



Module 13

Session Hijacking

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Basics Concepts

1. Introduction



Basics Concepts

- Session hijacking refers to an attack where an attacker **takes over** a valid **TCP** communication **session** between two computers.
- Since most **authentication** only occurs at the **start** of a **TCP** session, this allows the attacker to gain access to a machine.
- Attackers can **sniff** all the **traffic** from the **established TCP sessions** and perform **identity theft**, **information theft**, **fraud**, etc.
- The attacker **steals** a valid **session ID** and use it to **authenticate** himself with the server.

2. Why Session Hijacking is Successful?



Basics Concepts

- No account lockout for invalid session IDs.
- Weak session ID generation algorithm or small session IDs.
- Insecure handling of session IDs.
 - ▷ DNS poisoning, XSS, exploiting a bug in browser
- Indefinite session expiration time.
- Most computers using TCP/IP are vulnerable.
- Most countermeasures do not work unless you use encryption.

3. Session Hijacking Process



Basics Concepts

■ **Stealing:** The attacker uses different techniques to steal session IDs.

- ▷ Some of the techniques used to steal session IDs:
 - ▷ Using the HTTP **referrer** header.
 - ▷ **Sniffing** the network **traffic**.
 - ▷ Using the **cross-site-scripting** attacks.
 - ▷ Sending **Trojans** on client machines.



Basics Concepts

Guessing: The attacker tries to guess the session IDs by **observing variable parts** of the session IDs.

- ▶ <http://www.hacksite.com/view/VW48266762824302>
- ▶ <http://www.hacksite.com/view/VW48266762826502>
- ▶ <http://www.hacksite.com/view/VW48266762828902>

Brute Forcing: The attacker **attempts different** IDs until he succeeds.

- ▶ Using brute force attacks, an attacker **tries to guess** a session ID until he finds the correct session ID.



Basics Concepts

Stealing Session IDs:

- ▷ Using a "referrer attack," an attacker tries to **lure a user** to **click** on a **link to malicious** site (say www.hacksite.com)
- ▷ For example, GET /index.html HTTP/1.0 Host: www.hacksite.com
Referrer:
www.webmail.com/viewmsg.asp?msgid=689645&SID=2556X54VA75
- ▷ The browser **directs the referrer** URL that **contains** the **user's session ID** to the **attacker's site** (www.hacksite.com), and now the attacker possesses the user's session ID.



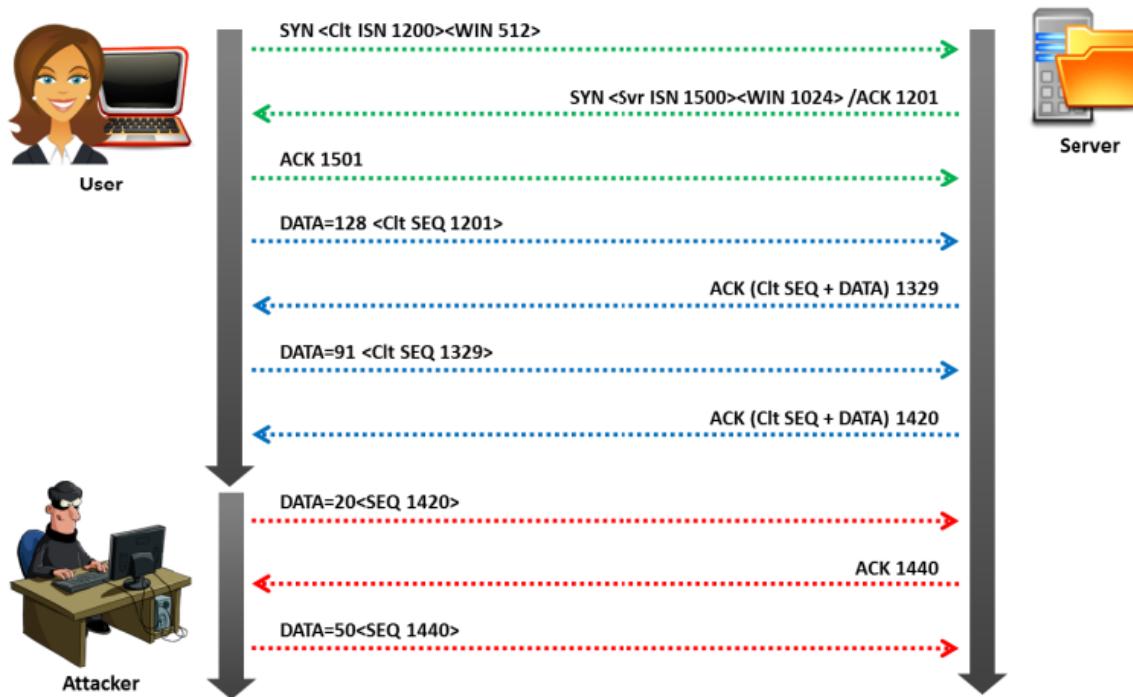
Basics Concepts

- **Note:** Session ID brute forcing attack is known as **session prediction attack** if the **predicted range** of values for a session ID is **very small**.
- **Command Injection:** Start injecting packets to the target server.
- **Session ID prediction:** **Take over** the session.
- **Session Desynchronization:** **Break the connection** to the victim's machine.
- **Monitor:** Monitor the **flow of packets** and predict the **sequence number**.
- **Sniff:** Place yourself **between** the **victim** and the **target** (you must be able to sniff the network).

4. Packet Analysis of a Local Session Hijack



Basics Concepts





Basics Concepts

- According to the diagram, the next expected sequence number would be 1420. If you can transmit that packet sequence number before the user does, you can desynchronize the connection between the user and the server.
- After establishing the connection between the attacker and the server, though the user sends the data with the correct sequence number, the server drops the data considering it as a resent packet.

5. Types of Session Hijacking



Basics Concepts

- **Active Attack:** In an active attack, an attacker finds an active session and takes over.
- **Passive Attack:** With a passive attack, an attacker hijacks a session but sits back and watches and records all the traffic that is being sent forth.

6. Session Hijacking in OSI Model



Basics Concepts

- **Network Level Hijacking:** Network level hijacking can be defined as the **interception** of the packets during the transmission between the client and the server in a **TCP and UDP session**.
- **Application Level Hijacking:** Application level hijacking is about gaining **control** over the **HTTP's user session** by obtaining the session IDs.

