

Section II: Heat Equation Solution using alpaka

PLASMA-PEPSC Workshop on Alpaka and OPENPMD

October 14, 2024





Overview

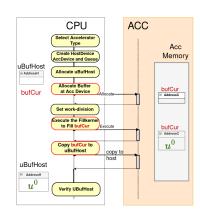
- 1 Steps filling a buffer in parallel
- 2 Memory Allocation and Passing to Kernel
- 3 Define InitilizeBuffer Kernel and Execute
- 4 Introduction for Heat Equation
- 5 Main Simulation Loop of Heat Eqn.
- 6 Heat Eqn. Domain and Stencil
- 7 Setting up the stage to run kernels
- 8 Optimization of Heat Eqn. Solution
- 9 Recap





Steps of Filling a Buffer in Parallel

- Select the accelerator
- Create host-device, acc-device and the queue
- 3 Allocate host and device memory
- Decide how to parallelize: set work-division
- Decide where will the parallel and non-parallel parts of the code run
- 6 Create the kernel instance and execute kernel
- Copy the result from Acc (e.g GPU) back to the host buffer.







Allocate memory at Host and at Device

Define number of dim and index type

```
using Dim = alpaka::DimInt<2u>; // Number of dim: 2 as a type
using Idx = std::size_t; // Index type of the threads and buffers
```

Define domain and halo extents

```
// alpaka::Vec is a static array similar to std::array.
// Dim is a compile—time constant, which is 2.
// Create a static array of size Dim.

constexpr alpaka::Vec < Dim, Idx > numNodes {64, 64};
constexpr alpaka::Vec < Dim, Idx > haloSize {2, 2};
constexpr alpaka::Vec < Dim, Idx > extent = numNodes + haloSize;
```

Allocate memories at host and accelerator

```
// Allocate memory for host—buffer
auto uBufHost = alpaka::allocBuf <double, Idx>(devHost, extent);

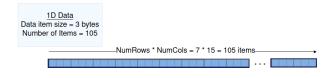
// Allocate memory for accelerator buffer
auto uBufAcc = alpaka::allocBuf <double, Idx>(devAcc, extent);
```

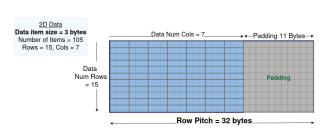




Allocated area at the memory

Let's assume that 105 item with 3-byte each will be allocated to pass to the kernel. The pitch value (actually the row-pitch) depends on the GPU or CPU type.



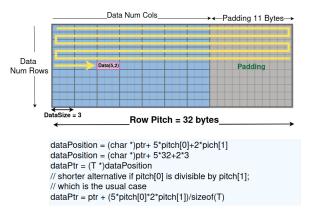






How to access data given the pointer and the pitch?

How to access data at index (5,2) given the pointer ptr and pitch? pitch = {32bytes,3 bytes} as {row-pitch, datasize}







Passing multi dimensional buffer to the kernel

■ Pass 3 variables for a buffer: pointer, row-pitch, and datasize

Multi-dimensional memory allocated in memory uses aligned rows. Hence, if a pointer of a 2D buffer is passed to the kernel as a pointer; 2 additional values **pitch** and item **data-size** should also be passed.

```
// Signature of function operator of the Kernel
templatetypename TAcc, typename TDim, typename TIdx>
ALPAKA_FN_ACC auto operator()(

TAcc const& acc,
double* const bufData,
// 2 variables row-pitch and data-type size
alpaka::Vec<TDim, TIdx> const pitch,
double dx,
double dy) const -> void
```

■ Simple Alternative: Pass an alpaka mdspan object

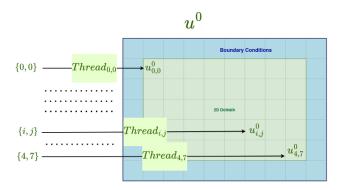
```
template<typename TAcc, typename TDim, typename TIdx, typename TMdSpan>
ALPAKA_FN_ACC auto operator()(
TAcc const& acc,
TMdSpan uAccMdSpan
...) const -> void
```





The Kernel to Initialize Heat Values

Calculate and set initial heat values, the u^0 matrix, by running a grid of threads.







