

Algorithm Fourth Homework

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1 Problem 3: Interval Scheduling Problem

1.

This problem can be formulated as follows:

$$\text{maximize} \quad \sum_{i,j} y_{ij}$$

y_{ij} means in the i th classroom, whether arranging the j th class.

s.t.

$$y_{ij} = 0 \quad \text{or} \quad 1$$

$$\sum_{i=1}^m y_{ij} \leq 1 \quad j = 1, \dots, n$$

$$y_{ij} + y_{ik} \leq 1 \quad i = 1, \dots, m$$

j th job can not arrange in several classrooms.

j th and k th job can not arrange in the same classroom.

I construct an instance as follows:

$$\begin{bmatrix} j & 1 & 2 & 3 & 4 & 5 & 6 \\ S & 0 & 1 & 1 & 3 & 4 & 5 \\ F & 5 & 3 & 4 & 8 & 6 & 7 \end{bmatrix}$$

My input as follows:

```
/* Variables */
var y_11 >= 0, integer;
var y_21 >= 0, integer;
var y_12 >= 0, integer;
var y_22 >= 0, integer;
var y_13 >= 0, integer;
var y_23 >= 0, integer;
var y_14 >= 0, integer;
var y_24 >= 0, integer;
var y_15 >= 0, integer;
var y_25 >= 0, integer;
var y_16 >= 0, integer;
var y_26 >= 0, integer;
```

```

/*var x >= 0, integer;
var y >= 0, integer;
var z >= 0, integer;*/

/* Object function */
maximize q: y_11 + y_12 + y_21 + y_22 + y_13 + y_23 +y_14+ y_24 + y_15 + y_25;
/*minimize q: -2*x + -3*y + -2*z; */

/* Constraints */
s.t. con1: y_11 + y_21 <= 1;
s.t. con2: y_12 + y_22 <= 1;
s.t. con3: y_13 + y_23 <= 1;
s.t. con4: y_14 + y_24 <= 1;
s.t. con5: y_15 + y_25 <= 1;
s.t. con6: y_16 + y_26 <= 1;
s.t. con7: y_11 + y_12 <= 1;
s.t. con8: y_21 + y_22 <= 1;
s.t. con9: y_11 + y_13 <= 1;
s.t. con10: y_21 + y_23 <= 1;
s.t. con11: y_11 + y_14 <= 1;
s.t. con12: y_21 + y_24 <= 1;
s.t. con13: y_11 + y_15 <= 1;
s.t. con14: y_21 + y_25 <= 1;
s.t. con15: y_12 + y_13 <= 1;
s.t. con16: y_22 + y_23 <= 1;
s.t. con17: y_13 + y_14 <= 1;
s.t. con18: y_23 + y_24 <= 1;
s.t. con19: y_14 + y_15 <= 1;
s.t. con20: y_24 + y_25 <= 1;
s.t. con21: y_14 + y_16 <= 1;
s.t. con22: y_24 + y_26 <= 1;
s.t. con23: y_15 + y_16 <= 1;
s.t. con24: y_25 + y_26 <= 1;
/*s.t. con1: -2*x + -y + -z >= -4;
s.t. con2: -x + -2*y + -z >= -7;
s.t. con3: -z >= -5; */

end;

```

The output solved by GLPK as follows:

```

Problem:    glpsolEx
Rows:       25
Columns:    12 (12 integer, 0 binary)
Non-zeros:  60
Status:     INTEGER OPTIMAL
Objective:  q = 4 (MAXimum)

```

No.	Row name	Activity	Lower bound	Upper bound
-----	----------	----------	-------------	-------------

1	q	4	
2	con1	1	1
3	con2	1	1
4	con3	0	1
5	con4	1	1
6	con5	0	1
7	con6	1	1
8	con7	1	1
9	con8	1	1
10	con9	0	1
11	con10	1	1
12	con11	1	1
13	con12	1	1
14	con13	0	1
15	con14	1	1
16	con15	1	1
17	con16	0	1
18	con17	1	1
19	con18	0	1
20	con19	1	1
21	con20	0	1
22	con21	1	1
23	con22	1	1
24	con23	0	1
25	con24	1	1

