

# Matrix-chain Multiplication

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## Problem

Given sequence of matrices

$$A_1 \times A_2 \times A_3 \times A_4,$$

with dimensions

$$A_{5 \times 4} \times A_{4 \times 6} \times A_{6 \times 2} \times A_{2 \times 7},$$

we want to minimize the cost of multiplying the matrices.

## Finding Cost

Given two matrices  $A$  and  $B$ , to find the cost or number of scalar multiplications of the matrix product  $A \times B = C$ , multiply the dimension of the rows and columns of  $A$  with the dimension of the columns of  $B$ . For example, given

$$A_{5 \times 4} \times B_{4 \times 3} = C_{5 \times 3}$$

The number of scalar multiplications would be:  $5 \cdot 4 \cdot 3 = 60$ . Using a bottom-up approach, we start with single matrices, then move up to multiply pairs, and then sets of three, and so on, to get the cost of the entire sequence.

Now, create tables  $m$  and  $s$  with row  $i$  and column  $j$ ,

- $m$ : at  $m[i, j]$ , the minimum cost is stored when multiplying matrices from  $A_i$  to  $A_j$ . We want to get to  $m[1, 4]$ , since we have 4 matrices.
- $s$ : at  $s[i, j]$ , row  $i$  will give the sequence of the optimal ordering; basically, it tells us how to parenthesize the sequence of matrices from  $i$  to  $j$ .

		1	2	3	4
	1				
$m$	2				
	3				
	4				

		1	2	3	4
	1				
$s$	2				
	3				
	4				

For the procedure, we only need to fill out half of the table along the diagonal going top-left to bottom-right, as cases below are symmetrical.

Start at  $m[i, j]$ , where  $i = j$ .

- The cost is zero, just having one matrix by itself would not cost anything to multiply. This is trivial, but for this diagonal, fill in all zeros.

		1	2	3	4
	1	0			
$m$	2		0		
	3			0	
	4				0

Move up to the next diagonal.

- $m[1, 2]$ : multiply matrices  $A_1$  and  $A_2$ .  
 $(A_{5 \times 4}) \times A_{4 \times 6}$  gives a cost of  $5 \cdot 4 \cdot 6 = 120$ .  
Thus,  $m[1, 2] = 120$ .  
Place value of 1 in  $s[1, 2]$  since we split the sequence at  $A_1$ .
- $m[2, 3]$ : multiply matrices  $A_2$  and  $A_3$ .  
 $(A_{4 \times 6}) \times A_{6 \times 2}$  gives a cost of  $4 \cdot 6 \cdot 2 = 48$ .  
Thus,  $m[2, 3] = 48$ .  
Place value of 2 in  $s[2, 3]$  since we split the sequence at  $A_2$ .
- $m[3, 4]$ : multiply matrices  $A_3$  and  $A_4$ .  
 $(A_{6 \times 2}) \times A_{2 \times 7}$  gives a cost of  $6 \cdot 2 \cdot 7 = 84$ .  
Thus,  $m[3, 4] = 84$ .  
Place value of 3 in  $s[3, 4]$  since we split the sequence at  $A_3$ .

This gives the updated tables  $m$  and  $s$ .

		1	2	3	4
$m$	1	0	120		
	2		0	48	
	3			0	84
	4				0

		1	2	3	4
$s$	1		1		
	2			2	
	3				3
	4				

Now we have calculated the cost of multiplying all pairs possible. Now let's calculate the cost of multiplying sequences with three matrices.

- $m[1, 3]$ : Consider two possibilities.

1.  $A_{5 \times 4} \times (A_{4 \times 6} \times A_{6 \times 2}) = A_{5 \times 4} \times A_{4 \times 2}$

Now, get the cost of  $A_1$ , a single matrix, plus the cost of  $A_2 \times A_3$  (the inner multiplication), and getting the overall cost with this specific sequence. Recall  $m[1, 1]$  and  $m[2, 3]$  have already been calculated.

$$m[1, 1] + m[2, 3] + 5 \cdot 4 \cdot 2 = 0 + 48 + 40 = 88$$

2.  $(A_{5 \times 4} \times A_{4 \times 6}) \times A_{6 \times 2} = A_{5 \times 6} \times A_{6 \times 2}$

$$m[1, 2] + m[3, 3] + 5 \cdot 6 \cdot 2 = 120 + 0 + 60 = 180$$

Now, take the smaller cost, which is 88, thus  $m[1, 3] = 88$ . Also,  $s[1, 3] = 1$  since  $A_1$  was our cut-off.

- $m[2, 4]$ : Consider two possibilities again.

1.  $A_{4 \times 6} \times (A_{6 \times 2} \times A_{2 \times 7}) = A_{4 \times 6} \times A_{6 \times 7}$

$$m[2, 2] + m[3, 4] + 4 \cdot 6 \cdot 7 = 0 + 84 + 168 = 252$$

2.  $(A_{4 \times 6} \times A_{6 \times 2}) \times A_{2 \times 7} = A_{4 \times 2} \times A_{2 \times 7}$

$$m[2, 3] + m[4, 4] + 4 \cdot 2 \cdot 7 = 48 + 0 + 56 = 104$$

Now, take the smaller cost, which is 104, thus  $m[2, 4] = 104$ . Also,  $s[2, 4] = 3$  since  $A_3$  separates  $A_2 \times A_3$  from  $A_4$ .

That gives us the updated tables.

		1	2	3	4
$m$	1	0	120	88	
	2		0	48	104
	3			0	84
	4				0

		1	2	3	4
$s$	1		1	1	
	2			2	3
	3				3
	4				

Now, calculate the cost of the entire sequence,  $m[1, 4]$ .

$$\begin{aligned}
 m[1, 4] &= \min\{m[1, 1] + m[2, 4] + 5 \cdot 4 \cdot 7, \\
 &\quad m[1, 2] + m[3, 4] + 5 \cdot 6 \cdot 7, \\
 &\quad m[1, 3] + m[4, 4] + 5 \cdot 2 \cdot 7 \\
 &= \min\{0 + 104 + 140, \\
 &\quad 120 + 84 + 210, \\
 &\quad 88 + 0 + 70\} \\
 &= \min\{244, 414, 158\} \\
 m[1, 4] &= 158
 \end{aligned}$$

Since our solution comes from splitting  $A_1 \cdot A_2 \cdot A_3$  from  $A_4$ ,  $s[1, 4] = 3$ . This gives our final tables  $m$  and  $s$ .

		1	2	3	4
$m$	1	0	120	88	158
	2		0	48	104
	3			0	84
	4				0

		1	2	3	4
$s$	1		1	1	3
	2			2	3
	3				3
	4				

## Solution

This is the general formula for finding the ideal sequence to multiply matrices  $i$  to  $j$ , with dimensions for rows and columns labeled with  $d$ .

$$m[i, j] = \min_{i \leq k \leq j} \{m[i, k] + m[k + 1, j] + d_{i-1} \cdot d_k \cdot d_j\}$$

Given  $n$  matrices, we will have  $n + 1$  dimensions. Here is the optimal parenthesization.

$$(A_1 \cdot A_2 \cdot A_3) \cdot (A_4)$$

Cut off the sequence at third matrix since  $s[1, 4] = 3$ .

$$((A_1) \cdot (A_2 \cdot A_3)) \cdot (A_4)$$

Cut off the sequence at first matrix since  $s[1, 3] = 1$ .

Thus, this is the sequence of matrices that requires the least number of scalar multiplications.

$$((A_{5 \times 4}) \times (A_{4 \times 6} \times A_{6 \times 2})) \times (A_{2 \times 7})$$