

A Generalization of the Cauchy–Davenport theorem

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November 1, 2024

Abstract

The Cauchy–Davenport theorem is a fundamental result in additive combinatorics. It was originally independently discovered by Cauchy [2] and Davenport [3] and has been formalized in the AFP entry [1] as a corollary of Kneser’s theorem. More recently, many generalizations of this theorem have been found. In this entry, we formalise a generalization due to DeVos [4], which proves the theorem in a non-abelian setting.

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1 Preliminaries on well-orderings, groups, and sum-sets

theory *Generalized-Cauchy-Davenport-preliminaries*
imports
Complex-Main
Jacobson-Basic-Algebra.Group-Theory

begin

1.1 Well-ordering lemmas

lemma *wf-prod-lex-fibers-inter:*

fixes $r :: ('a \times 'a) \text{ set}$ **and** $s :: ('b \times 'b) \text{ set}$ **and** $f :: 'c \Rightarrow 'a$ **and** $g :: 'c \Rightarrow 'b$
and
 $t :: ('c \times 'c) \text{ set}$
assumes $h1: wf ((inv\text{-}image\ r\ f) \cap t)$ **and**
 $h2: \bigwedge a. a \in range\ f \implies wf (\{x. f\ x = a\} \times \{x. f\ x = a\} \cap (inv\text{-}image\ s\ g))$
 $\cap t)$ **and**
 $h3: trans\ t$
shows $wf ((inv\text{-}image\ (r\ <*\text{lex}*\>\ s)\ (\lambda\ c. (f\ c,\ g\ c))) \cap t)$
 $\langle proof \rangle$

lemma *wf-prod-lex-fibers:*

fixes $r :: ('a \times 'a) \text{ set}$ **and** $s :: ('b \times 'b) \text{ set}$ **and** $f :: 'c \Rightarrow 'a$ **and** $g :: 'c \Rightarrow 'b$
assumes $h1: wf (inv\text{-}image\ r\ f)$ **and**
 $h2: \bigwedge a. a \in range\ f \implies wf (\{x. f\ x = a\} \times \{x. f\ x = a\} \cap (inv\text{-}image\ s\ g))$
shows $wf (inv\text{-}image\ (r\ <*\text{lex}*\>\ s)\ (\lambda\ c. (f\ c,\ g\ c)))$
 $\langle proof \rangle$

context *monoid*

begin

1.2 Pointwise set multiplication in a monoid: definition and key lemmas

inductive-set $smul :: 'a \text{ set} \Rightarrow 'a \text{ set} \Rightarrow 'a \text{ set}$ **for** $A\ B$

where

$smulI[intro]: \llbracket a \in A; a \in M; b \in B; b \in M \rrbracket \implies a \cdot b \in smul\ A\ B$

syntax $smul :: 'a \text{ set} \Rightarrow 'a \text{ set} \Rightarrow 'a \text{ set}$ $((-\ \cdots -))$

lemma *smul-eq:* $smul\ A\ B = \{c. \exists a \in A \cap M. \exists b \in B \cap M. c = a \cdot b\}$
 $\langle proof \rangle$

lemma *smul:* $smul\ A\ B = (\bigcup a \in A \cap M. \bigcup b \in B \cap M. \{a \cdot b\})$
 $\langle proof \rangle$

lemma *smul-subset-carrier*: $smul\ A\ B \subseteq M$
 $\langle proof \rangle$

lemma *smul-Int-carrier* [simp]: $smul\ A\ B \cap M = smul\ A\ B$
 $\langle proof \rangle$

lemma *smul-mono*: $\llbracket A' \subseteq A; B' \subseteq B \rrbracket \implies smul\ A'\ B' \subseteq smul\ A\ B$
 $\langle proof \rangle$

lemma *smul-insert1*: $NO-MATCH\ \{\} \ A \implies smul\ (insert\ x\ A)\ B = smul\ \{x\}\ B \cup smul\ A\ B$
 $\langle proof \rangle$

lemma *smul-insert2*: $NO-MATCH\ \{\} \ B \implies smul\ A\ (insert\ x\ B) = smul\ A\ \{x\} \cup smul\ A\ B$
 $\langle proof \rangle$