7. Spoofing vs Hijacking



Basics Concepts

Spoofing Attack:

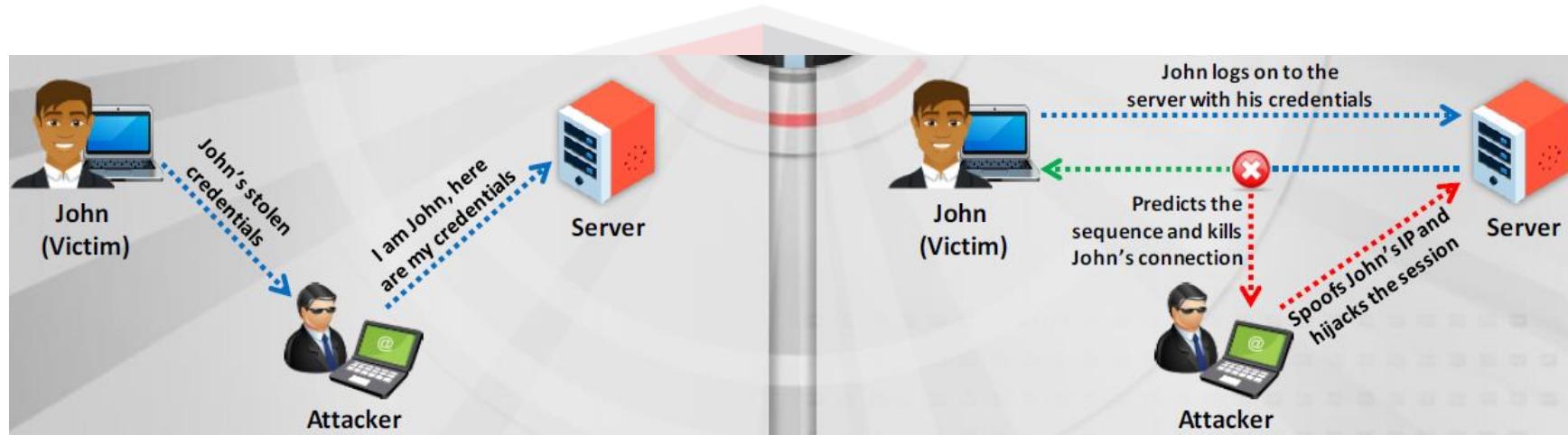
- ▶ Attack pretends to be another user or machine (victim) to gain access.
- ▶ Attacker does not take over an existing active session. Instead he initiates a new session using the victim's stolen credentials.

Hijacking:

- ▶ Session hijacking is the process of taking over an existing active session.
- ▶ Attacker relies on the legitimate user to make a connection and authenticate.



Basics Concepts



Application Level Session Hijacking

Module 13





Application Level Session Hijacking

- A session token is **stolen** or valid session token is **predicted** to gain unauthorized access to the web server.
- A session token can be compromised in various ways:
 - ▷ Session **sniffing**, Session **replay** attack, Session **fixation**
 - ▷ **Predictable** session **token**
 - ▷ **Man-in-the-middle** attack
 - ▷ **Man-in-the-browser** attack
 - ▷ **Cross-site script** attack
 - ▷ Cross-site request forgery attack (**CSRF**)

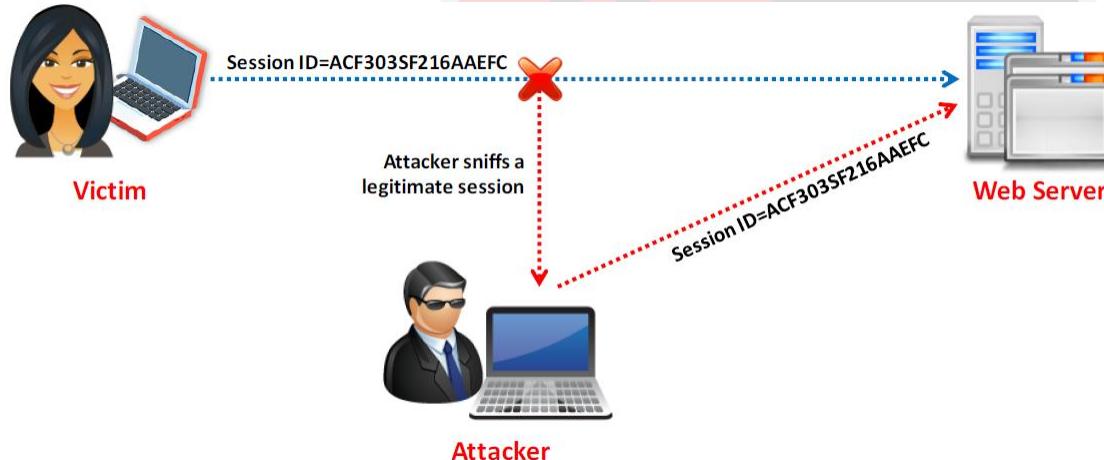
1. Sniffing



Application Level Session Hijacking

Compromising Sessions IDs using Sniffing

- ▶ Attacker uses a **sniffer** to **capture** a valid session **token** or session ID.
- ▶ Attacker then uses the valid token session to gain unauthorized access to the web server.



2. Predicting



Application Level Session Hijacking

Compromising Session IDs by Predicting Session Token

- ▶ Attacker can predict session IDs generated by weak algorithms and impersonate a web site user.
- ▶ Attackers perform analysis of variable section of session IDs to determine the existence of a pattern.
- ▶ The analysis is performed manually or by using various cryptanalytic tools.
- ▶ Attackers collect a high number of simultaneous session IDs in order to gather samples in the same time window and keep the variable constant.



Application Level Session Hijacking

How to Predict a Session Token

- ▶ Most of the web servers use **custom algorithms** or a **predefined pattern** to generate sessions IDs.
- ▶ Attacker **guess** the **unique** session value or **deduce** the session ID to hijack the sessions.
- ▶ **Captures:** Attacker **captures several** session IDs and **analyzes** the pattern.
 - ▶ <http://www.juggyboy.com/view/JBEX25022014152820>
 - ▶ <http://www.juggyboy.com/view/JBEX25022014153020>
 - ▶ <http://www.juggyboy.com/view/JBEX25022014160020>
 - ▶ <http://www.juggyboy.com/view/JBEX25022014164020>



Application Level Session Hijacking

- **Predicts:** At **16:25:55** on **Feb-25, 2014**, the attacker can successfully **predict** the session ID to be <http://www.juggyboy.com/view/JBEX25022014162555>
 - ▷ **JBEX:** Constant
 - ▷ **25022014:** Date
 - ▷ **162555:** Time

