$TUM\ ModSim,\ SoSe\ 2023$

Mitschriften basierend auf der Vorlesung von Dr. Hans-Joachim Bungartz

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Introduction

About

Hier sind die wichtigsten Konzepte der ModSim Vorlesung von Dr. Hans-Joachim Bungartz im Sommersemester 2023 zusammengefasst.

Die Mitschriften selbst sind in Markdown geschrieben und werden mithilfe einer GitHub-Action nach jedem Push mithilfe von Pandoc zu einem PDF konvertiert.t

Eine stets aktuelle Version der PDFs kann über modsim_SS23_IN2010_merge.pdf heruntergeladen werden.

Implementation

Außerdem befindet sich eine Implementation von verschiedenen Algorithmen im Ordner /algorithms auf GitHub. Diese sind in Python und unter der Verwendung von NumPy geschrieben.

How to Contribute

- 1. Fork this Repository
- 2. Commit and push your changes to your forked repository
- 3. Open a Pull Request to this repository
- 4. Wait until the changes are merged

Contributors

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Focus Analysis / Calculus

Foundations

Functions and their representations

• One-Dimensional

$$f: D \subseteq \mathbb{R}^n \to \mathbb{R}^m, x \mapsto f(x)$$

• Multidimensional

$$f: D \subseteq \mathbb{R} \to \mathbb{R}, x = \begin{pmatrix} x_1 \\ \vdots \\ x_n \end{pmatrix} \mapsto f(x) = \begin{pmatrix} f_1(x_1, \dots, x_n) \\ \vdots \\ f_m(x_1, \dots, x_n) \end{pmatrix}$$

Names for special types of functions

- Curves: n = 1 and $m \in \mathbb{N}$
 - plane curves (2D): n = 1 and m = 2
 - space curves (3D): n = 1 and m = 3
- Surfaces: n=2 and m=3
- Scalar fields: $n \in \mathbb{N}$ and m = 1
- Vector fields: n = m

Topology concepts in higher dimensions

There is an analogous concept to open and closed intervals in multi-dimensional spaces. Given a domain $D \subset \mathbb{R}^n$ and its complement $D^c = \mathbb{R}^n \setminus D$

- A point x is called *inner point* if there exists an arbitrarily small ball around this points that fullly lies inside D.
- The set of all inner points of D is called the *interior* of D and is denoted as \mathring{D} .
- The domain is called open if $D = \mathring{D}$
- A point $x_0 \in \mathbb{R}^n$ is called boundary point if any arbitrarily small ball around this point intersects with both D and its complement D^c
- The set of all boundary points of D is called the boundary of D, denoted ∂D
- The set $\overline{D} = D \cup \partial D$ is called the *closure* of D

Using these definitions there are multiple attributes assignable to domains.

A domain D is called:

- closed if $\partial D \subseteq D$, i.e. $\bar{D} = D$
- bounded if $\exists K \in \mathbb{R} : ||x|| < K, \forall x \in D$
- compact if it is closed and bounded
- convex if all points on a straight line between to points in D are themselves element of D

Continuity

We define continuity in multi-dimensional spaces using converging vector sequences. A sequence $(x^{(k)})$ converges to the limit x if

$$\lim_{k \to \infty} ||x^{(k)} - x|| = 0$$

Converges of a vector sequence is also equivalent to the convergence of all components.

A vector function is then called continuous at $a \in D$ if for all sequences $(x^{(k)})_{k \in \mathbb{N}_0}$ in D converging to a the corresponding sequence $(f(x^{(k)}))_{k \in \mathbb{N}_0}$ in \mathbb{R}^m converges to f(a) and continuous on D if this holds for all points $a \in D$

Partial Differentiation

Gradient

The Gradient of a function gives the direction of the steepest ascent of the function. It requires that f represents a scalar field.

When applying the limit definition of the derivative to a function in higher dimesions it is not clear from which direction the derivative should be taken.

Using

$$\frac{\partial f}{\partial v}(a) = \lim_{h \to 0} \frac{f(a+hv) - f(a)}{h}$$

we can define the directional derivative of a function $f: \mathbb{R}^n \to \mathbb{R}$ along a vector $v \in \mathbb{R}^n$ at a point $a \in \mathbb{R}^n$.

