

# Tree Dynamic Programming

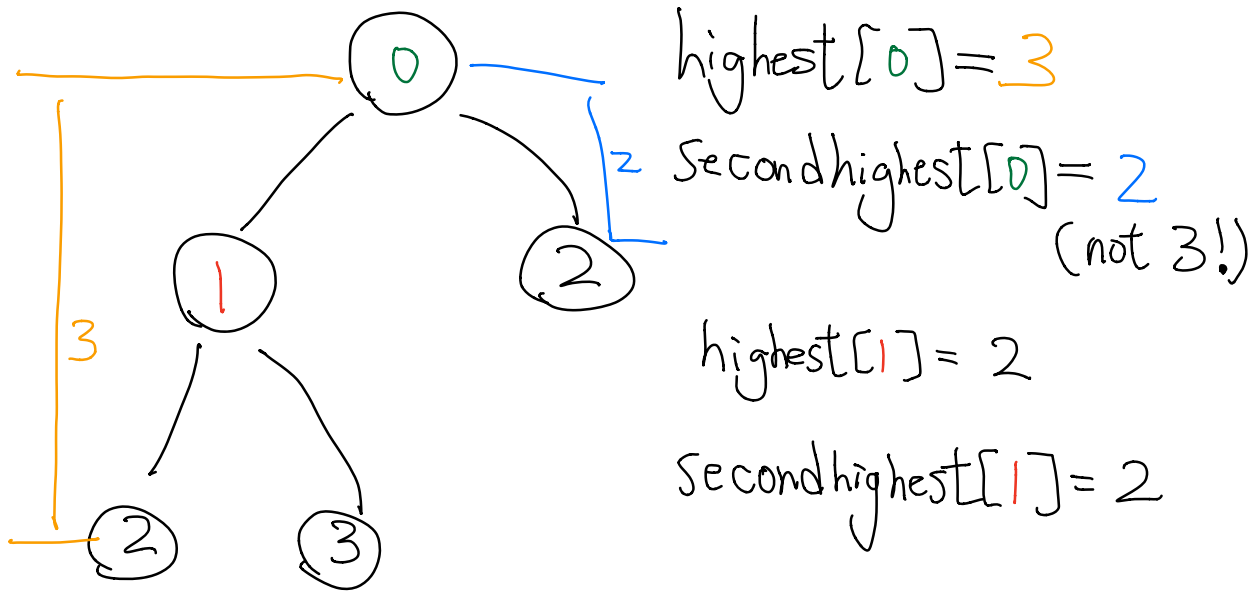
Setting: A tree roots at 0.  $\Rightarrow$  Compute  $dp[i]$

Trick: for each node  $i$ , we store two values

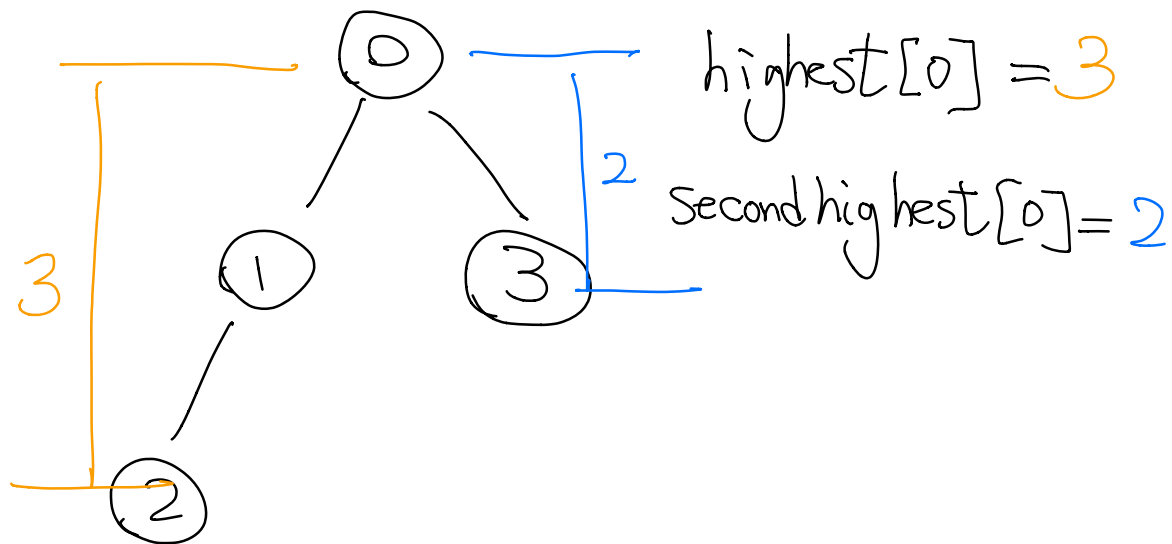
$highest[i]$

$secondhighest[i]$

Ex 1.



