ABAB

Naive: for each i, go as right as possible until k is used up

Seach A, B, C = 2 separately

000,020305 1 1

> (l,r) is the longest starting from l

When we move I by 1,
we know that it has at least
[[11] r) is a valid one!

Directly apply a sliding window

000102013 TH 0506 record a freg [26] if in [l,r], max(freq)+k>r-l+1 -) We can keep go right if [l, r) is the longest one, we will need to move I right

Let say max_count = max (freq)

We will find out that when we

Move I to the right, we don't know

how to update max_count to

reflect the (true) max(freq)

within a window.

However, closs it really matter that

Max_count reflects max(freq)
?

Suppose:

max_count = 3 (: # of A)

AAA BCDEF it we don't update max-count, [l', r] at most can be expanded to max_count + K Observe that if max_count reflects the wrong max (freq.), the window length will not exceed max-count + k I) not affect our answer!

Claim: Suppose [l,r] is the optimal solution and max (freq, [l,r]) = true_max count and k ≥0 is the # operations we allow to do

Case 1: r-l+1 = true - max count + KCase 2: r-l+1 < true - max count + K

For Case 1:

when we expand using a fake max count which does not reflect max (freq, window),

it is ok be cause fake max count < true max count

The window created by fake maxcount will not exceed true-maxcount+k

we still have operations to use

>> keep expanding

-> contradicts that this is

the optimal solution