$s_0 s_1 s_2 s_3 \cdots \quad s_m$

$t_0 t_1 t_2 t_3 \cdots t_n$

$freq(t) = [\underset{0}{a} \quad \underset{2}{b} \quad \underset{1}{c} \quad \underset{0}{d} \quad e \cdots]$

$s_i \; s_{i+1} \; s \; s \cdots s_j$

$freq(s_i : s_j) \geq freq(t)$

for all alphabets

(Vector comparison

A naive solution is

$O(m^2)$: for every (i, j) pair $(i \le j)$

we check if $freq(S_i \sim S_j) \ge freq(t)$

Observe that

if $freq(S_i \sim S_j) \ge freq(t)$

there is no need to check

$freq(S_i \sim S_{j+1})$

because it must $\ge freq(t)$:

And $len(S_i \sim S_j) < len(S_i \sim S_{j+1})$

so it is useless to search

$S_i \sim S_k$ where $k \geq j+1$!