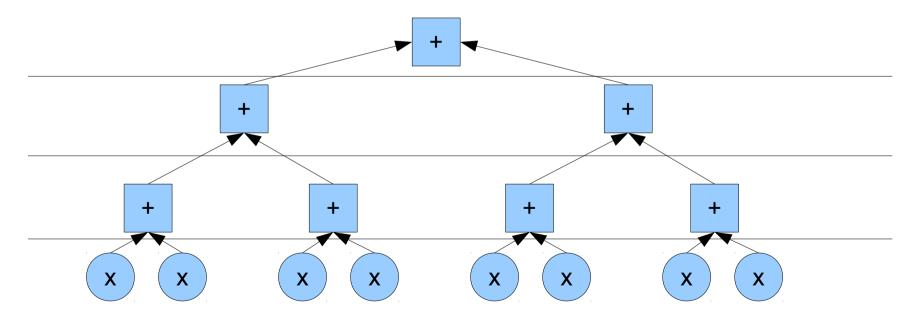
```
CPU:
void vector sum(float* A, float* B
                                          Many threads
   float* gA=GPU get reference(A)
   float* gB=GPU get reference(B);
   float* gC=GPU allocate mem(n):
  GPU set num threads(n);
     // GPU will invoke n threads
  GPU run(vector sum kernel(gA,gB,gC));
  GPU retrieve(C,gC);
                                 Fine-grained parallelism
GPU:
                                     (1 op per thread)
void vector sum kernel(float
     int my=Hardware_areadID;
    gC[my]=gA[my]+gC[my];
```

#### **BUT!**

- Vector sum is simple purely data parallel
- What if we need coordination between tasks
   Example: parallel dot product

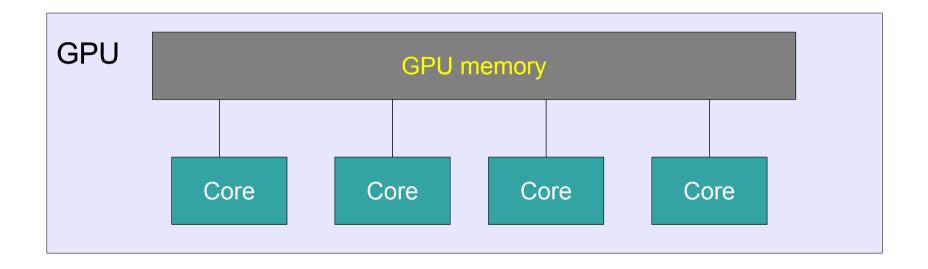
#### **BUT!**

- Vector sum is simple purely data parallel
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   Example: parallel dot product