lemma *smul-subset-Un1*: $smul\ (A \cup A')\ B = smul\ A\ B \cup smul\ A'\ B$
 $\langle proof \rangle$

lemma *smul-subset-Un2*: $smul\ A\ (B \cup B') = smul\ A\ B \cup smul\ A\ B'$
 $\langle proof \rangle$

lemma *smul-subset-Union1*: $smul\ (\bigcup A)\ B = (\bigcup a \in A. smul\ a\ B)$
 $\langle proof \rangle$

lemma *smul-subset-Union2*: $smul\ A\ (\bigcup B) = (\bigcup b \in B. smul\ A\ b)$
 $\langle proof \rangle$

lemma *smul-subset-insert*: $smul\ A\ B \subseteq smul\ A\ (insert\ x\ B) \wedge smul\ A\ B \subseteq smul\ (insert\ x\ A)\ B$
 $\langle proof \rangle$

lemma *smul-subset-Un*: $smul\ A\ B \subseteq smul\ A\ (B \cup C) \wedge smul\ A\ B \subseteq smul\ (A \cup C)\ B$
 $\langle proof \rangle$

lemma *smul-empty* [simp]: $smul\ A\ \{\} = \{\} \wedge smul\ \{\} \ A = \{\}$
 $\langle proof \rangle$

lemma *smul-empty'*:
assumes $A \cap M = \{\}$
shows $smul\ B\ A = \{\} \wedge smul\ A\ B = \{\}$
 $\langle proof \rangle$

lemma *smul-is-empty-iff* [simp]: $smul\ A\ B = \{\} \longleftrightarrow A \cap M = \{\} \vee B \cap M = \{\}$
 $\langle proof \rangle$

lemma *smul-D* [simp]: $smul\ A\ \{\mathbf{1}\} = A \cap M \wedge smul\ \{\mathbf{1}\}\ A = A \cap M$

$\langle \text{proof} \rangle$

lemma *smul-Int-carrier-eq* [simp]: $\text{smul } A (B \cap M) = \text{smul } A B \text{ smul } (A \cap M) B$
 $= \text{smul } A B$
 $\langle \text{proof} \rangle$

lemma *smul-assoc*:
shows $\text{smul } (\text{smul } A B) C = \text{smul } A (\text{smul } B C)$
 $\langle \text{proof} \rangle$

lemma *finite-smul*:
assumes *finite* A *finite* B **shows** *finite* $(\text{smul } A B)$
 $\langle \text{proof} \rangle$

lemma *finite-smul'*:
assumes *finite* $(A \cap M)$ *finite* $(B \cap M)$
shows *finite* $(\text{smul } A B)$
 $\langle \text{proof} \rangle$

1.3 Exponentiation in a monoid: definitions and lemmas

primrec *power* :: 'a \Rightarrow nat \Rightarrow 'a (infix \wedge 100)
where
power0: $\text{power } g \ 0 = 1$
| *power-suc*: $\text{power } g \ (\text{Suc } n) = \text{power } g \ n \cdot g$

lemma *power-one*:
assumes $g \in M$
shows $\text{power } g \ 1 = g$ $\langle \text{proof} \rangle$

lemma *power-mem-carrier*:
fixes n
assumes $g \in M$
shows $g \wedge n \in M$
 $\langle \text{proof} \rangle$

lemma *power-mult*:
assumes $g \in M$
shows $g \wedge n \cdot g \wedge m = g \wedge (n + m)$
 $\langle \text{proof} \rangle$

lemma *mult-inverse-power*:
assumes $g \in M$ **and** *invertible* g
shows $g \wedge n \cdot ((\text{inverse } g) \wedge n) = 1$
 $\langle \text{proof} \rangle$

lemma *inverse-mult-power*:
assumes $g \in M$ **and** *invertible* g
shows $((\text{inverse } g) \wedge n) \cdot g \wedge n = 1$ $\langle \text{proof} \rangle$

lemma *inverse-mult-power-eq*:

assumes $g \in M$ **and** *invertible* g

shows $\text{inverse } (g \wedge n) = (\text{inverse } g) \wedge n$

<proof>

definition *power-int* :: $'a \Rightarrow \text{int} \Rightarrow 'a$ (**infixr** *powi* 80) **where**