Ex 2.



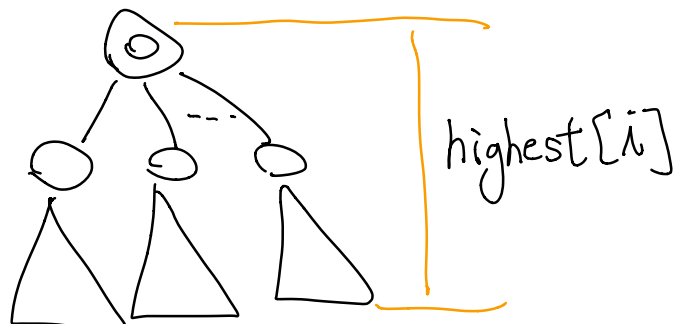
Intuition:

Let  $dp[i]$  be the height of a tree rooted at  $i$ .  $O(n)$

And assume we have already computed

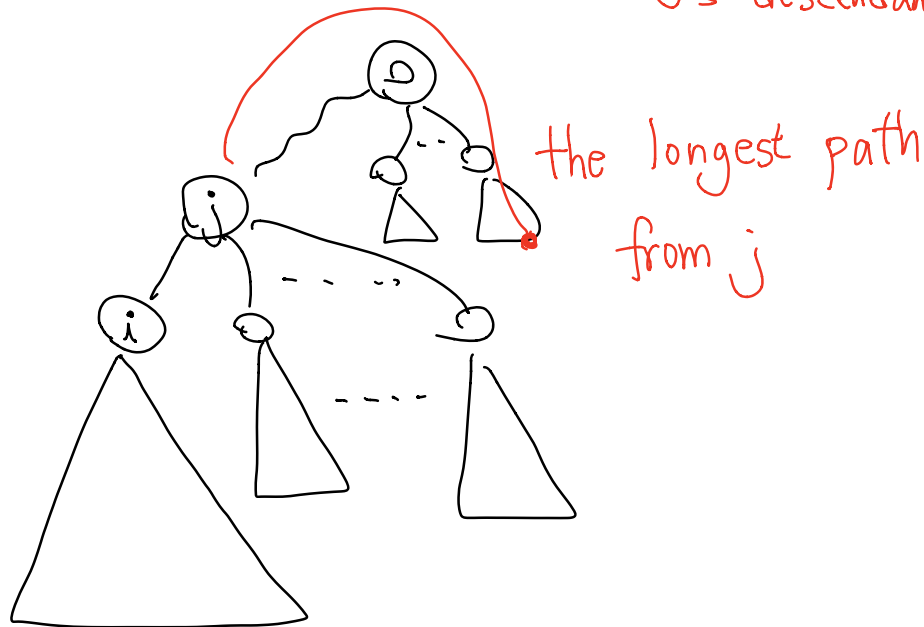
highest[i]  $i \in [0, n)$   
second highest[i] in advance by DFS

It will be easy to know  $dp[0]$



How about  $dp[i]$ ?