No.	Column name	Activity	Lower bound	Upper bound
1	y_11	*	0	0
2	y_21	*	1	0
3	y_12	*	1	0
4	y_22	*	0	0
5	y_13	*	0	0
6	y_23	*	0	0
7	y_14	*	1	0
8	y_24	*	0	0
9	y_15	*	0	0
10	y_25	*	0	0
11	y_16	*	0	0
12	y_26	*	1	0

Integer feasibility conditions:

KKT.PE: max.abs.err = 0.00e+00 on row 0
max.rel.err = 0.00e+00 on row 0
High quality

```
KKT.PB: max.abs.err = 0.00e+00 on row 0
        max.rel.err = 0.00e+00 on row 0
        High quality
```

End of output

We can see the optimal solution is 4, which arrange class 1 and class 6 on room 2, and arrange class 2 and class 4 on room 1. After analyzing the instance, we know the result is right, but we can also find other suitable arrangements.

2.

No, if we change the requirement above, set $s_6 = 10$ and $f_6 = 12$, apparently we will get answer is 5, we can get the answer as follows:

```
Problem:      glpsolEx
Rows:         21
Columns:      12
Non-zeros:    52
Status:       OPTIMAL
Objective:    q = 6 (MAXimum)
```

No.	Row name	St	Activity	Lower bound	Upper bound	Marginal
1	q	B	6			
2	con1	NU	1		1	
3	con2	NU	1		1	
4	con3	NU	1		1	
5	con4	NU	1		1	<
6	con5	B	1		1	
7	con6	NU	1		1	
8	con7	NU	1		1	<
9	con8	B	1		1	
10	con9	NU	1		1	<
11	con10	B	1		1	
12	con11	B	1		1	
13	con12	B	1		1	
14	con13	NU	1		1	<
15	con14	B	1		1	
16	con15	B	1		1	
17	con16	NU	1		1	<
18	con17	B	1		1	
19	con18	B	1		1	
20	con19	NU	1		1	
21	con20	NU	1		1	

No.	Column name	St	Activity	Lower bound	Upper bound	Marginal
1	y_11	B	0.5	0		
2	y_21	B	0.5	0		
3	y_12	B	0.5	0		

4	y_22	B	0.5	0
5	y_13	B	0.5	0
6	y_23	B	0.5	0
7	y_14	B	0.5	0
8	y_24	B	0.5	0
9	y_15	B	0.5	0
10	y_25	B	0.5	0
11	y_16	B	1	0
12	y_26	NL	0	0

Karush-Kuhn-Tucker optimality conditions:

KKT.PE: max.abs.err = 0.00e+00 on row 0
max.rel.err = 0.00e+00 on row 0
High quality

KKT.PB: max.abs.err = 0.00e+00 on row 0
max.rel.err = 0.00e+00 on row 0
High quality

KKT.DE: max.abs.err = 0.00e+00 on column 0
max.rel.err = 0.00e+00 on column 0
High quality

KKT.DB: max.abs.err = 0.00e+00 on row 0
max.rel.err = 0.00e+00 on row 0
High quality

End of output

But in ILP, the answer as follows:

Problem: glpsolEx
Rows: 21
Columns: 12 (12 integer, 0 binary)
Non-zeros: 52
Status: INTEGER OPTIMAL
Objective: q = 5 (MAXimum)

No.	Row name	Activity	Lower bound	Upper bound
1	q	5		
2	con1	0		1
3	con2	1		1
4	con3	1		1
5	con4	1		1
6	con5	1		1
7	con6	1		1
8	con7	1		1