The Kernel to Initialize Heat Values

The InitializeBufferKernel fills the given buffer at the accelerator device (e.g GPU) Prepare kernel to set initial heat values

- Thread Index: Find thread index in the kernel to be used as index to set 2D buffer.
- Initial Condition at the point: Find analytically the heat value at the point which has coordinates equal to the 2D thread index.
- Memory Adress in Buffer: Calculate the corresponding memory adress in buffer using thread index. Take into account row-pitch and data-size
- Set Value at the Adress: Set the data at the memory position to the calculated initial condition.



Hands-On Session

Hands-on Session3: Filling an accelerator buffer paralelly







The Heat Equation

■ The heat equation models the Heat Diffusion over time in a given medium.

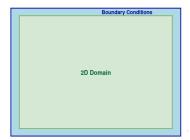
$$\frac{\partial u(x, y, t)}{\partial t} = \alpha \left(\frac{\partial^2 u(x, y, t)}{\partial x^2} + \frac{\partial^2 u(x, y, t)}{\partial y^2} \right)$$

Difference approximations for Time and Spatial Derivatives:

$$\left. \frac{\partial u(x,y,t)}{\partial t} \right|_{t=t^n} \approx \frac{u_{i,j}^{n+1} - u_{i,j}^n}{\Delta t} \qquad \left. \frac{\partial^2 u(x,y,t)}{\partial x^2} \right|_{x=x_i,y=y_j} \approx \frac{u_{i+1,j}^n - 2u_{i,j}^n + u_{i-1,j}^n}{\Delta x^2}$$

Resulting difference equation:

$$u_{i,j}^{n+1} = u_{i,j}^n + \alpha \Delta t \left(\frac{u_{i+1,j}^n - 2 u_{i,j}^n + u_{i-1,j}^n}{\Delta x^2} + \frac{u_{i,j+1}^n - 2 u_{i,j}^n + u_{i,j-1}^n}{\Delta y^2} \right)$$





The Heat Equation- Cont.

■ The difference equation:

$$u_{i,j}^{n+1} = u_{i,j}^n + \alpha \Delta t \left(\frac{u_{i+1,j}^n - 2u_{i,j}^n + u_{i-1,j}^n}{\Delta x^2} + \frac{u_{i,j+1}^n - 2u_{i,j}^n + u_{i,j-1}^n}{\Delta y^2} \right)$$

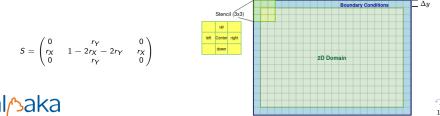
■ Substitute: $\alpha = 1$, $r_X = \frac{\Delta t}{\Delta x^2}$, $r_Y = \frac{\Delta t}{\Delta y^2}$

Then $u_{i,j}^{n+1}$ is:

$$u_{i,j}^{n+1} = u_{i,j}^{n} + r_{X} \left(u_{i+1,j}^{n} - 2u_{i,j}^{n} + u_{i-1,j}^{n} \right) + r_{Y} \left(u_{i,j+1}^{n} - 2u_{i,j}^{n} + u_{i,j-1}^{n} \right)$$

By regrouping the terms related to $u_{i,j}^n$, the equation can be rewritten as:

$$u_{i,j}^{n+1} = u_{i,j}^{n} \left(1 - 2r_{X} - 2r_{Y}\right) + r_{X} \left(u_{i+1,j}^{n} + u_{i-1,j}^{n}\right) + r_{Y} \left(u_{i,j+1}^{n} + u_{i,j-1}^{n}\right)$$





Main Simulation Loop: Leveraging Parallelism

Initialization:

- Define the "host device" and "accelerator device". The "Host" and "Device" in short.
- Set initial conditions and boundary conditions.
- Allocate data buffers to host and device.
- Copy data from host to device buffer to pass to the kernel.
- Define parallelisation strategy (determine block size).