If we use the coordinate vectors e_i as basis vectors for \mathbb{R}^n we can define the *Gradient* of f at a as

$$\nabla f(a) = \operatorname{grad} f(a) = \begin{pmatrix} \frac{\partial f}{\partial x_1}(a) \\ \vdots \\ \frac{\partial f}{\partial x_n}(a) \end{pmatrix}$$

For continuous functions the directional derivative at the point a along a vector v can be computed as

$$\frac{\partial f}{\partial v}(a) = \langle \nabla f(a), v \rangle$$

Example:

$$f(x,y) = x^2 + y^2 \rightarrow \nabla f(a) = \begin{pmatrix} 2x \\ 2y \end{pmatrix}$$

Hessian Matrix

The Hessian matrix of a function $f: \mathbb{R}^n \to \mathbb{R}$ at a point $a \in \mathbb{R}^n$ is the matrix of all second partial derivatives of f at a.

$$H_f(a) = \begin{pmatrix} \frac{\partial^2 f}{\partial x_1^2}(a) & \frac{\partial^2 f}{\partial x_1 \partial x_2}(a) & \dots & \frac{\partial^2 f}{\partial x_1 \partial x_n}(a) \\ \frac{\partial^2 f}{\partial x_2 \partial x_1}(a) & \frac{\partial^2 f}{\partial x_2^2}(a) & \dots & \frac{\partial^2 f}{\partial x_2 \partial x_n}(a) \\ \vdots & \vdots & \ddots & \vdots \\ \frac{\partial^2 f}{\partial x_n \partial x_1}(a) & \frac{\partial^2 f}{\partial x_n \partial x_2}(a) & \dots & \frac{\partial^2 f}{\partial x_n^2}(a) \end{pmatrix}$$

Example:

$$f(x,y) = x^2 + y^2 \to H_f(a) = \begin{pmatrix} 2 & 0 \\ 0 & 2 \end{pmatrix}$$

Jacobian Matrix

The Jacobian matrix of a function $f: \mathbb{R}^n \to \mathbb{R}^m$ at a point $a \in \mathbb{R}^n$ is the matrix of all partial derivatives of f at a.

In contrast to the Gradient the Jacobian matrix works for vector fields. It gives an analogue to the gradient for vector fields.

$$Df(a) = J_f(a) = \begin{pmatrix} \frac{\partial f_1}{\partial x_1}(a) & \frac{\partial f_1}{\partial x_2}(a) & \dots & \frac{\partial f_1}{\partial x_n}(a) \\ \frac{\partial f_2}{\partial x_1}(a) & \frac{\partial f_2}{\partial x_2}(a) & \dots & \frac{\partial f_2}{\partial x_n}(a) \\ \vdots & \vdots & \ddots & \vdots \\ \frac{\partial f_m}{\partial x_1}(a) & \frac{\partial f_m}{\partial x_2}(a) & \dots & \frac{\partial f_m}{\partial x_n}(a) \end{pmatrix} = \begin{pmatrix} \nabla f_1(a)^T \\ \nabla f_2(a)^T \\ \ddots \\ \nabla f_m(a)^T \end{pmatrix}$$

Example:

$$f(x,y) = \begin{pmatrix} x^2 + y \\ x^2 + y^2 \end{pmatrix} \rightarrow J_f(a) = \begin{pmatrix} 2x & 1 \\ 2x & 2y \end{pmatrix}$$

Calculation rules for the Jacobian

• Addition rule: $J(f+g) = J_f + J_g$

• Homogeneous rule: $J(cf) = cJ_f$ • Product rule: $J(f^T \cdot g) = f(x)^T J_q(x) + g(x)^T J_f(x)$

Laplace Operator

The Laplace operator is a second order partial derivative operator. It is defined on Scalar fields and is used to compute the rate of change of a scalar field.

$$\Delta f = \nabla^2 f = \sum_{i=1}^n \frac{\partial^2 f}{\partial x_i^2}$$

Example:

$$f(x,y) = x^2 + y^2 \to \Delta f(a) = 2 + 2 = 4$$

Divergence

The Divergence of a vector field is the rate of shrinkage or expansion around a point. It is defined as the sum of the partial derivatives of the components of the vector field.

$$\operatorname{div} f = \sum_{i=1}^{n} \frac{\partial f_i}{\partial x_i} = \nabla \cdot f$$

Example:

$$f(x,y) = \begin{pmatrix} x^2 \\ y^2 \end{pmatrix} \to \operatorname{div} f(a) = \frac{\partial f_1}{\partial x} + \frac{\partial f_2}{\partial y} = 2x + 2y$$