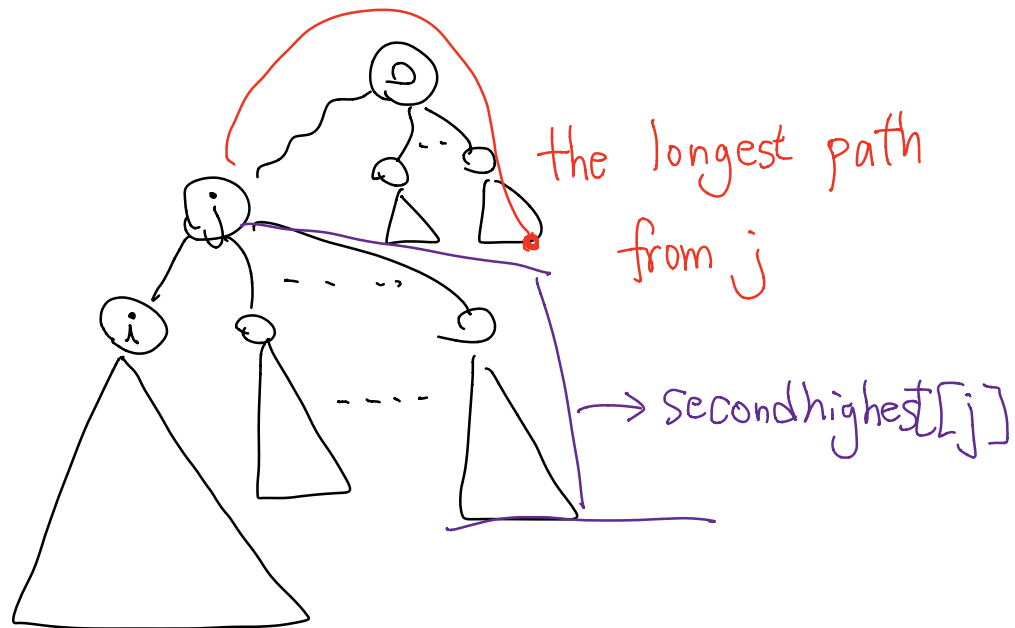
Assume we know the longest path from  $j$  to other nodes other than  $j$ 's descendants



It will be easy to see:

$$dp[j] = \max(\text{highest}[j], \text{the longest path from } j)$$

How about the longest path from  $i$  ?

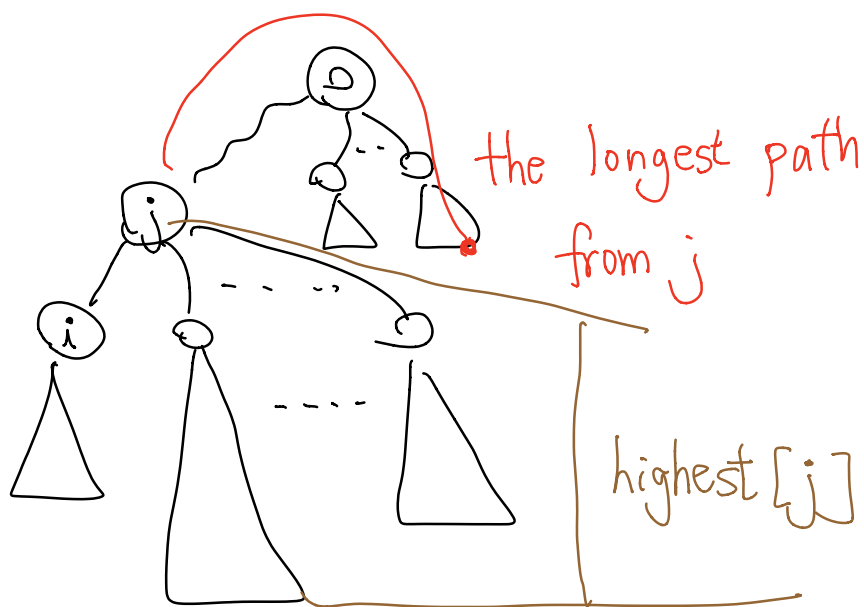


It turns out it is easy to see :

① If  $\text{highest}[i] + 1 \Rightarrow \text{highest}[j]$  :