Simulation Loop:

- Step 1: Execute StencilKernel to compute next values.
- Step 2: Apply boundary conditions using BoundaryKernel.
- Step 3: Swap buffers for the next iteration so that calculated $u_{i,j}^{n+1}$ becomes the $u_{i,j}^{n}$ for the next step.

Parallel Efficiency:

- Subdomains are processed in parallel, with halos ensuring data consistency and correct boundary conditions.
- Optimization: Shared memory optimizes memory access within each block using chunks of data.

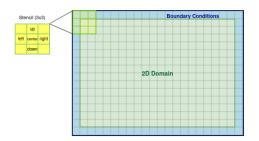
Validation





Parallel Heat Equation Solution

- Data Parallelism: Each point on the grid can be updated independently based on its neighbors, enabling parallel computation.
- Stencil Operations: Stencil is a core computational pattern in PDE solvers. Updates a grid point in time
 using its immediate neighbors (left, right, up, down) according to the difference equation. A 5-point stencil
 is needed.



- Halo Region for BC: A layer of grid cells surrounding the problem domain for Boundary Conditions.
 - Facilitates stencil operations at the boundaries of subdomains.



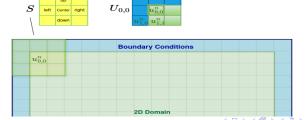


Calculation of $u_{i,j}^{n+1}$ from $u_{i,j}^n$

- lacksquare Each kernel execution by alpaka calculates $u_{i,j}^{n+1}$ using $u_{i,j}^n$
- Each heat point is separately calculated by a thread using Frobenious Inner Product (FIP)
- The Frobenius Inner Product between matrix S and matrix $U_{i,j}$ is:

$$u_{i,j}^{n+1} = \langle S, U_{i,j}^n \rangle_F = \sum_{m=1}^M \sum_{k=1}^K s_{m,k} u_{m,k}$$

■ S and $U_{0,0}^n$ is used by a thread to calculate $u_{0,0}^{n+1}$ using FIP







■ Another thread calculates $u_{0,1}^{n+1}$ using S and $U_{0,1}^{n}$

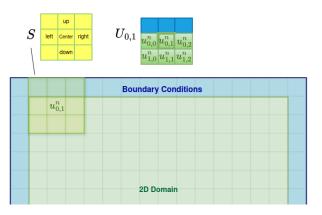
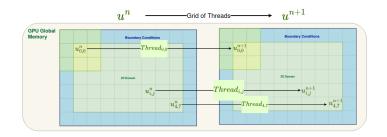


Figure: Second thread calculates $u_{0,1}^{n+1}$ using





Stencil Kernel Execution by a grid of threads

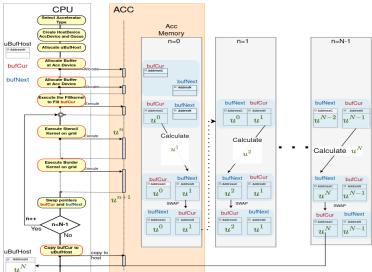


- Stencil kernel will update only core nodes not the border
- The workdiv for stencil kernel can be calculated by setting gridthread extent to nodes domain
- The workdiv for the borders kernel would need extended extents, because halo is going to updated as well





Complete Heat Equation solution







The Stencil Kernel

What kind of parallelization needed to calculate $u_{i,j}^{n+1}$ using $u_{i,j}^{n}$

- StencilKernel needs a Work-division to work on domain of all nodes (without halo)
- Boundary kernel needs a Workdivision which covers nodes + halo region
- Boundary kernel will only use threads corresponding to halo region