$\text{power-int } g \ n = (\text{if } n \geq 0 \text{ then } g \wedge (\text{nat } n) \text{ else } (\text{inverse } g) \wedge (\text{nat } (-n)))$

definition *nat-powers* :: $'a \Rightarrow 'a \text{ set}$ **where** $\text{nat-powers } g = ((\lambda n. g \wedge n) \text{ ` UNIV})$

lemma *nat-powers-eq-Union*: $\text{nat-powers } g = (\bigcup n. \{g \wedge n\})$ *<proof>*

definition *powers* :: $'a \Rightarrow 'a \text{ set}$ **where** $\text{powers } g = ((\lambda n. g \text{ powi } n) \text{ ` UNIV})$

lemma *nat-powers-subset*:

$\text{nat-powers } g \subseteq \text{powers } g$

<proof>

lemma *inverse-nat-powers-subset*:

$\text{nat-powers } (\text{inverse } g) \subseteq \text{powers } g$

<proof>

lemma *powers-eq-union-nat-powers*:

$\text{powers } g = \text{nat-powers } g \cup \text{nat-powers } (\text{inverse } g)$

<proof>

lemma *one-mem-nat-powers*: $\mathbf{1} \in \text{nat-powers } g$

<proof>

lemma *nat-powers-subset-carrier*:

assumes $g \in M$

shows $\text{nat-powers } g \subseteq M$

<proof>

lemma *nat-powers-mult-closed*:

assumes $g \in M$

shows $\bigwedge x \ y. x \in \text{nat-powers } g \implies y \in \text{nat-powers } g \implies x \cdot y \in \text{nat-powers } g$

<proof>

lemma *nat-powers-inv-mult*:

assumes $g \in M$ **and** *invertible* g

shows $\bigwedge x \ y. x \in \text{nat-powers } g \implies y \in \text{nat-powers } (\text{inverse } g) \implies x \cdot y \in$

$\text{powers } g$

<proof>

lemma *inv-nat-powers-mult*:

assumes $g \in M$ **and** *invertible* g

shows $\bigwedge x \ y. x \in \text{nat-powers } (\text{inverse } g) \implies y \in \text{nat-powers } g \implies x \cdot y \in$

powers g
⟨proof⟩

lemma *powers-mult-closed*:
assumes $g \in M$ and invertible g
shows $\bigwedge x y. x \in \text{powers } g \implies y \in \text{powers } g \implies x \cdot y \in \text{powers } g$
⟨proof⟩

lemma *nat-powers-submonoid*:
assumes $g \in M$
shows *submonoid* (*nat-powers g*) $M (\cdot) \mathbf{1}$
⟨proof⟩

lemma *nat-powers-monoid*:
assumes $g \in M$
shows *monoid* (*nat-powers g*) $(\cdot) \mathbf{1}$
⟨proof⟩

lemma *powers-submonoid*:
assumes $g \in M$ and invertible g
shows *submonoid* (*powers g*) $M (\cdot) \mathbf{1}$
⟨proof⟩

lemma *powers-monoid*:
assumes $g \in M$ and invertible g
shows *monoid* (*powers g*) $(\cdot) \mathbf{1}$
⟨proof⟩

lemma *mem-nat-powers-invertible*:
assumes $g \in M$ and invertible g and $u \in \text{nat-powers } g$
shows *monoid.invertible* (*powers g*) $(\cdot) \mathbf{1} u$
⟨proof⟩

lemma *mem-nat-inv-powers-invertible*:
assumes $g \in M$ and invertible g and $u \in \text{nat-powers } (\text{inverse } g)$
shows *monoid.invertible* (*powers g*) $(\cdot) \mathbf{1} u$
⟨proof⟩

lemma *powers-group*:
assumes $g \in M$ and invertible g
shows *group* (*powers g*) $(\cdot) \mathbf{1}$
⟨proof⟩

