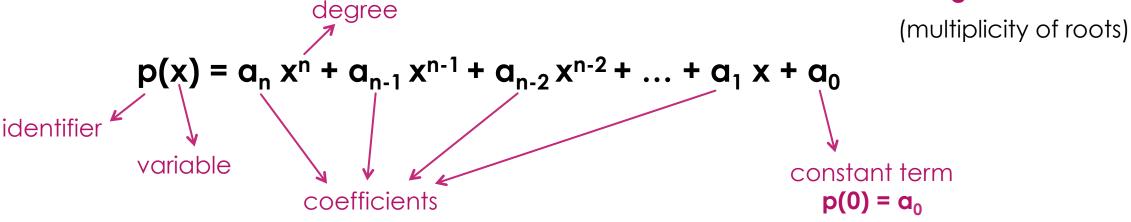
- ✓ Algebraic expressions (single variable or multi-variable)
- ✓ Domain of input variable: real numbers (more generally: complex numbers)
  - Range of <u>output</u> values: real numbers (more generally: complex numbers)
- ✓ Only **non-negative** integer powers (n = 0,1,2,3,4,...)
- $\checkmark$  n is degree: deg(p(x)) = n

A polynomial of deg = n has maximum n real roots (exactly n complex roots)

Fundamental theorem of algebra



(**Real** or **Complex** numbers)

#### Fundamental theorem of algebra

(https://en.wikipedia.org/wiki/Fundamental\_theorem\_of\_algebra)

From Wikipedia, the free encyclopedia

Not to be confused with Fundamental theorem of arithmetic.

The fundamental theorem of algebra states that every non-constant single-variable polynomial with complex coefficients has at least one complex root. This includes polynomials with real coefficients, since every real number can be considered a complex number with its imaginary part equal to zero.

Equivalently (by definition), the theorem states that the field of complex numbers is algebraically closed.

$$p(x) = (x-1)^2 = x^2 - 2x + 1 \rightarrow two identical roots$$

The theorem is also stated as follows; every non-zero, single-variable, degree *n* polynomial with complex coefficients has, counted with multiplicity, exactly *n* complex roots. The equivalence of the two statements can be proven through the use of successive polynomial division.

In spite of its name, there is no purely algebraic proof of the theorem, since any proof must use some form of the analytic completeness of the real numbers, which is not an algebraic concept.<sup>[1]</sup> Additionally, it is not fundamental for modern algebra; its name was given at a time when algebra was synonymous with theory of equations.

A polynomial of deg = n has maximum n real roots (exactly n complex roots)

Fundamental theorem of algebra

identifier 
$$p(x) = a_n x^n + a_{n-1} x^{n-1} + a_{n-2} x^{n-2} + ... + a_1 x + a_0$$

$$constant term$$

$$p(0) = a_0$$

(Real or Complex numbers)

- ✓ What we need to know:
  - Construct a polynomial (single variable) where an array of real or complex numbers
  - Degree Highest power (integer number ≥ 0)
  - Value & string representation  $p(x) = a_n x^n + ... + a_1 x + a_0$  where  $a_n \neq 0$
  - Composition p(q(x))
  - Algebra: addition, subtraction  $p(x) \pm q(x) = s(x)$ : polynomial
  - Algebra: multiplication p(x) \* q(x) = s(x): polynomial
  - Division is more complicated p(x) / q(x) = s(x) + r(x)/q(x) where deg(r(x)) < deg(q(x))
  - Derivative p'(x) = s(x) = d(p(x))/dx where deg(p'(x)) = deg(p(x))-1
  - Integral  $\longrightarrow$  s'(x) = p(x) where deg(s(x)) = deg(p(x))+1
  - Roots (zeros)  $\triangleright$  (x) = 0
    - May be real or complex
    - Even if the coefficients are real, some roots may be complex numbers

