

# Computational Algebra. Lecture 4

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# Traditional Euclidean Algorithm

- 1  $r_0 \leftarrow f, r_2 \leftarrow g.$
- 2  $i \leftarrow 1$   
**while**  $r_i \neq 0$  **do**  $r_{i+1} \leftarrow r_{i+1} \text{ rem } r_i, i \leftarrow i + 1$
- 3 **return**  $r_{i-1}$

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- Which is the algebraic complexity of "MCD" ?

$$c(MCD) = c(While) = c(s[N]) = Fib(N) - 1,$$

where N is the number of steps (ie. divisions).

$$c(s[0]) = c(s[1]) = 0$$

$$c(s[i]) = c(s[i - 2]) + c(s[i - 1]) + 1$$

Writing a table of values, we conclude that:

$$c(s[i]) = Fib(i) - 1,$$

where  $Fib(i)$  is the i-th Fibonacci number. Moreover, if  $a > b > 0$ , it can be seen that  $N < 5\log_{10}(b) + 1$  (where  $\log_{10}(b)$  is the number of digits of b in basis 10).

# Homework: Stein Algorithm

This algorithm (1962, Silver and Tersian) (1967, Stein) does not involve any division in  $\mathbb{Q}$ , only subtractions, parity tests and divisions by two of even numbers. So, it is well adapted to binary calculus.

a) Show that if  $a$  and  $b$  are nonnegative integers, then

- $\gcd(2a, 2b) = 2\gcd(a, b)$
- $\gcd(2a + 1, 2b) = \gcd(2a + 1, b)$
- $\gcd(2a + 1, 2b + 1) = \gcd(\min(2a + 1, 2b + 1), |a - b|)$ .

# Homework: Stein Algorithm

- b) Design and implement a recursive algorithm which allows the computation of the gcd of any pair of integers by using the three properties mentioned above. Show that the algorithm concludes in a finite number of steps.

# Homework: Stein Algorithm

- c) Show that the number of bits operations of this algorithm is at most  $O((\log N)^2)$ , with  $N = \max(a, b)$ .

# Homework: Stein Algorithm

- c) Design and implement an iterative version of this algorithm.

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## Algorithm 1 Iterbin

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```
procedure Iterbin( $a, b$ )
     $d \leftarrow 1, t;$ 
    while EvenQ[ $a$ ] and EvenQ[ $b$ ] do
         $a \leftarrow \frac{a}{2}; b \leftarrow \frac{b}{2}; d \leftarrow 2d;$ 
    while  $a \neq 0$  do
        while EvenQ[ $a$ ] do  $a \leftarrow \frac{a}{2}$ 
        while EvenQ[ $b$ ] do  $b \leftarrow \frac{b}{2}$ 
         $t \leftarrow \frac{|a-b|}{2}$ 
        if  $a \geq b$  then  $a \leftarrow t$ 
        else  $b \leftarrow t$ 
    return  $d + t$ 
```

# Homework: Stein Algorithm

- e) Modify either of the versions (recursive or iterative) of your code in such a way that the output is a relation of the form  $r, s, u$ , with  $r.a + s.b = u = \gcd(a, b)$ .

**Hint:** Encyclopedia of Cryptography and Security