9	con8	0	1
10	con9	0	1
11	con10	1	1
12	con11	1	1
13	con12	0	1
14	con13	0	1
15	con14	1	1
16	con15	1	1
17	con16	1	1
18	con17	1	1
19	con18	1	1
20	con19	1	1
21	con20	1	1

No.	Column name		Activity	Lower bound	Upper bound

1	y_11	*	0	0	
2	y_21	*	0	0	
3	y_12	*	1	0	
4	y_22	*	0	0	
5	y_13	*	0	0	
6	y_23	*	1	0	
7	y_14	*	1	0	
8	y_24	*	0	0	
9	y_15	*	0	0	
10	y_25	*	1	0	
11	y_16	*	0	0	
12	y_26	*	1	0	

Integer feasibility conditions:

KKT.PE: max.abs.err = 0.00e+00 on row 0
max.rel.err = 0.00e+00 on row 0
High quality

KKT.PB: max.abs.err = 0.00e+00 on row 0
max.rel.err = 0.00e+00 on row 0
High quality

End of output

According to lemma, we can not always set each column of A has at most one +1 entry and at most one -1 entry.

2 Problem 5: Stable Matching Problem

I select condition 1 as my condition, the problem can be formulated as follows:

$$\begin{aligned}
& \text{maximize} \quad \sum_{i,j,l,k} y_{i,j,l,k} \\
& \text{s.t.} \quad \sum_{i,l \in p} y_{i,j(i),l,k(l)} S_{i,j(i),l,k(l)} = 0 \quad p \in P
\end{aligned}$$

P means permutations including all pairs between man and woman, we can conclude the number of permutations is $(n!)^2$. So we wanna find a permutation can meet the constraint and in the mean time, its optimal function can also be max of all p in P. In an instance p of P, pairs are fixed, all we need to do is to meet the constraints and get the maximum.

I give an instance as follows:

$$\begin{bmatrix} 0 & 1 & 2 & 3 \\ A & 3 & 2 & 1 \\ B & 2 & 1 & 3 \\ C & 3 & 1 & 2 \end{bmatrix}$$

$$\begin{bmatrix} 0 & A & B & C \\ 1 & C & B & A \\ 2 & A & B & C \\ 3 & C & B & A \end{bmatrix}$$

My input as follows:

```

/* Variables */
var y_1_B_2_A >= 0, binary;
var y_1_A_2_B >= 0, binary;
var y_1_C_2_B >= 0, binary;
var y_1_B_2_C >= 0, binary;
var y_1_A_3_C >= 0, binary;
var y_1_C_3_A >= 0, binary;
var y_1_C_2_A >= 0, binary;
var y_1_A_2_C >= 0, binary;
var y_1_B_3_A >= 0, binary;
var y_1_A_3_B >= 0, binary;
var y_1_B_3_C >= 0, binary;
var y_1_C_3_B >= 0, binary;
var y_2_B_3_A >= 0, binary;
var y_2_A_3_B >= 0, binary;
var y_2_A_3_C >= 0, binary;
var y_2_C_3_A >= 0, binary;
var y_2_B_3_C >= 0, binary;
var y_2_C_3_B >= 0, binary;

/* Object function */
/*maximize q: y_1_A_2_B + y_1_A_3_C + y_2_B_3_C;*/
/*maximize r: y_1_A_2_C + y_1_A_3_B + y_2_C_3_B;*/
maximize s: y_1_B_2_A + y_1_B_3_C + y_2_A_3_C;
/*maximize t: y_1_B_2_C + y_1_B_3_A + y_2_C_3_A;*/
/*maximize u: y_1_C_2_A + y_1_C_3_B + y_2_A_3_B;*/

