Parallel Functional Programming with Interaction Nets

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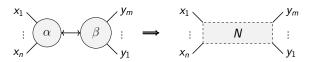
What are interaction nets?

Graph rewriting system (Lafont 1990). "A new kind of programming language"

Finite set of *user-defined* agents:



Finite set of *user-defined* rewrite rules:



Maximum one rule per agent pair.

Set of auxiliary ports preserved.



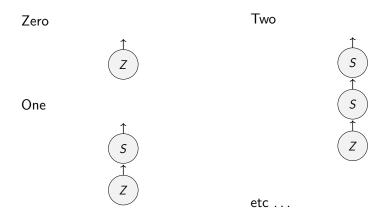
Properties as programming language

- Turing complete
- ▶ Pattern matching
- Constant time rewrite operations
- Visual debugging
- Local reductions
- Diamond property



- Explicit mandatory memory management
 - → Natural parallel execution

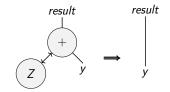
Example constructor - Unary numbers



Example function - Unary number addition

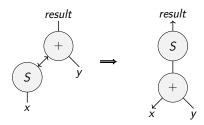
$$Z + y = y$$

add Z y = y



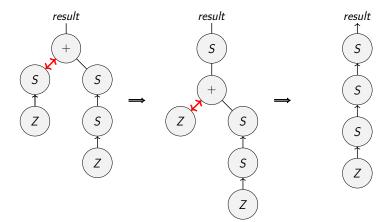
$$(S x) + y = S (x + y)$$

add (S x) y = S (add x y)

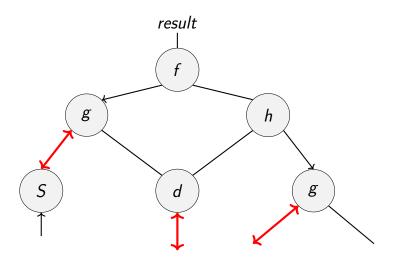


Example function - Unary number addition

$$1 + 2 = 3$$

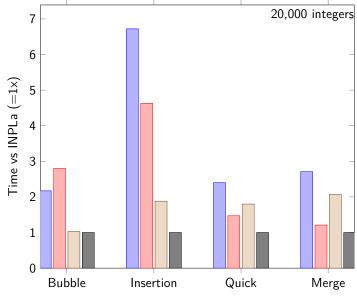


Parallel evaluation



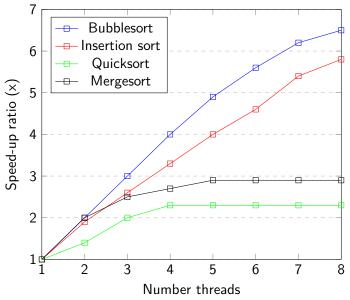
Sequential algorithm \equiv parallel algorithm.

Impact of parallelism - benchmark results

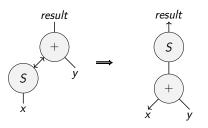


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Impact of parallelism - benchmark results



Towards a programming language

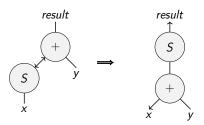


Flatten net¹:

$$add(result,y)> result-S(aux), add(aux,y)-x$$

¹Sato, 2014; https://github.com/inpla/inpla

Towards a programming language



Agents whose principal port acts as input \rightarrow *functions*.

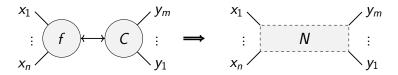


Agents whose principal port acts as output \rightarrow constructors.



FLIN - a Functional Language for Interaction Nets

If f is a function and C is a constructor:



then:

f (C
$$\vec{y}$$
) $\vec{x}' = N \vec{x}' \vec{y}$

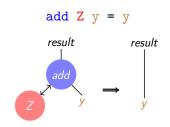
where:

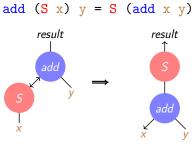
$$N = f \dots \mid C \dots \mid \vec{y} \mid \dots$$

and

 $\vec{x}' = \vec{x}$ adjusted for output.

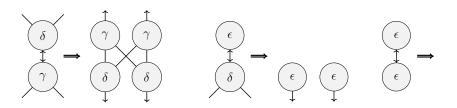
FLIN ≅ Interaction Nets





But not all interaction rules are functions!

Interaction combinators (Lafont 1997)



 δ has two outputs.

 ϵ consumes its input!

None are function-constructor systems!!

FLIN syntax for non-functions

Multiple outputs:

```
delta Z = Z,Z
delta (S x) = let x1,x2 = delta x in <math>(S x1),(S x2)
```

No output or not function-constructor:

```
{eps><delta(a1,a2) => eps~a1, eps~a2}
{eps><eps => }
```

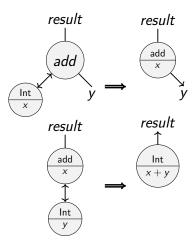
Or rewrite the algorithm!

FLIN examples

```
add Z y = y
add (S x) y = S(add x y)
mult Z y = Z {erase~y}
mult (S x) y = let y1, y2 = dup x in add y1 (mult x y2)
fib Z = Z
fib (S x) = fibS x
fibS Z = S Z
fibS (S x) = let x1,x2=dup x in
                 add (fibS x1) (fib x2)
append [] ys = ys
append (x:xs) ys = x:(append xs ys)
```

Extension - Attributes

Hold values within agents - ints, bools, strings etc & tuples of. (Fernández, Mackie, Pinto 2001)

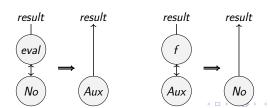


cf. λ -calculus \rightarrow PCF.

Monads

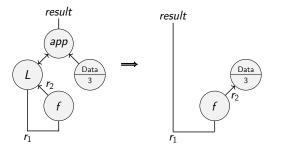
Translate: bind m f \rightarrow f (eval m) e.g. Maybe monad (following Jiresch 2010)

eval Nothing = Aux ; f Aux = Nothing

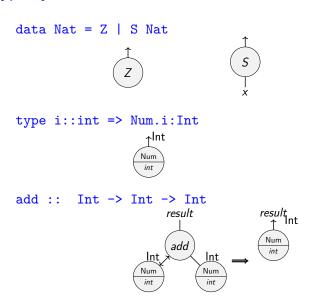


Higher order functions

"Package up" function in a constructor (λ -abstraction):



Type system



Conclusions

- Interaction nets provide inherently parallel evaluation.
- ▶ INPLa implementation has encouraging benchmarks.
- ► FLIN a simple function-constructor language maps 1:1 to interaction nets.
- ► FLIN can encode standard functional programming structures.
- ► FLIN programs run sequentially or in parallel based on resources.
- lacktriangle We have implemented a FLIN ightarrow INPLa transpiler.
- ► Language can be used directly for programming or as an intermediate language for a more complete language.