## PCA algorithms proof

Friday, June 26, 2020 1:00 AM

Suppose we have the following collection of mpoints  $\{x^{(1)}, x^{(2)}, \dots, x^{(m)}\}$  in  $\mathbb{R}^m$ 

We like to reduce each 2" in the collection, i.e,

 $f: \mathbb{R}^n \longrightarrow \mathbb{R}^\ell$ 

 $f'(\alpha) = c$ 

So f is a function that map x to a ley olimons or veels space, so f is our woling function. That means we'll neo a decode fondon for f g

9: IR = IR 1 En and 9 (-5 (2)) 2 20 the most simple transformation is define of (c) = Dc, Dci to keep the obocooling function easy let's oleting on properties for D

1) the columns of D define an orthonormal busis Let be d(1), 6/(2), ..., d(4) 

2) DTD = 11

this properties give the problem a unuque solution Now we need a definition for fix) = e, and we remem JULIANO DO ME CAN OLE TIME

problem

c = avgmin 11 >c - g (c) 1/2

because Lz norm is alway positive then optimize the before is the same that appropriate the square of is because square Punction are monotonic for positive are

C = argmin 11 or - g cer 1/2

 $\|x - g(c)\|_{L^{2}} = (x - g(c))^{7} \cdot (x - g(c))$ =  $x^{7} x - x^{7} g(c) - g(c)^{7} x + g(c)^{7} g(c)$ 

 $z^{\dagger}g(c) \in \mathbb{R}$   $\Rightarrow z^{\dagger}g(c) = (xc^{\dagger}g(c))^{\dagger} \Rightarrow z^{\dagger}x - 2x^{\dagger}g(c) + g$ 

= g(cot x => x1x doalnot depend 30 15 content

Using the definition of g(c) = Dc

 $e^{*} = \operatorname{ovgmin} \left( -2 \times^{\mathsf{T}} g(\mathcal{O}) \right) +$