Curl / Rotation

The Curl of a vector field is the rate of rotation around a point. It is defined as the cross product of the partial derivatives of the components of the vector field.

$$\operatorname{rot} f = \nabla \times f = \begin{pmatrix} \frac{\partial f_3}{\partial x_2} - \frac{\partial f_2}{\partial x_3} \\ \frac{\partial f_1}{\partial x_3} - \frac{\partial f_3}{\partial x_1} \\ \frac{\partial f_2}{\partial x_1} - \frac{\partial f_1}{\partial x_2} \end{pmatrix}$$

Example:

$$f(x,y,z) = \begin{pmatrix} x^2y \\ y^2x \\ yz \end{pmatrix} \to \operatorname{rot} f(a) = \begin{pmatrix} z \\ 0 \\ y^2 - x^2 \end{pmatrix}$$

Taylor Expansion

It is also possible to approximate functions of multiple variables by Taylor expansions, by using the analog for higher order derivatives for functions of multiple variables.

Coordinate Transformations

A bijection between two coordinate systems is called a coordinate transformation. It is a function $\phi: \mathbb{R}^n \to \mathbb{R}^m$ that maps points in one coordinate system to points in another coordinate system and vice versa.

Jacobian Matrix of a Coordinate Transformation

The Jacobian matrix of this transformation is called *coordinate transformation matrix* and is defined as the matrix of all partial derivatives of the coordinate transformation.

Its determinant is called the Jacobian determinant.

Example:

We define the Transformation ϕ from polar coordinates to cartesian coordinates as follows:

$$\phi(r,\phi) = \begin{pmatrix} r\cos\phi\\ r\sin\phi \end{pmatrix} := \begin{pmatrix} x\\ y \end{pmatrix}$$

$$\implies J_{\phi}(r,\phi) = \begin{pmatrix} \cos\phi & -r\sin\phi\\ \sin\phi & r\cos\phi \end{pmatrix}$$

$$\implies \det J_{\phi}(r,\phi) = r$$

Roots and Optima

Newton's Method

Newtons method in higher dimensions works similar to the one dimensional case. The only difference is that we have to compute the gradient and the Hessian matrix instead of the derivative.

1D Newton's Method

$$x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}$$

Higher Dimensional Newton's Method

$$\mathbf{x_{n+1}} = \mathbf{x_n} - \left(\mathbf{J_f}(\mathbf{x_n})\right)^{-1} \cdot \mathbf{f}(\mathbf{x_n})$$

It konverges quadratically to a root in the neighborhood of the initial guess.

Optima

If f is a scalar field and $\mathbf{x_0}$ is a point in the domain of f:

- $\mathbf{x_0}$ is a global maximum if $f(\mathbf{x_0}) \geq \mathbf{f}(\mathbf{x})$ for all $\mathbf{x} \in \mathbf{D}$.
- $\mathbf{x_0}$ is a global minimum if $f(\mathbf{x_0}) \leq \mathbf{f}(\mathbf{x})$ for all $\mathbf{x} \in \mathbf{D}$.
- $\mathbf{x_0}$ is a local maximum if $f(\mathbf{x_0}) \ge \mathbf{f}(\mathbf{x})$ for all $\mathbf{x} \in \mathbf{B}_{\epsilon}(\mathbf{x_0})$.
- $\mathbf{x_0}$ is a local minimum if $f(\mathbf{x_0}) \leq \mathbf{f}(\mathbf{x})$ for all $\mathbf{x} \in \mathbf{B}_{\epsilon}(\mathbf{x_0})$.

All local extrema can be found from critical-points ($\nabla f(\mathbf{x}) = \mathbf{0}$) or at the boundary points.

But analog to the one-dimensional case, its also possible that a critical-point is neither a minum or a maximum point, because it is a "Saddle-Point"

Criteria

A critical Point $\mathbf{x_0} \in \mathbf{D}$ is:

- A local maximum if $H_f(\mathbf{x_0})$ is negative definite
- A local minum if $H_f(\mathbf{x_0})$ is positive definite
- A saddle point if $H_f(\mathbf{x_0})$ is indefite

Curves and Surfaces

Curve

A curve is a mapping $\gamma: I \subseteq \mathbb{R} \to \mathbb{R}^n$. In other words it maps a line into higher dimensional space. Example:

$$\gamma(t) = \begin{pmatrix} \cos(2t) \\ \sin(t) \end{pmatrix}$$

Surface

A surface is a mapping $\gamma: I \subseteq \mathbb{R}^2 \to \mathbb{R}^n$. In other words it maps a 2d-surface into higher dimensional space.

