# Advanced R Programming - Lecture 6 Parallel programming

Krzysztof Bartoszek (slides based on Leif Jonsson's and Måns Magnusson's)

Linköping University

krzysztof.bartoszek@liu.se

7 October 2019 (U1)



- Parallelism
- Theoretical limits
- Parallelism in R
- Balance and subsetting





### Classical introduction: dining philosophers problem



wait until left fork available wait until right fork available eat release left fork release right fork

Solutions: synchronize: pick lower number (Dijkstra) wait for permission (mutex) release and random wait

https://en.wikipedia.org/wiki/Dining\_philosophers\_problem

Mostly we will not (a.s.) be concerned with such problems.



# What is parallelism?

Parallelism

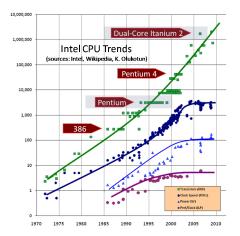
### Multiple cores

Each core work with its own part

Cores can exchange information



### Why parallelism?



# Why parallelism?

Single core limits

Handling larger data

Solving problems faster

More and more important

Is there any **but** ...?



# Types of parallelism

Parallelism

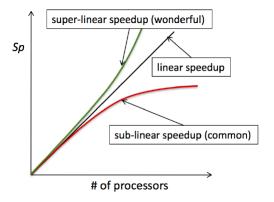
Multicore systems

Distributed systems

Graphical processing units (GPU)



# Speedup



https://portal.tacc.utexas.edu/c/document\_library/get\_file?uuid=

e05d457a-0fbf-424b-87ce-c96fc0077099&groupId=13601



### Theoretical limits

### Strong scaling: Amdahl's law

Deals with fixed problem size, increasing resources

Weak scaling: Gustafsons law

Deals with increasing size problem along with increasing resources

### Amdahl's law

Speedup: 
$$S_p = \frac{\text{execution time on 1 processor}}{\text{execution time on } P \text{ processors}}$$

$$S_p \le \frac{1}{f_s + \frac{f_p}{P}}$$

Where:

 $f_s$  = serial fraction of computations

 $f_p$  = parallel fraction of computations

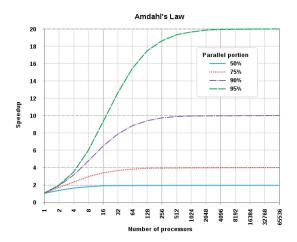
P = number of cores

For a *fixed size problem*, single core computation time is fixed = 1!



lism Theoretical limits Parallelism in R Balance and subsetting

### Amdahl's law



https://en.wikipedia.org/wiki/Amdahl's\_law



### Gustafsons law

$$\text{Speedup}: \quad S_p = \frac{\text{execution time on 1 processor}}{\text{execution time on } P \text{ processors}}$$

$$S_p \leq \frac{\alpha + (1-\alpha)P}{\alpha + (1-\alpha)} = P - \alpha * (P-1)$$

Where:

 $\alpha =$  fraction of time dedicated to serial computations

P = number of cores

Problem size scales with P, parallel execution time is fixed = 1!if we only had one core, then the P parallel computations would have to be done on that core with time  $(1 - \alpha)P$ 



### Practical problems

Costs of parallelism communication

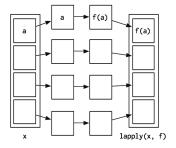
load balancing (NP-hard) scheduling (NP-hard) but 4/3 = k-approximate algorithm

fine-grained vs embarrassingly parallel (EP)



# Parallelism in R (embarassingly parallel)

Based on lapply() iterations inside a loop are indepenendent of each other



(H. Wickham, Advanced R, p. 201)

What about: fctl<-1;for(i in 1:n){fctl<-i\*fctl}?



# parallel package

### Two approaches:

- 1. mclapply()
- 2. parLapply()



# mclapply()

#### **Pros**

Simple to use Low overhead (startup)

#### Cons

Does not work on Windows
Only multi core



# parLapply(type="psock")

### Pros

Works everywhere Good for testing/developing

#### Cons

Slow on multiple nodes



#### **Pros**

Good for multiple computers Good for production

### Cons

Can be used interactively Needs Rmpi package



### Load balance

```
load balancing is NP-hard
often we do not know running time for list[[i]]
submit to cores in order?
split into consecutive equal chunks?
Example:
all 2^p regression models on p predictors
```

generate all combinations:  $\emptyset$ ,  $\{1\}$ , ...,  $\{p\}$ ,  $\{1,2\}$ , ...

lapply(list,...) does in order of list

randomize order in list

R's par family has load balancing capabilities



## Subset methods (non-EP)

glm(): large number of observations

Chunk averaging (estimation)

break data into chunks of rows to each chunk (in parallel) apply glm() average the results to obtain single estimate

Observations i.i.d. and model is *decent* Asymptotic equivalance for large samples

Chunks are not identically distributed: first cases then controls randomly permute observations (will not harm an initial random arrangement)



# Subset methods (non-EP)

glm(): large number of predictors

**Subsetting variables** (prediction, *decent* model)

create random subsets of predictors to each subset apply glm() do prediction for each subset combine predictions e.g. average, majority rule

prediction difficult in high dimensions—curse of dimensionality



### Example

 $\label{lem:https://github.com/STIMALiU/AdvRCourse/blob/master/Code/parallel_example.R$ 

Parallel code example



The End... for today. Questions? See you next time!