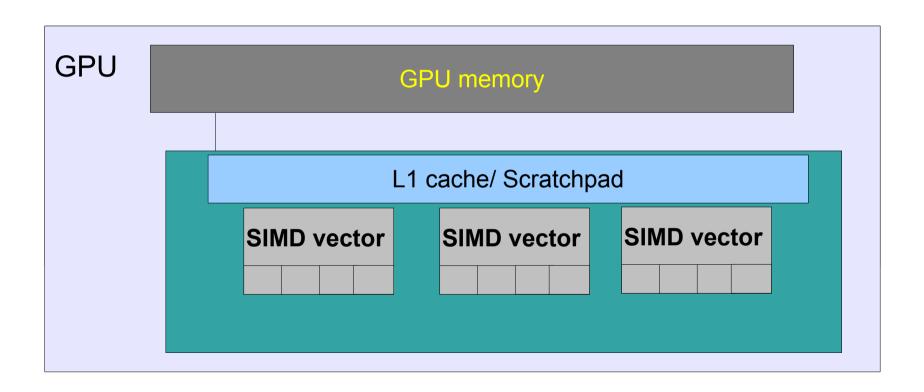


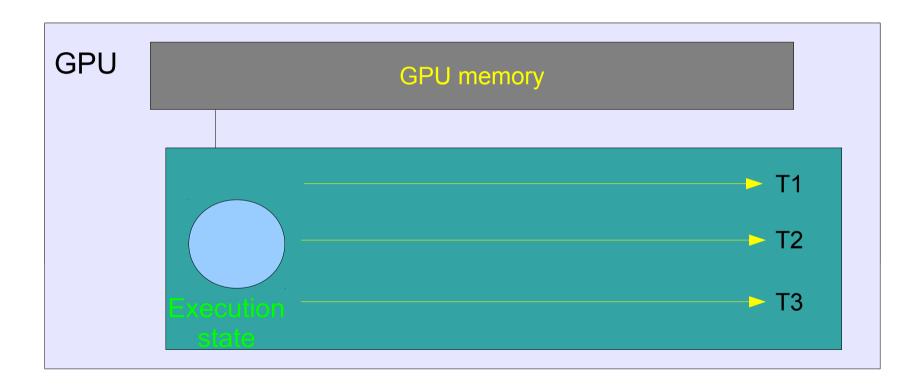
#### **GPU** hardware

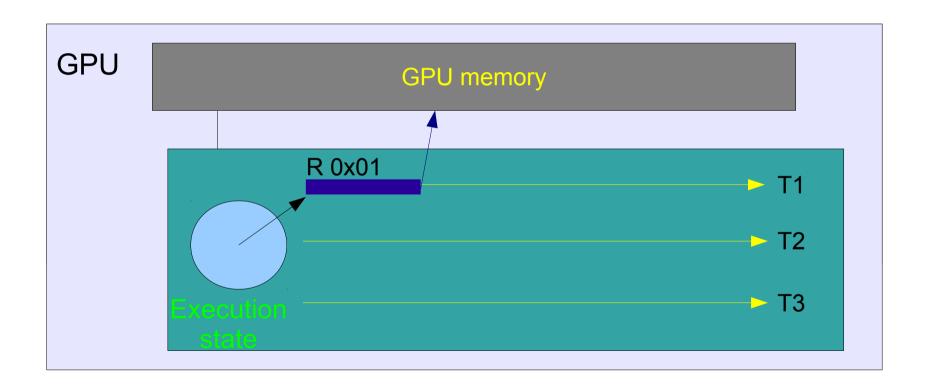
## GPU hardware parallelism 1. Multi-core

