

Trees

Special Binary Trees

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Outline

Basic Tree Concepts

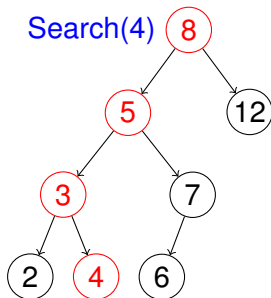
Binary Trees

Special Binary Trees

- Binary Search Trees

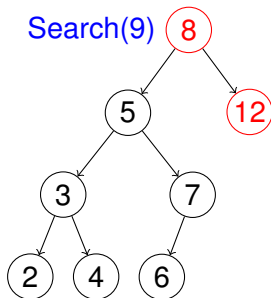
- Heap and Priority Queues

Binary Search Trees (BST)



- For a node whose value is K ,
 - all nodes in its left subtree have key values $< K$, and
 - all nodes in its right subtree have key values $\geq K$
- LNR Traversal \Rightarrow ascending order
- RNL Traversal \Rightarrow descending order
- Search a value along the path of the tree

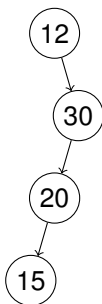
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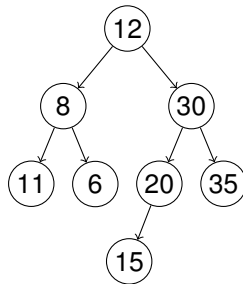
Types of BST

- **Degenerate**: only one child
- **Balanced**: mostly two children



Degenerate BST

Searching: $O(n)$



Balanced BST

Searching: $O(\log n)$

BST Insertion

Algorithm 1: Insert Node whose value is K into BST

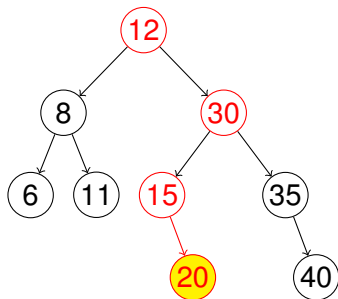
Input: New node whose value is K

Output: a BST with new node inserted

```
1  N ← root
2  while true do
3      if N's value ≤ K then
4          if N's right child is NULL then
5              N's right child ← New node
6              Break
7          else
8              N ← N's right child
9          end
10     else
11         if N's left child is NULL then
12             N's left child ← New node
13             Break
14         else
15             N ← N's left child
16         end
17     end
18 end
19 return
```

BST Insertion Example

- Insert(20)



- $12 < 20$, right
- $30 > 20$, left
- $15 < 20$, right

BST Node Deletion

Algorithm 2: Delete Node whose value K in BST

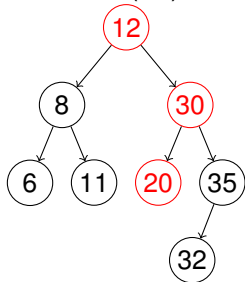
Input: Value K

Output: a BST node whose value is K is removed

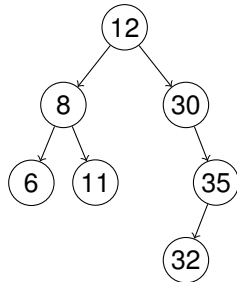
```
1 Find a node whose value is K
2 if found then
3   N  $\leftarrow$  found node
4   if N is leaf then
5     delete node N
6   else
7     L  $\leftarrow$  leftmost of N's right subtree
8     N's value  $\leftarrow$  L's value
9     replace L with its right child and delete node L
10  end
11 return
```

Deleted Node is a leaf

Delete(20)



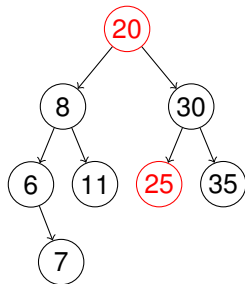
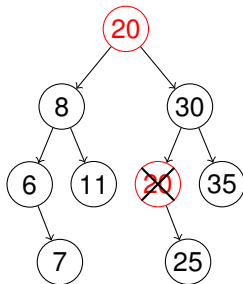
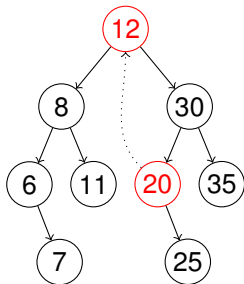
- $12 < 20$, right
- $30 > 20$, left
- $20 = 20$, delete



