

The QR iteration method for quasiseparable matrices

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1 Introduction

2 Definitions

Let $\{a_k\}, k = 1, \dots, N$ be a family of matrices of sizes $r_k \times r_{k-1}$. For positive integers i, j , $i > j$ define the operation a_{ij}^\times as follows: $a_{ij}^\times = a_{i-1} \cdots a_{j+1}$ for $i > j + 1$, $a_{j+1,j}^\times = I_{r_j}$.

Let $\{b_k\}, k = 1, \dots, N$ be a family of matrices of sizes $r_{k-1} \times r_k$. For positive integers i, j , $j > i$ define the operation b_{ij}^\times as follows: $b_{ij}^\times = b_{i+1} \cdots b_{j-1}$ for $j > i + 1$, $b_{i,i+1}^\times = I_{r_i}$.

It is easy to see that

$$a_{ik}^\times = a_{ij}^\times a_{j+1,k}^\times, \quad i > j \geq k \quad (2.1)$$

and

$$b_{kj}^\times = b_{k,i+1}^\times b_{i,j}^\times, \quad k \leq i < j. \quad (2.2)$$

Let $R = \{R_{ij}\}_{i,j=1}^N$ be a matrix with block entries R_{ij} of sizes $m_i \times n_j$. Assume that the entries of this matrix are represented in the form

$$R_{ij} = \begin{cases} p_i a_{ij}^\times q_j, & 1 \leq j < i \leq N, \\ d_i, & 1 \leq i = j \leq N, \\ g_i b_{ij}^\times h_j, & 1 \leq i < j \leq N. \end{cases} \quad (2.3)$$

Here p_i ($i = 2, \dots, N$), q_j ($j = 1, \dots, N - 1$), a_k ($k = 2, \dots, N - 1$) are matrices of sizes $m_i \times r'_{i-1}$, $r'_j \times n_j$, $r'_k \times r'_{k-1}$ respectively; these elements are said to be *lower generators of the matrix R with orders r'_k* ($k = 1, \dots, N - 1$). The elements g_i ($i = 1, \dots, N -$

1), h_j ($j = 2, \dots, N$), b_k ($k = 2, \dots, N-1$) are matrices of sizes $m_i \times r''_i$, $r''_{j-1} \times n_j$, $r''_{k-1} \times r''_k$ respectively; these elements are said to be *upper generators of the matrix R* with orders r''_k , ($k = 1, \dots, N-1$). The matrices d_k ($k = 1, \dots, N$) of sizes $m_k \times n_k$ are said to be *diagonal entries* of the matrix R . We define also orders of generators r'_k , r''_k for $k = 0$, N setting them to be zeros. For scalar matrices the generators p_i, g_i and q_j, h_j are rows and columns of the corresponding sizes. Set $n_L = \max_{1 \leq k \leq N-1} r'_k$, $n_U = \max_{1 \leq k \leq N-1} r''_k$, the matrix R is said to be *lower quasiseparable* of order n_L and *upper quasiseparable* of order n_U or *quasiseparable* of order (n_L, n_U) .

Formally, we use some calculation rules with matrices that have blocks with dimension zero. Aside from obvious rules, the product of an “empty” matrix of dimension $m \times 0$ and an empty matrix of dimension $0 \times n$ is a matrix of dimension $m \times n$ with all elements equal to 0. All further rules of block matrix multiplication remain consistent. Such operations are used in MATLAB.

3 The QR factorization

Let $R = \{R_{ij}\}_{i,j=1}^N$ be a matrix with entries from \mathbb{C} with given generators. We present here an algorithm for computing generators and diagonal entries of unitary matrix Q and upper triangular matrix S such that $R = QS$. The main part of the algorithm is based on the following result from [1].