Quadrature

Integral over rectangular domains

If $f:[a,b]\times[c,d]\to\mathbb{R}$ is a scalar field, the "sum" of scalar values along the area of the Domain can be computed as:

$$\int \int_{D} f(x,y) dx dy = \int_{a}^{d} \int_{c}^{d} f(x,y) dx dy$$

Integration over simple domains

A 2d-standard Domain is defined as $D = \{(x,y)|x \in [a,b], l(x) \le y \le u(x)\}$. This means such a domain has one-varing parameter, which determines the lower- and upper bound of the surface at that point.

We can integrate over such domains using:

$$\int \int_D f(x,y) dx dy = \int_{x=a}^b \int_{l(x)}^{u(x)} f(x,y) dx dy$$

Integration under coordinate transformations

If $B,D\subseteq\mathbb{R}^n$ are Domains and $\phi:B\to D$ is a transformation between these domains. The Domain-Integral can be rewritten as:

$$\int_D f(x_1, \dots, x_n) dx_1 \cdots dx_2 = \int_B f(\phi(t_1, \dots, t_n)) \cdot |\det(J_\phi(t_1, \dots, t_n))| dt_1 \cdots dt_2$$

Partial differential equations

A PDE is a differential equation with multiple changing, variables.

Examples

- First order:
 - Traffic flow: $u_t + vu_x = f(x,t)$, where v is velocity, t is time and x is a point along the road. The resulting u(t,x) describes the cars at point x and time t
- - Heat equation: $u_t t c^2 \Delta u = 0$ Laplace equaiton: $-\Delta u = 0$

Introduction to Mathematical Modeling

Terminology

Model

A model is a simplified image of a partial reality.

- Practical Models
 - Wind tunnel
- Scale model
- Abstract Models
 - Mathematical models

Derivation

Questions to ask when modeling:

- What exactly should be modeled?
 - Population growth
 - Rocket trajectory
- Which attributes play a role in the model?
 - Population size, children per family, death rate...
 - Rocket mass, thrust, air resistance...
- What relations exist between the attributes?
 - Population growth is proportional to the population size
 - Rocket thrust is proportional to the fuel consumption
- What mathematical tools are needed to describe the relations?
 - Differential equations, probability theory, statistics, algebraic equations and inequalities, automata theory, graph theory, etc.

Simulation Tasks

- What is the goal of the simulation?
 - Find an arbitrary solution
 - Find the only solution
 - Show that a solution exists
 - Solve a constrained optimization problem
 - Find a critical point

Analysis of the Model

- What is the behavior of the model?
 - Is the model stable?
 - Does it converge to a steady state?
 - Is the solution point, really an optimum?
 - Is the solution unique or do exist better solutions?

Problems are **well-posed** if the following conditions are met:

- The solution exists
- The solution is unique
- It is stable

Aplicability of the Model

- Is enough input data available, to run the model?
- Is the hardware available to run the model?
- Is it fast enough, to be useful?
- Is it sensitive to small changes in the input data?

Mathematical Modeling

Process of formal derivation and analysis of mathematical models.

- 1. Informal description of the problem
- 2. Semi-formal description of the problem, using tools of the specific discipline
- 3. A strict formal description of the problem, (consistent)

Simulation

A virtual, computer based experiment with a mathematical model.

The goals of simulation are:

- To understand the behavior of the system
 - Why earthquakes occur
 - Why buildings collapse
- To optimize a system
 - Better flight schedules
 - higher throughput
- To predict the behavior of the system
 - Climate change
 - Characteristics of a new drug

Simulation Pipeline

- 1. Modeling the system
- 2. Numerical methods needed to solve the model
- 3. Implementation of the numerical methods in an efficient way
- 4. Visualization of the results
- 5. Validation of the model
- 6. Embedding the model in a larger system i.e a wheater forecast

Discrete Modeling and Simulation

Decision Model

Decisions can be made based of:

• certanty: all the information is known

• risk: the probability of each outcome is known

• uncertainty: the probability of each outcome is not known

Model

This way of making decision can be modeled using a Payoff Matrix.