Coding StencilKernel

- Input: Check the size of the input buffer, it should include halo region as well
- Thread Index: Find thread index in the kernel. This index will be used as the center of 3x3 stencil.
- Memory Adress in Buffer: Calculate the corresponding memory adress in buffer using thread index. Take into account pitch and data-size
- **Calculate new heat value:** Calculate $u_{i,i}^{n+1}$ using Frobenious Inner Product of 3x3 matrices
- Set Value at the Adress: Set the data at the memory adress.

```
struct StencilKernel
{
    template < typename TAcc, typename TDim, typename TIdx>
    ALPAKA_FN_ACC auto operator()(
        TAcc constk acc,
        double const* const uCurrBuf, double* const uNextBuf,
        alpaka::Vec<TDim, TIdx> const chunkSize,
        alpaka::Vec<TDim, TIdx> const pitchCurr, alpaka::Vec<TDim, TIdx> const pitchCurr, alpaka::Vec<TDim, Tidx> const pitchCurr, alpaka::Vec<TDim, Tidx> const pitchNext,
        double const dx, double const dy, double const dt) const -> void
    {
            ...
        }
}
```



Hands-On Session

Hands-on Session4: Stencil Kernel







alpaka Basics

Setting up the stage to run kernels

- Selecting the accelerator and host devices
- Allocating and setting host and accelerator device memory
- 3 Alpaka Vector, Buffer and View?
- Passing data to the accelarator
- 5 WorkDiv
- 6 Define Queue





Accelerator, Device and Host

Define number of dim and index type

```
using Dim = alpaka::DimInt<2u>; // Number of dim: 2 as a type
using Idx = std::size_t; // Index type of the threads and buffers
```

Define the accelerator

```
// AccGpuCudaRt, AccGpuHipRt, AccCpuThreads, AccCpuSerial,
// AccCpuOmp2Threads, AccCpuOmp2Blocks, AccCpuTbbBlocks
using Acc = alpaka::AccGpuHipRt < Dim, Idx >;
using DevAcc = alpaka::Dev < Acc >;
```

Select a device from platform of Acc

```
auto const platform = alpaka::Platform < Acc > {};
auto const devAcc = alpaka::getDevByIdx(platform, 0);
```

Select a host and hostype to allocate memory for data

```
// Get the host device for allocating memory on the host.
auto const platformHost = alpaka::PlatformCpu{};
auto const devHost = alpaka::getDevByIdx(platformHost, 0);
// Host device type is needed, still not known
using DevHost = alpaka::DevCpu;
```





What is Accelerator

Accelerator hides hardware specifics behind alpaka's abstract API

 On Host: Accelerator is a type. A Meta-parameter for choosing correct physical device and dependent types

```
using Acc = acc::AccGpuHipRt < Dim, Idx>;
```

- Inside Kernel: Accelerator is a variable. Contains thread state, provides access to alpaka's device-side API
 - The Accelerator provides the means to access to the indices

```
// get thread index on the grid
auto gridThreadIdx = alpaka::getIdx<Grid, Threads>(acc);
// get block index on the grid
auto gridBlockIdx = alpaka::getIdx<Grid, Blocks>(acc);
```

 The Accelerator gives access to alpaka's shared memory (for threads inside the same block)

```
// allocate a variable in block shared static memory
auto& sdata = alpaka::declareSharedVar<double[T_SharedMemSize1D], __COUNTER__>
acc);

// get pointer to the block shared dynamic memory
auto* const sharedN = alpaka::getDynSharedMem<float>(acc);
```

■ Enables synchronization on the block level

alpaka::syncBlockThreads(acc);

synchronize all threads within the block

```
al/saka
```



What is alpaka Buffer, Vector and View

- **alpaka::Buf** is multi-dimensional dynamic (runtime sized) container.
 - Contains memory adress, extent, datatype and the device that memory belongs to!
 - Since buffer already knows the it's device and extent; device to device copy is easy in alpaka
 - Supports [] operator but not [][].

```
// Allocate buffers
auto bufCpu = alpaka::allocBuf<float, Idx>(devCpu, extent);
auto bufGpu = alpaka::allocBuf<float, Idx>(devGpu, extent);
....
// Copy buffer from CPU to GPU — destination comes first
alpaka::memcopy(gpuQueue, bufGpu, bufCpu);
// cuda way: cudaMemcpy(b_d, b_host, sizeof(float)*N, cudaMemcpyHostToDevice)
```