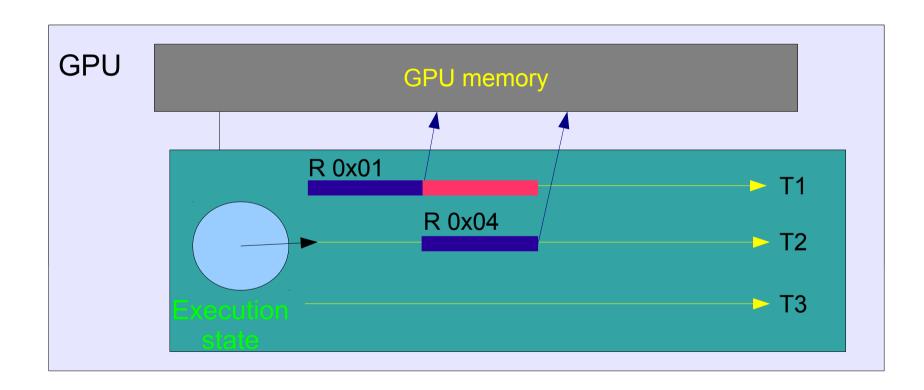


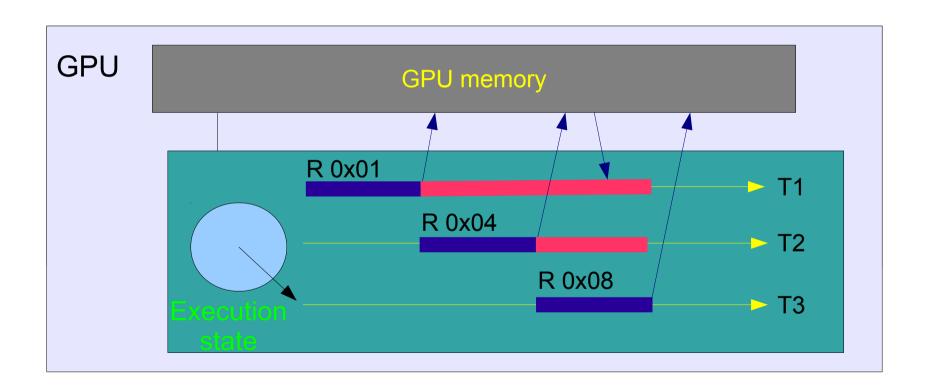
## GPU hardware parallelism 2. SIMD

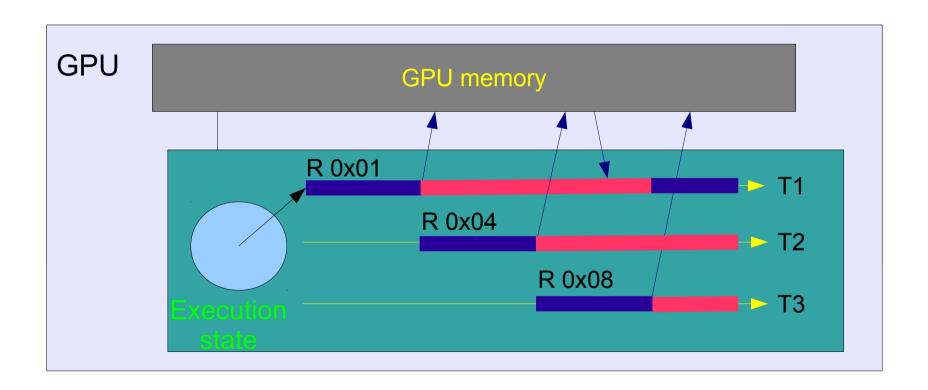




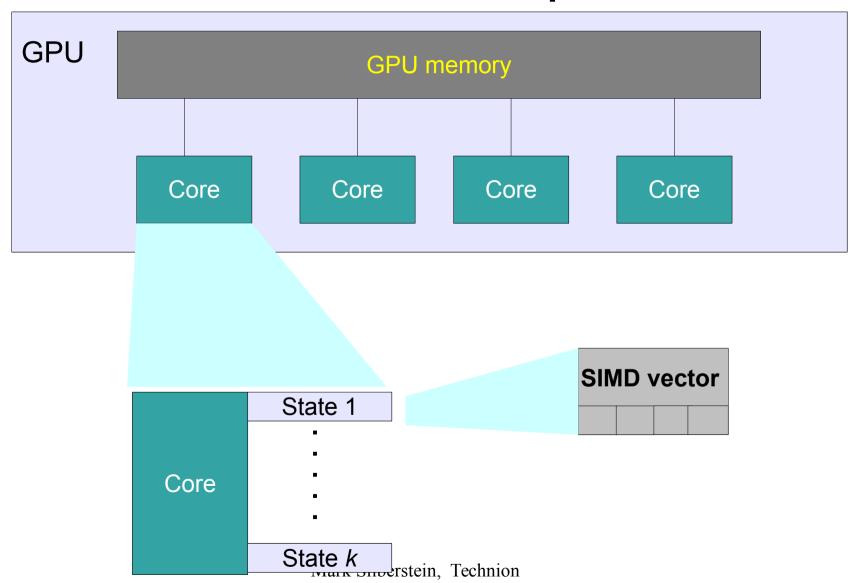




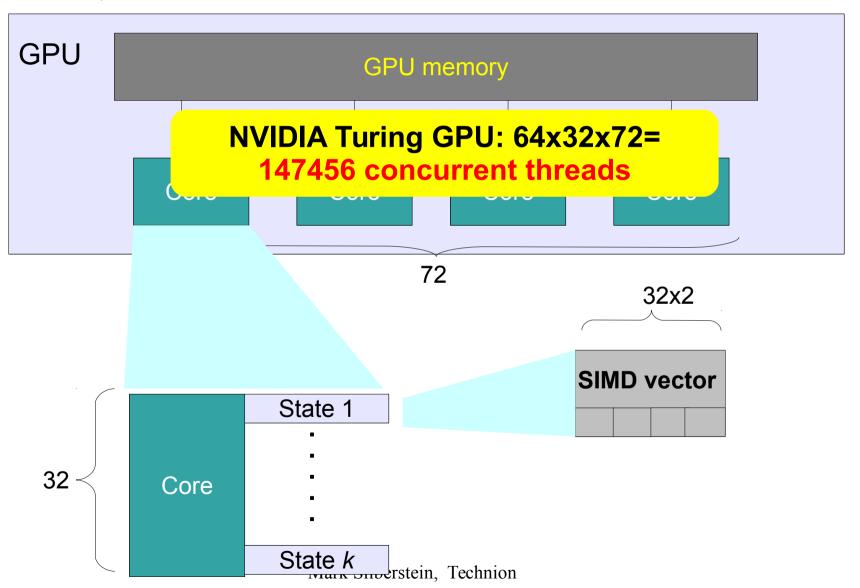




## Putting it all together: 3 levels of hardware parallelism



## Takeaway 1: 100,000-s of concurrent threads!



## Requirements for allowing fast GPU execution

We have enough parallelism

 We have enough space to store state per thread

We have enough bandwidth to memory

We have efficient scheduler to manage threads

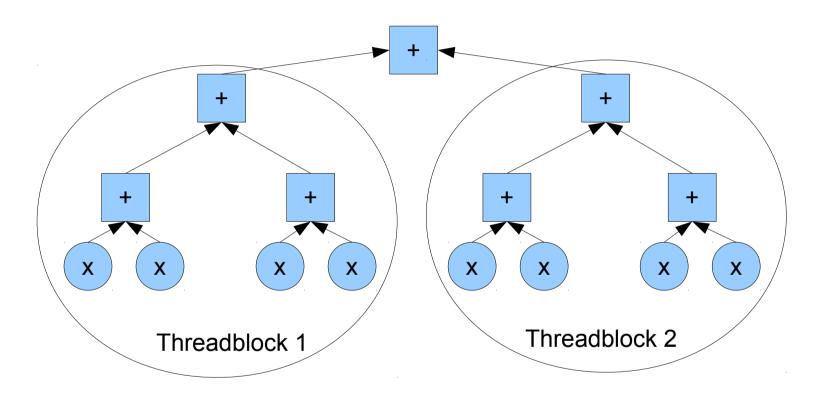