	State 1	State 2		State n
Action 1	a_{11}	a_{12}		a_{1n}
Action 2	a_{21}	a_{22}		a_{2n}
:	:	:	٠	:
Action m	a_{m1}	a_{m2}		a_{mn}

The **payoff** is the value of the Action based on the State.

Example Payoff Matrix - Prisoner's Dilemma

If both players cooperate, they will both get a sentence of 5 years in prison. If one of them cooperates and the other doesn't, the one that cooperates will be free and the other will get a sentence of 10 years in prison. If both of them don't cooperate, they will both get a sentence of 1 year in prison.

(Reward A, Reward B)	Player A cooperates	Player A doesn't cooperate
Player B cooperates	-5, -5	-10,0
Player B doesn't cooperate	0, -10	-1, -1

Startegies

- **Certainity**: if you are certain about the initial state, your action will simply be the one that gives you the highest payoff.
 - 1. Maximum
 - Find Action i as $\arg \max_i N_{ij}$. In other words, find the action that gives you the highest payoff.
 - * In the example above, if you would somehow know that Player A will cooperate, it is of your best interest to also cooperate. Since this will reduce your sentence from 10 years to 5 years.
- Risk: There are multiple Strategies you can take:
 - 1. Caution (max-min payoff)
 - Find Action i as $\arg \max_i \min_j N_{ij}$. In other words, find the action that gives you the highest minimum payoff.

* In the example above, Player B would choose to cooperate, since chosing this way, will guarantee him a sentence of maximum 5 years.

2. Full Risk - (max-max payoff)

- Find Action i as $\arg \max_i \max_j N_{ij}$. In other words, find the action that gives you the highest maximum payoff.
 - * In the example above, Player B would choose to cooperate, since the best case scenario for him is to get a sentence of 0 years in prison, if Player A doesn't cooperate.
- 3. Alternative Caution (min-max payoff)
 - Find Action i as $\arg \min_i \max_j R_{ij}$. Where $R_{ij} = \max_i N_{ij} N_{ij}$.

4. Pessimism-Optimism

- Let $m_i = \min_j N_{ij}$ and $M_i = \max_j N_{ij}$, and $\alpha \in [0, 1]$.
- Find Action i as $\arg \max_{i} (\alpha m_i + (1 \alpha) M_i)$.
- This allows you to choose how much risk you want to take. If $\alpha = 0$, you will choose the action that gives you the highest maximum payoff. If $\alpha = 1$, you will choose the action that gives you the highest minimum payoff.

Two-Player Zero-Sum Games

Idea: Both players choose their actions simultaneously.

Assumptions:

- Both players know the payoff matrix.
- Both players act with caution. If they alternate turns:
 - S1 tries to maximize the minimum payoff.
 - S2 tries to minimize the maximum loss.

Equilibrium

If you find a state where both players are happy with their payoff, you have found an equilibrium. No player will want to change their action.

$$\max_{i} \min_{j} N_{ij} = a_{\hat{i}\hat{j}} = \min_{j} \max_{i} N_{ij}$$

Group Decision Making

Question: How to combine decisions from multiple people in a **democtratic** way, to find the best collective decision?

- A is the set of all candidates.
- $r: A \to \mathbb{R}$ is the ranking function.
 - Each voter i orders the candidates to his preference with numbers from 1 to |A|, where a lower number means a higher preference. He does this by defining a ranking function r_i .
- The set $P_A = \{ \rho \subset A \times A | \rho \text{ is a transitive, and asymetric relation} \}$ is the set of all possible rankings.
 - This means in P_A a voter can have equal preferences, between two candidates.

A colletive choice function $K: P_A^{\otimes n} \to P_A$ is a function that takes all n rankings from the voters and returns a combined ranking, which is the collective decision.

Majority Decision (Condorcet Method)

For each voter let N(x,y) be the number of voters that prefer x over y.

	$r_i(x)$	$r_i(y)$	$r_i(z)$
i = 1	1	2	3
i=2	3	1	2
i = 3	2	3	1

Therefore:

N(a, b)	x	y	z
\overline{x}	0	2	1
\overline{y}	1	0	2
\overline{z}	2	1	0

Define:

• $a\rho b$ if N(a,b) > N(b,a)

It follows:

- $x \rho y$ because N(x, y) > N(y, x)
- $y\rho z$ because N(y,z) > N(z,y)
- $z\rho x$ because N(z,x) > N(x,z)

So: $x >_{\rho} y >_{\rho} z >_{\rho} x$. This is a cycle, and therefore there we cannot find a collective decision.