lemma *nat-powers-ne-one*:
assumes $g \in M$ and $g \neq \mathbf{1}$
shows $\text{nat-powers } g \neq \{\mathbf{1}\}$
⟨proof⟩

lemma *powers-ne-one*:

assumes $g \in M$ and $g \neq 1$
 shows powers $g \neq \{1\}$ $\langle proof \rangle$

end

context *group*

begin

lemma *powers-subgroup*:
 assumes $g \in G$
 shows subgroup (powers g) G (\cdot) 1
 $\langle proof \rangle$

end

context *monoid*

begin

1.4 Definition of the order of an element in a monoid

definition *order*

where order $g = (\text{if } (\exists n. n > 0 \wedge g^n = 1) \text{ then } \text{Min } \{n. g^n = 1 \wedge n > 0\} \text{ else } 0)$

definition *min-order* where $\text{min-order} = \text{Min } ((\text{order } ' M) - \{0\})$

end

1.5 Sumset scalar multiplication cardinality lemmas

context *group*

begin

lemma *card-smul-singleton-right-eq*:
 assumes finite A shows $\text{card } (\text{smul } A \{a\}) = (\text{if } a \in G \text{ then } \text{card } (A \cap G) \text{ else } 0)$
 $\langle proof \rangle$

lemma *card-smul-singleton-left-eq*:
 assumes finite A shows $\text{card } (\text{smul } \{a\} A) = (\text{if } a \in G \text{ then } \text{card } (A \cap G) \text{ else } 0)$
 $\langle proof \rangle$

lemma *card-smul-sing-right-le*:
 assumes finite A shows $\text{card } (\text{smul } A \{a\}) \leq \text{card } A$
 $\langle proof \rangle$

lemma *card-smul-sing-left-le*:
 assumes *finite A* **shows** $\text{card } (\text{smul } \{a\} A) \leq \text{card } A$
 $\langle \text{proof} \rangle$

lemma *card-le-smul-right*:
 assumes *A: finite A* $a \in A$ $a \in G$
 and *B: finite B* $B \subseteq G$
 shows $\text{card } B \leq \text{card } (\text{smul } A B)$
 $\langle \text{proof} \rangle$

lemma *card-le-smul-left*:
 assumes *A: finite A* $b \in B$ $b \in G$
 and *B: finite B* $A \subseteq G$
 shows $\text{card } A \leq \text{card } (\text{smul } A B)$
 $\langle \text{proof} \rangle$

lemma *infinite-smul-right*:
 assumes $A \cap G \neq \{\}$ **and** *infinite (B ∩ G)*
 shows *infinite (A ⋯ B)*
 $\langle \text{proof} \rangle$

lemma *infinite-smul-left*:
 assumes $B \cap G \neq \{\}$ **and** *infinite (A ∩ G)*
 shows *infinite (A ⋯ B)*
 $\langle \text{proof} \rangle$

1.6 Pointwise set multiplication in a group: auxiliary lemmas

lemma *set-inverse-composition-commute*:
 assumes $X \subseteq G$ **and** $Y \subseteq G$
 shows $\text{inverse } ' (X \cdots Y) = (\text{inverse } ' Y) \cdots (\text{inverse } ' X)$
 $\langle \text{proof} \rangle$

lemma *smul-singleton-eq-contains-nat-powers*:
 fixes $n :: \text{nat}$
 assumes $X \subseteq G$ **and** $g \in G$ **and** $X \cdots \{g\} = X$
 shows $X \cdots \{g^{\wedge n}\} = X$
 $\langle \text{proof} \rangle$

lemma *smul-singleton-eq-contains-inverse-nat-powers*:
 fixes $n :: \text{nat}$
 assumes $X \subseteq G$ **and** $g \in G$ **and** $X \cdots \{g\} = X$
 shows $X \cdots \{(\text{inverse } g)^{\wedge n}\} = X$
 $\langle \text{proof} \rangle$

lemma *smul-singleton-eq-contains-powers*:
 fixes $n :: \text{nat}$
 assumes $X \subseteq G$ **and** $g \in G$ **and** $X \cdots \{g\} = X$

