

Assignment 6.

EE 10039.

Manj Uniyal (D)

Date / /

Q1. even parity checker

(i) Data 1000000, parity 0.

1 → parity error

(ii) Do 11111110

0 → parity No error

(iii) 1100111

1 → No error.

(b) full adder,

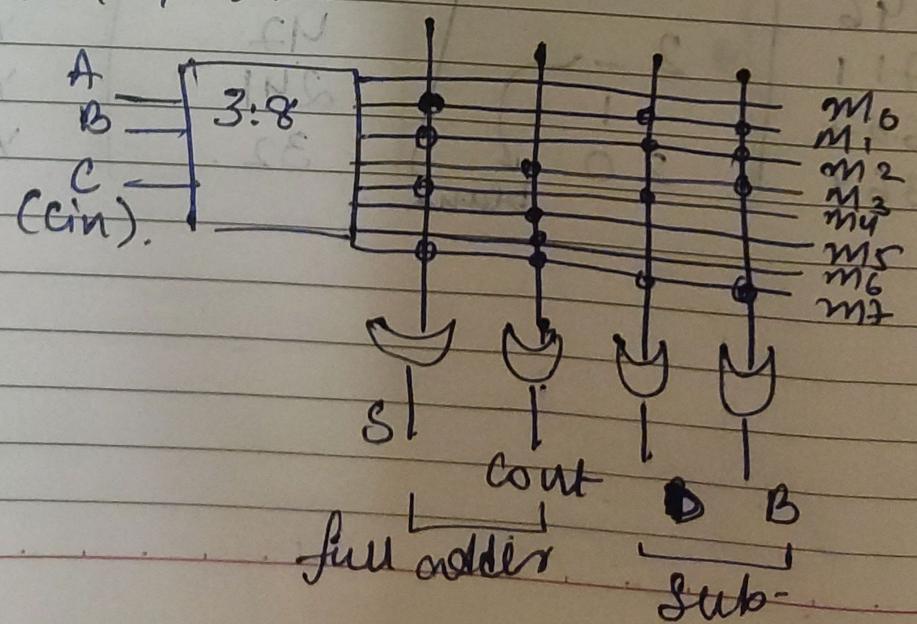
A	B	C	S	Cout	D	B.
0	0	0	0	0	0	0
0	0	1	1	0	1	1
0	1	0	1	0	1	0
0	1	1	0	1	0	1
1	0	0	1	0	1	0
1	0	1	0	1	0	0
1	1	0	0	1	0	0
1	1	1	1	1	1	1

$$Cout = \sum m(3, 5, 6, 7)$$

$$S = \sum m(1, 2, 4, 6, 7)$$

$$D = \sum m(1, 2, 4, 7)$$

$$B = \sum m(1, 2, 3, 7)$$



(2)

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- Q2.
- a) Select lines in
- $$1: 8192 \quad \text{DPMWx} = \log_{10} 8192 = 18 \quad (18)$$
- (b) $Y_2 = D_5 + D_7 = (D_5 + D_7)'$
- $$D_0 = 1, 1' 1_0' = (P+Q)'(PQ)' = (P'+Q+PQ)' = (P'+Q)'$$
- $$D_2 = 1, 1_0' = (P'+Q)(PQ')' = (P'+Q)(P'+Q') = P' + Q'Q = P'.$$
- $$Y_1 = D_0 + D_2 = PQ' + P = P' + Q'$$
- $$D_5 = 1_3' 1_2 = (P'+Q')'P = PQP' = PQ,$$
- $$D_7 = 1_3 1_2 = (P+Q')P = PQ'.$$
- $\therefore Y_2 = (D_5 + D_7)' = (PQ + PQ')' = P'$
- $\therefore Y_1 Y_2 = (P' + Q')'(P')' = ((PQ))'P' = PQP' = 0.$
- $\therefore Y_1 Y_2 = 0$
- Q3.
- (a)
-
- Inputs: $A_1, A_0, B_1, B_0, C_{in}$
- Outputs: $S_1, S_0, C_{out}, Y_1, Y_0$
- Truth table:
- | A_1 | A_0 | B_1 | B_0 | C_{in} | S_1 | S_0 | C_{out} | Y_1 | Y_0 |
|-------|-------|-------|-------|----------|-------|-------|-----------|-------|-------|
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 |
| 0 | 0 | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 0 |
| 0 | 0 | 0 | 1 | 1 | 0 | 1 | 1 | 0 | 1 |
| 0 | 1 | 0 | 0 | 0 | 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 |
| 0 | 1 | 0 | 1 | 0 | 1 | 1 | 0 | 0 | 1 |
| 0 | 1 | 0 | 1 | 1 | 0 | 0 | 1 | 0 | 1 |
| 1 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 |
| 1 | 0 | 0 | 0 | 1 | 0 | 1 | 1 | 0 | 0 |
| 1 | 0 | 0 | 1 | 0 | 0 | 1 | 0 | 1 | 0 |
| 1 | 0 | 0 | 1 | 1 | 1 | 0 | 1 | 0 | 1 |
| 1 | 1 | 0 | 0 | 0 | 1 | 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 0 | 0 |
| 1 | 1 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 1 |
| 1 | 1 | 0 | 1 | 1 | 1 | 0 | 0 | 0 | 1 |
| 1 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 0 | 1 | 0 | 1 | 1 | 1 | 1 | 0 | 0 |
| 1 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 |
| 1 | 0 | 1 | 1 | 1 | 0 | 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 |
| 1 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |
| 1 | 1 | 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 1 | 1 | 1 | 0 | 0 | 0 | 1 |