### How GPU manages threads?

- Application threads are grouped into threadblocks
- Programmer defines the number of threads / threadblocks for each run

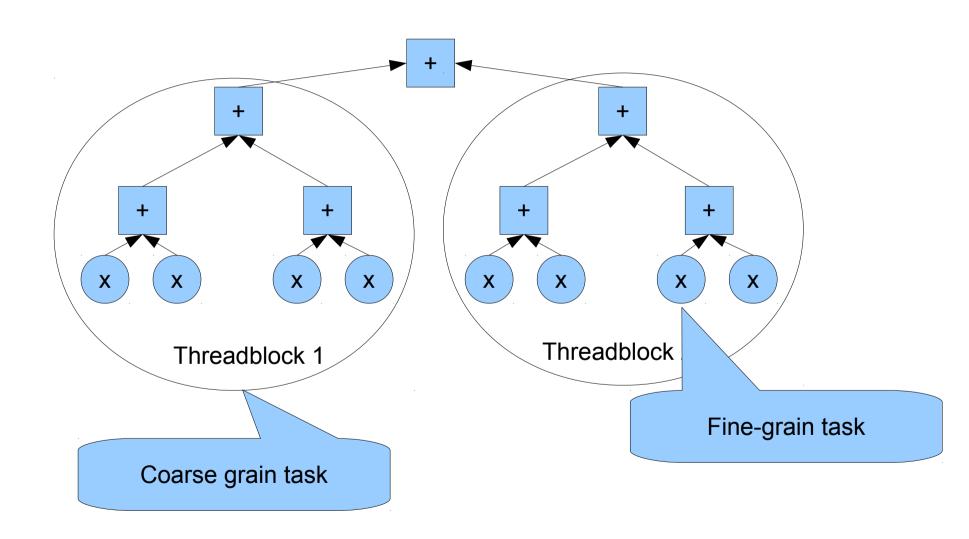
## Threadblock is a building block of GPU algorithms

- Threads inside a threadblock can communicate efficiently!
  - Share a small fast scratchpad memory
  - Can be synchronized via barriers
- Threadblocks are independent
- A program consists of many threadblocks

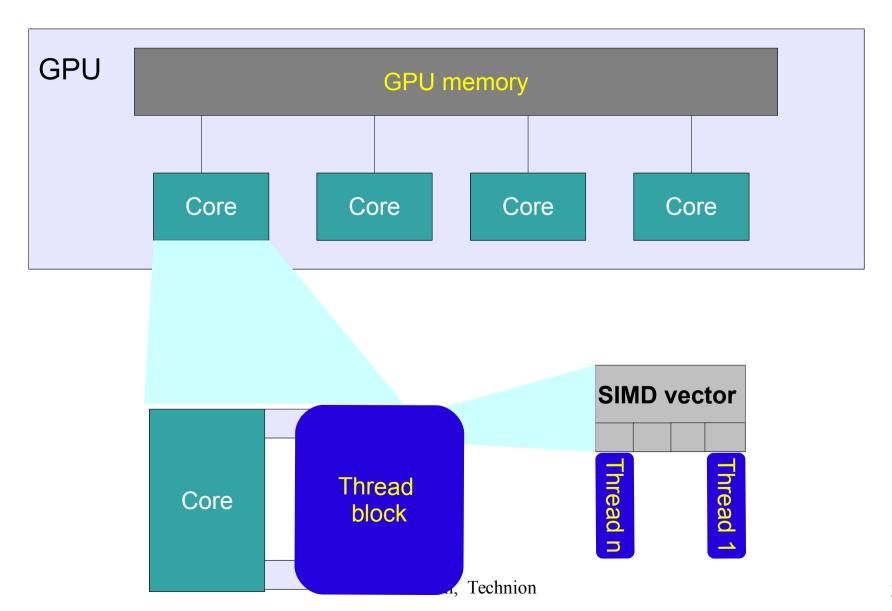
### Dot-product: hierarchical parallelization Decomposing into threadblocks



