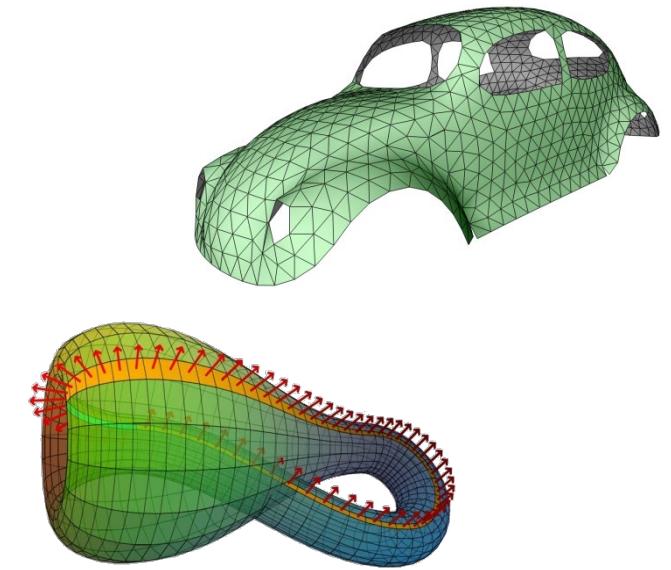
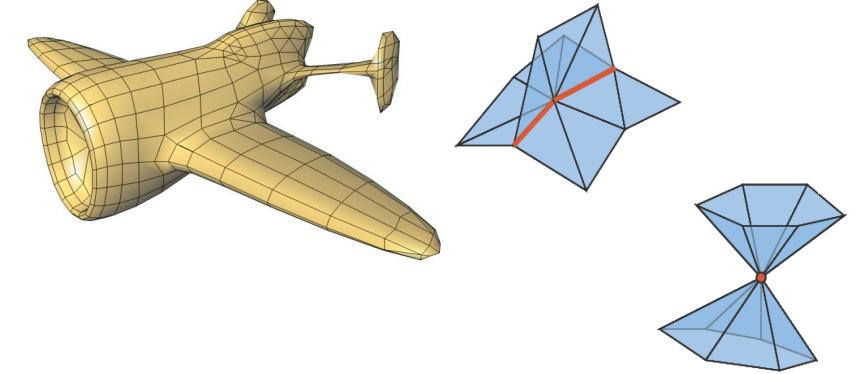
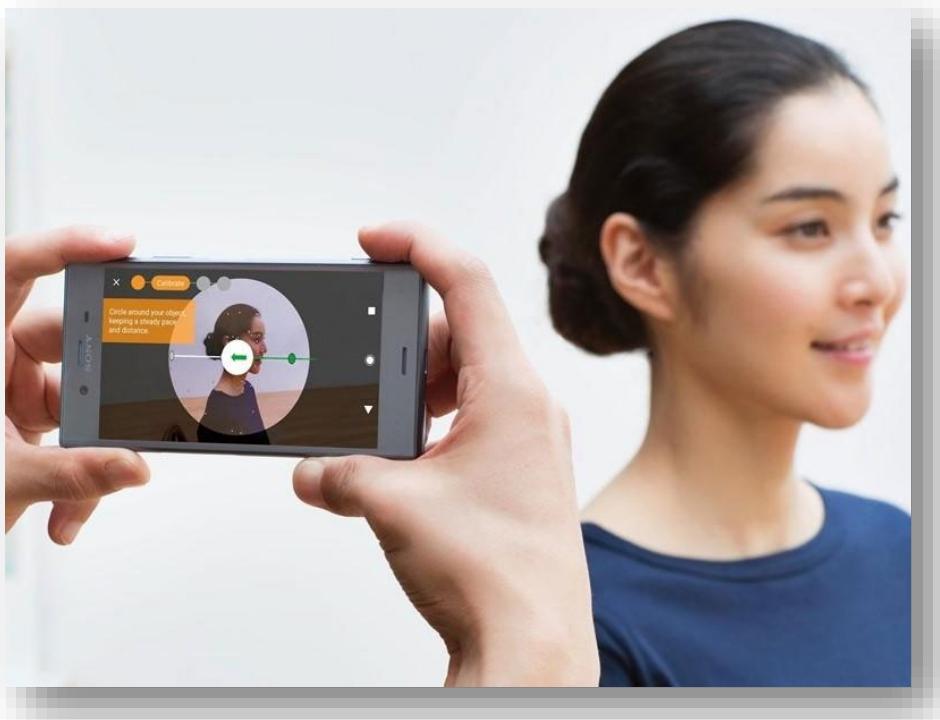


# Basic Concepts



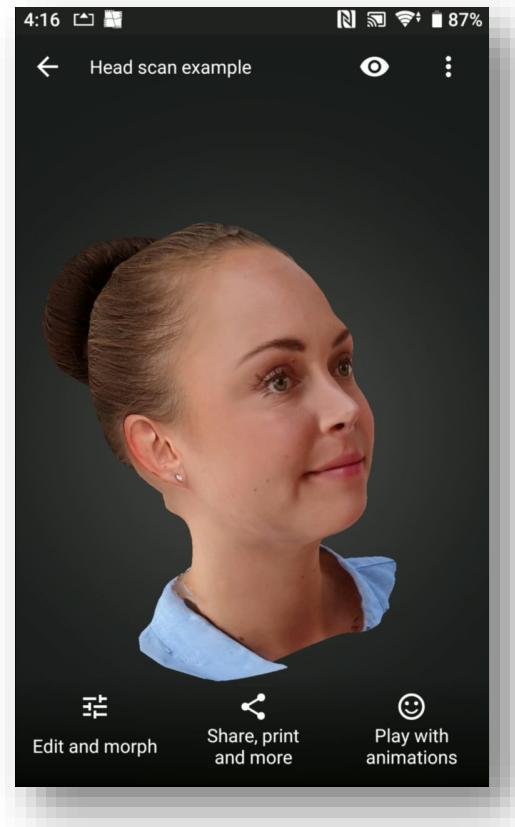
# Last Time

## Cool applications



# Today

## How is the 3D data represented computationally?



# Motivation

## Requirements:

- Efficient = surface only
- arbitrary shapes
  - Holes, boundaries, sharp creases
- more data = better accuracy
- adaptive



# Triangle Meshes

- Connectivity: vertices, edges, triangles

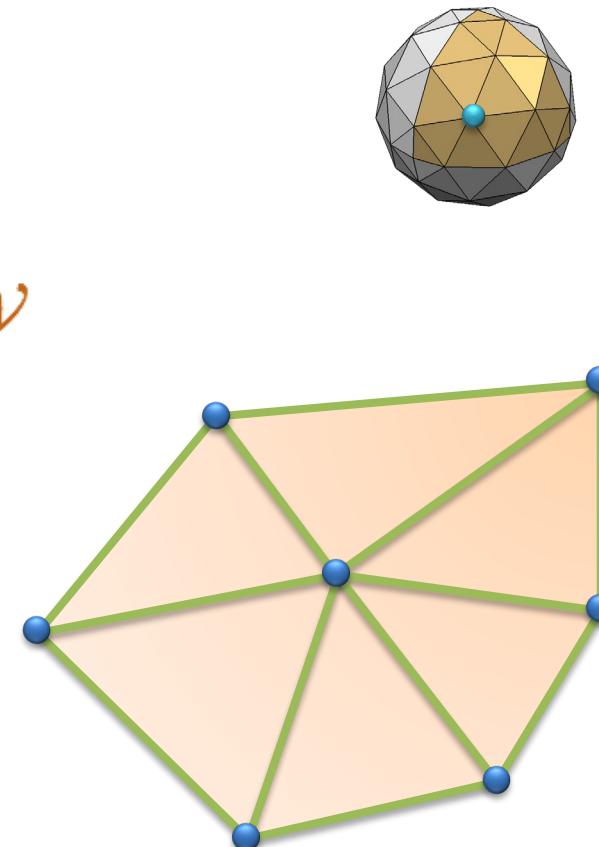
$$\mathcal{V} = \{v_1, \dots, v_n\}$$

$$\mathcal{E} = \{e_1, \dots, e_k\}, \quad e_i \in \mathcal{V} \times \mathcal{V}$$

$$\mathcal{F} = \{f_1, \dots, f_m\}, \quad f_i \in \mathcal{V} \times \mathcal{V} \times \mathcal{V}$$

- Geometry: vertex positions

$$\mathcal{P} = \{\mathbf{p}_1, \dots, \mathbf{p}_n\}, \quad \mathbf{p}_i \in \mathbb{R}^3$$