This is bad, because even tough even tough every voter had a preference, we couldn't find a collective decision.

This violates rule 2 of a democratic decision. Because sometimes we cannot find a collective decision.

Rank Addition

This method simply adds the rankings of all voters together, and finds the collective decision from that

Define $x \rho y$ if $\sum_{i=1}^{n} r_i(x) < \sum_{i=1}^{n} r_i(y)$.

	$r_i(x)$	$r_i(y)$	$r_i(z)$
i = 1	1	2	3
i=2	2	1	3
$\sum r_i$	3	3	6

So by this method x and y are equally good, and z is the worst. This means we cannot find a unnique collective decision.

But this method yields another problem:

If we vote again:

	$r_i(x)$	$r_i(y)$	$r_i(z)$
i = 1	1	2	3
i=2	3	1	2
Σr_i	4	3	5

We get a different result. This times y is the clear winner. But every candidate still has the same preference between x and y but this time y is the winner.

This is bad, because the collective decision should not change if the relative order between candidates doesn't change.

This violates rule 4 of a democratic decision.

Rules of Democraty

- 1. For every set of Ballots, it must be possible to find a collective decision. (K is defined for all P_{Λ}^{n})
- 2. The result of the collective decision must be included in A. No external candidate can win.
- 3. If all voters decide unanimously, the collective decision must be the same. (Pareto Principle)
- 4. Ballots with the same order of candidates for every pair (x, y) must yield the same collective decision. (Independence of Irrelevant Alternatives)

5. No voter can always determine the collective decision. (Non-Dictatorship)\$\$

Arrow's Theorem

If |A| > 2, and more than 1 voter, then there is no collective choice function K that satisfies all 5 rules of democracy.

Scheduling

Problem:

- A Process consists of n tasks A_1, A_2, \ldots, A_n .
- There exist machines M_1, M_2, \ldots, M_m .
- Each task A_i needs a machine M_j to be executed. With an execution time of $t_i^{(j)}$

The goal is to find a schedule, that minimizes the total execution time.

Example: Process Scheduling

- Execution times is purely task-dependent.
- Start time s_i , completion time $c_i = s_i + t_i$.
- $A_i \to A_j$ means that A_i must be executed before A_j .
 - There might be no cycle in this graph.

Algorithm

```
Idea: Start the tasks as soon as possible.
add_initial_vertex()
#forward
while there are vertices left:
    find a vertex where all predecessors are already scheduled
    mark this vertex as scheduled
    set its start time as s=max(c of all predecessors)
    set its completion time as c=s+t
#backward
while there are vertices left:
    find a vertex where all successors are already processed
    mark this vertex as processed
    set its latest_end_time as l=min(s of all successors)
    set its latest_start time as latest_s=l-s
#critical path
critical_path=[]
for all vertices:
    if s==latest_s:
        add vertex to critical path
```

Job Shop Scheduling

- Every Job A_i consists of mutiple subjobs $A_{i1}, A_{i2}, \ldots, A_{ik}$.
 - Each subjob needs a machine M_j to be executed.
 - No circulation. Every Job requires M_j at most once. # Continuos Modeling

Population Dynamics

Model of Malthus

• There exists only one species P

- constant birth rate γ
- constant death rate δ
 - constant growth rate $\lambda = \gamma \delta$

This leads to the following equation:

$$p(t + \Delta t) = p(t) + \lambda p(t) \Delta t$$

in the limit $\Delta t \to 0$:

$$\dot{p}(t) = \lambda p(t)$$

This has the solution:

$$p(t) = p_0 e^{\lambda t}$$

Verhulst Model

- The model of Malthus is not realistic, since the population cannot grow indefinitely
- At some point the population will be saturated
- Ideas:
 - The population growth rate is proportional to the population size
 - linear birth rate $\gamma(t) = \gamma_0 \gamma_1 p(t)$
 - * Larger population size leads to smaller birth rate
 - linear death rate $\delta(t) = \delta_0 + \delta_1 p(t)$
 - * Larger population size leads to larger death rate

This leads to the following equation:

$$\dot{p}(t) = \gamma(t) - \delta(t) = \gamma_0 - \gamma_1 p(t) - \delta_0 - \delta_1 p(t) = -m \cdot (p(t) - p_\infty)$$

where $m = \gamma_1 + \delta_1$ and $p_{\infty} = \frac{\gamma_0 + \delta_0}{m}$

This has the solution:

$$p(t) = p_{\infty} + (p_0 - p_{\infty})e^{-mt}$$

This models starts to saturate right at the beginning, meaning that the growth rate shrinks right from the start. This is not realistic, since the population needs some time to run into resource limitations.