(B)

(b) 3 bit comparator:

Used 2³ combinations $\Rightarrow 2^3 = 64$ (i) # A=B at 8 combination
left over = $64 - 8 = 56$.(ii) # $A > B = 56/2 = 28$ (iii) # $A < B = 56/2 = 28$.

$$\text{Q4. } (168.65)_0 + (253.9)_0 \rightarrow (?)_{x_3-3}$$

(i)

Carry $\rightarrow 1111$

$$\begin{array}{r}
 1.111 \\
 1010 + 0101 \\
 \hline
 10000
 \end{array}
 \quad
 \begin{array}{r}
 1111 \\
 1001 + 1011 \\
 \hline
 0010
 \end{array}
 \quad
 \begin{array}{r}
 1111 \\
 0110 + 1100 \\
 \hline
 0010
 \end{array}
 \quad
 \begin{array}{r}
 1111 \\
 1001 + 1100 \\
 \hline
 0101
 \end{array}
 \quad
 \begin{array}{r}
 1111 \\
 0011 + 0011 \\
 \hline
 0000
 \end{array}
 \quad
 \begin{array}{r}
 1111 \\
 0011 + 0011 \\
 \hline
 0000
 \end{array}
 \quad
 \begin{array}{r}
 1111 \\
 0011 + 0011 \\
 \hline
 0000
 \end{array}
 \quad
 \begin{array}{r}
 1111 \\
 0011 + 0011 \\
 \hline
 0000
 \end{array}
 \quad
 \begin{array}{r}
 1111 \\
 0011 + 0011 \\
 \hline
 0000
 \end{array}$$

$$F = (4355.88)_{x_3-3}$$

$$(ii) (58)_0 - (49)_0 = (?)_{x_3-3}$$

$$\begin{array}{r}
 1000 \\
 0101 \\
 \hline
 1110
 \end{array}
 \quad
 \begin{array}{r}
 1011 \\
 1000 \\
 \hline
 0011
 \end{array}
 \quad
 \begin{array}{r}
 0011 \\
 + 0011 \\
 \hline
 0000
 \end{array}
 \quad
 \begin{array}{r}
 0011 \\
 + 0011 \\
 \hline
 0000
 \end{array}
 \quad
 \begin{array}{r}
 0011 \\
 + 0011 \\
 \hline
 0000
 \end{array}
 \quad
 \begin{array}{r}
 0011 \\
 + 0011 \\
 \hline
 0000
 \end{array}
 \quad
 \begin{array}{r}
 0011 \\
 + 0011 \\
 \hline
 0000
 \end{array}
 \quad
 \begin{array}{r}
 0011 \\
 + 0011 \\
 \hline
 0000
 \end{array}
 \quad
 \begin{array}{r}
 0011 \\
 + 0011 \\
 \hline
 0000
 \end{array}$$

$$\begin{array}{r}
 0100 \\
 1 \\
 \hline
 1001
 \end{array}
 \quad
 \begin{array}{r}
 1001 \\
 6 \\
 \hline
 0110
 \end{array}$$

$$F = (99)_{x_3-3}$$

(4)

(a) $(547)_8$.

(i) $(101100111)_2$

$$\text{grey code} = 111010100 = (764)_8.$$

(ii) $(\text{DEC})_{10} \rightarrow (10111101100)_2$

$$\text{grey code} = 101100011010$$

(b) grey code to binary

$$111100111001 \rightarrow (10100010110)_2$$

(c) ASCII

American Standard code for Information Interchange.

use 8bit for each character.