# Polygonal ~~Triangle~~ Meshes

- Connectivity: vertices, edges, triangles

$$\mathcal{V} = \{v_1, \dots, v_n\}$$

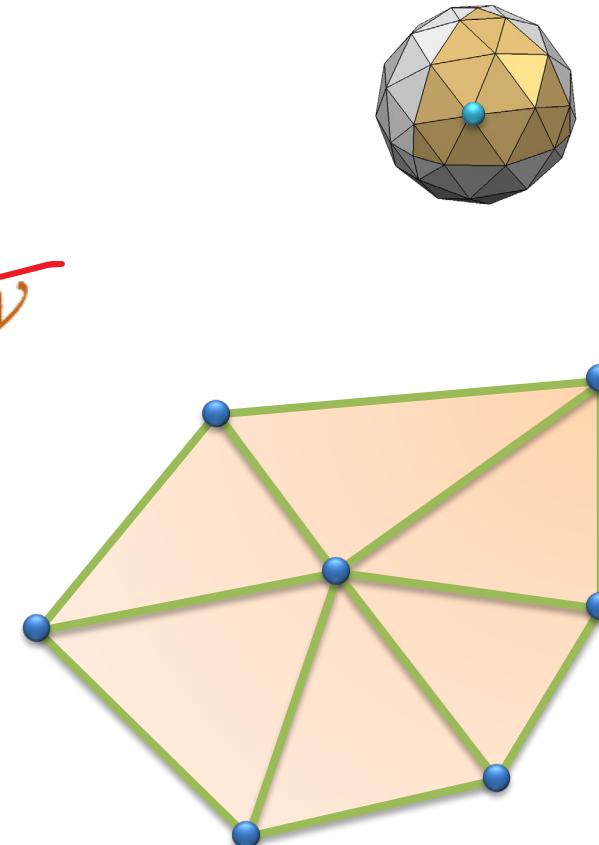
$$\mathcal{E} = \{e_1, \dots, e_k\}, \quad e_i \in \mathcal{V} \times \mathcal{V}$$

$$\mathcal{F} = \{f_1, \dots, f_m\}, \quad f_i \in \mathcal{V} \times \mathcal{V} \times \mathcal{V}$$

$$f_i \in \mathcal{V} \times \mathcal{V} \times \dots \times \mathcal{V}$$

- Geometry: vertex positions

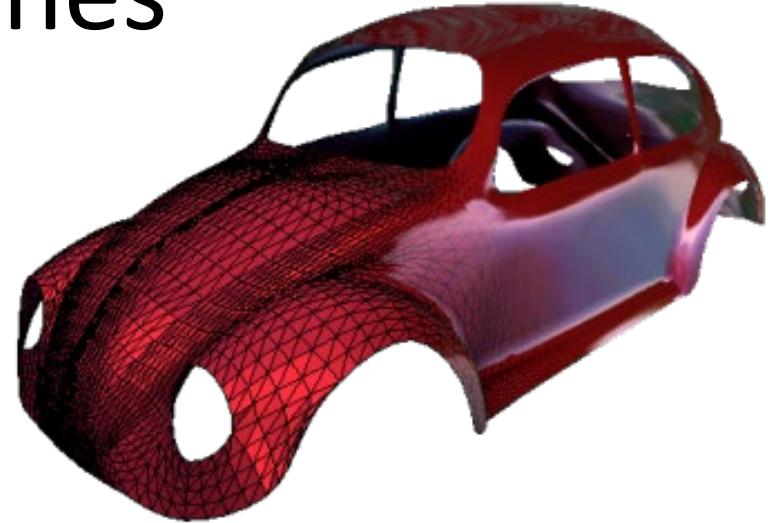
$$\mathcal{P} = \{\mathbf{p}_1, \dots, \mathbf{p}_n\}, \quad \mathbf{p}_i \in \mathbb{R}^3$$



# Polygonal Meshes

A good representation:

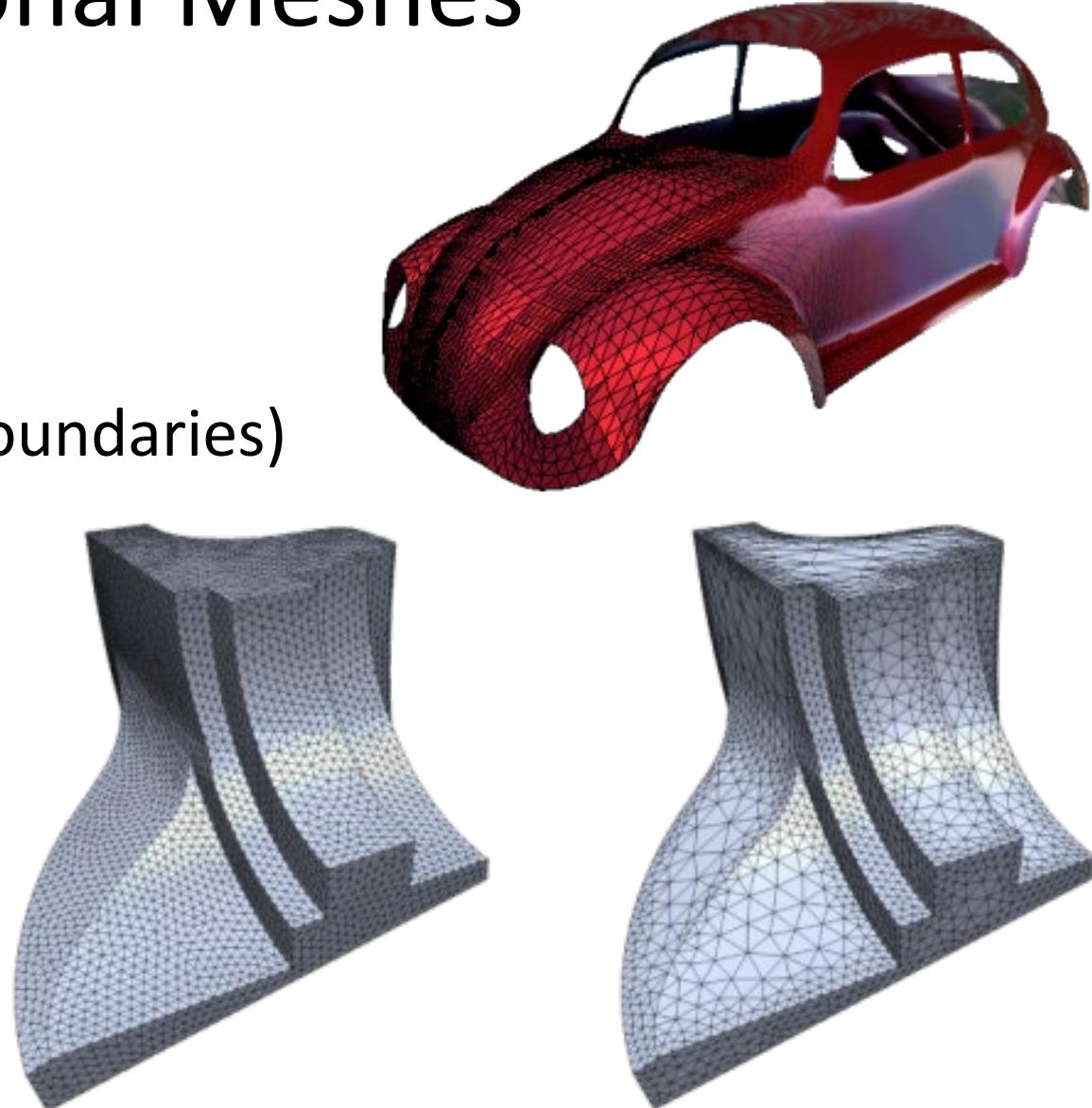
- arbitrary topology (holes, boundaries)
- piecewise smooth surfaces



# Polygonal Meshes

A good representation:

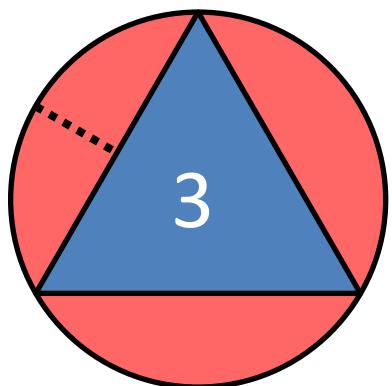
- arbitrary topology (holes, boundaries)
- piecewise smooth surfaces
- adaptive refinement
- efficient rendering



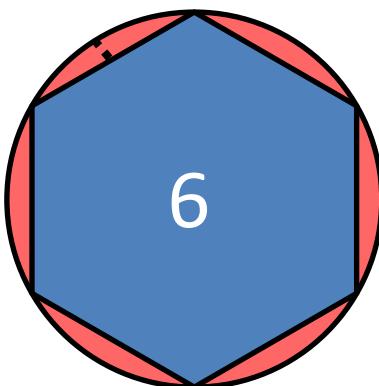
# Polygonal Meshes

Piecewise linear approximation

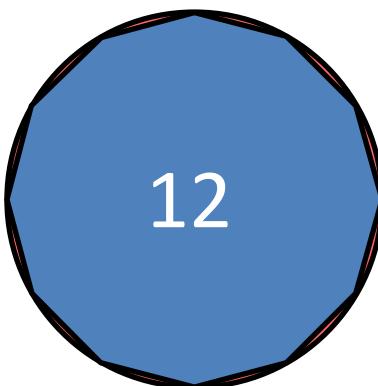
- More data = better accuracy
- Error is  $O(h^2)$ 
  - Edge length divided by 2 → Error divided by 4