Logistic Model

- The model of Verhulst is not realistic, since the growth rate shrinks right from the start
- When using a quadratic term, it allows for an inflection point

$$\dot{p}(t) = (a - bp(t))p(t) = ap(t) - bp^2(t)$$

This has the solution:

$$p(t) = \frac{ap_0}{bp_0 + (a - bp_0)e^{-at}}$$

If $p_0 < \frac{a}{b}$, then the population will grow exponentially until it reaches the inflection point. Then it will start to saturate.

If a >> b, then the quadratic term will kick in very late, meaning that the population will grow exponentially for a long time.

Oscillations

• The logistic model does not allow for oscillations

$$\ddot{p}(t) + \mu \dot{p}(t) + \omega^2 (p(t) - p_{\infty}) = 0$$

where μ is the damping factor and ω_0 is the natural frequency.

This has the solution:

$$p(t) = (p_0 - p_\infty)e^{-\frac{\mu}{2}t}\cos(\sqrt{\omega^2 - \frac{\mu^2}{4}t}) + p_\infty$$

Lotka-Volterra Model

- Model of two species P and Q
- P is the prey and Q is the predator

$$\dot{p}(t) = f(p(t), q(t)) \cdot p(t)$$
$$\dot{q}(t) = g(p(t), q(t)) \cdot p(t)$$

If both $\dot{p}(t)$ and $\dot{q}(t)$ are zero, then the system is in equilibrium. This means that the population sizes are constant from that point on.

The values at this point are called fixed points.

Attractive Equilibrium Points

• The fixed point is called attractive, if the system will converge to this point, if it starts close to it

$$F(p,q) = \begin{bmatrix} f(p,q) \cdot p \\ g(p,q) \cdot q \end{bmatrix}$$

The Jacobian matrix is defined as:

$$J_F = \begin{bmatrix} \frac{\partial f}{\partial p} & \frac{\partial f}{\partial q} \\ \frac{\partial g}{\partial p} & \frac{\partial g}{\partial q} \end{bmatrix}$$

If we look at the eigenvalues of the Jacobi Matrix at the fixed point $[\bar{p}, \bar{q}]^T$ and their real parts are both negative, then the fixed point is attractive.

Numerical ODE Solvers

There exist two problem categories:

- Initial value problem
 - The initial value is given
 - The goal is to find the solution at a later time $p(0) = p_0$
- Boundary value problem
 - Both the initial and the final value are given
 - The goal is to find the solution in between $p(0) = p_0$ and $p(T) = p_T$

Definitions

- Local error
 - Maximum error between the exact solution and the numerical solution after one step

$$- e_{loc} = \max_{n \in [0, N-1]} \left| \frac{f(t_{n+1}) - f(t_n)}{\Delta t} - \dot{y}(t_n, y(t_n)) \right|$$

- Global error
 - Maximum error accumulated over all steps

$$-e_{glob} = \max_{n=0,\dots,N} |y_n - f(t_n)|$$

- Consistency
 - The method is consistent, if the local error goes to zero, if the step size goes to zero
 - $-\lim_{\Delta t \to 0} e_{loc} = 0$
- Convergence
 - The method is convergent, if the global error goes to zero, if the step size goes to zero * $\lim_{\Delta t \to 0} e_{glob} = 0$
 - Convergence is stronger than consistency
 - Convergence = Consistency + Stability
- Condition
 - Property of the problem
 - Measure for sensitifity of the solution to small changes in the initial value
 - Does not depend on the algorithm

$$\kappa_{rel} = \frac{|x \cdot \dot{y}(t, x)|}{|y(t, x)|}$$

- Stability
 - Property of the algorithm
 - Measure for the sensitivity of the solution to small changes in the initial value
 - Depends on the algorithm (ϵ -stability)
 - Sometimes implcit methods are more stable than explicit methods
- Stiffness
 - A problem is stiff, when the solution of a Differential Equation only converges, if the step size is very small
 - Even if the method is consistent and stable, the solution will not converge, if the step size is too large
 - Solution: Implicit methods