the longest path from  $i$  =

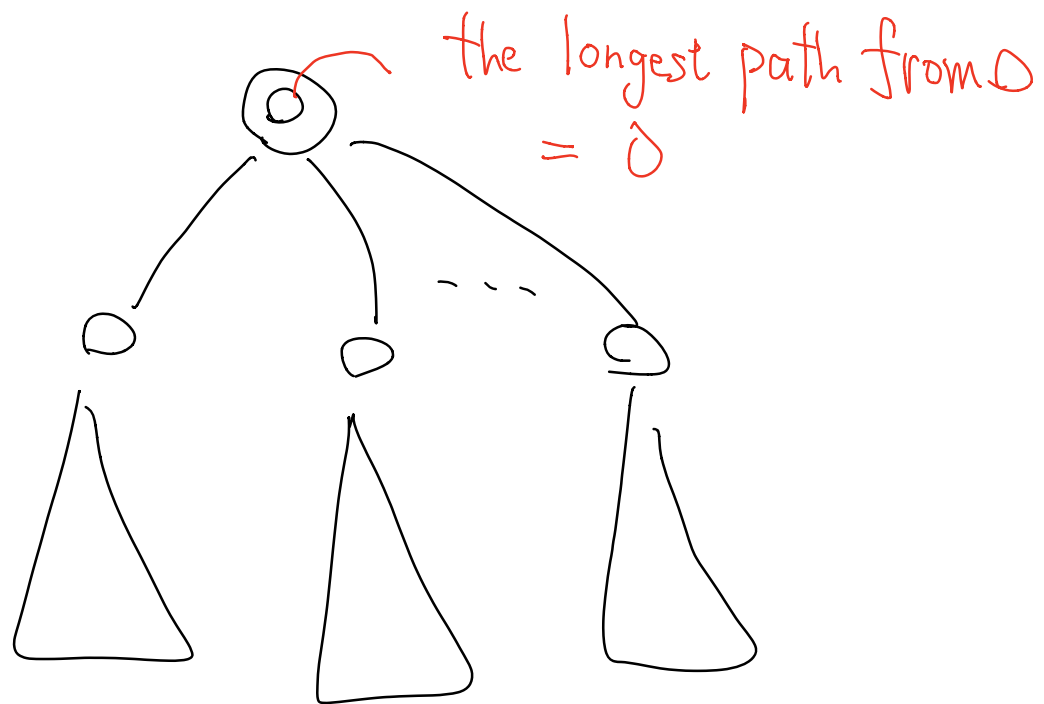
$\max(\text{secondhighest}[j], \text{the longest path from } j) + 1$



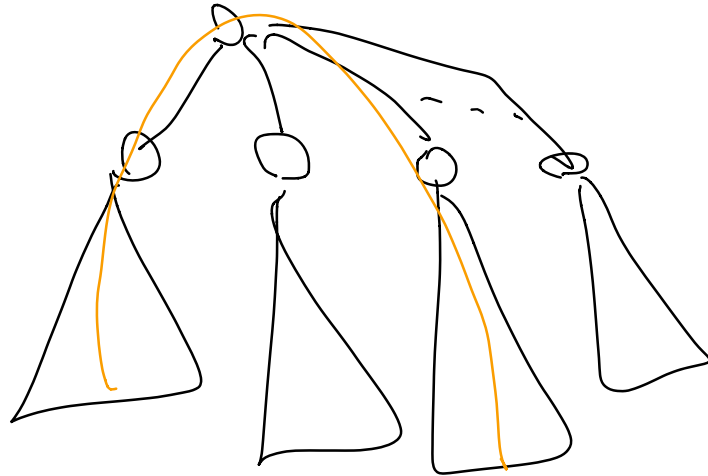
② If  $\text{highest}[i] + 1 < \text{highest}[j]$ :

the longest path from  $i$  =

$\max(\text{highest}[j], \text{the longest path from } j) + 1$



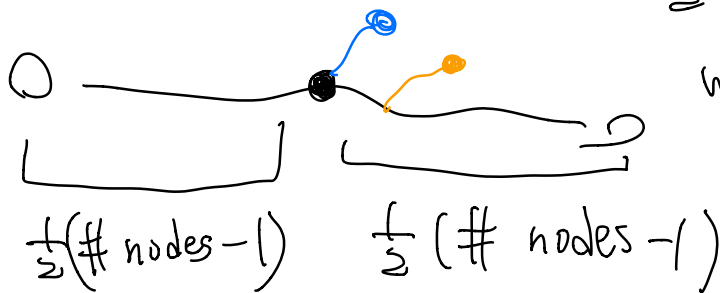
Assume: there is only one longest path



The path contains odd nodes


When we choose the middle point, the height of it is  $\frac{1}{2}(\# \text{ of nodes} - 1) + 1$

If we choose , it will have

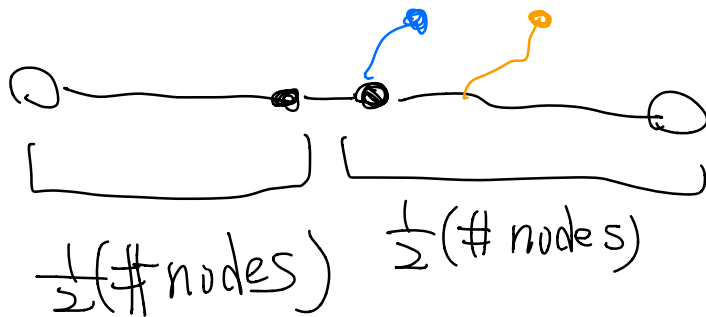


$$\frac{1}{2}(\# \text{ nodes} - 1) + 1 + 1$$

which is greater

Choose  will be even longer.