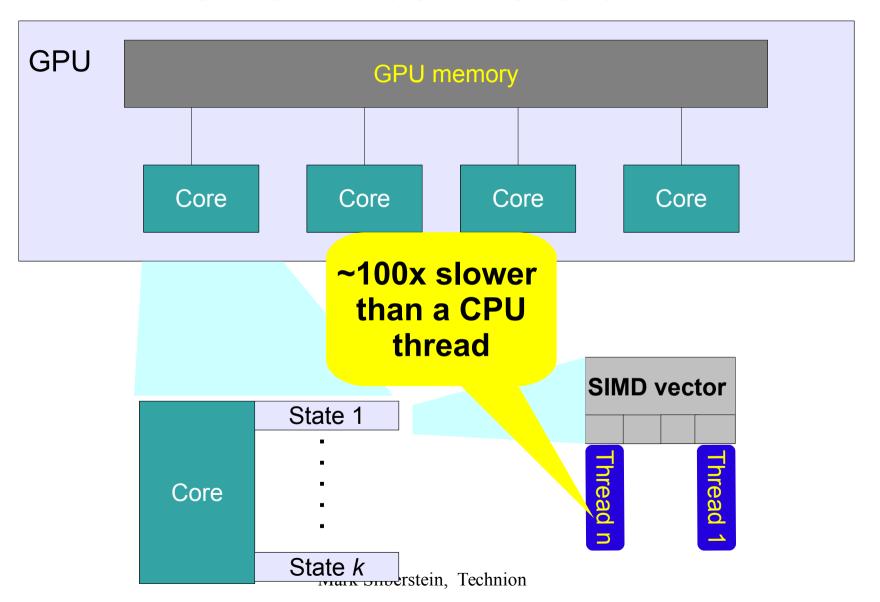
### Takeaway: parallelism hierarchy



### Software-Hardware mapping



#### Takeaway 2: One thread is slow



#### Dot product

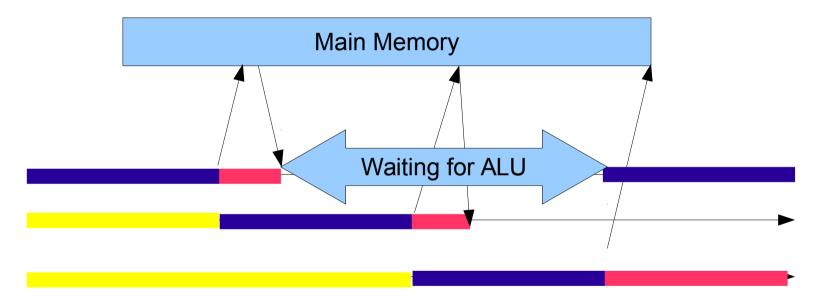
```
void vector dotproduct kernel(float* gA, float* gB, float* gOut)
{
     local float l res[TB SIZE]; //local core memory
     int thid=LocalThreadID;
     int tbid=ThreadBlockID;
   int offset=tbid*WG SIZE+thid;
   l res[tid]=qA[offset]*qB[offset];
     BARRIER(); // wait for all products
     for(int i=TB SIZE/2; i>0;i/=2)
      if (thid<i) l res[thid]=l res[thid]+l res[i+thid];</pre>
          BARRIER(); // wait for all partial sums
   if (thid==0) gOut[tbid]=l res[0];
}
```

# Parallelism structure of GPU programs

- Having many independent tasks is not enough
- Parallel structure should map well on hardware hierarchy

#### Estimating application performance on highthroughput processors

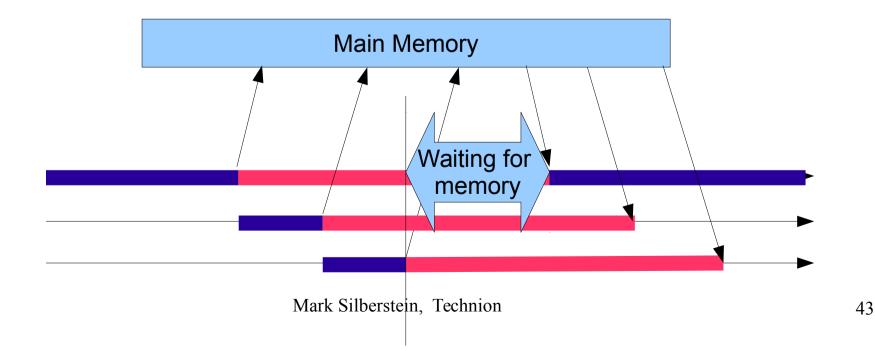
#### Compute-bound tasks



## Performance bounded by maximum ALU capacity

#### Memory-bound tasks

- We can fully utilize a processor only if data is available in an ALU on time
- How fast can the data be made available?
  - Assume infinite number of threads, ideal parallelization



## Measure of ALU/memory ratio: Arithmetic intensity

- Number of OPs per memory(\*\*) access
  - Vector sum: 1 operation per 3 accesses. A=1/3

### Upper bound on performance

 For memory bound algorithms only (why?)