3. Man in the Middle

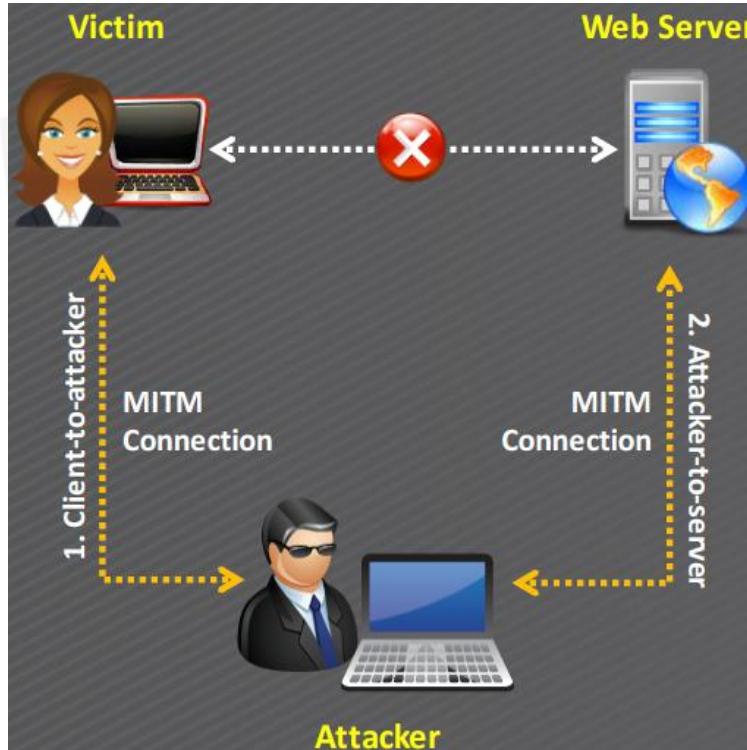


Application Level Session Hijacking

- The man-in-the-middle attack is used to **intrude** into an **existing connection** between systems and to **intercept messages** being exchanged.
- Attackers use different techniques and split the TCP connection into two connections.
 - ▶ **Client-to-attacker** connection
 - ▶ **Attacker-to-server** connection
- After the successful interception of TCP connection, an attacker can **read**, **modify**, and **insert** fraudulent data into the intercepted communication.
- In the case of an **http** transaction, the TCP connection between the client and the server becomes the target.



Application Level Session Hijacking



4. Man in the Browser



Application Level Session Hijacking

- Man-in-the-browser attack uses a **Trojan Horse** to **intercept** the **calls** between the **browser** and its **security mechanisms or libraries**.
- It works with an **already installed** Trojan horse and acts between the browser and its security mechanisms.
- Its main objective is to **cause financial deceptions** by manipulating transactions of **Internet Banking systems**.



Application Level Session Hijacking

Steps to Perform Man-in-the-Browser Attack

- The Trojan first **infects** the computer's software (OS or application).
- The Trojan **installs** malicious code (extension files) and **saves** it into the browser configuration.
- After the user **restarts** the **browser**, the malicious **code** in the form of extension files is **loaded**.
- The extension files **register** a **handler** for every visit to the **webpage**.
- When the page is loaded, the **extension uses the URL** and matches it with a **list of known sites** targeted for attack.



Application Level Session Hijacking

- ▷ The user logs in securely to the website.
- ▷ It registers a button event handler when a specific page load is detected for a specific pattern and compares it with its targeted list.
- ▷ When the user clicks on the button, the extension uses DOM interface and extracts all the data from all form fields and modifies the values.
- ▷ The browser sends the form and modified values to the server.



Application Level Session Hijacking

- The server receives the modified values but cannot distinguish between the original and the modified values.
- After the server performs the transaction, a receipt is generated.
- Now, the browser receives the receipt for the modified transaction.
- The browser displays the receipt with the original details.
- The user thinks that the original transaction was received by the server without any interceptions.

5. Client Side Attacks



Application Level Session Hijacking

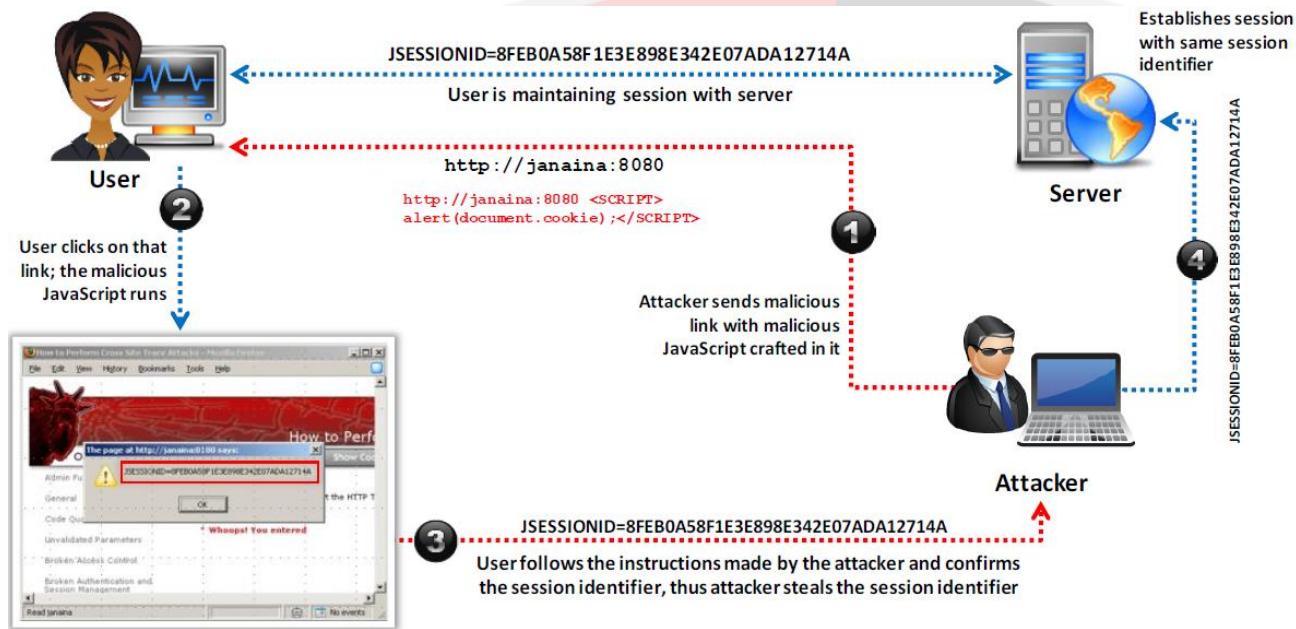
Compromising Session IDs Using Client-side Attacks

- ▷ **Cross-Site Scripting (XSS)**: XSS enables attackers to inject malicious client side scripts into the web pages viewed by other users.
- ▷ **Malicious JavaScript Codes**: A malicious script can be embedded in a web page that does not generate any warning but it captures session tokens in the background and send it to the attacker.
- ▷ **Trojans**: A Trojan horse can change the proxy settings in user's browser to send all the sessions through the attackers machine.