■ alpaka::Vec is a static 1D array.

```
alpaka::Vec<SizeOfArrayAsType,DataT> myVec;
```

 alpaka::View is a non-owning view to an already allocated memory, so that it can be used in alpaka::memcpy





What is alpaka::Queue

- alpaka::Queue is "a queue of tasks".
- Used for sycnhronization of tasks like memcpy or kernel-execution
- Queue is always FIFO, everything is sequential (in-order) inside the alpaka::queue
- If there is a second queue, queue feature "blocking" and "non-blocking" becomes important
- Different queues can run in parallel for many devices
- Within a single queue accelerator back-ends can be mixed (used in interleaves)

```
using QueueProperty = alpaka::NonBlocking;
// Create queue
using QueueAcc = alpaka::Queue<Acc, QueueProperty>;

QueueAcc computeQueue(devAcc);
// Copy host -> device, use the queue
alpaka::memcpy(computeQueue, uCurrBufAcc, uBufHost);
alpaka::wait(computeQueue, uCurrBufAcc, uBufHost);
// Create kernel instance
stencilKernel
// Create kernel instance
StencilKernel
// Execute kernel using queue
// // Execute kernel using queue
alpaka::exec
// Execute kernel using queue
// Create kernel using Account (ComputeQueue, workDiv_manual, stencilKernel...)
```



Queue Operations

Queues are used for synchronization

```
// block caller until all operations have completed
alpaka::wait(myQueue);
// block myQueue until myEvent has been reached
alpaka::wait(myQueue, myEvent);
// block myQueue until otherQueue has completed
alpaka::wait(myQueue, otherQueue);
```

Queues can be checked for completion of all tasks

```
bool done = alpaka::empty(myQueue);
```





Tasks and Events

- Device-side related operations (kernels, memory operations) can be wrapped in tasks.
- Tasks are executed by **enqueue()** function.
- Tasks on the same queue are executed in order (FIFO principle)

```
alpaka::enqueue(queueA, task1);
alpaka::enqueue(queueA, task2);
// task2 starts after task1 has finished, even if queueA is non-blocking
```

Order of tasks in different queues is unspecified

```
alpaka::enqueue(queueA, task1);
alpaka::enqueue(queueB, task2);
// task2 starts before, after or in parallel to task1 if queueA is non blocking
```

For easier synchronization, alpaka Events can be inserted before, after or between Tasks:

```
auto myEvent = alpaka::Event<alpaka::Queue>(myDev);
alpaka::enqueue(queueA, myEvent);
alpaka::wait(queueB, myEvent);
// queueB will only resume after queueA reached myEvent
alpaka:wait(queueB, myEvent);
// queueB will only resume after queueA reached myEvent
alpaka:wait(queueB, myEvent);
// queueB will only resume after queueA reached myEvent
alpaka:wait(queueB, myEvent);
// queueB will only resume after queueA reached myEvent
```



Execute the Kernel and copy data back to host

First create the queue

```
// Create queue,
// queue is needed for kernel execution and copies to/from accelerator
alpaka::Queue<Acc, alpaka::NonBlocking> queue{devAcc};
```

■ Execute the kernel using the queue, the workdiv and kernel arguments:

```
alpaka::exec<Acc>(queue, workDiv, initBufferKernel, uBufAcc.data(),
    pitchCurrAcc, dx, dy);
```

Copy the filled buffer back to the host

```
// Copy device -> host
// Since buffers know their corresponding devices (host or acc) memcopy does
not need any device variable
alpaka::memcpy(queue, uBufHost, uBufAcc);
alpaka::wait(queue);
```