Euler Method

$$y_{n+1} = y_n + \Delta t \cdot f(t_n, y_n)$$

- Consistency
 - $-e_{loc}(\Delta t) = \mathcal{O}(\Delta t)$
- Convergence
 - $-e_{qlob}(\Delta t) = \mathcal{O}(\Delta t)$

Heun's Method (Runge-Kutta 2)

$$y_{n+1} = y_n + \frac{\Delta t}{2} \cdot (f(t_n, y_n) + f(t_{n+1}, y_n + \Delta t \cdot f(t_n, y_n)))$$

- Consistency
 - $-e_{loc}(\Delta t) = \mathcal{O}(\Delta t^2)$
- Convergence
 - $-e_{alob}(\Delta t) = \mathcal{O}(\Delta t^2)$

Multi-Step Methods

• Since Runge-Kutta methods are costly to compute, we can use multi-step methods, those method use previous steps to compute the next step

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Adams-Bashforth Method

$$y_{n+1} = y_n + \frac{\Delta t}{2} \cdot (3f(t_n, y_n) - f(t_{n-1}, y_{n-1}))$$

Since this method needs two previous steps, we need to use a different method for the first step
 Calculate the first step with the Euler method with a small step size

Implicit Methods

• Implicit methods are more stable than explicit methods

Implicit Euler Method

$$y_{n+1} = y_n + \Delta t \cdot f(t_{n+1}, y_{n+1})$$

You need to solve this equation for y_{n+1} and find a solution for y_{n+1} .

Predictor-Corrector Methods

- Predictor
 - Use an explicit method to predict the next step
- Corrector
 - Use an implicit method to correct the prediction

$$y_{n+1}^{pred} = y_n + \Delta t \cdot f(t_n, y_n)$$

$$y_{n+1} = y_n + \frac{\Delta t}{2} \cdot (f(t_n, y_n) + f(t_{n+1}, y_{n+1}^{pred}))$$

Higher Order ODEs

- Higher order Derivatives can be transformed into a system of first order ODEs
- Introduce new Variables $y_1 = y, y_2 = \dot{y}, y_3 = \ddot{y}, \dots$
- Transform the ODE $y^{(n)} = f(t, y, y', y'', \dots, y^{(n-1)})$ into a system of first order ODEs

$$y'_1 = y_2$$

 $y'_2 = y_3$
 \vdots
 $y'_{n-1} = y_n$
 $y'_n = f(t, y_1, y_2, \dots, y_n)$

Boundary Value Problems

• Initial value problems are easy to solve, but boundary value problems are hard to solve

$$\ddot{y} = b \cdot y + c$$

Finite Difference Method

• Calculate the finite difference approximation of the ODE

$$\ddot{y}(t_n) = \frac{y_{n+1} - 2y_n + y_{n-1}}{\Delta t^2}$$

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• Solve the system of equations via Gauss-Seidel or Jacobi

$$\frac{y_{i+1} - 2y_i + y_{i-1}}{\Delta t^2} - b_i \cdot y_i = c_i$$

$$\begin{pmatrix} 2 + b_1 \cdot \Delta t^2 & -1 & 0 & \dots & 0 \\ -1 & 2 + b_2 \cdot \Delta t^2 & -1 & \dots & 0 \\ 0 & -1 & 2 + b_3 \cdot \Delta t^2 & \dots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \dots & 2 + b_{N-1} \cdot \Delta t^2 \end{pmatrix} \begin{pmatrix} y_1 \\ y_2 \\ y_3 \\ \vdots \\ y_{N-1} \end{pmatrix} = \begin{pmatrix} -\Delta t^2 \cdot c_1 + y_0 \\ -\Delta t^2 \cdot c_2 \\ -\Delta t^2 \cdot c_3 \\ \vdots \\ -\Delta t^2 \cdot c_{N-1} + y_N \end{pmatrix}$$

• Solve the system of equations via Gauss-Seidel or Jacobi

Shooting Method

- Transform the boundary value problem into an initial value problem
- Solve the initial value problem with an ODE solver
- Adjust the initial value, until the boundary conditions are met
- Very expensive, since the ODE solver needs to be called multiple times

$$\ddot{y} = f(t, y, \dot{y}), \qquad y(t_0) = y_0, y(t_1) = y_1$$