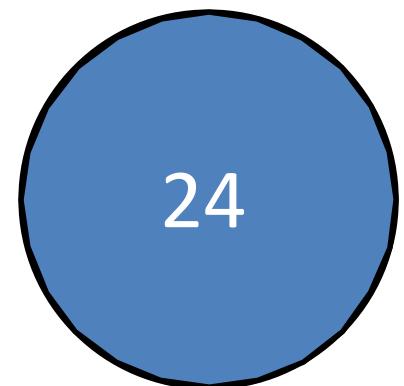
25%



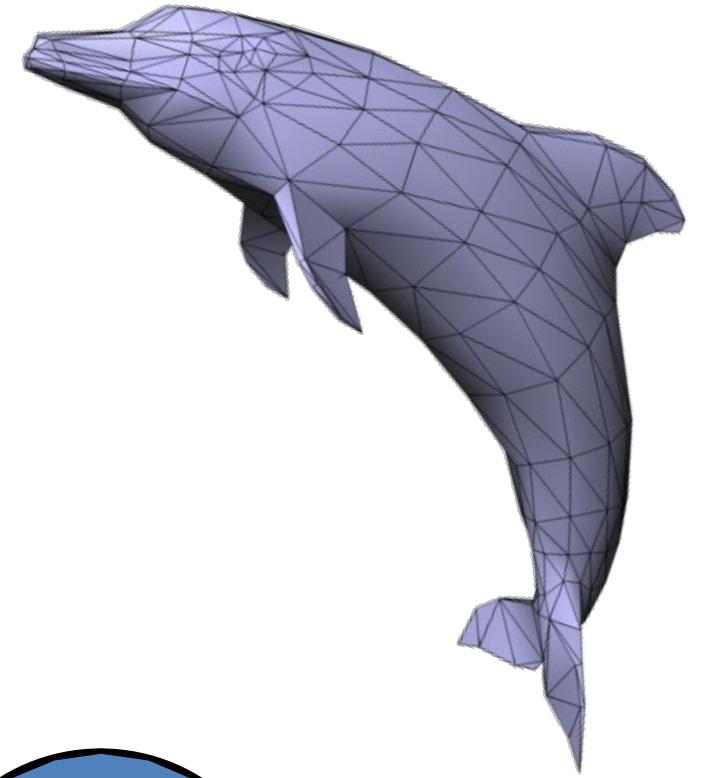
6.5%



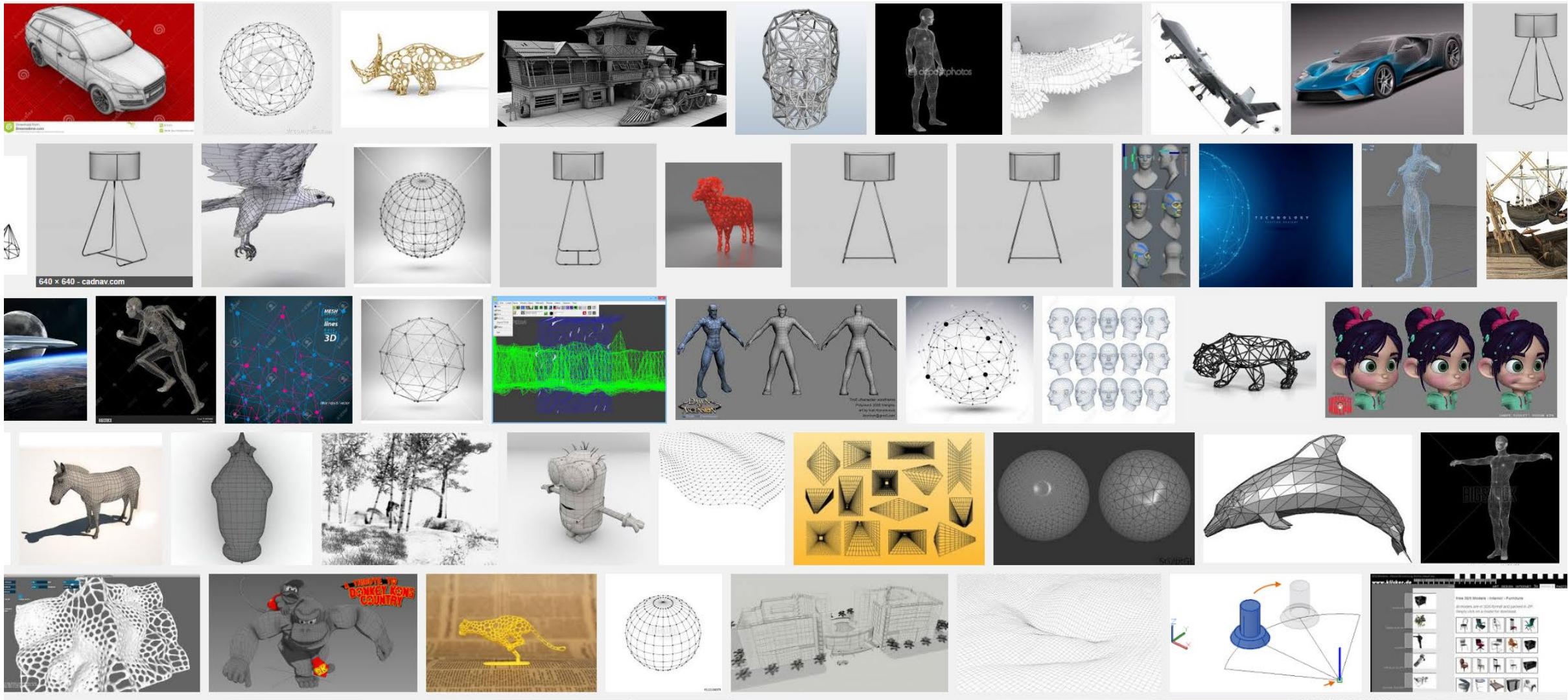
1.7%



0.4%



# Available Meshes



# “Good” Meshes

## Real time compression of triangle mesh connectivity

S Gumhold, W Straßer - Proceedings of the 25th annual conference on ..., 1998 - dl.acm.org

... To allow the encoding of an arbitrarily connected and **oriented triangle mesh** in one run, several basic ... of the current cut-border part is known during compression and decompression, the “**close** cut-bo

## The ball-pivoting algorithm for surface reconstruction

Cited F Bernardini, J Mittleman, H Rushmeier... - IEEE transactions on ..., 1999 - ieeexplore.ieee.org

... By checking that **triangle** and data point normals are consistently **oriented**, we avoid generating a **triangle** in this ... In all cases, the BPA outputs an **orientable**, triangulated **manifold**. ... other potential seed **triangles**, leading to the construction of small sets of **triangles** lying **close** to the ...

Cit

## Fair morse functions for extracting the topological structure of a surface **mesh**

X Ni, M Garland, JC Hart - ACM Transactions on Graphics (TOG), 2004 - dl.acm.org

... Its values are defined on the vertices of an **oriented 2-manifold triangle mesh**  $M$ , and ... Such **manifolds** contain one or more boundaries each represented by a simple **closed** loop of ... properties of the field, and depend on a smooth and well-shaped function over the **manifold**. ...

Cited by 201 Related articles All 13 versions Web of Science: 61 Cite Save

# Tasks / Applications

- Determine inside/outside

- Collision detection
    - Animation
    - Mesh design
    - Fabrication

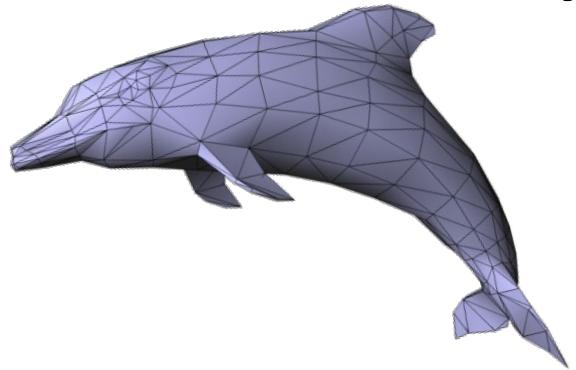


From: [Computational Design of Reconfigurables, Garg et al., 2016]

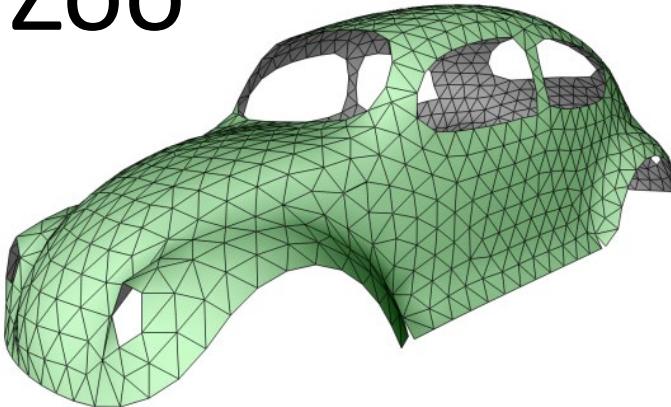
- Find boundaries

- Preserve boundaries during processing
  - Set boundary values for numerically solving differential equations

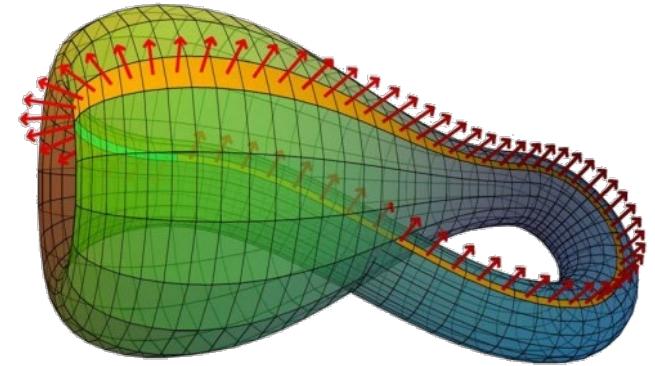
# Mesh Zoo



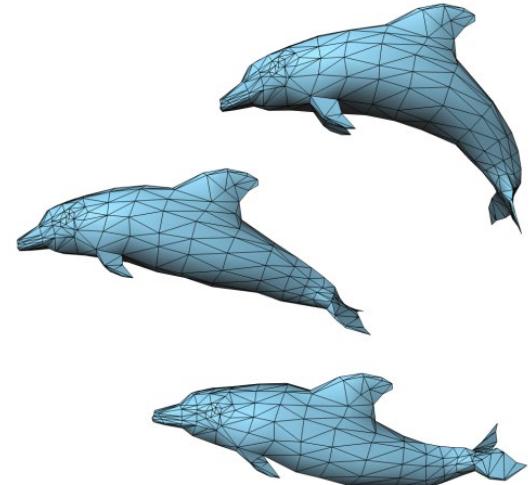
**Single component,  
closed, triangular,  
orientable manifold**