EBCDIC

Extended Binary Coded Decimal Interchange Code

use 8bit for each character.

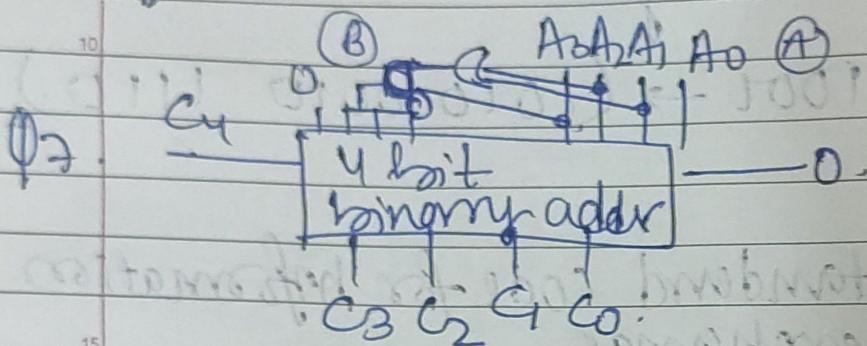
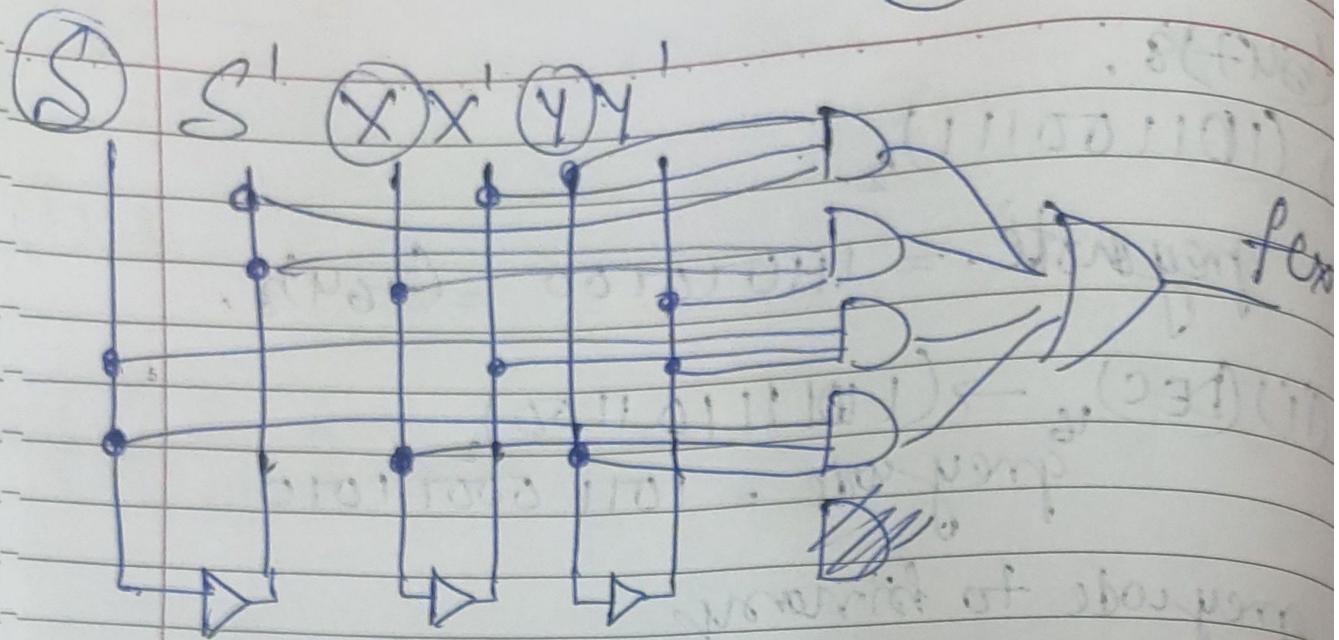
Q6. S=0, f = XOR (x,y)

S=1, f = XNOR (x,y)

S	x	y	f	S	x'y'	xy'	xy	xy'
0	0	0	0	0	1	1	0	0
0	0	1	1	1	0	0	1	0
0	1	0	1	0	0	1	0	1
0	1	1	0	1	1	1	1	1
1	0	0	1					
1	0	1	0					
1	1	0	0					
1	1	1	1					

$$f = S'x'y + S'xy' + Sx'y + Sxy$$

(5)



10. unique s/p.

