

Data Structure

Monday, June 19, 2023 7:49 AM

Section B - Data Structure - No. Of Questions - 10

Searching

- Linear Search

- Binary Search

Sorting

- Selection Sort

- Bubble Sort

- Insertion Sort

- Merge Sort

- Quick Sort

Stack

- Implementation

- Applications of stack - Expression conversion and evaluation

Queue

- Implementation Linear Queue Circular Queue

Linked List

- All four types - implementation

Tree and Graph - Terminologies

Hash Table

Data Structure:

Algorithm:

Logical flow of given problem statement

Pseudocode - can be represented using pseudocode or

Flowchart - logical flow represent in pictorial format

There are different algorithms are present:

1. Divide and Concur Algorithm
2. In place Algorithm
3. Greedy Algorithm

Program:

Set of instructions

Efficiency:

Can be decided based on timespan and memory consumed by a specific algorithm.

Timespan can be calculated using concept Time Complexity

Memory consumed by any algorithm can be calculated using a concept of Space Complexity

Time Complexity is measured in 2 different ways

1. Asymptotic Time Complexity
 1. Best Case - Omega notation
 2. Average Case - Theta notation
 3. Worst Case - Big Oh notation
2. Symptotic Time Complexity

Space Complexity is calculated based on 2

factors

1. Fixed Space
2. Variable Space

↓

0	1	2	3	4
11	22	33	44	55

key = 11

$a[0] == 11$

comparision count = 1
best time

$O(1)$

↓ ↓ ↓

0	1	2	3	4
11	22	33	44	55

key = 33

$a[0] == 33$ X

$a[1] == 33$ X

$a[2] == 33$ ✓

comparision count = 3 $(n+1)/2$
Average.

① ↓ ② ↓ ③ ↓ ④ ↓ ⑤ ↓

0	1	2	3	4
11	22	33	44	55

key = 55

$a[0] == 55$ X

$a[1] == 55$ X

$a[2] == 55$ X

$a[3] == 55$ X

$a[4] == 55$ ✓

comparision cent = 5 $O(n)$

↓ ↓ ↓ ↓ ↓ X

0	1	2	3	4
11	22	33	44	55

key = 28

$a[0] == 28$ X

$a[1] == 28$ X

$a[2] == 28$ X

$a[3] == 28$ X

$a[4] == 28$ X

Binary Search uses divide and conquer algorithm

One of the fastest searching algorithm

Pre-requisite of binary search algorithm is data has to be in sorted order

If key is present in collection it will be always at mid location

Divide and Conquer algorithm is used by

1. Binary Search
2. Merge Sort
3. Quick Sort

low=0 *mid=4* *high=9* *Key=55*

0	1	2	3	4	5	6	7	8	9
11	22	33	44	55	66	77	88	99	110

l=0 *h=3* *mid=4* *h=9* *Key=33*

0	1	2	3	4	5	6	7	8	9
11	22	33	44	55	66	77	88	99	110

l=0 *h=1* *l=2* *h=3*

0	1	2	3
11	22	33	44

l=0 *h=2* *l=3* *h=3*

0	1	2
11	22	33

int low=0, high=9, mid

```

while(low<=high)
{
    mid = (low+high)/2;
    if(a[mid] == key)
        return mid;
    else if(key < a[mid])
        high = mid -1;
    else if(key > a[mid])
        low = mid + 1;
}
    
```

2	3
33	44

}

key=110

int low=0, high=9, mid

~~low=0~~ mid=4 low=5 high=9

0	1	2	3	4	5	6	7	8	9
11	22	33	44	55	66	77	88	99	110

~~low=5~~ mid=7 low=8 high=9

5	6	7	8	9
66	77	88	99	110

~~mid=8~~ low=9 high=9

8	9
99	110

mid=9
high=9
low=9

9
110

```
while(low<=high)
{
    mid = (low+high)/2;
    if(a[mid] == key)
        return mid;
    else if(key < a[mid])
        high = mid -1;
    else if(key > a[mid])
        low = mid + 1;
}
```

In case of selection sort lowest element of collection is placed in order after completion of first pass

0	1	2	3	4
33	22	11	25	10

Selection sort algorithm needs to execute n-1 passes

1st pass ===== n-1 comparisons == 5-1 = 4 ✓
 2nd pass ===== n-2 comparisons == 5-2 = 3
 3rd pass ===== n-3 comparisons == 5-3 = 2
 4th pass ===== n-4 comparisons == 5-4 = 1