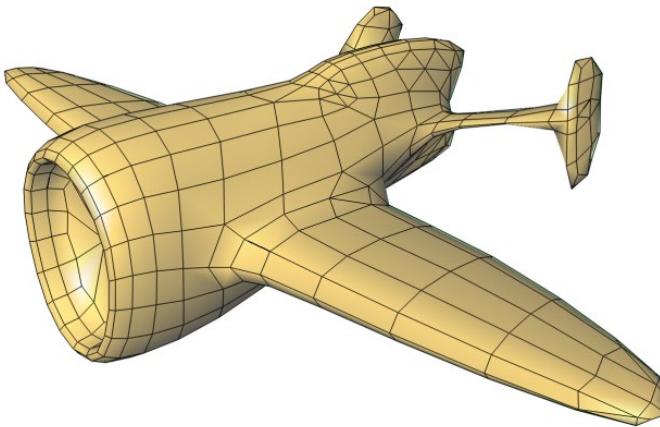
**With boundaries**



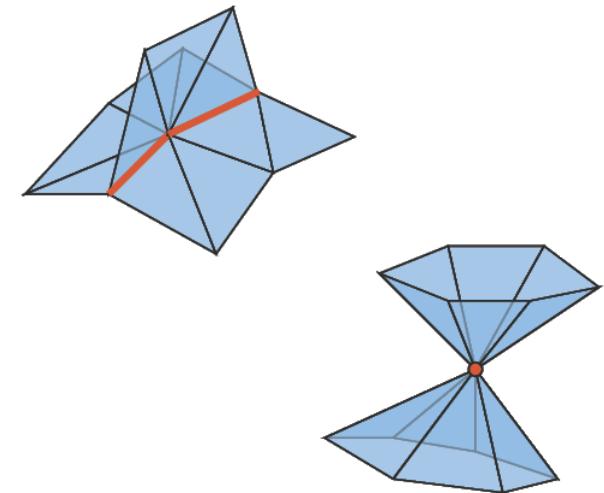
**Not orientable**



**Multiple components**



**Not only triangles**

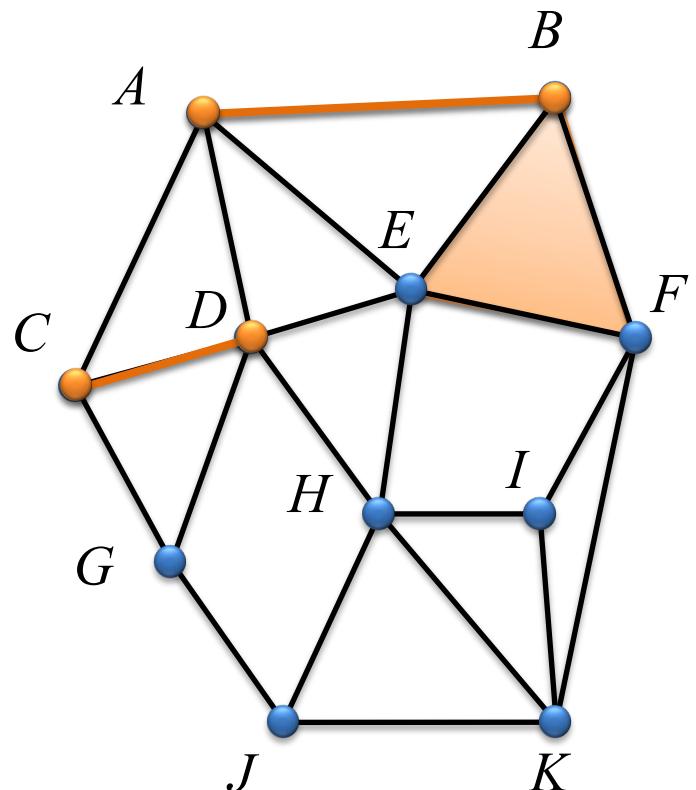


**Non manifold**

# Mesh Definitions

- Need vocabulary to describe zoo meshes
- The connectivity of a mesh is just a graph
- We'll start with some graph theory

# Graph Definitions



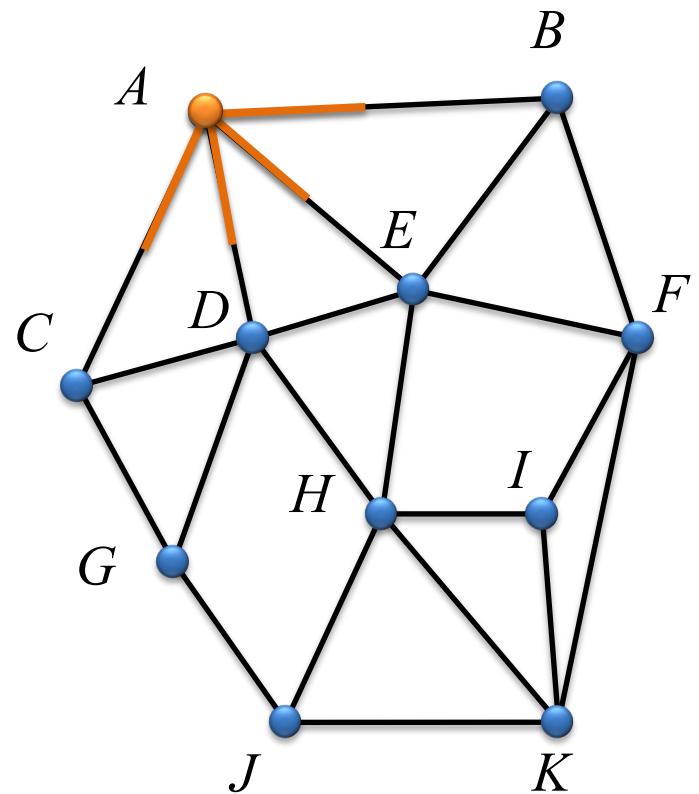
$G = \text{graph} = \langle V, E \rangle$

$V = \text{vertices} = \{A, B, C, \dots, K\}$

$E = \text{edges} = \{(AB), (AE), (CD), \dots\}$

$F = \text{faces} = \{(ABE), (DHJG), \dots\}$

# Graph Definitions



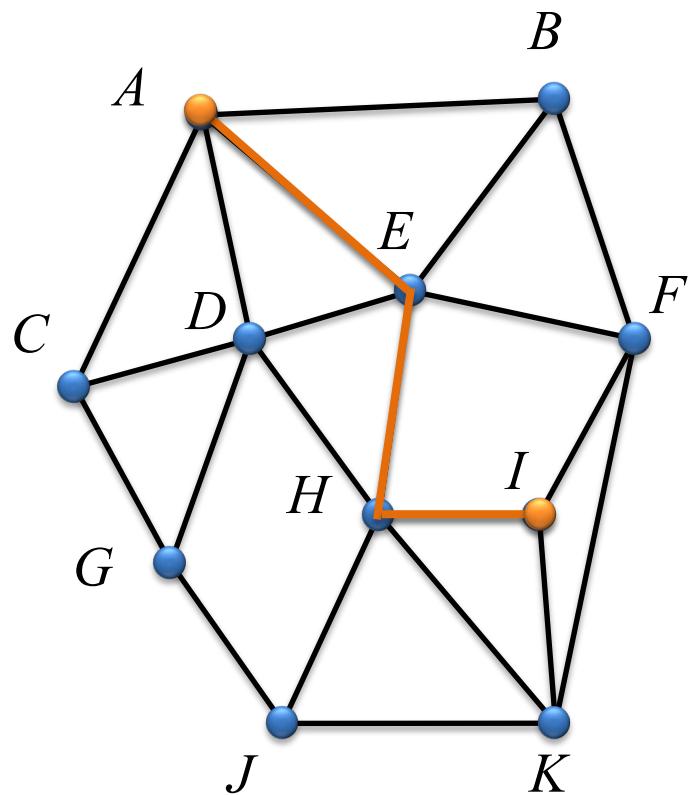
**Vertex degree or valence** =  
number of incident edges