```

```

/*maximize v: y_1_C_2_B + y_1_C_3_A + y_2_B_3_A;*/

/* Constraints */
s.t. con1: y_1_A_2_B * 1 + y_1_A_3_C * 0 + y_2_B_3_C * 0 = 0;
s.t. con2: y_1_A_2_C * 1 + y_1_A_3_B * 1 + y_2_C_3_B * 1 = 0;
s.t. con3: y_1_B_2_A * 0 + y_1_B_3_C * 0 + y_2_A_3_C * 0 = 0;
s.t. con4: y_1_B_2_C * 1 + y_1_B_3_A * 0 + y_2_C_3_A * 1 = 0;
s.t. con5: y_1_C_2_A * 0 + y_1_C_3_B * 1 + y_2_A_3_B * 1 = 0;
s.t. con6: y_1_C_2_B * 0 + y_1_C_3_A * 1 + y_2_B_3_A * 0 = 0;
/*s.t. con1: S_1_B_2_A = 0;
s.t. con2: S_1_A_2_B = 1;
s.t. con3: S_1_C_2_B = 0;
s.t. con4: S_1_B_2_C = 1;
s.t. con5: S_1_A_3_C = 0;
s.t. con6: S_1_C_3_A = 1;
s.t. con7: S_1_C_2_A = 0;
s.t. con8: S_1_A_2_C = 1;
s.t. con9: S_1_B_3_A = 0;
s.t. con10: S_1_A_3_B = 1;
s.t. con11: S_1_B_3_C = 0;
s.t. con12: S_1_C_3_B = 1;
s.t. con13: S_2_B_3_A = 0;
s.t. con14: S_2_A_3_B = 1;
s.t. con15: S_2_A_3_C = 0;
s.t. con16: S_2_C_3_A = 1;
s.t. con17: S_2_B_3_C = 0;
s.t. con18: S_2_C_3_B = 1;*/

end;

```

My output as follows:

```

Problem:      StableMatch
Rows:         7
Columns:      12 (12 integer, 12 binary)
Non-zeros:    12
Status:       INTEGER OPTIMAL
Objective:    s = 3 (MAXimum)

```

No.	Row name	Activity	Lower bound	Upper bound
1	s	3		
2	con1	0	-0	=
3	con2	0	-0	=
4	con3	0	-0	=
5	con4	0	-0	=
6	con5	0	-0	=
7	con6	0	-0	=

No.	Column name		Activity	Lower bound	Upper bound
1	y_1_B_2_A	*	1	0	1
2	y_1_A_2_B	*	0	0	1
3	y_1_B_2_C	*	0	0	1
4	y_1_C_3_A	*	0	0	1
5	y_1_A_2_C	*	0	0	1
6	y_1_A_3_B	*	0	0	1
7	y_1_B_3_C	*	1	0	1
8	y_1_C_3_B	*	0	0	1
9	y_2_A_3_B	*	0	0	1
10	y_2_A_3_C	*	1	0	1
11	y_2_C_3_A	*	0	0	1
12	y_2_C_3_B	*	0	0	1

Integer feasibility conditions:

```
KKT.PE: max.abs.err = 0.00e+00 on row 0
        max.rel.err = 0.00e+00 on row 0
        High quality
```

```
KKT.PB: max.abs.err = 0.00e+00 on row 0
        max.rel.err = 0.00e+00 on row 0
        High quality
```

End of output

So we can see that, the answer means 1 and B, 2 and A, 3 and C are the stable pairs, we can verify it.