Given n elements if we apply Selection Sort/
 BubbleSort/MergeSort/Insertion Sort/ Quick sort then after
 completion of specific pass what will be state of an
 collection

i=0 j=1

0	1	2	3	4
33	22	11	25	10

i < 4 ✓
 j < 5 ✓

0	1	2	3	4
22	33	11	25	10

i=0 j=2

0	1	2	3	4
22	33	11	25	10

2 < 5 ✓

0	1	2	3	4
11	33	22	25	10

i=0 j=3

0	1	2	3	4
11	33	22	25	10

3 < 5 ✓

i=0 j=4

0	1	2	3	4
11	33	22	25	10

4 < 5

0	1	2	3	4
10	33	22	25	11

j=5

5 < 5

i=1 j=2

0	1	2	3	4
10	33	22	25	11

1 < 5-1

2 < 5 ✓

i=1 j=3

0	1	2	3	4
10	22	33	25	11

3 < 5 ✓

i=1 j=4

0	1	2	3	4
10	22	33	25	11

4 < 5 ✓

i=1 j=5

0	1	2	3	4
10	11	33	25	22

j=5

5 < 5 X

```
for(i=0; i < size-1; i++)
{
    for(j=i+1; j < size; j++)
    {
        if(p[i] > p[j])
        {
            p[i] = p[i] + p[j];
            p[j] = p[i] - p[j];
            p[i] = p[i] - p[j];
        }
    }
}
```

pass. select element
 compare with remaining els.
 swap
 a = a + b
 b = a - b
 a = a - b

a = 5 b = 7
 a = 5 + 7 = a = 12
 b = 12 - 7 = b = 5 ✓
 a = 12 - 5 = a = 7 ✓

0	1	2	3	4
10	11	33	25	22

i=2 j=3

$2 < 5 - 1$
 $3 < 5$

0	1	2	3	4
10	11	25	33	22

0	1	2	3	4
10	11	25	33	22

i=2 j=4

$4 < 5 \checkmark$

0	1	2	3	4
10	11	22	33	25

$j=5$

$5 < 5 \times$

$i=3 \quad 3 < 5 - 1$

0	1	2	3	4
10	11	22	33	25

i=3 j=4

$j=4 \quad 4 < 5 \checkmark$

0	1	2	3	4
10	11	22	25	33

$j=5$

$5 < 5 \times$

$i=4 \quad 4 < 5 - 1 \times$

In case of bubble sort highest element from collection will be place in order

Bubble sort algorithm executes n-1 passes

Time Complexity = $(n-1) + (n-2) + (n-3) \dots$

```
int i, j;
for(i=0; i<size-1; i++)
{
    for(j=0; j<size-1-i; j++)
    {
        if(p[j] > p[j+1])
        {
            p[j] = p[j] + p[j+1];
            p[j+1] = p[j] - p[j+1];
            p[j] = p[j] - p[j+1];
        }
    }
}
```

0	1	2	3	4
33	22	11	25	10

$j=0, j+=1$

0	1	2	3	4
33	22	11	25	10

$i=0$ 1st pass
 $0 < 5-1$

$j < 5-1-0$
 $0 < 5-1-0 \checkmark$

0	1	2	3	4
22	33	11	25	10

$j=1, j+=1$

0	1	2	3	4
22	33	11	25	10

$1 < 5-1-0 \checkmark$

$j=2, j+=1$

0	1	2	3	4
22	11	33	25	10

$2 < 5-1-0 \checkmark$

$j=3, j+=1$

0	1	2	3	4
22	11	25	33	10

$3 < 5-1-0 \checkmark$

$j=4$

0	1	2	3	4
22	11	25	10	33

$5 < 5-1-0 \times$



```
int i, j;
int temp;
for(i=1; i<n; i++)
{
    ✓temp = arr[i];
    for(j=i-1; j>=0 && arr[j] > temp; j--)
        arr[j+1] = arr[j];
    → arr[j+1] = temp;
}
```