$$\deg(A) = 4$$

$$\deg(E) = 5$$

**Regular mesh** =  
all vertex degrees are equal

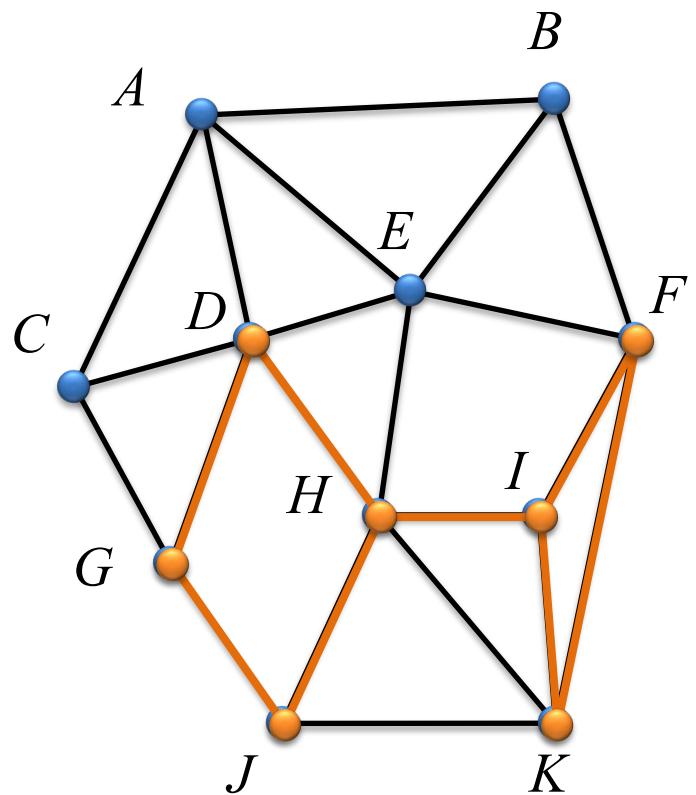
# Connectivity



**Connected =**

path of edges connecting every two vertices

# Connectivity



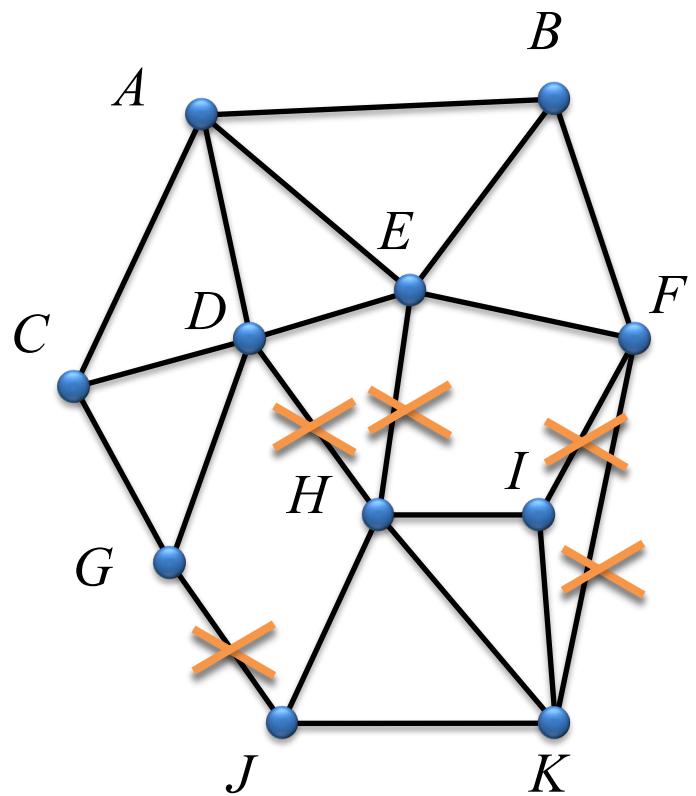
**Connected =**

path of edges connecting every two vertices

**Subgraph =**

$G' = \langle V', E' \rangle$  is a subgraph of  $G = \langle V, E \rangle$  if  
 $V'$  is a subset of  $V$  and  
 $E'$  is the subset of  $E$  incident on  $V'$

# Connectivity



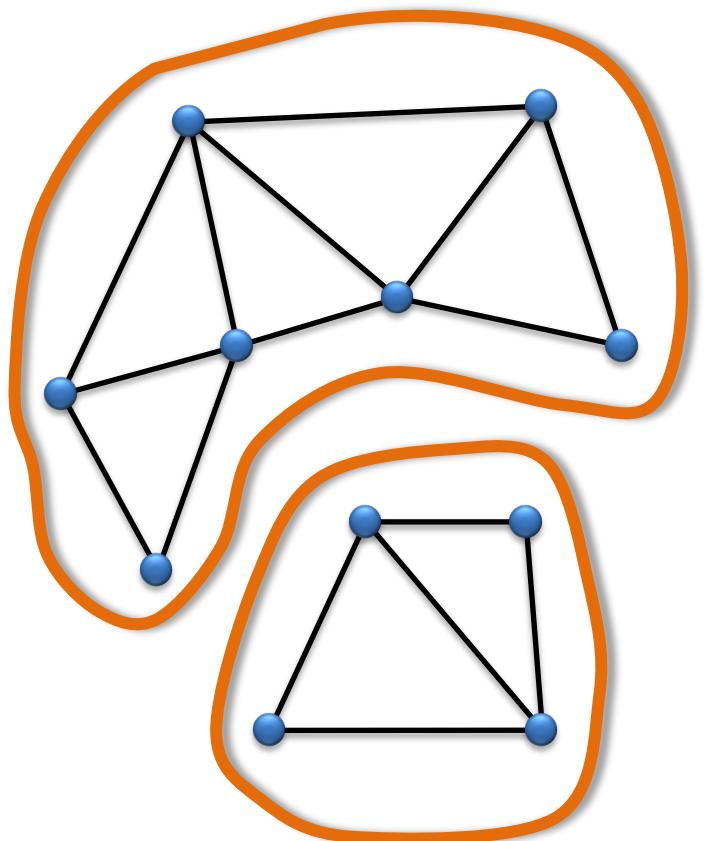
**Connected =**

path of edges connecting every two vertices

**Subgraph =**

$G' = \langle V', E' \rangle$  is a subgraph of  $G = \langle V, E \rangle$  if  
 $V'$  is a subset of  $V$  and  
 $E'$  is the subset of  $E$  incident on  $V'$

# Connectivity



**Connected =**

path of edges connecting every two vertices

**Subgraph =**

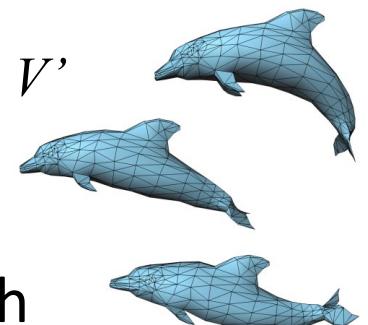
$G' = \langle V', E' \rangle$  is a subgraph of  $G = \langle V, E \rangle$  if

$V'$  is a subset of  $V$  and

$E'$  is the subset of  $E$  incident on  $V'$

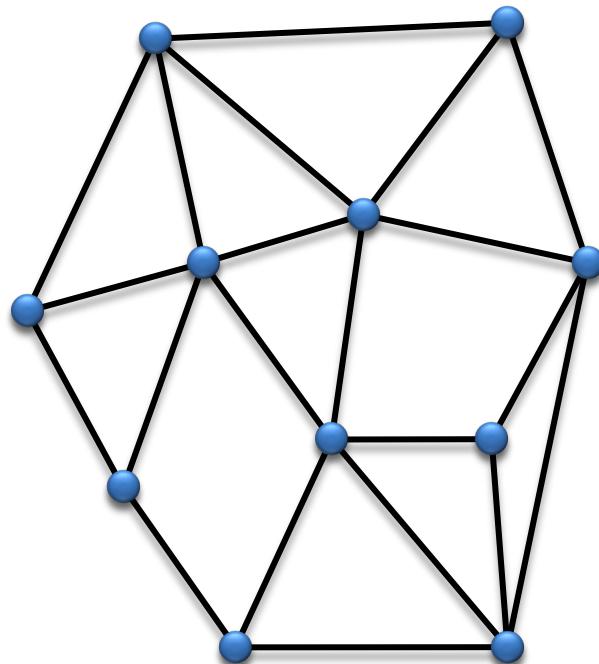
**Connected Component =**

maximally connected subgraph

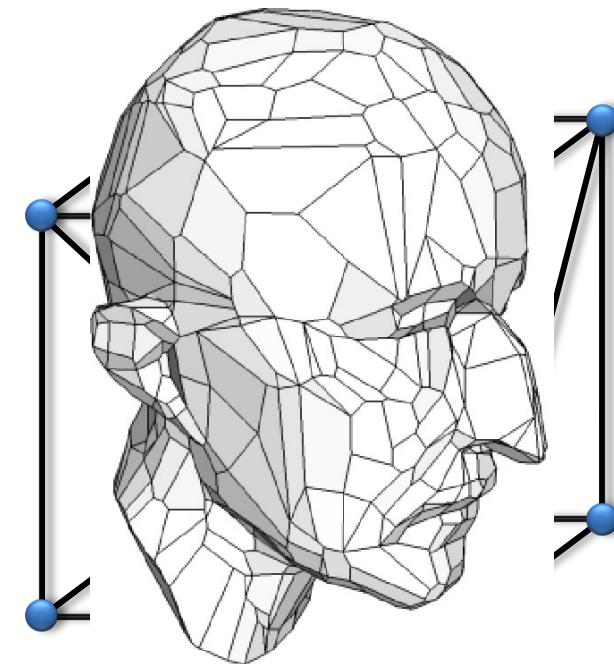


# Graph Embedding

**Embedding:**  $G$  is *embedded* in  $R^d$ , if each **vertex** is assigned a position in  $R^d$



Embedded in  $R^2$



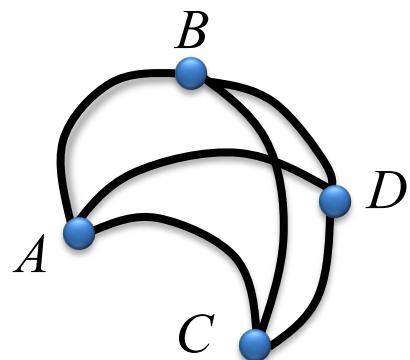
Embedded in  $R^3$

# Planar Graphs

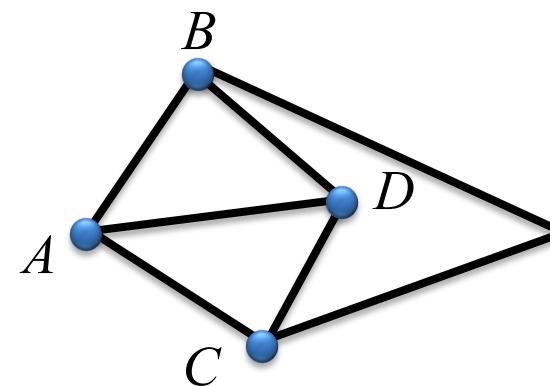
**Planar Graph** =

Graph whose vertices and edges *can be* embedded in  $\mathbb{R}^2$   
such that its edges *do not intersect*