Application Level Session Hijacking

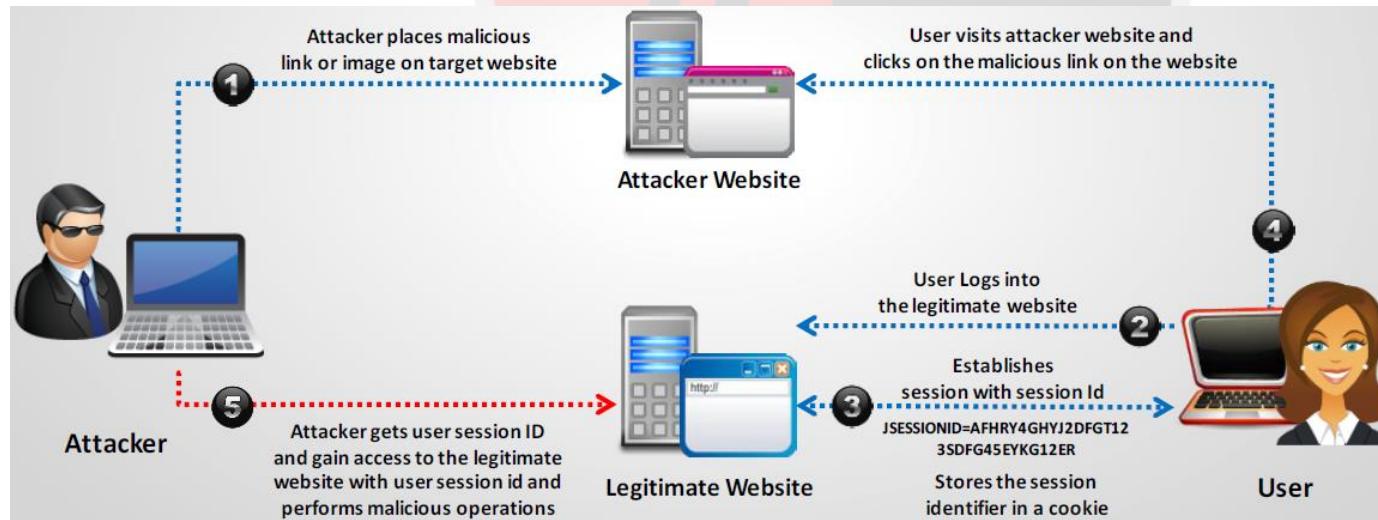
Cross-site Script Attack





Application Level Session Hijacking

Cross-site Request Forgery Attack

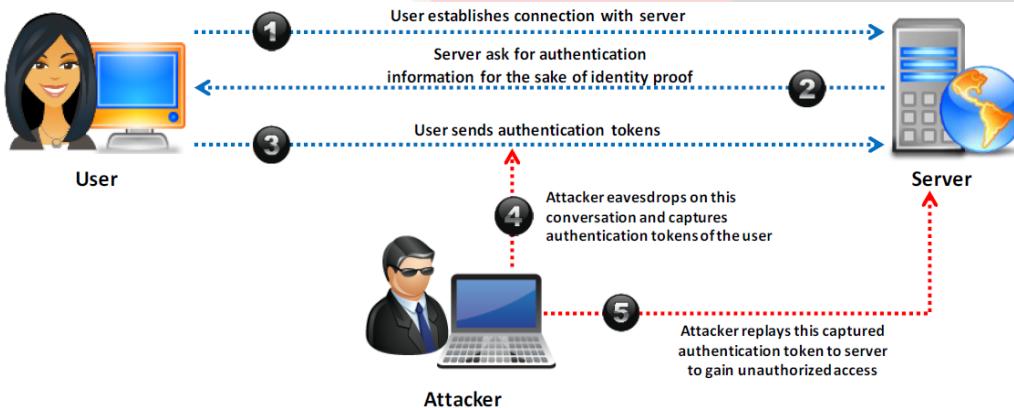


6. Replay Attacks



Application Level Session Hijacking

- In a session replay attack, the attacker **listens** to the conversation between the user and the server and **captures** the authentication token of the user.
- Once the authentication token is captured, the attacker **replays** the **request** to the **server** with the captured authentication token and gains unauthorized access to the server.



7. Session Fixation

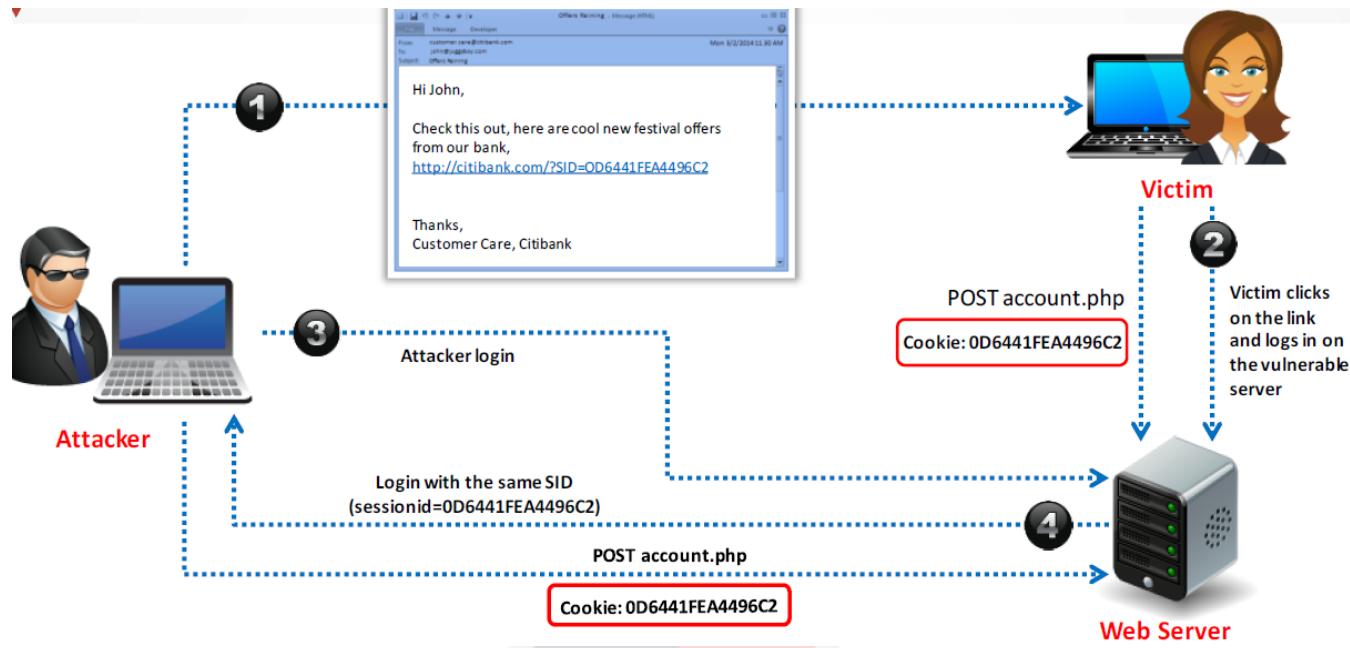


Application Level Session Hijacking

- The attack tries to **lure** a user to authenticate himself with a **known** session ID and then **hijacks** the **user-validated** session by the **knowledge** of the **used** session ID.
- The attacker has to **provide** a **legitimate** web application session ID and try to lure victim browser to use it.
- Several techniques to execute Session Fixation attack are:
 - ▶ Session **token** in the **URL** argument
 - ▶ Session **token** in a **hidden form** field
 - ▶ Session **ID** in a **cookie**