$A_3 A_2 A_1 A_0$ $B_2 B_1$. $C_3 - C_0$

0 0 0 0	0	0 0 0 0
---------	---	---------

0 0 0 1	0	0 0 0 1
---------	---	---------

0 0 1 0	0	0 0 1 0
---------	---	---------

0 0 1 1	0	0 0 1 1
---------	---	---------

0 1 0 0	0	0 1 0 0
---------	---	---------

0 1 0 1	0	0 1 0 1
---------	---	---------

0 1 1 0	0	0 1 1 0
---------	---	---------

0 1 1 1	0	0 1 1 1
---------	---	---------

1 0 0 0	1	1 0 0 0
---------	---	---------

1 0 0 1	1	1 1 1 1
---------	---	---------

1 0 1 0	1	0 0 0 0
---------	---	---------

1 0 1 1	1	0 0 0 1
---------	---	---------

1 1 0 0	1	0 0 1 0
---------	---	---------

1 1 0 1	1	0 0 1 1
---------	---	---------

1 1 1 0	1	0 1 0 0
---------	---	---------

1 1 1 1	1	0 1 0 1
---------	---	---------

1 1 1 1	1	0 1 0 1
---------	---	---------

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(a).

Q8. m parity bit. n=16 data bit.

$$2^m > m+n.$$

$$2^m > m+16.$$

$$m=5 \quad 2^5 > 5+16 \\ 32 > 21.$$

So 5 parity bit are required to transfer 16 bit data.

(b) $x_1, x_0, y_1, y_0,$

$$\begin{array}{ll} p_f & x < y \\ 1 & \\ 0 & x \geq y. \end{array}$$

 $x_1, x_0, y_1, y_0, F.$

$$\begin{array}{cccc} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{array}$$

$$\begin{array}{cccc} 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 \end{array}$$

$$\begin{array}{cccc} 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 \end{array}$$

$$\begin{array}{cccc} 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 \end{array}$$

$$\begin{array}{cccc} 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 \end{array}$$

$$\begin{array}{cccc} 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 \end{array}$$

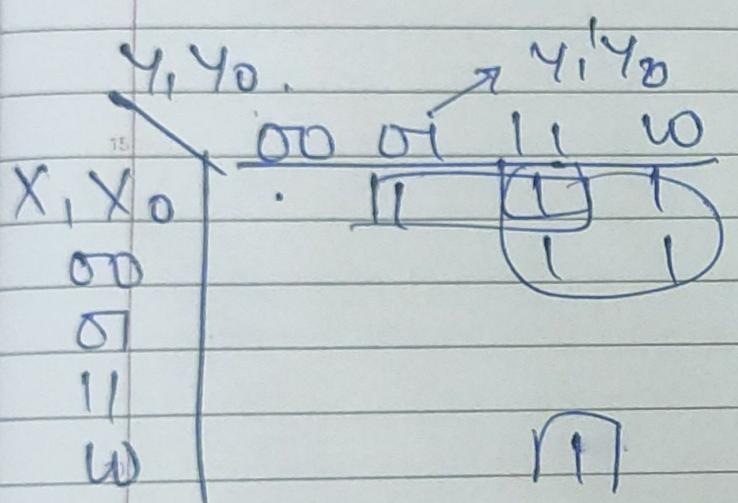
$$\begin{array}{cccc} 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 \end{array}$$

$$\begin{array}{cccc} 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 \end{array}$$

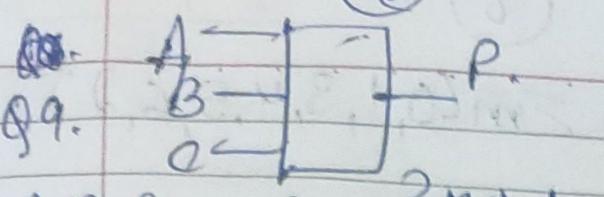
$$\begin{array}{cccc} 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 \end{array}$$

$$\begin{array}{cccc} 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 \end{array}$$

$$\begin{array}{cccc} 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 \end{array}$$



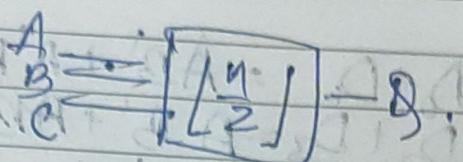
$$F = x_1' y_1 + x_1' x_0' y_0 + x_0' y_1 y_0.$$



$P = \sum m(?)$
when odd no. of i/p is high.

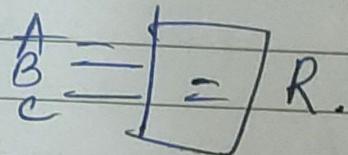
$$ABC \quad P \quad 2n+1.$$

000	0							
001	1							
010	0							
011	0							
100	1							
101	0							
110	0							
111	1							



$Q = \sum m(?)$
when either 0 i/p or 1 i/p high.