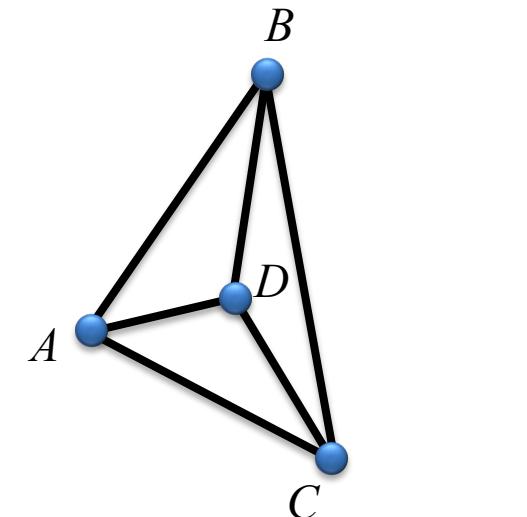
**Planar** Graph



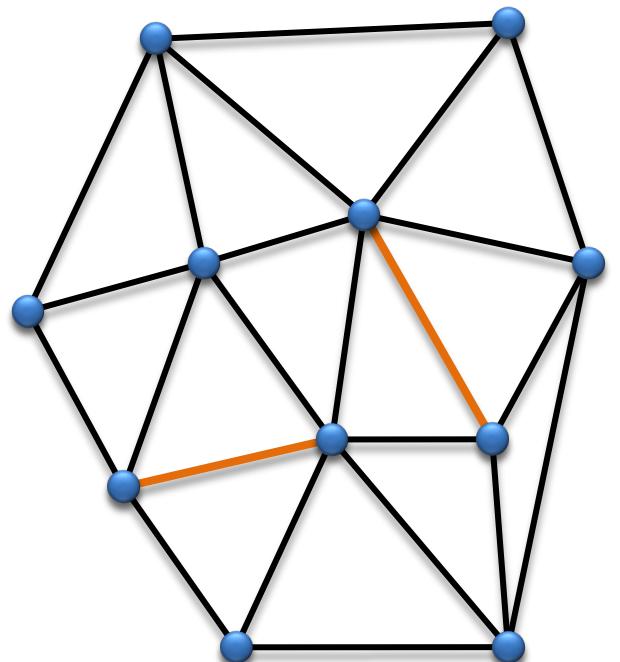
**Plane** Graph



**Straight Line** Plane Graph

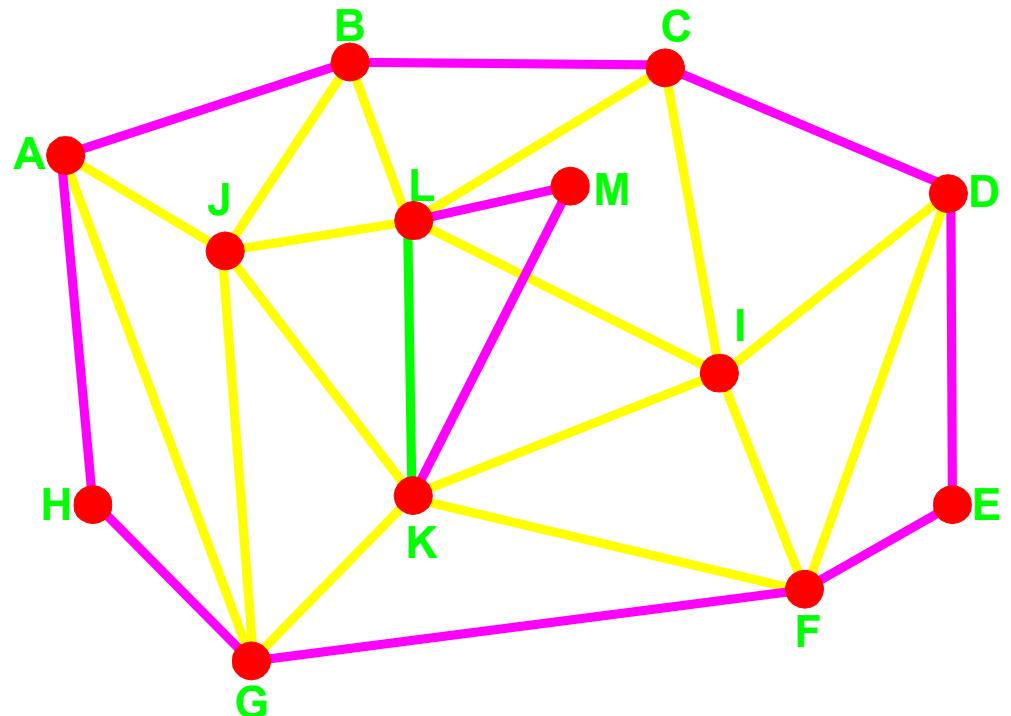


# Triangulation



**Triangulation:**  
Straight line plane graph where  
every face is a *triangle*.

# Mesh

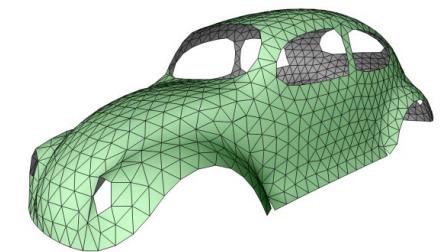


## Mesh:

straight-line graph embedded in  $R^3$

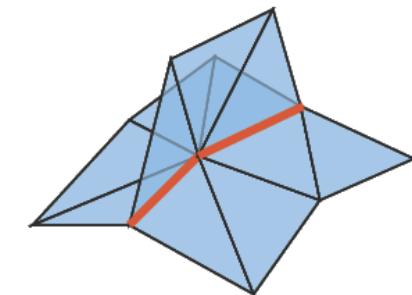
## Boundary edge:

adjacent to exactly *one* face



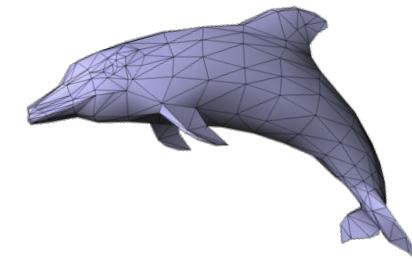
## Regular edge:

adjacent to exactly *two* faces



## Singular edge:

adjacent to more than two faces

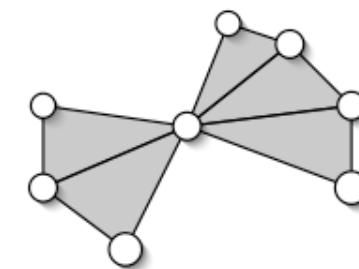
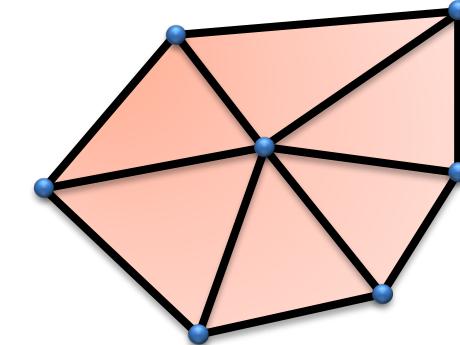
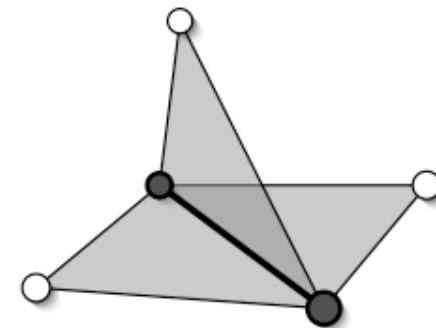
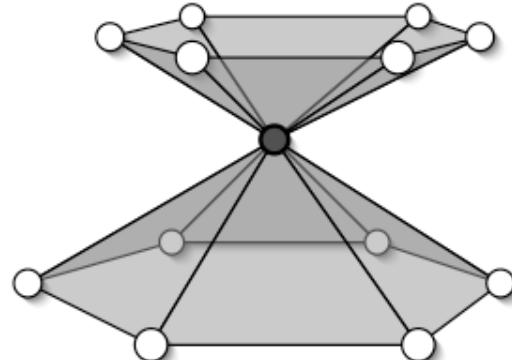