Determine WorkDiv

Let alpaka calculate work division for you:

```
// All kernel inputs are needed because work—division depends on the kernel
// Create kernel instance
InitializeBufferKernel initBufferKernel;
// Elements per thread needed to determine work—div
constexpr alpaka::Vec<Dim, Idx> elemPerThread(1, 1);
// Bundle the extent and elements per thread
alpaka::KernelCfg<Acc> const kernelCfg = {extent, elemPerThread};
// Kernel input row—pitch and data—type—size
auto const pitchCurrAcc{alpaka::getPitchesInBytes(uBufAcc)};
// Determine the work—div
auto workDiv = alpaka::getValidWorkDiv(kernelCfg, devAcc, initBufferKernel, uBufAcc.data(), pitchCurrAcc, dx, dy);
```





Determine the work-division manually

```
// Set Dim and Index type
using Dim1D = alpaka::DimInt<1u>; // 1 as a type
using Idx = uint32_t;
// M blocks each has 4 threads, each level is 1D
alpaka::WorkDivMembers<Dim1D, Idx> workdiv1D{M, 4, 1};
// Set Dim2D and use same index type
using Dim2D = alpaka::DimInt<2u>; // 2 as a type
alpaka::Vec<Dim2D, Idx> gridBlockExtent{M,N}; // 2D grid
alpaka::Vec<Dim2D, Idx> blockThreadExtent{32,32}; // 2D block
alpaka::Vec<Dim2D, Idx> elementExtentPerThread{1,1};
// MxN blocks each has 32x32 threads, each level is 2D
alpaka::WorkDivMembers<Dim2D, Idx> workdiv2D{gridBlockExtent,
blockThreadExtent, elementExtentPerThread}:
```

```
using DimiD = alpaka::DimInt<|2;//Set number of dims to 1
using Vec1D = alpaka::Vec<DimID, Idx>;//Define alias
auto workbiv1D = alpaka::MorkbivMembers(Vec1D{M}, Vec1D{Au}, Vec1D{Au});
// alternatively
using Dim3D = alpaka::DimInt<3;//Set number of dims to 3
using Vec3D = alpaka::Vec<Dim3D, Idx>; //Define alias
auto workbiv3D = alpaka::WorkDivMembers(Vec3D{1,1,M}, Vec3D{1,1,4u}, Vec3D{1,1,1u});
```





Optimization and usability

alpaka Usability and Optimization Features

- Use alpaka mdspan to set,get, pass buffers easily
- Use Domain Decomposition: Divide the domain in chunks
- Use 2 asynch queues for performance increase
- 4 Use **shared memory** for performance increase





alpaka::experimental::mdspan

Mdspan a multi-dimensional and non-owning view

- Part of C++23 standard. Can be used with C++17.
- Consists Data Pointer, Data Pitch and Data item size
- Has member functions to get/set data and to get extents

Mdspan Installation

- Set alpaka_USE_MDSPAN cmake variable to FETCH while installing alpaka
- Alternatively, set alpaka_USE_MDSPAN cmake variable to FETCH while configuring example if it is not already set while installation

```
// in build directory
cmake -Dalpaka_USE_MDSPAN=FETCH ..
```

Create an mdspan view of a buffer, then pass to the kernel

```
// Host code: Allocate device memory
auto bufDevA = alpaka::allocBuf < DataType, Idx > (devAcc, extentA);
// Create mdspan views for device buffers using alpaka::experimental::getMdSpan
auto mdDevA = alpaka::experimental::getMdSpan(bufDevA);
// Execute the kernel
alpaka::exec < Acc > (queue, workDiv, kernel, mdDevA, mdDevB, mdDevC);
```