Multiply

## Polynomials

#### ✓ Horner's method

- **Most efficient way** of evaluating a polynomial
- Only **n summation** and **n multiplications** (n = deg(p(x)))

$$p(x) = a_0 + a_1 x + a_2 x^2 + \dots + a_n x^n = a_0 + x(a_1 + x(a_2 + x(a_3 + \dots + x(a_{n-1} + xa_n)\dots)))$$

Evaluating  $p(x_0)$ 

• 
$$b_n = a_n$$

• 
$$b_{n-1} = a_{n-1} + b_n x_0$$

$$b_{n-2} = a_{n-2} + b_{n-1} x_0$$

• 
$$b_0 = a_0 + b_1 x_0$$

This is just another way of writing the polynomial! (pretty clever way!!)

n multiplications & n summation

•  $b_{n-2} = a_{n-2} + b_{n-1} x_0$  | Evaluate this sequence  $p(x_0) = b_0$ 



Convolution sum

remainder

### Polynomials

#### ✓ Algebra

- Addition, subtraction:  $p(x) \pm q(x) = s(x)$ 
  - Result is <u>always</u> a polynomial
  - Add coefficients term by term:  $s_n = p_n \pm q_n$
  - $deg(s(x)) = max(deg(p(x)), deg(q(x))) \rightarrow general case$
- Multiplication:  $p(x) * q(x) = s(x) \rightarrow deg(s) = deg(p) + deg(q)$ 
  - Result is <u>always</u> a polynomial
  - Convolution of coefficients:  $s_n = p_n$  convolves  $q_n \rightarrow s_n = p_n q_0 + p_{n-1} q_1 + ... + p_0 q_n$
- Division
  - Result is **not** always a polynomial
  - Result is a rational expression
  - Fundamental theorem of division: p(x)/q(x) = s(x) + r(x)/q(x) where deg(r(x)) < deg(q(x))

-

Tip: implement algebra as part of the "operator overloading".

- Other useful libraries
  - Michael Flanagan's Java Scientific Library
  - Archive (jar) available on this personal website:
    - ✓ https://www.ee.ucl.ac.uk/~mflanaga/java/
  - Source code available on github
    - https://github.com/bgithub1/flanagan
  - You can write adapter methods from your polynomial class to Flanagan's polynomial class
    - ✓ Adapters are very useful when using methods of different java libraries
      - o They allow interfacing with other java libraries
    - ✓ We will see this in more details in "linear algebra" topic.

Adapter method: public flanagan.Polynomial toFlanaganPolynomial()

method name

### Polynomial Division

- Result of division is not a polynomial in general
  - Instead it should return two polynomials
  - $P(x) = B(x) Q(x) + R(x) \rightarrow Q(x)$ : quotient, R(x): remainder, P(x): dividend, B(x): divisor
  - Remember: this is an "identity" that holds for any real or complex number
- Fundamental Lemma: deg(R(x)) < deg(B(x))</li>
  - Equality is not allowed
- Special cases
  - $P(x) = (x-a) Q(x) + R \rightarrow remainder is a number \rightarrow R = P(a)$
  - $P(x) = (x-a)(x-b) Q(x) + (mx+n) \rightarrow$  remainder is first order
    - $\mathbf{m} = (P(a)-P(b))/(a-b), \mathbf{n} = (a P(b)-b P(a))/(a-b)$
- What to do in the general case?
  - Must find both Q(x) and R(x)

### Polynomial Division

- Polynomial Long Division
  - Sort coefficients from highest degree to lowest
  - Make sure the coefficient of the highest degree is NOT zero
  - If  $deg(P(x)) < deg(B(x)) \rightarrow Nothing to do: Q(x)=0, R(x)=P(x)$

$$P(x) = B(x) Q(x) + R(x)$$

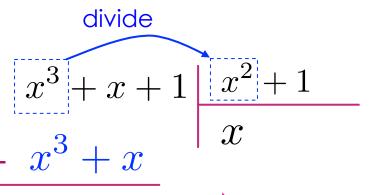
$$\begin{cases} P(x) = p_n x^n + ... + p_1 x + p_0 \\ B(x) = b_m x^m + ... + b_1 x + b_0 \end{cases}$$