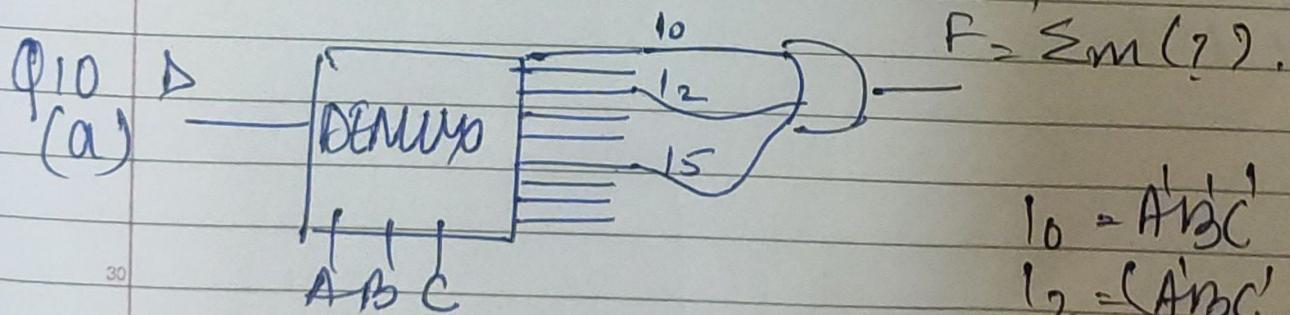
ABC	Q
000	1
001	1
010	1
011	0
100	1
101	1
110	0
111	0



$R = 0$.

gate high
when all
i/p are equal

ABC	R
000	0
001	1
010	1
011	1
100	1
101	1
110	1
111	0



$$F = \sum m(?)$$

$$I_0 = ABC'$$

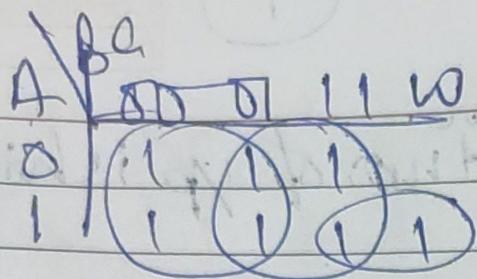
$$I_2 = (A'BC')'$$

$$I_5 = AB'C$$

$$F = AB'C + A'B'C + ABC' + AB'C = A + B + C$$

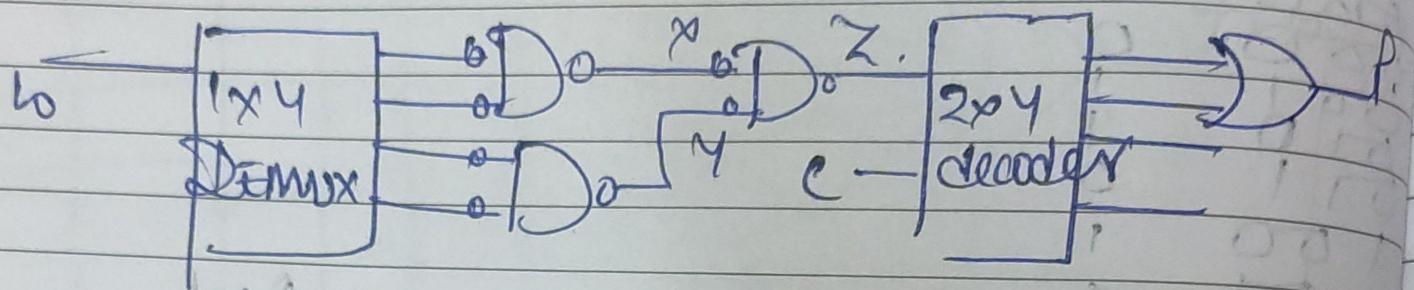
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$$f = \sum m(0, 1, 3, 4, 5, 6, 7)$$

(b)



$$X = (D_0' D_1')^l = D_0 + D_1 = AB' + A'B = A'$$

$$Y = (D_2' D_3')^l = D_2 + D_3 = A'B' + AB = A$$

$$Z = (X' Y' Z)^l = X + Y = A + A' = 1$$

$$\begin{aligned} A_0 &= 1 & A_1 &= C \\ P &= A_1 A_0' + A_1' A_0 = A' \\ \boxed{P = C'} \end{aligned}$$