3 Problem 7: Simplex Algorithm

I use the big-M method to get the initial situation, and use simplex algorithm to solve the problem.

```
import numpy as np
```

```
class Simplex_Table:
```

```
    def __init__(self, obj):
        self.obj=obj
        self.rows=[]
        self.columns=[]
```

```
    def addConstraint(self, expression, value):
        self.rows.append(expression)
        self.columns.append(value)
```

```

def check(self):
    if max(self.obj[0:-1])<= 0: return 1
    else:return 0

def display(self):
    print '\n',np.matrix([self.obj]+self.rows)

def select_column(self):
    high=0
    index=[]
    for i in range(0,len(self.obj)-1):
        if self.obj[i]>high:
            index.append(i)
    if index==[]: return -1
    return index[0]

def select_row(self,col):
    right=[self.rows[i][-1] for i in range(len(self.rows))]
    left=[self.rows[i][col] for i in range(len(self.rows))]
    ratio=[]
    for i in range(len(right)):
        if left[i]==0:
            ratio.append(999*abs(max(right)))
            continue
        ratio.append(right[i]/left[i])
    if ratio[np.argmin(ratio)]<0:
        return -1
    return np.argmin(ratio)

def pivot(self,row,col):
    e=self.rows[row][col]
    self.rows[row]=self.rows[row]/e
    for r in range(len(self.rows)):
        if r==row:continue
        self.rows[r]=self.rows[r]-self.rows[r][col]*self.rows[row]
    self.obj-=self.obj[col]*self.rows[row]

def solve(self):
    # build a simplex table
    for i in range(len(self.rows)):
        self.obj+=[-99999]
        col=[0 for r in range(len(self.rows))]
        col[i]=-1
        self.rows[i]+=col+[self.columns[i]]
        self.rows[i]=np.array(self.rows[i],dtype=float)
    self.obj=np.array(self.obj+[0],dtype=float)
    self.display()

```

```

checkNum=0
value_value=0
while not self.check():
    checkNum+=1
    if checkNum>9999:
        print "No answer!"
        break
    columns=self.select_column()
    rows=self.select_row(columns)
    if rows==-1:
        print "No optimal answer!"
        break
    self.pivot(rows,columns)
    print '\nPivot column: %s\n Pivot row: %s'%(columns+1,rows+1)
    self.display()

if __name__=='__main__':
    """
        In the program, we use the big-M method to set the initial situation
        min z=-2x+-3y+-2z
        s.t.
        -2x+-y+-z-a=-4
        -x+-2y+-z-b=-7
        -z-c      =-5
        x,y,z,a,b,c>=0
    """
    t=Simplex_Table([2,3,2,0,0,0])
    t.addConstraint([-2,-1,-1,-1,0,0],-4)
    t.addConstraint([-1,-2,-1,0,-1,0],-7)
    t.addConstraint([0,0,-1,0,0,-1],-5)
    t.solve()

```

Compared the GLPK, we use two program solving the same problem, the answer shows below:
GLPK:

```

Problem:      glpsolEx
Rows:         4
Columns:      3
Non-zeros:    10
Status:       OPTIMAL
Objective:    q = -11 (MINimum)

```

No.	Row name	St	Activity	Lower bound	Upper bound	Marginal
1	q	B	-11			
2	con1	NL	-4	-4		
3	con2	NL	-7	-7		
4	con3	B	-1	-5		

No.	Column name	St	Activity	Lower bound	Upper bound	Marginal
1	x	NL	0	0		
2	y	B	3	0		
3	z	B	1	0		

Karush-Kuhn-Tucker optimality conditions:

KKT.PE: max.abs.err = 0.00e+00 on row 0
max.rel.err = 0.00e+00 on row 0
High quality

KKT.PB: max.abs.err = 0.00e+00 on row 0
max.rel.err = 0.00e+00 on row 0
High quality

KKT.DE: max.abs.err = 0.00e+00 on column 0
max.rel.err = 0.00e+00 on column 0
High quality

KKT.DB: max.abs.err = 0.00e+00 on row 0
max.rel.err = 0.00e+00 on row 0
High quality