Perf=GPUMemBandwidth\*A

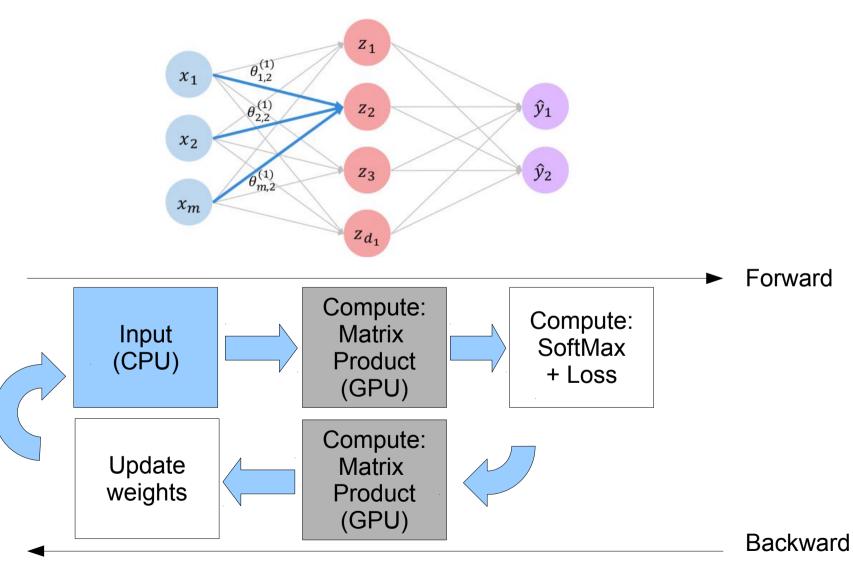
Bytes/sec\*Ops/Bytes=Ops/sec

#### Vector sum performance estimate

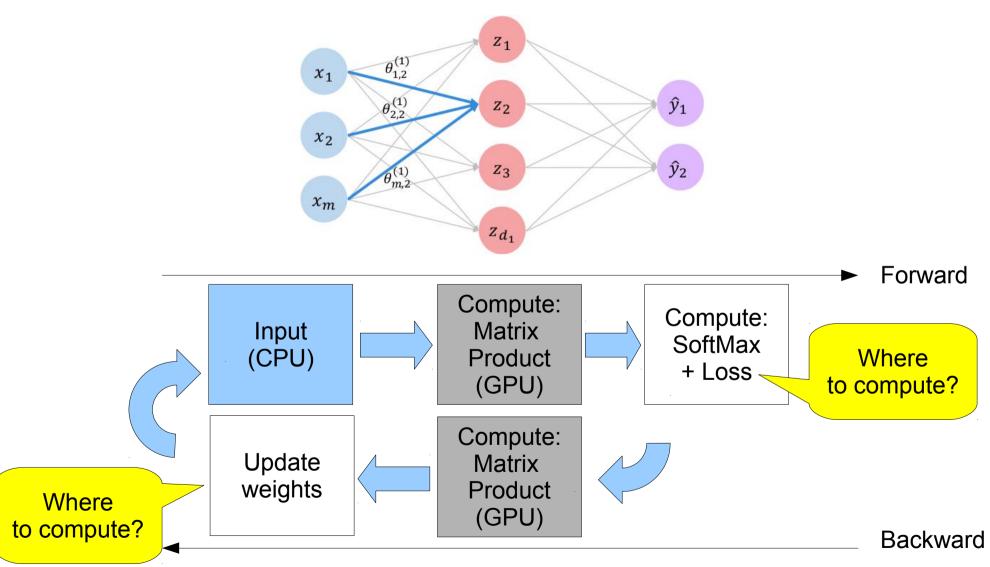
- Example: sum of vectors, GPU GTX Titan
- A=1/(3\*(4 bytes)), MemBW=~70GFloat/s:
   Performance= ~23GFLOP/s
- For comparison: raw capacity: 4.5 TFLOPs
- Only 0.5% of computing capacity utilized!!

