

$[w_0, w_1, w_2, w_3 \dots w_{n-1}]$

$\uparrow$   
 $w_l$

$\uparrow$   
 $w_r$

Naïve: use hashmap  $[w] = \text{a list of indices}$

$\Rightarrow O(n^2)$  at worst

Ex.

$[A, B, A, B, A, B, A, B]$

$\text{hashmap}[w_l] \rightarrow \text{sorted}$

$\text{hashmap}[w_r] \rightarrow \text{sorted}$

Ex.  $[1, 3, 5, 7]$   
 $[0, 2, 4, 6]$

or

$\text{hashmap}[w_l] = [i_1, i_2, i_3, i_4, \dots]$

$\text{hashmap}[w_r] = [j_1, j_2, j_3, \dots]$

Observe that:

for  $j_1$ ,  $i$  minimize  $\text{abs}(i - j_1)$

can only be the  $\max(i)$   
 $i < j_1$

and  $\min(i)$   
 $i > j_1$

Visualization

$i_{\max} (j_1) \quad i_{\min}$

Thus, when we encounter  $i > j_1$ ,  
we know that every  $i$  later in the  
array will never give us smaller  
 $\text{abs}(i - j_1)$

Thus, we should move  $j_1$  to  $j_2$ ,

$\Rightarrow O(n) !$