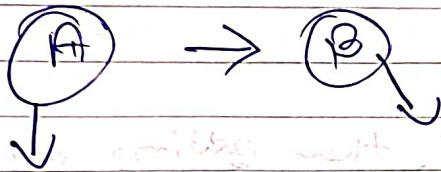


Normalisation

- It is an step towards Database optimisation.
- Functional dependency (FD)

When value of B is dependent on A. Like from taking value of A we can determine B.



Determinant

dependent

→ Types of FD

① Trivial FD

$A \rightarrow B$; B is ^{part} subset of A.

$\{ \text{Emp-ID, name} \} \rightarrow \{ \text{Emp-ID} \}$

② Non-trivial FD

$A \rightarrow B$; B is not a subset of A.

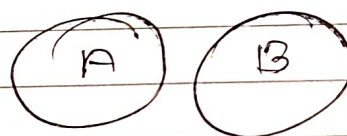
$\{ \text{Emp-ID, name} \} \rightarrow \{ \text{Emp-address} \}$

- Venn diagram of the both types of FD above:-

Trivial FD



Non-trivial FD



$$A \cap B = \text{NULL}$$

→ Rules of FD

① Reflexive

∴ If "A" is a set of attribute and "B" is a subset of "A", then $A \rightarrow B$.

② Augmentation

∴ If 'A' can determine 'B', then adding an attribute to this functional dependency won't change anything.

$$A \rightarrow B$$

$$Ax \rightarrow Bx.$$

③ Transitivity

If A determines B and B determines C, then A determines C.

$$A \rightarrow B, B \rightarrow C, A \rightarrow C.$$

④ Why Normalization?

∴ To avoid redundancy from DB.

⑤ What happen if we have redundant data?
∴ Insertion, deletion and updation anomalies arises.

→ Anomalies / or abnormality

① Insertion anomaly

- When certain data (attribute) can not be inserted into the DB without the presence of other data.

② Deletion anomaly

- The delete anomaly, a situation where the deletion of data results in the unintended loss of some other important data.

③ Updation anomaly

- A single data update requires multiple rows of data to be updated.

→ Due to these anomalies, DB size increases and DB performance become very slow.

→ To rectify these anomalies, we use database optimisation technique called Normalization.

→ Normalisation divides the composite attributes into individual attributes or larger table into smaller and link them using relationships.

⇒ We decompose the table until we achieve "SRP"

SRP ⇒ Single Responsibility Principle,

Like one table contains info of one entity only.

→ Types of Normal forms

① 1NF

- Every relation cell must have atomic value.
- Relation must not have multi-valued attributes.

② 2NF

- Relation must be in 1NF.
- There should not be any partial dependency.
- All non-prime attributes must be fully dependent on Primary key.
- Non prime attributes can not depend on the part of the primary.

③ 3NF

- Relation must be in 2NF.
- No transitivity dependency exists.
- Non-prime attribute should not find a non-prime attribute.

④ BCNF (Boyce - Codd Normal Form)

- Relation must be in 3NF.
- FD: - $A \rightarrow B$, A must be a Super key.
- We must not derive prime attribute from any prime or non-prime attribute.