Taking use of

multiplication and

subtraction

Strictly less than



Deg(R(x)) < deg(B(x)) : STOP

reduce() method

#### **Flowchart**

Initialize 
$$Q(x) = 0$$
,  $R(x) = 0$ 
 $+ \dots + p_1 \times + p_0$ 
 $-1 + \dots + b_1 \times + b_0$ 

Sort  $P(x)$  coefficients

 $Q(x) + = (p_n/b_m) \times p_{-m}$ 

Taking use of multiplication and subtraction

Strictly less than (equality is not allowed)

STOP

Initialize  $Q(x) = 0$ ,  $R(x) = 0$ 
 $Q(x) + = (p_n/b_m) \times p_{-m}$ 
 $Q(x) + = (p_n/b_m) \times p_{-m}$ 
 $Q(x) + = (p_n/b_m) \times p_{-m}$ 

Initialize  $Q(x) = 0$ ,  $Q(x) = 0$ 
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 $Q(x) + = (p_n/b_m) \times p_{-m}$ 

Initialize  $Q(x) = 0$ ,  $Q(x) = 0$ 

Initialize  $Q(x) = 0$ 

Initi

# Taylor Series

#### $\checkmark$ A polynomial series representing a smooth function f(x)

- Polynomial expansion at a given point

Based on derivatives of the function 
$$f(x) = f(x_0) + \frac{1}{1!} \frac{df}{dx}|_{x_0} (x - x_0) + \frac{1}{2!} \frac{d^2f}{dx^2}|_{x_0} (x - x_0)^2 + \dots$$
 Polynomial Point of interest

coefficient = 1/n!

- Locally approximates the function
- **Example**: Taylor expansion of exp(x) at x = 0  $e^x = \sum_{n=0}^{\infty} \frac{x^n}{n!} = 1 + x + \frac{x^2}{2} + \frac{x^3}{6} + \dots$  Different ways of implementing
- Different ways of implementing
  - Use the symbolic math library we developed before (call **diff()** multiple times)
  - Numerically evaluate higher order derivatives at each point
    - Richardson extrapolation

# Taylor Series

- Taylor series has a **uniform** convergence
- You can apply term-by-term algebra
  - Add, subtract
  - Multiply, divide by a number
  - Differentiate
  - Integrate
- Necessary condition for the convergence
  - $\lim |a_n(x-x_0)^n| \rightarrow 0 \text{ as } n \rightarrow \infty$

$$R = \frac{1}{\lim_{n \to \infty} \sqrt[n]{|a_n|}} \longrightarrow$$

Cauchy-Hadamard theorem

$$f(x) = \sum_{n=0}^{\infty} f_n(x - x_0)^n$$
  $g(x) = \sum_{n=0}^{\infty} g_n(x - x_0)^n$ 

$$f(x) \pm g(x) = \sum_{n=0}^{\infty} (f_n \pm g_n)(x - x_0)^n$$

$$c f(x) = \sum_{n=0}^{\infty} (c f_n)(x - x_0)^n$$

$$f'(x) = \sum_{n=1}^{\infty} (n f_n)(x - x_0)^{n-1}$$

If  $a_n$ 's have the same sign as  $n \rightarrow \infty$ 

$$R = \lim_{n \to \infty} \frac{a_n}{a_{n+1}}$$

If  $a_n$ 's have alternating sign as  $n \rightarrow \infty$ 

$$b_n = \sqrt{|a_n a_{n-2} - a_{n-1}^2|} \implies R = \lim_{n \to \infty} \frac{b_n}{b_{n+1}}$$

R can be infinite



f(x): entire