## Closed mesh:

mesh with no boundary edges

# 2-Manifolds Meshes

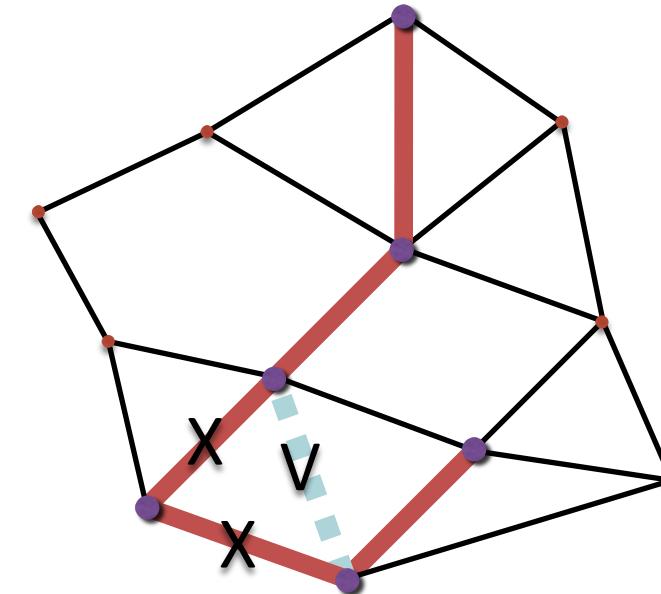
Disk-shaped neighborhoods



non-manifolds

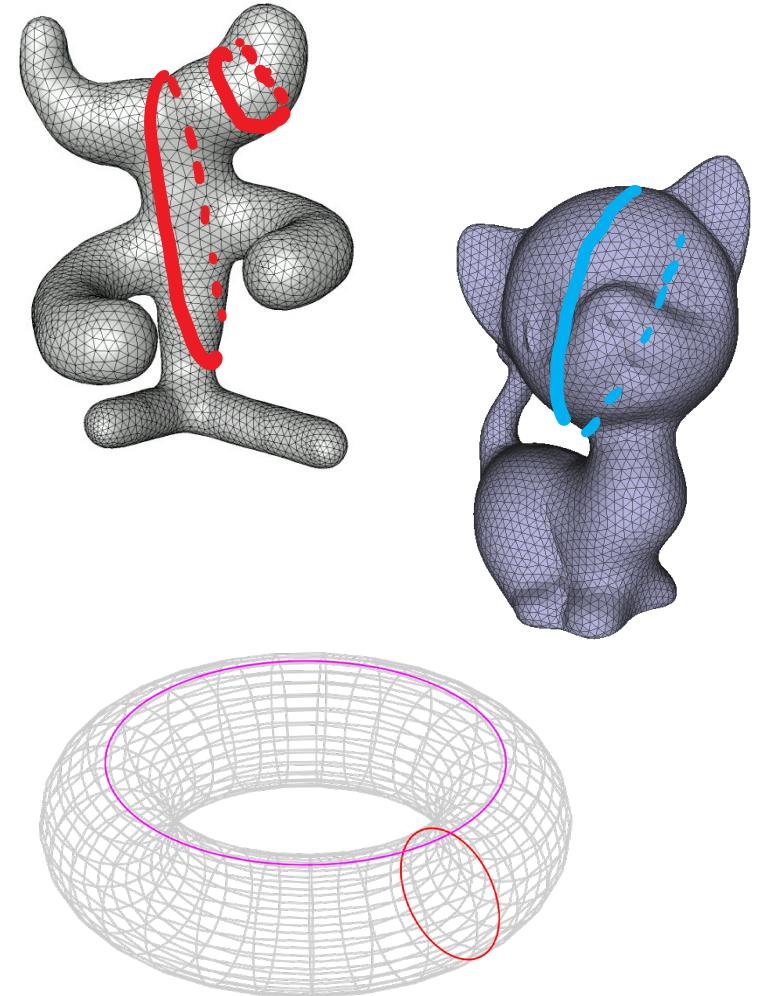
# Topology

- **Path morphing:** local editing without breaking the path connectivity.
- **Independent loops:** cannot be edited to each other.



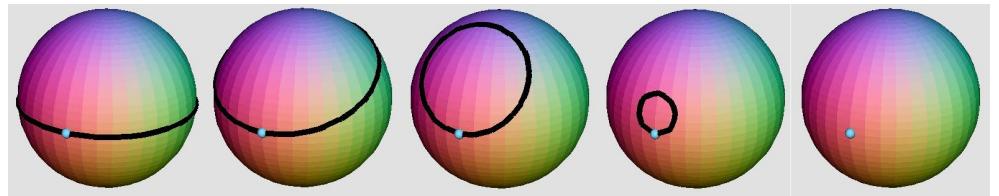
# Topology

- **Contractible loop**: closed path that can be **contracted** to a single vertex.
- Otherwise, **incontractible** loops.
- **Handles**: incontractible loops that do not disconnect the mesh.
  - Which incontractible loops are not handles?



# Topology

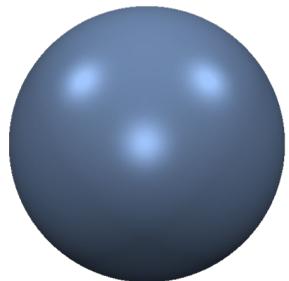
- **Simply-connected patch**: all loops are contractible.
- **Genus**, Informal: #independent handles.
- Genus 0 **does not** imply simply connected! Why?



# Global Topology: Genus

**Genus:**

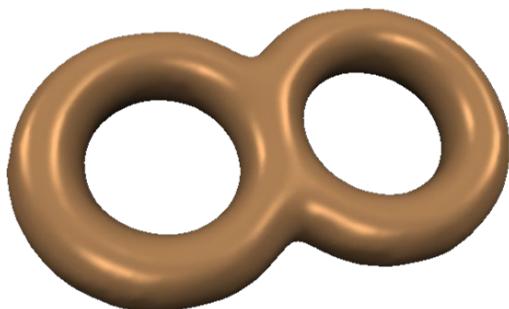
Half the maximal number of closed paths that do not disconnect the mesh  
= the number of holes



Genus 0



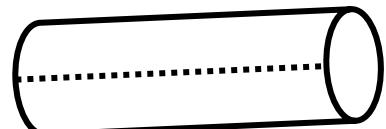
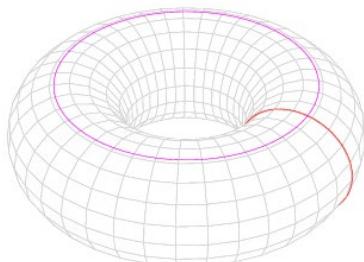
Genus 1



Genus 2



Genus ?



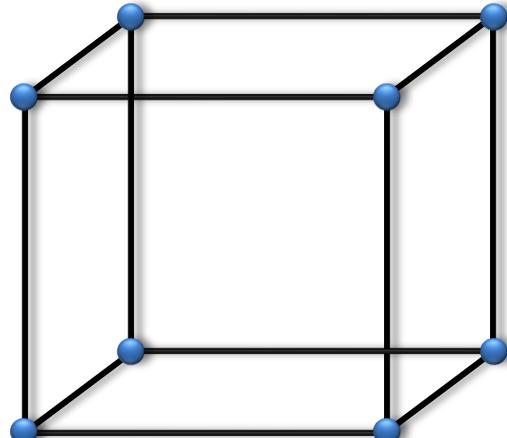
# Closed 2-Manifold Polygonal Meshes

Euler-Poincaré formula

[\*\*20 proofs\*\*](#)

[Video of proof](#)

$$V + F - E = \boxed{\chi} \quad \text{Euler characteristic}$$

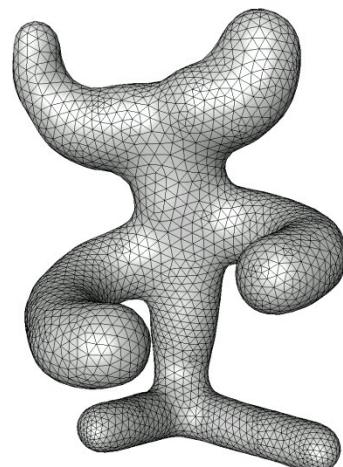


$$V = 8$$

$$E = 12$$

$$F = 6$$

$$\chi = 8 + 6 - 12 = \mathbf{2}$$



$$V = 3890$$

$$E = 11664$$

$$F = 7776$$

$$\chi = \mathbf{2}$$

# Closed 2-Manifold Polygonal Meshes

Euler-Poincaré formula

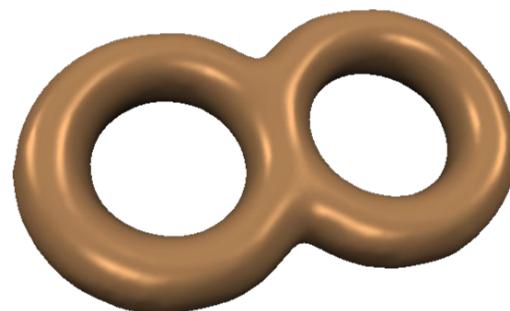
$$V + F - E = \chi = 2$$



$$V = 1500, E = 4500$$

$$F = 3000, g = 1$$

$$\chi = 0$$



$$g = 2$$

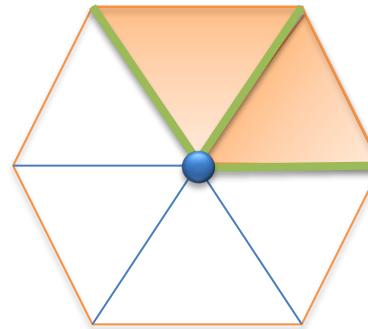
$$\chi = -2$$

# Closed 2-Manifold Triangle Meshes

- *Triangle* mesh statistics

$$E \approx 3V$$

$$F \approx 2V$$



- Avg. valence  $\approx 6$

*Show using Euler Formula*



- When can a closed triangle mesh be 6-regular?