Application Level Session Hijacking



8. Proxy Servers



Application Level Session Hijacking

- Attacker lure victim to click on bogus link which looks legitimate but redirect user to attacker server.
- Attacker forwards request to the legitimate server on behalf of victim and serve as a proxy for the entire transaction.
- Attacker then captures the sessions information during interaction of legitimate server and user.

Network Level Session Hijacking

Module 13



Network Level Session Hijacking

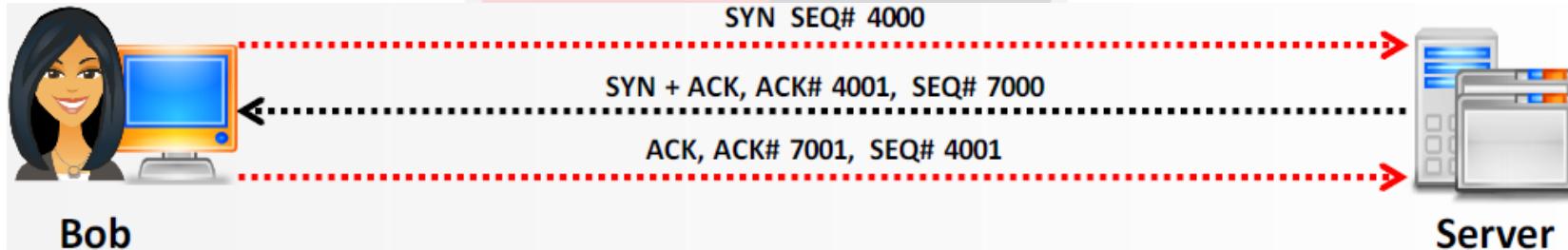
- The network-level hijacking relies on **hijacking transport** and **Internet protocols** used by web applications in the application layer.
- Attacker gathers some critical information **used to attack application level**.
- Network-level hijacking includes:
 - ▷ **Blind Hijacking**
 - ▷ **TCP/IP Hijacking, UDP Hijacking**
 - ▷ **RST Hijacking**
 - ▷ Man-in-the-Middle: **Packet Sniffer**
 - ▷ **IP Spoofing**: Source Routed Packets



Network Level Session Hijacking

The 3-Way Handshake

- If the attacker can anticipate the next sequence and ACK number that Bob will send, he/she will spoof Bob's address and start a communication with the server.



1. TCP/IP Hijacking

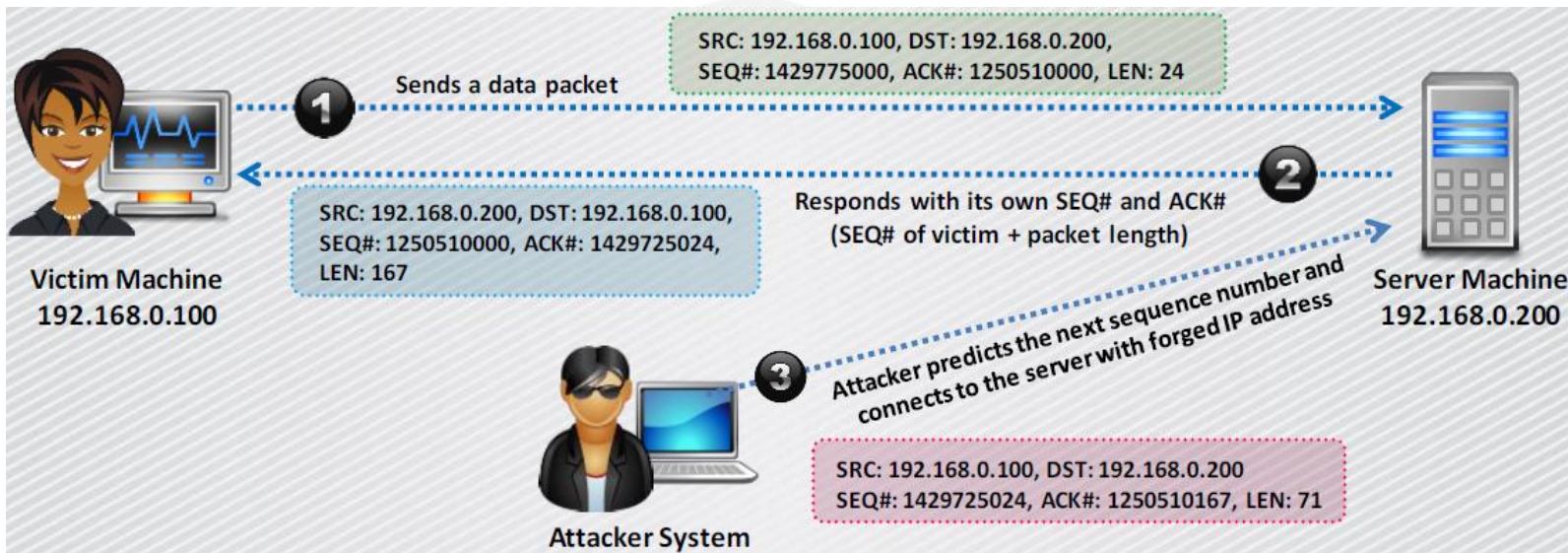


Network Level Session Hijacking

- TCP/IP hijacking is a hacking technique that uses **spoofed packets** to take over a connection between a victim and a target machine.
- The **victim's connection hangs** and the **attacker** is then able to communicate with the **host's machine** as if the **attacker is the victim**.
- To launch a TCP/IP hijacking attack, the attacker **must** be on **the same network** as the **victim**.
- The target and the victim machines can be **anywhere**.