temp
1

j=0 i=1

0	1	2	3	4	5	6	7
7	1	15	6	3	34	12	23

j=1 j=0

0	1	2	3	4	5	6	7
7	1	15	6	3	34	12	23

temp
15

j=1 i=2

0	1	2	3	4	5	6	7
1	7	15	6	3	34	12	23

temp
6

j=0 j=1 j=2 i=3

0	1	2	3	4	5	6	7
1	7	15	6	3	34	12	23

1 6 7 15 3 34 12 23

temp
3

j=0 j=1 j=2 j=3 i=4

0	1	2	3	4	5	6	7
1	6	7	15	3	34	12	23

1 3 6 7 15 34 12 23

temp
34

j=4 i=5

0	1	2	3	4	5	6	7
1	3	6	7	15	34	12	23

temp
12

j=3 j=4 j=5 i=6

0	1	2	3	4	5	6	7
1	3	6	7	15	34	12	23

1 3 6 7 12 15 34 23

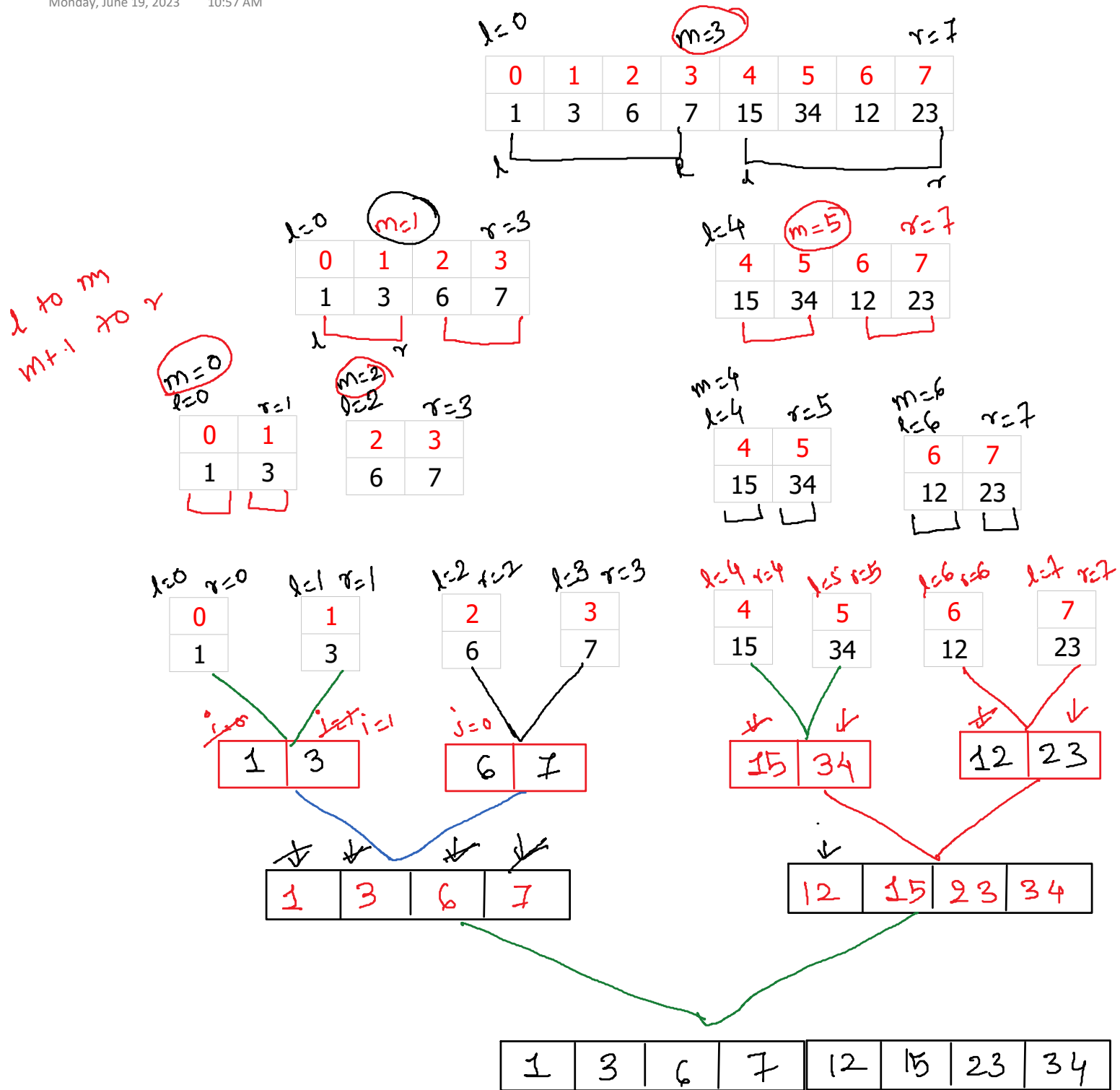
temp
23

j=5 j=6 i=7

0	1	2	3	4	5	6	7
1	3	6	7	12	15	34	23

1 3 6 7 12 15 23 34

i=8



$i=1$ $i=2$ $j=4$ $j=5$ $j=6$ $j=7$

0	1	2	3	4	5	6	7
7	1	15	6	3	34	12	23
\uparrow	\uparrow	\uparrow		\uparrow	\uparrow	\uparrow	\uparrow

$i=2$ $j=6$ $i=4$ $j=4$

0	1	2	3	4	5	6	7
7	1	3	6	15	34	12	23
\uparrow		\uparrow	\uparrow	\uparrow			

$j=6$

0	1	2	3	4	5	6	7
7	1	3	6	15	34	12	23

0	1	2	3	4	5	6	7
6	1	3	7	15	34	12	23