Deleted node is an internal

Delete(12)

- find node whose value is 12 => internal
- find leftmost of its right subtree
- replace value
- delete node/replace with right child of leftmost node

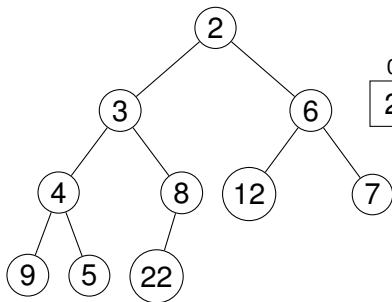


Heaps and Priority Queues

- Normal queues: FIFO
- Priority queues: Highest priority removed first
 - Most critical patient is treated first
 - Highest priority task is executed first
- Heaps are used to implement priority queues

Heaps

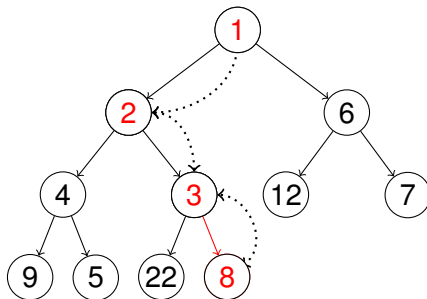
- Complete binary tree
- Partially order
 - Max-heap: parent's value \geq its children's value
 - Min-heap: parent's value \leq its children's value



0	1	2	3	4	5	6	7	8	9	10
2	3	6	4	8	12	7	9	5	22	

Heap Insertion

- Insert(1)



Heap Insertion

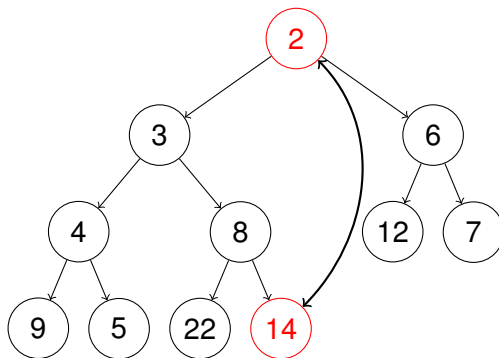
Algorithm 3: Heap Insertion

Input: New value V

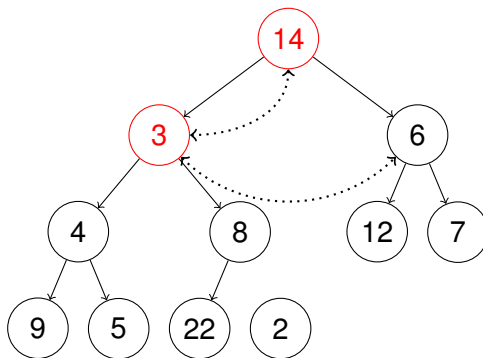
Output: New element of V is inserted into heap

- 1 Increase heap size by 1 to have room for new value;
 - 2 Assign V to new element ;
 - 3 $i \leftarrow$ index of new element ;
 - 4 **while** i is not index of root and $V <$ value at parent of i **do**
 - 5 | swap V and value at parent of i ;
 - 6 | $i \leftarrow$ index of parent of i ;
 - 7 **end**
-

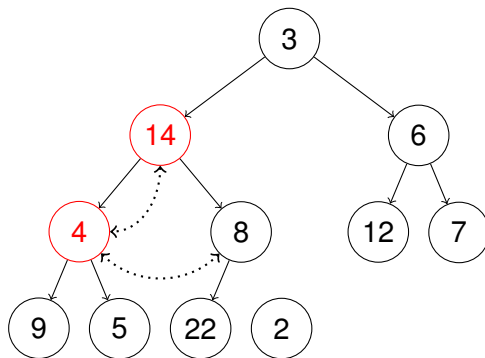
Heap Deletion



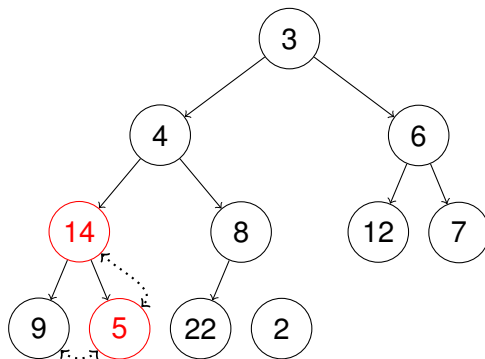
Heap Deletion



Heap Deletion



Heap Deletion



Heap Deletion