Theorem 3.1 *Let $R = \{R_{ij}\}_{i,j=1}^N$ be a scalar matrix with lower generators p_i ($i = 2, \dots, N$), q_j ($j = 1, \dots, N-1$), a_k ($k = 2, \dots, N-1$) of orders r'_k ($k = 1, \dots, N-1$), upper generators g_i ($i = 1, \dots, N-1$), h_j ($j = 2, \dots, N$), b_k ($k = 2, \dots, N-1$) of orders r''_k ($k = 1, \dots, N-1$) and diagonal entries d_k ($k = 1, \dots, N$). Let us define the numbers ρ_k via recursive relations $\rho_N = 0$, $\rho_{k-1} = \min\{1 + \rho_k, r'_{k-1}\}$, $k = N, \dots, 2$, $\rho_0 = 0$ and the numbers $m_k = 1, n_k = 1, \nu_k = 1 + \rho_k - \rho_{k-1}$, $\rho'_k = r'_k + \rho_k$, $k = 1, \dots, N$.*

The matrix R admits the factorization

$$R = VUS,$$

where V is a unitary matrix represented in the block lower triangular form with blocks of sizes $m_i \times \nu_j$ ($i, j = 1, \dots, N$), lower generators $(p_V)_i$ ($i = 2, \dots, N$), $(q_V)_j$ ($j = 1, \dots, N-1$), $(a_V)_k$ ($k = 2, \dots, N-1$) of orders ρ_k ($k = 1, \dots, N-1$) and diagonal entries $(d_V)_k$ ($k = 1, \dots, N$), U is a unitary matrix represented in the block upper triangular form with blocks of sizes $\nu_i \times n_j$ ($i, j = 1, \dots, N$), upper generators $(g_U)_i$ ($i = 1, \dots, N-1$), $(h_U)_j$ ($j = 2, \dots, N$), $(b_U)_k$ ($k = 2, \dots, N-1$) of orders ρ_k ($k = 1, \dots, N-1$) and diagonal entries $(d_U)_k$ ($k = 1, \dots, N$) and S is an upper triangular matrix with upper generators $(g_S)_i$ ($i = 1, \dots, N-1$), $(h_S)_j$ ($j = 2, \dots, N$), $(b_S)_k$ ($k = 2, \dots, N-1$) of orders ρ'_k ($k = 1, \dots, N-1$) and diagonal entries $(d_S)_k$ ($k = 1, \dots, N$).

The generators and the diagonal entries of the matrices V, U, S are determined using the following algorithm.

1.1. If $r'_{N-1} > 0$ set

$$X_N = p_N, \quad (p_V)_N = 1, \quad (h_S)_N = \begin{bmatrix} h_N \\ d_N \end{bmatrix},$$

$(d_v)_N$ to be 1×0 empty matrix, Δ_N to be 0×1 empty matrix;
if $r'_{N-1} = 0$ set X_N to be 0×0 empty matrix, $(p_V)_N$ to be 1×0 empty matrix,

$$(d_V)_N = 1, \quad (h_S)_N = h_N, \quad \Delta_N = d_N.$$

1.2. For $k = N - 1, \dots, 2$ perform the following. Compute the QR factorization

$$\begin{bmatrix} p_k \\ X_{k+1}a_k \end{bmatrix} = V_k \begin{pmatrix} X_k \\ 0 \end{pmatrix},$$

where V_k is a unitary matrix of sizes $(1 + \rho_k) \times (1 + \rho_k)$, X_k is a matrix of sizes $\rho_{k-1} \times r'_{k-1}$. Determine matrices $(p_V)_k$, $(a_V)_k$, $(d_V)_k$, $(q_V)_k$ of sizes $1 \times \rho_{k-1}$, $\rho_k \times \rho_{k-1}$, $1 \times \nu_k$, $\rho_k \times \nu_k$ from the partition