Network Level Session Hijacking





Network Level Session Hijacking

- The attacker **sniffs** the victim's connection and **uses** the victim's IP to send a **spoofed packet** with the predicted sequence number.
- The receiver **processes** the **spoofed packet**, **increments** the **sequence number**, and **sends acknowledgement** to the victim's IP.
- The victim machine is **unaware** of the **spoofed packet**, so it **ignores** the receiver machine's **ACK** packet and **turns sequence number count off**.
- Therefore, the receiver **receives** packets with the **incorrect sequence number**.



Network Level Session Hijacking

- The attacker forces the victim's connection with the receiver machine to a desynchronized state.
- The attacker tracks sequence numbers and continuously spoofs packets that comes from the victim's IP.
- The attacker continues to communicate with the receiver machine while the victim's connection hangs.



2. IP Spoofing



Network Level Session Hijacking

- Packet **source routing** technique is used for gaining unauthorized access to a computer with the help of a **trusted host's IP address**.
- The attackers **spoofs** the **host's IP** address so that the server managing a session with the host, accepts the packets from the attacker.
- When the session is established, the attacker **injects forged packets** before the **host responds** to the server.
- The **original packet** from the **host is lost** as the server **gets the packet** with a sequence number already used by the attacker.
- The packets are **source-routed** where the **path** to the **destination IP** can be **specified** by the attacker.

3. RST Hijacking

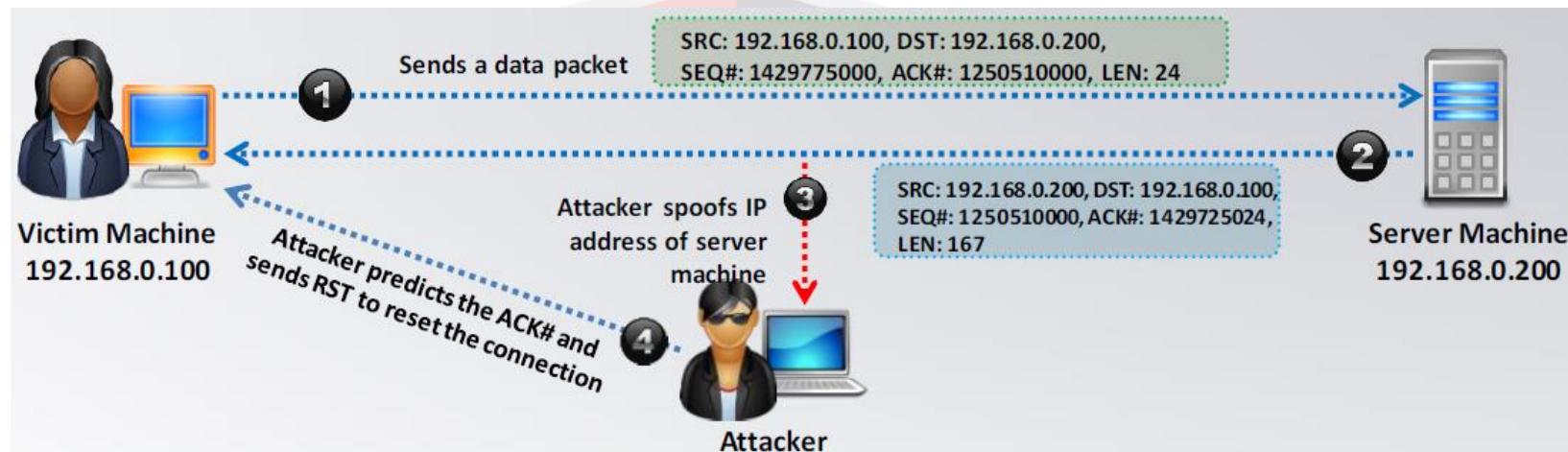


Network Level Session Hijacking

- RST hijacking involves injecting an authentic-looking reset (RST) packet using spoofed source address and predicting the acknowledgment number.
- The hacker can reset the victim's connection if it uses an accurate acknowledgement number.
- The victim believes that the source actually sent the reset packet and resets the connection.
- RST Hijacking can be carried out using a packet crafting tool such as Colasoft's Packet Builder and TCP/IP analysis tool such as tcpdump.



Network Level Session Hijacking

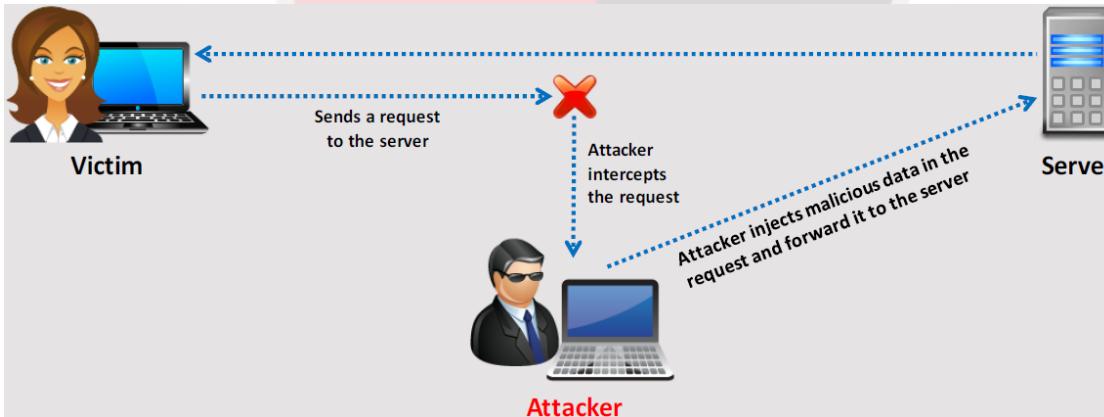


4. Blind Hijacking



Network Level Session Hijacking

- The attacker can inject the **malicious data or commands** into the intercepted communications in the TCP session even if the **source-routing** is **disabled**.
- The **attacker** can send the data or commands but **has no access** to see the response.