shows $X \cdots (\text{powers } g) = X \langle \text{proof} \rangle$

end

end

2 Generalized Cauchy–Davenport Theorem: main proof

theory *Generalized-Cauchy-Davenport-main-proof*

imports

Generalized-Cauchy-Davenport-preliminaries

begin

context *group*

begin

2.1 The counterexample pair relation in [4]

definition *devos-rel* **where**

$\text{devos-rel} = (\lambda (A, B). \text{card}(A \cdots B)) <_{\text{mlex}} (\text{inv-image } (\{(n, m). n > m\}) <_{\text{lex}})$
 $\text{measure } (\lambda (A, B). \text{card } A)) (\lambda (A, B). (\text{card } A + \text{card } B, (A, B)))$

lemma *devos-rel-iff*:

$((A, B), (C, D)) \in \text{devos-rel} \iff \text{card}(A \cdots B) < \text{card}(C \cdots D) \vee$
 $(\text{card}(A \cdots B) = \text{card}(C \cdots D) \wedge \text{card } A + \text{card } B > \text{card } C + \text{card } D) \vee$
 $(\text{card}(A \cdots B) = \text{card}(C \cdots D) \wedge \text{card } A + \text{card } B = \text{card } C + \text{card } D \wedge \text{card } A < \text{card } C)$
 $\langle \text{proof} \rangle$

lemma *devos-rel-le-smul*:

$((A, B), (C, D)) \in \text{devos-rel} \implies \text{card}(A \cdots B) \leq \text{card}(C \cdots D)$
 $\langle \text{proof} \rangle$

Lemma stating that the above relation due to DeVos is well-founded

lemma *devos-rel-wf* : *wf (Restr devos-rel*

$\{(A, B). \text{finite } A \wedge A \neq \{\} \wedge A \subseteq G \wedge \text{finite } B \wedge B \neq \{\} \wedge B \subseteq G\}$ **(is** *wf*
 $(\text{Restr devos-rel } ?\text{fin}))$
 $\langle \text{proof} \rangle$

2.2 $p(G)$ – the order of the smallest nontrivial finite subgroup of a group : definition and lemmas

definition *p* **where** $p = \text{Inf } (\text{card } ' \{H. \text{subgroup } H \ G \ (\cdot) \ \mathbf{1} \wedge \text{finite } H \wedge H \neq \{\mathbf{1}\}\})$

lemma *subgroup-finite-ge*:
 assumes *subgroup* H G (\cdot) **1** and $H \neq \{1\}$ and *finite* H
 shows $\text{card } H \geq p$
 $\langle \text{proof} \rangle$

lemma *subgroup-infinite-or-ge*:
 assumes *subgroup* H G (\cdot) **1** and $H \neq \{1\}$
 shows *infinite* $H \vee \text{card } H \geq p$ $\langle \text{proof} \rangle$

end

2.3 Proof of the Generalized Cauchy–Davenport Theorem for (non-abelian) groups

Generalized Cauchy–Davenport Theorem for (non-abelian) groups due to Matt DeVos [4]

theorem (*in group*) *Generalized-Cauchy-Davenport*:
 assumes $hAne: A \neq \{\}$ and $hBne: B \neq \{\}$ and $hAG: A \subseteq G$ and $hBG: B \subseteq G$ and
 $hAfin: \text{finite } A$ and $hBfin: \text{finite } B$ and
 $hsub: \{H. \text{ subgroup-of-group } H \ G \ (\cdot) \ \mathbf{1} \wedge \text{finite } H \wedge H \neq \{1\}\} \neq \{\}$
 shows $\text{card } (A \cdots B) \geq \min p (\text{card } A + \text{card } B - 1)$
 $\langle \text{proof} \rangle$

end

References

- [1] M. Bakšys and A. Koutsoukou-Argraki. Kneser’s theorem and the Cauchy–Davenport Theorem. *Archive of Formal Proofs*, November 2022. https://isa-afp.org/entries/Kneser_Cauchy_Davenport.html, Formal proof development.
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