## Integrating GPUs with applications

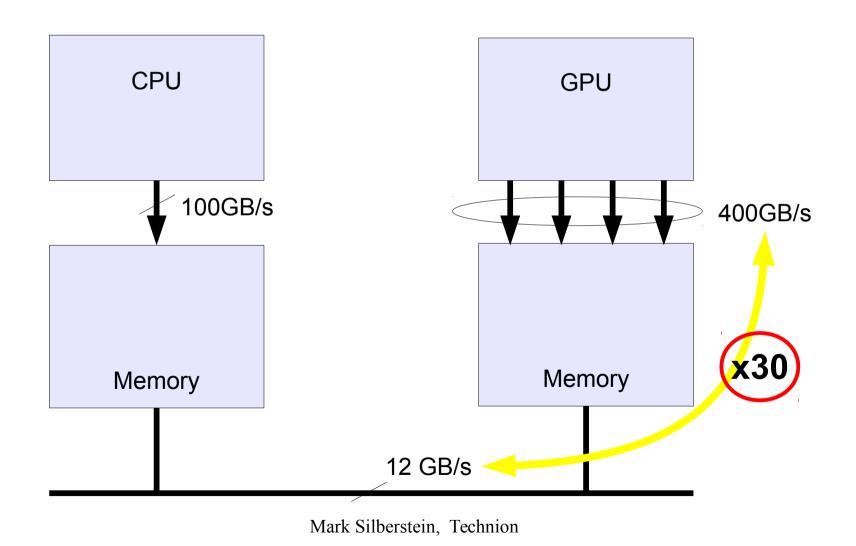
### GPUs and data locality



#### GPUs and data locality

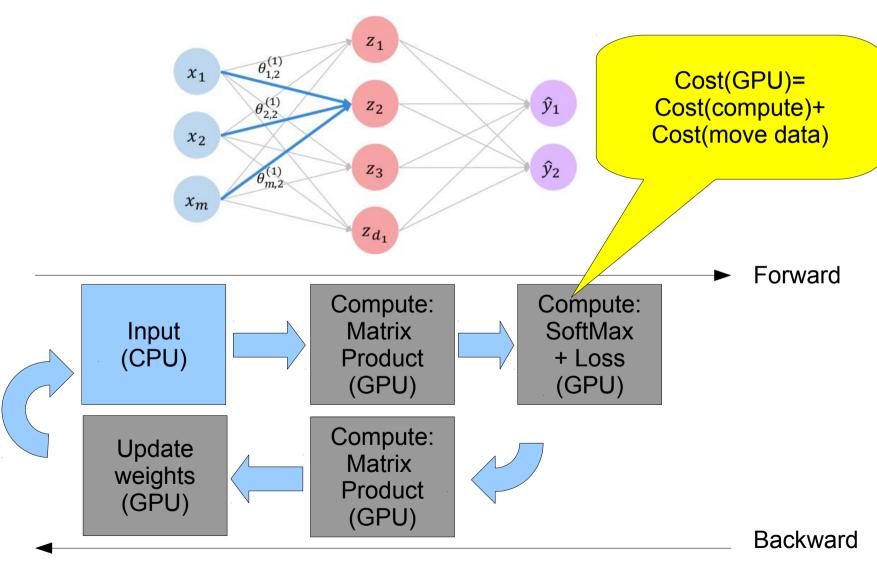


### Problem: separate GPU memory

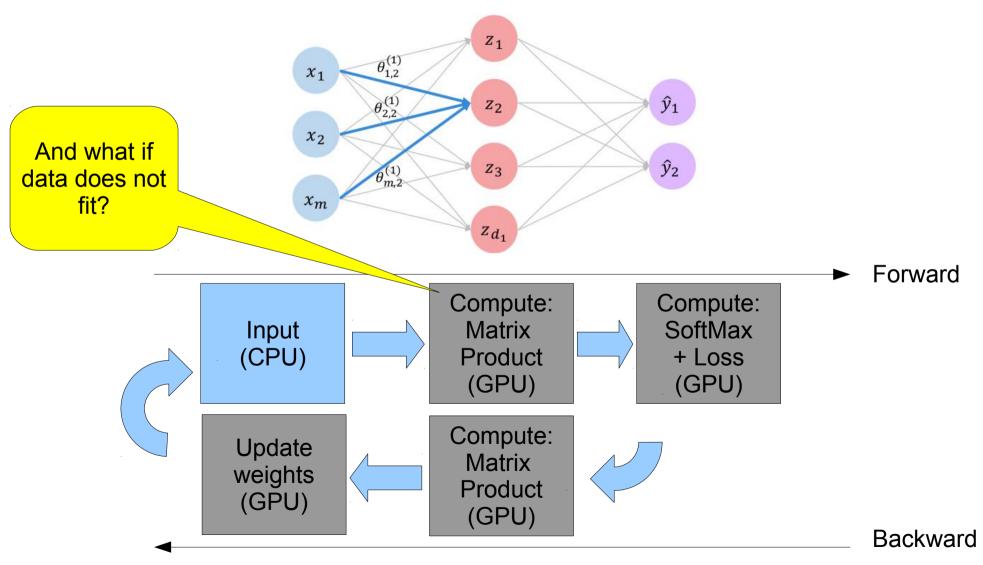


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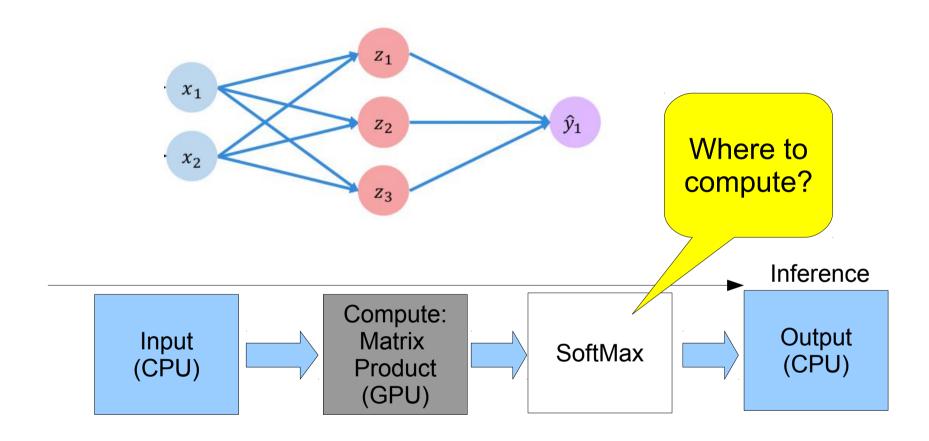
#### Data must be on a GPU



### Data must stay on a GPU



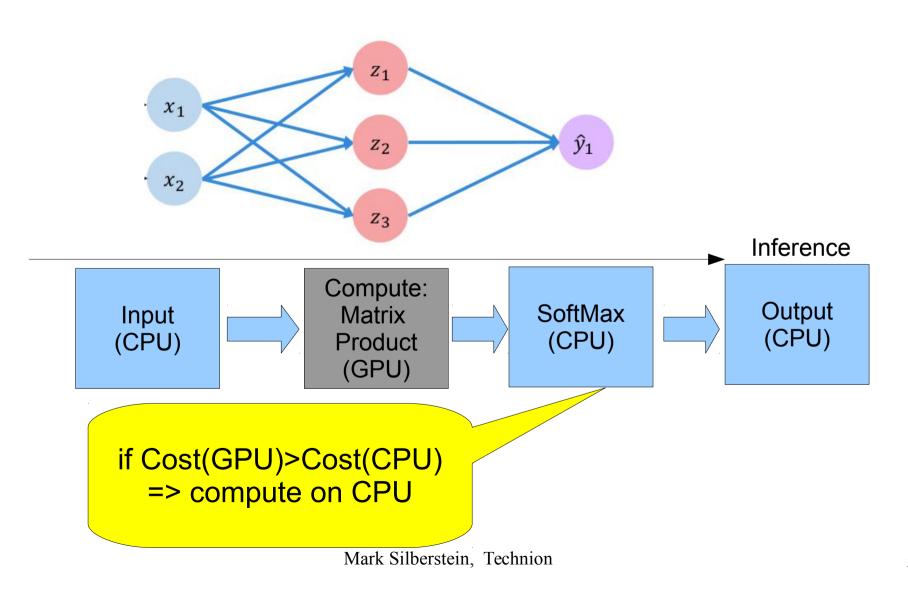
#### GPUs and invocation cost



### Management overhead

- GPU invocation takes at least 30,000 single CPU core cycles
- Short GPU invocations do not pay off
  - Batch multiple invocations

## Sometimes CPU is faster even when it's slower



#### GPUs are used automatically, but...

- cuBLAS, cuDNN,cuFFD
  - provide powerful GPU-accelerated functions
- TensorFlow, MxNet, Caffe,...
  - run on GPUs
- So why learn how to use GPUs?

### Other interesting projects in my lab

- GPU networking
- GPU access to storage
- GPU security
- GPU interaction with Smart NICs and other accelerators

#### Want to hear more?

- Accelerators and accelerated systems: 236278
- Undergraduate/master projects @ Accelerated Computing Systems Lab (ACSL)

mark@ee.technion.ac.il Fishbach 408, TCE