5. Forged ICMP and ARP Spoofing



Network Level Session Hijacking

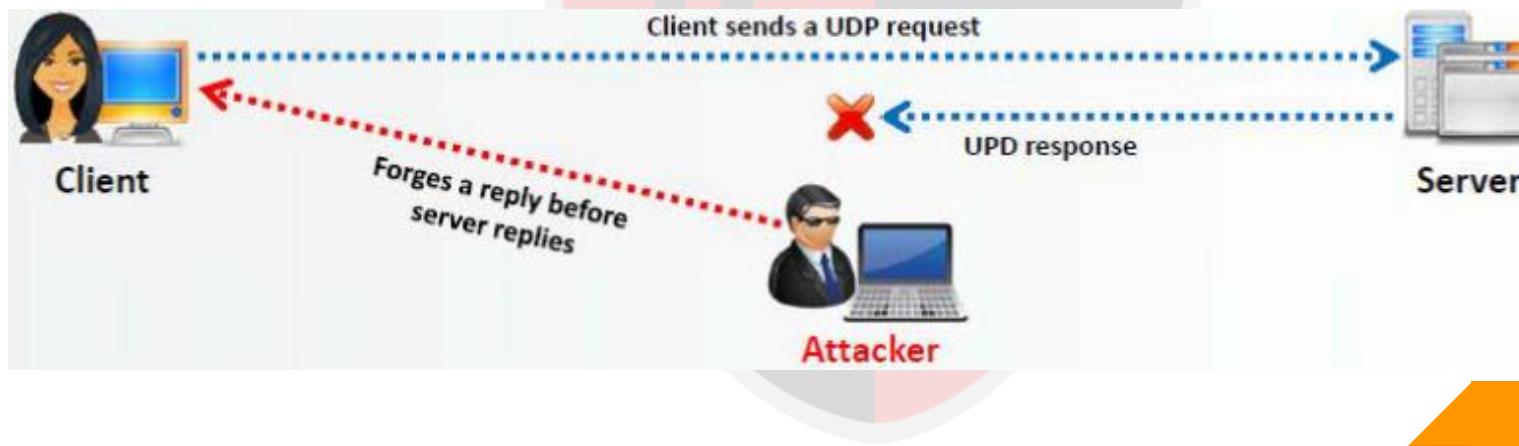
- Packet sniffer is used as an interface between the client and the server.
- ARP spoofing involves fooling the host by broadcasting the ARP request and changing its ARP tables by sending the forged ARP replies.
- The packets between the client and the server are routed through the hijacker's host by using two techniques:
 - ▷ Using **Forged Internet Control Message Protocol (ICMP)**: It is an extension of IP to send error messages where the attacker can send messages to fool the client and the server.
 - ▷ **Using Address Resolution Protocol (ARP) Spoofing**: ARP is used to map the network layer address (IP address) to link layer addresses (MAC address).

6. UDP Hijacking



Network Level Session Hijacking

- A network-level session hijacking where the attacker sends **forged server reply** to a victim's UDP request before the **intended server** replies to it.
- The attacker uses **man-in-the-middle** attack to **intercept** server's **response** to the client and sends its **own forged reply**.



Countermeasures

1. Detection



Countermeasures

■ Detection Method

- ▷ **Manual** Method
 - ▷ Using **Packet Sniffing** Software
 - ▷ Normal **Telnet** Session
 - ▷ Forcing an ARP Entry
- ▷ **Automatic** Method
 - ▷ Intrusion Detection Systems (IDS)
 - ▷ Intrusion Prevention Systems (IPS)

2. Protection



Countermeasures

- Use **Secure Shell (SSH)** to create a secure communication channel.
- Pass the authentication cookies over **HTTPS** connection.
- Implement the **log-out** functionality for user to **end the session**.
- Generate the session **ID** after **successful login** and **accept sessions IDs generated by server only**.
- Ensure data in transit is **encrypted** and implement **defense-in-depth** mechanism.

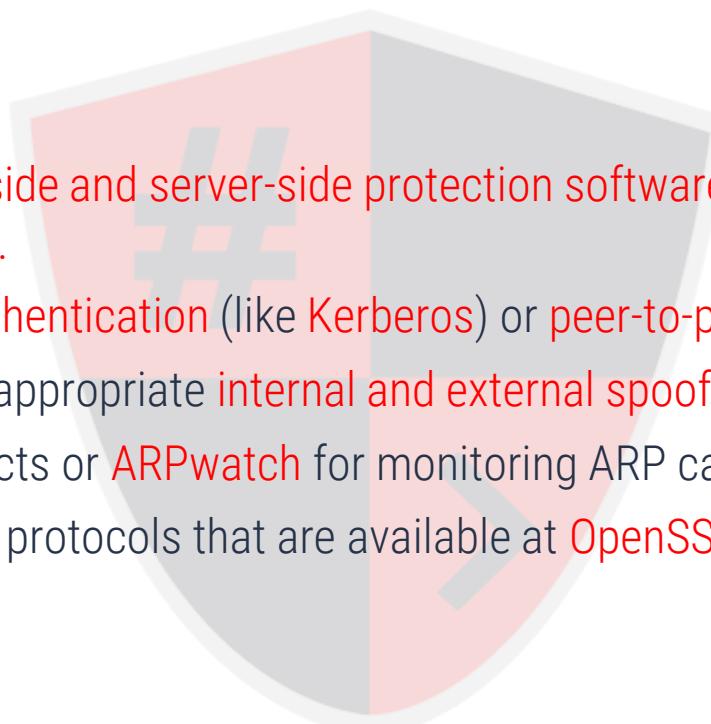


Countermeasures

- Use string or long random number as a session key.
- Use different user name and passwords for different accounts.
- Educate the employees and minimize remote access.
- Implement timeout() to destroy the session when expired.
- Do not transport session ID in query string.
- Use switches rather than hubs and limit incoming connections.



Countermeasures

- 
- Ensure client-side and server-side protection software are in active state and up to date.
 - Use strong authentication (like Kerberos) or peer-to-peer VPN's.
 - Configure the appropriate internal and external spoof rules on gateways.
 - Use IDS products or ARPwatch for monitoring ARP cache poisoning.
 - Use encrypted protocols that are available at OpenSSH suite.



Countermeasures

For Web Developers:

- Create session keys with **lengthy strings or random number** so that it is **difficult** for an attacker **to guess** a valid session key.
- Regenerate the session ID after a successful login to **prevent session fixation** attack.
- Encrypt the **data and session** key that is transferred between the user and the web servers.
- Expire the session as soon as the user **logs out**.
- Prevent Eavesdropping within the network.
- Reduce the **life span** of a session or a cookie.



Countermeasures

For Web Users:

- ▷ Do not click on the **links** that are received through mails or IMs.
- ▷ Use **Firewalls** to prevent the malicious content from entering the network.
- ▷ Use **firewall and browser** settings to **restrict cookies**.
- ▷ Make sure that the **website** is certified by the certifying authorities.
- ▷ Make sure you **clear history, offline content**, and **cookies** from your browser **after every** confidential and sensitive **transaction**.
- ▷ Prefer **https**, a secure transmission, **rather than http**
- ▷ **Logout** from the browser **instead of closing** the browser.