Kernel using mdspan instead of data buffer pointer and pitch

Example usage to access/set multi-dim data at host or in kernel

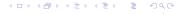
```
struct MatrixAddKernel
        template < typename TAcc, typename TMdSpan >
         ALPAKA_FN_ACC void operator()(TAcc const& acc, TMdSpan A, TMdSpan B, TMdSpan
               C) const
         ſ
             auto const i = alpaka::getIdx<alpaka::Grid, alpaka::Threads>(acc)[0];
             auto const j = alpaka::getIdx<alpaka::Grid, alpaka::Threads>(acc)[1];
             if(i < C.extent(0) && i < C.extent(1))
8
             ſ
                      C(i, j) = A(i, j) + B(i, j);
9
         } }:
1
   struct StencilKernel
3
       template < typename TAcc, typename TDim, typename TIdx, typename TMdSpan>
        ALPAKA_FN_ACC auto operator()(
            TAcc const& acc,
            TMdSpan uCurrBuf.
            TMdSpan uNextBuf.
            alpaka::Vec < TDim , TIdx > const chunkSize ,
            double const dx.
            double const dy,
            double const dt) const -> void
```



Hands-On Session

Hands-on Session5: Use mdspan for the kernel using shared memory

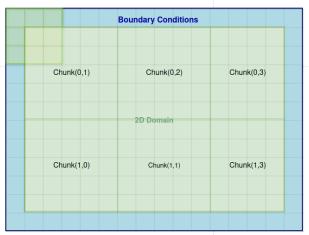






Chunk Definition

Chunk: Subdomains needed for latency management of block level parallelisation

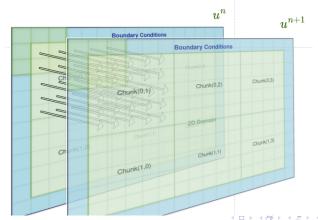






Calculation by a block of grids of stencil kernel

- Chunking is a domain decomposition method
- A block of threads update a chunk of heat data
- A grid of threads updates the whole domain

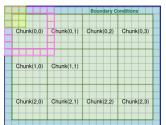


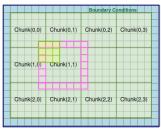




Chunks in Parallel Grid Computations - I

- Halo Region around chunk: A layer of grid cells surrounding the subdomains. In order to use the heat value beside the current chunk
- Halo Size: Typically 1 for a 5-point stencil.
- Chunks might include more than one blocks depending on the blocksize

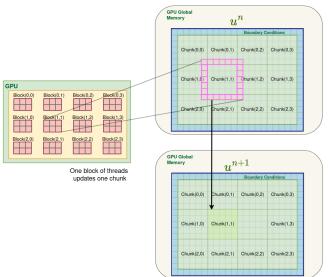








Chunks in Parallel Grid Computations - II







Determine WorkDiv for Chunked Solution

Set work division fields directly:

```
// Define a workdiv for the shared memory based heat egn solution
    constexpr alpaka::Vec < Dim , Idx > elemPerThread {1, 1};
    // Get max threads that can be run in a block for this kernel
    auto const kernelFunctionAttributes = alpaka::getFunctionAttributes<Acc>(
        devAcc.
        stencilKernel.
        uCurrBufAcc.data(), uNextBufAcc.data(),
        chunkSize.
        pitchCurrAcc, pitchNextAcc,
9
        dx, dy, dt);
10
    auto const maxThreadsPerBlock = kernelFunctionAttributes.maxThreadsPerBlock:
    auto const threadsPerBlock
        = maxThreadsPerBlock < chunkSize.prod() ? alpaka::Vec < Dim , Idx > {
             maxThreadsPerBlock, 1} : chunkSize:
    alpaka::WorkDivMembers < Dim. Idx > workDiv manual \( \) numChunks. threadsPerBlock.
14
         elemPerThreadl:
```





Hands-On Session

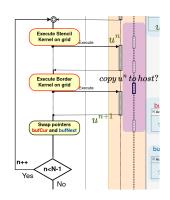
Hands-on Session6: Optimized Heat Eqn. solution by Domain Decomposition