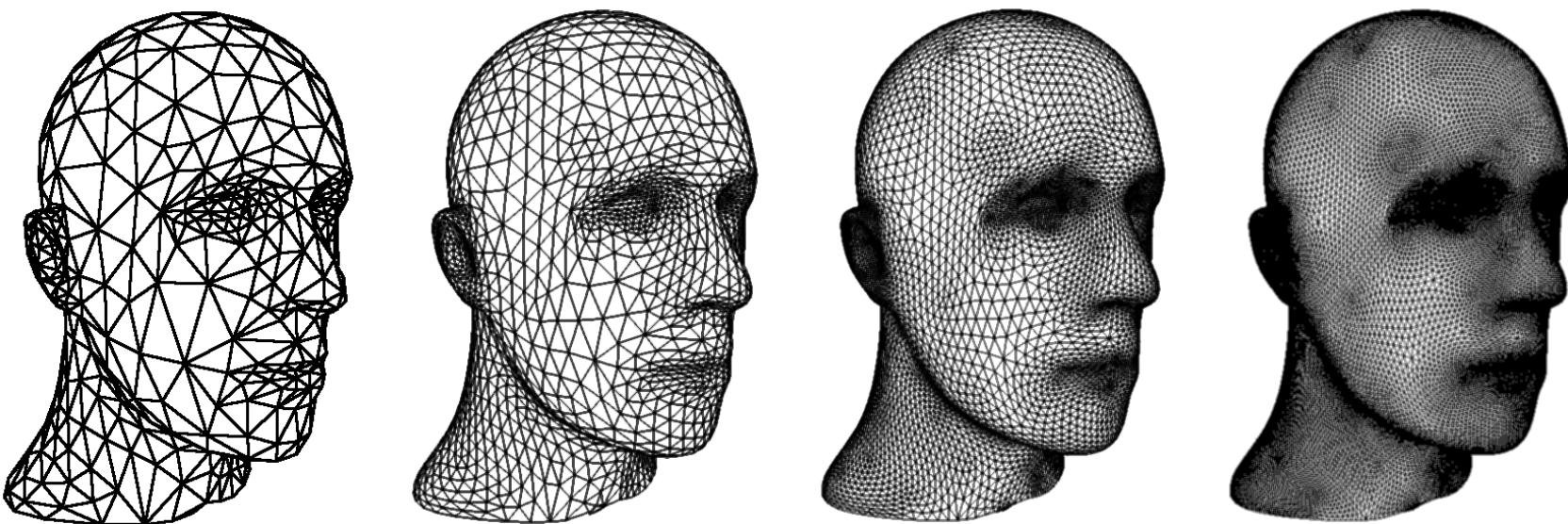
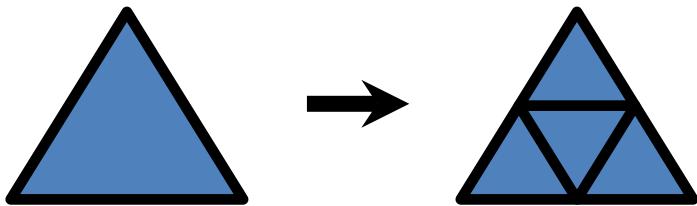
# How Many Pentagons?

- Closed genus 0 surface
- Only hexagons and pentagons
- 3-Regular vertices
- How many pentagons?

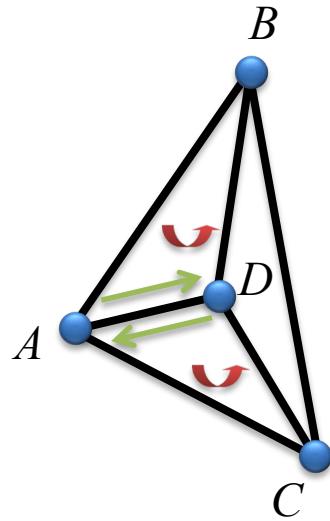


# Regularity

- semi-regular



# Orientability



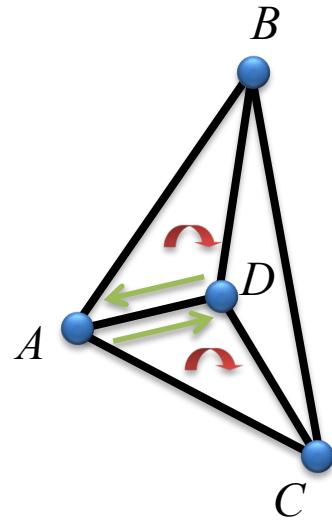
**Face Orientation =**  
clockwise or anticlockwise order in  
which the vertices listed

defines direction of face **normal**

Oriented CCW:  
 $\{(C, \text{D}, \text{A}), (\text{A}, \text{D}, \text{B}), (\text{C}, \text{B}, \text{D})\}$

Oriented CW:  
 $\{(\text{C}, \text{A}, \text{D}), (\text{D}, \text{A}, \text{B}), (\text{B}, \text{C}, \text{D})\}$

# Orientability



**Face Orientation =**  
clockwise or anticlockwise order in  
which the vertices listed

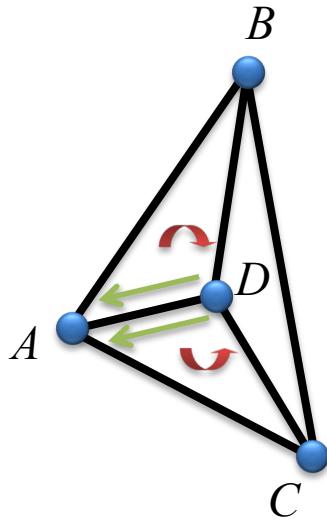
defines direction of face **normal**

Oriented **CCW**:  
 $\{(C,D,A), (A,D,B), (C,B,D)\}$

Oriented **CW**:  
 $\{(C,\text{A},\text{D}), (\text{D},\text{A},B), (B,C,D)\}$

Not oriented:  
 $\{(C,D,A), (D,A,B), (C,B,D)\}$

# Orientability



**Face Orientation =**  
clockwise or anticlockwise order in  
which the vertices listed

defines direction of face **normal**

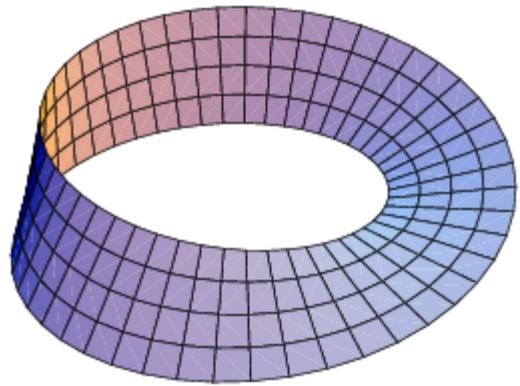
Oriented **CCW**:  
 $\{(C,D,A), (A,D,B), (C,B,D)\}$

Oriented **CW**:  
 $\{(C,A,D), (D,A,B), (B,C,D)\}$

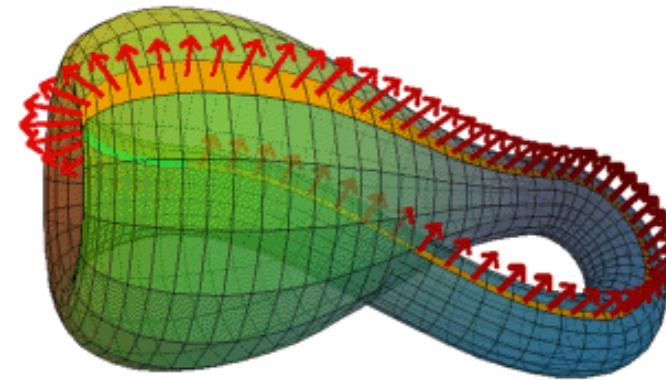
**Not** oriented:  
 $\{(C,\text{D},\text{A}), (\text{D},\text{A},\text{B}), (C,B,D)\}$

**Orientable Plane Graph =**  
orientations of faces can be chosen  
so that each non-boundary edge is  
oriented in *both* directions

# Non-Orientable Surfaces



Möbius Strip

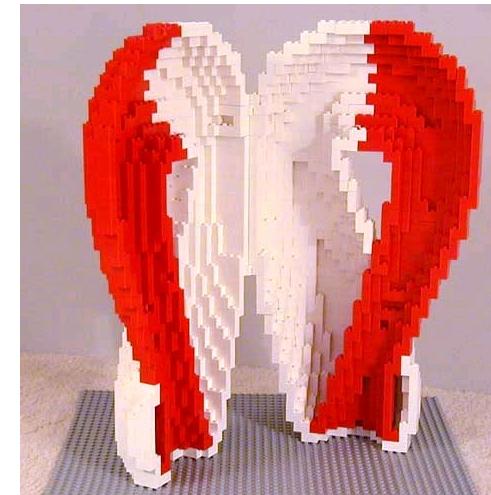


Klein Bottle

# Garden Variety Klein Bottles

Glass Klein Bottles for sale - inquire within

Need a zero-volume bottle?  
Searching for a one-sided surface?  
Want the ultimate in non-orientability?  
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