Countermeasures

Issue	Solution	Notes
Telnet, rlogin	OpenSSH or ssh (Secure Shell)	It sends encrypted data and makes it difficult for attacker to send the correctly encrypted data if session is hijacked
FTP	sFTP	It reduces the chances of successful hijacking
HTTP	SSL (Secure Socket Layer)	It reduces the chances of successful hijacking
IP	IPSec	It prevents hijacking by securing IP communications
Any Remote Connection	VPN	Implementing encrypted VPN such as PPTP, L2PT, IPSec, etc. for remote connection prevents session hijacking
SMB (Server Message Block)	SMB signing	It improves the security of the SMB protocol and reduces the chances of session hijacking
Hub Network	Switch Network	It mitigates the risk of ARP spoofing and other session hijacking attacks

3. IPSec



Countermeasures

- IPSec is a **protocol suite** developed by the **IETF** for securing IP communications by authenticating and encrypting each IP packet of a communication session.
- It is deployed widely to implement **virtual private networks** (VPNs) and for remote user access through dial-up connection to **private networks**.
- **Benefits:**
 - ▷ Network-level peer authentication
 - ▷ Data origin authentication
 - ▷ Data **integrity**
 - ▷ Data **confidentiality** (encryption)
 - ▷ Replay protection

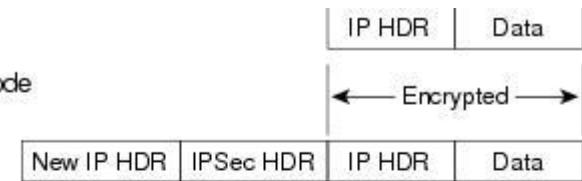


Countermeasures

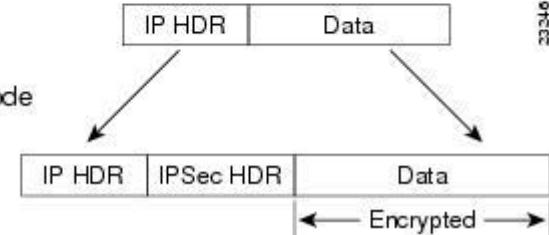
Transport Mode:

- ▷ Authenticates two connected computers
- ▷ Has an option to encrypt data transfer
- ▷ Compatible with NAT
- ▷ Encrypts only IP data

Tunnel mode



Transport mode

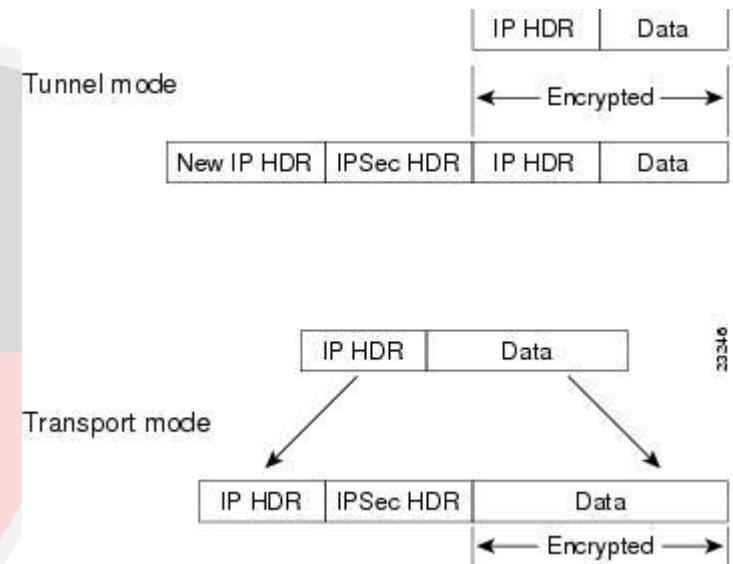




Countermeasures

Tunnel Mode:

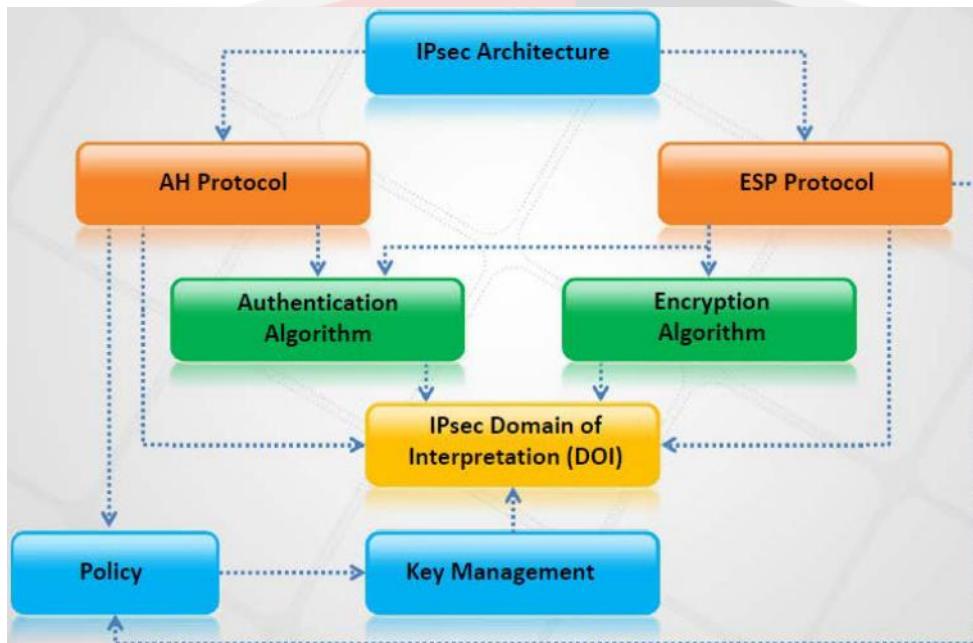
- ▷ Encapsulates packets being transferred
- ▷ Has an option to encrypt data transfer
- ▷ Encrypts TCP Data + Header
- ▷ Not compatible with NAT





Countermeasures

IPSec Architecture





Countermeasures

- IPsec uses two different security services for authentication and confidentiality:
 - ▷ **Authentication Header** (AH): Provide data authentication of the sender.
 - ▷ **Encapsulation Security Payload** (ESP): Provides both data authentication and encryption (confidentiality) of the sender.



Countermeasures

Components of IPsec (?)

- **IPsec driver:** A **software**, that performs **protocol-level functions** that are required to **encrypt** and **decrypt** the packets.
- **Internet Key Exchange (IKE):** IPsec protocol that **produces security keys** for IPsec and other protocols.
- **Internet Security Association Key Management Protocol:** Software that **allows** two computers to **communicate by encrypting** the data that is exchanged between them.



Countermeasures

Components of IPsec (?)

- ▶ **Oakley**: A protocol, which uses the Diffie-Heilman algorithm to **create master key**, and a key that is **specific to each session** in IPsec data transfer.
- ▶ **IPsec Policy Agent**: A **service** of the Windows 2000, **collects IPsec policy settings** from the **active directory** and sets the **configuration** to the system at start up.

HACKING

Is an art, practised through a creative mind.