Running 2 parallel queues: Additional queue to dump temporary results

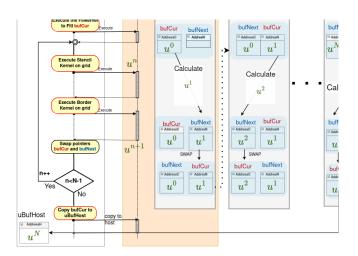
- Create an additional alpaka::queue instance at accelerator to run parallely
- The temporary heat result uⁿ will be copied to host from accelerator at the end of each iteration
- Copying can start while the stencil and boundary kernels are running
- In order to run 2 queues paralelly they should be a NonBlocking queues
- The copied heat data will be used to create an animation of images







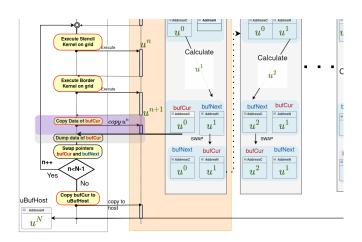
Current Loop







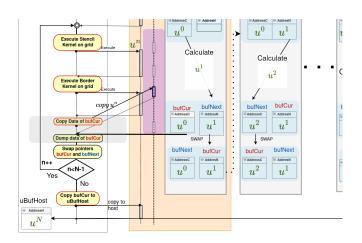
Copy u^n back at each iteration sequentially







Copy u^n back at each iteration in parallel







Hands-On Session

Hands-On Session7: Running 2 parallel queues







Efficient Stencil Application with Shared Memory

Shared Memory at GPUs

- A fast, limited-size memory accessible by all threads within a block.
- Used to store data locally in Compute Unit(or SM), reducing the need to access slower global memory.
- Shared Memory allocation can be static or dynamic
 - Static (compile time determined extent)
 - Dynamic (runtime determined extent)
- Filling shared memory is done by the same kernel calculating the stencil
- Threads in a block must synchronize to ensure all data is loaded into shared memory before computation begins.

Benefits:

 Reduces memory latency by storing the working set of data (halo + core) in shared memory.





Shared memory data: chunk with halo

- Halo Region around chunk: A layer of grid cells surrounding the subdomains. In order to use the heat value beside the current chunk
- Halo Size: Typically 1 for a 5-point stencil.
- Chunks might include more than one blocks depending on the blocksize
- Kernel will install the data to shared memory then use the data from shared memory

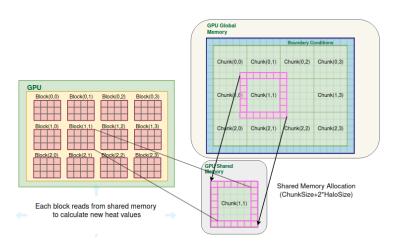








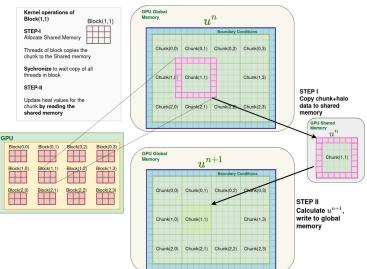
Updating chunks using shared memory







Kernel Operations of a Block to Find u^{n+1} for block data







Steps of Stencil Kernel using shared memory

Allocate shared memory inside kernel

- Calculate thread index
- Fill the shared memory by block of threads
- Wait for shared memory to be filled by all block threads

```
1 alpaka::syncBlockThreads(acc);
```

- Calculate new heat value using the data from the shared memory
- Set the new heat value





Hands-On Session

Hands-on Session8: Optimized Heat Eqn. solution by using shared memory







Conclusion: Parallel Techniques For Solving Heat Equation

Kernel Definition

- Kernel to fill a buffer in parallel
- Stencil Kernel for calculating the next set of heat values
- Boundary Kernel

Work division

- Getting a valid work division according to accelerator
- Setting work-division manually

Allocating and Setting Memory at Host and Accelerator

- Using alpaka::buffer
- Using alpaka::memcpy

Alpaka Structures

- Accelerator, Device, Queue, Task
- Optimizations and Usability
 - Using alpaka Mdspan
 - Domain Decomposition
 - Using Multiple Async Queues
 - Using GPU's Shared Memory





Questions?