When the path contain even nodes:



Choose ②:

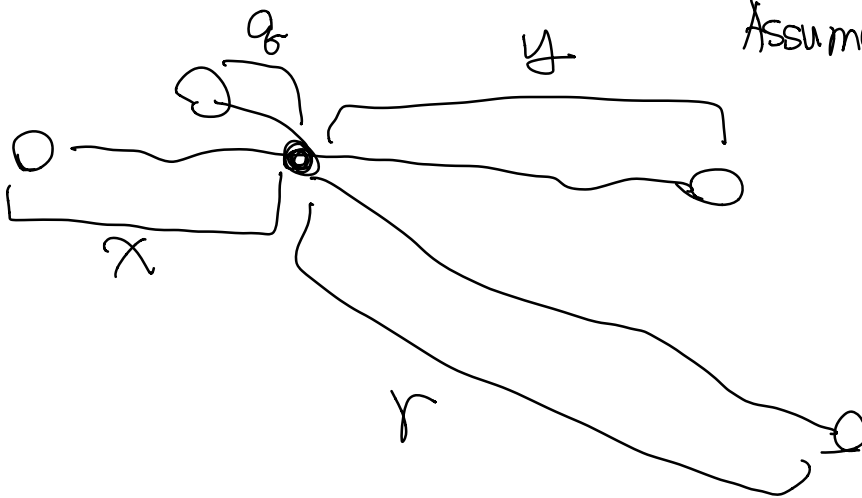
$$\text{height} = \frac{1}{2}(\# \text{ nodes}) + 1$$

Choose ③:

$$\text{height} = \frac{1}{2}(\# \text{ nodes})$$

Choose ④ will be even larger.

When there are multiple longest path



Assume  $\#$  nodes on the path  
=  $p$

Assume  $x + 1 + y = q + r + 1 = p$



WLOG, we first assume  $r > y$

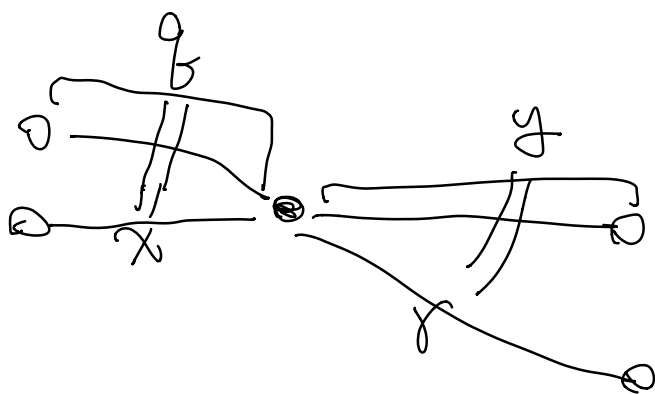
If  $r > y$ , then  $x + r + 1$  can form a  
even longer path,

contradicts that  $x + 1 + y$  or  
 $q + r + 1$  are  
the longest path

So apparently,  $r = y$ ,

if  $r = y$ , then  $x = q$ ,

The tree should be like this

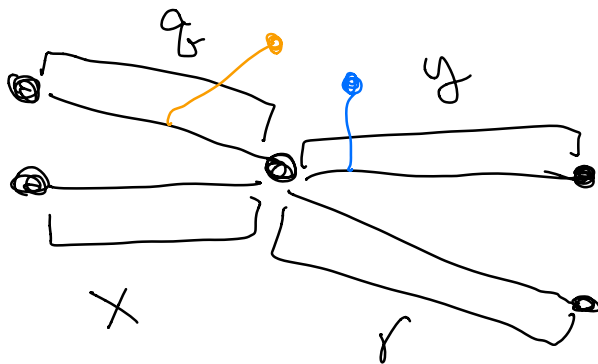


WLOG,  $x \leq y$  ( $q \leq r$ )

However, if  $x = q \leq y = r$ ,



$y + r + 1$  can form a longer path!

It should look like this instead.



where  $x = q = y = r$

⇒ We can just find the middle point.

Because choosing  or  will result a longer height