End of output

My Program:

```
[[ 2.00000000e+00  3.00000000e+00  2.00000000e+00  0.00000000e+00
  0.00000000e+00  0.00000000e+00 -9.99990000e+04 -9.99990000e+04
 -9.99990000e+04  0.00000000e+00]
 [-2.00000000e+00 -1.00000000e+00 -1.00000000e+00 -1.00000000e+00
  0.00000000e+00  0.00000000e+00 -1.00000000e+00  0.00000000e+00
  0.00000000e+00 -4.00000000e+00]
 [-1.00000000e+00 -2.00000000e+00 -1.00000000e+00  0.00000000e+00
 -1.00000000e+00  0.00000000e+00  0.00000000e+00 -1.00000000e+00
  0.00000000e+00 -7.00000000e+00]
 [ 0.00000000e+00  0.00000000e+00 -1.00000000e+00  0.00000000e+00
  0.00000000e+00 -1.00000000e+00  0.00000000e+00  0.00000000e+00
 -1.00000000e+00 -5.00000000e+00]]
```

Pivot column: 1

Pivot row: 1

```
[[ 0.00000000e+00  2.00000000e+00  1.00000000e+00 -1.00000000e+00
  0.00000000e+00  0.00000000e+00 -1.00000000e+05 -9.99990000e+04
```

```

-9.99990000e+04 -4.00000000e+00]
[ 1.00000000e+00 5.00000000e-01 5.00000000e-01 5.00000000e-01
-0.00000000e+00 -0.00000000e+00 5.00000000e-01 -0.00000000e+00
-0.00000000e+00 2.00000000e+00]
[ 0.00000000e+00 -1.50000000e+00 -5.00000000e-01 5.00000000e-01
-1.00000000e+00 0.00000000e+00 5.00000000e-01 -1.00000000e+00
0.00000000e+00 -5.00000000e+00]
[ 0.00000000e+00 0.00000000e+00 -1.00000000e+00 0.00000000e+00
0.00000000e+00 -1.00000000e+00 0.00000000e+00 0.00000000e+00
-1.00000000e+00 -5.00000000e+00]]

```

Pivot column: 2

Pivot row: 2

```

[[ 0.00000000e+00 0.00000000e+00 3.33333333e-01 -3.33333333e-01
-1.33333333e+00 0.00000000e+00 -9.99993333e+04 -1.00000333e+05
-9.99990000e+04 -1.06666667e+01]
[ 1.00000000e+00 0.00000000e+00 3.33333333e-01 6.66666667e-01
-3.33333333e-01 0.00000000e+00 6.66666667e-01 -3.33333333e-01
0.00000000e+00 3.33333333e-01]
[ -0.00000000e+00 1.00000000e+00 3.33333333e-01 -3.33333333e-01
6.66666667e-01 -0.00000000e+00 -3.33333333e-01 6.66666667e-01
-0.00000000e+00 3.33333333e+00]
[ 0.00000000e+00 0.00000000e+00 -1.00000000e+00 0.00000000e+00
0.00000000e+00 -1.00000000e+00 0.00000000e+00 0.00000000e+00
-1.00000000e+00 -5.00000000e+00]]

```

Pivot column: 3

Pivot row: 1

```

[[ -1.00000000e+00 0.00000000e+00 0.00000000e+00 -1.00000000e+00
-1.00000000e+00 0.00000000e+00 -1.00000000e+05 -1.00000000e+05
-9.99990000e+04 -1.10000000e+01]
[ 3.00000000e+00 0.00000000e+00 1.00000000e+00 2.00000000e+00
-1.00000000e+00 0.00000000e+00 2.00000000e+00 -1.00000000e+00
0.00000000e+00 1.00000000e+00]
[ -1.00000000e+00 1.00000000e+00 0.00000000e+00 -1.00000000e+00
1.00000000e+00 -0.00000000e+00 -1.00000000e+00 1.00000000e+00
-0.00000000e+00 3.00000000e+00]
[ 3.00000000e+00 0.00000000e+00 0.00000000e+00 2.00000000e+00
-1.00000000e+00 -1.00000000e+00 2.00000000e+00 -1.00000000e+00
-1.00000000e+00 -4.00000000e+00]]

```

We can see the results above, the answers are same, after testing several cases, I verify the fact that my program is correct, but wastes more time due to python's high time cost.