$$V_k = \begin{bmatrix} (p_V)_k & (d_V)_k \\ (a_V)_k & (q_V)_k \end{bmatrix}.$$

Compute

$$\begin{aligned} h'_k &= (p_V)_k^* d_k + (a_V)_k^* X_{k+1} q_k, \quad (h_S)_k = \begin{bmatrix} h_k \\ h'_k \end{bmatrix}, \quad (b_S)_k = \begin{pmatrix} b_k & 0 \\ (p_V)_k^* g_k & (a_V)_k^* \end{pmatrix}, \\ \Theta_k &= \begin{bmatrix} (d_V)_k^* g_k & (q_V)_k^* \end{bmatrix}, \quad \Delta_k = (d_V)_k^* d_k + (q_V)_k^* X_{k+1} q_k. \end{aligned}$$

1.3. Set $V_1 = I_{\nu_1}$ and define matrices $(d_V)_1$, $(q_V)_1$ of sizes $1 \times \rho_1$, $\rho_1 \times \nu_1$ from the partition

$$V_1 = \begin{bmatrix} (d_V)_1 \\ (q_V)_1 \end{bmatrix};$$

compute

$$\Delta_1 = \begin{pmatrix} d_1 \\ X_2 q_1 \end{pmatrix}, \quad \Theta_1 = \begin{pmatrix} g_1 & 0 \\ 0 & I_{\rho_1} \end{pmatrix}.$$

Thus we have computed generators and diagonal entries of the matrix V and generators $(b_S)_k$, $(h_S)_k$ of the matrix S .

2.1. Compute the QR factorization

$$\begin{bmatrix} \Delta_1 & \Theta_1 \end{bmatrix} = U_1 \begin{bmatrix} (d_S)_1 & (g_S)_1 \\ 0 & Y_1 \end{bmatrix},$$

where U_1 is a unitary matrix of sizes $\nu_1 \times \nu_1$, $(d_S)_1$ is a number, $(g_S)_1$ is a row of size ρ'_1 , Y_1 is a matrix of sizes $\rho_1 \times \rho'_1$. Determine matrices $(d_U)_1$, $(g_U)_1$ of sizes $\nu_1 \times 1$, $\nu_1 \times \rho'_1$ from the partition

$$U_1 = \begin{bmatrix} (d_U)_1 & (g_U)_1 \end{bmatrix}.$$

2.2. For $k = 2, \dots, N - 1$ perform the following. Compute the QR factorization

$$\begin{bmatrix} Y_{k-1}(h_S)_k & Y_{k-1}(b_S)_k \\ \Delta_k & \Theta_k \end{bmatrix} = U_k \begin{bmatrix} (d_S)_k & (g_S)_k \\ 0 & Y_k \end{bmatrix},$$

where U_k is a unitary matrix of sizes $(1 + \rho_k) \times (1 + \rho_k)$, $(d_S)_k$ is a number, $(g_S)_k$ is a row of size ρ'_k , Y_k is a matrix of sizes $\rho_k \times \rho'_k$.

2.3. If $r'_{N-1} > 0$ set $(d_U)_N = 1$ and $(h_U)_N$ to be 0×1 empty matrix;

if $r'_{N-1} = 0$ set $(h_U)_N = 1$ and $(d_U)_N$ to be 0×1 empty matrix;

compute

$$(d_S)_N = \begin{bmatrix} Y_{N-1}(h_S)_N \\ \Delta_N \end{bmatrix}.$$

Thus we have computed generators and diagonal entries of the matrix U and generators $(g_S)_k$ and diagonal entries $(d_S)_k$ of the matrix S .

Theorem 3.1 yields the QR-factorization of the matrix R , i.e. representation of R in the form $R = QS$ with the unitary matrix $Q = UV$ and the upper triangular matrix S . For the next considerations we should obtain generators of the matrix Q explicitly.

Theorem 3.2 Let $R = \{R_{ij}\}_{i,j=1}^N$ be a scalar matrix with lower generators p_i ($i = 2, \dots, N$), q_j ($j = 1, \dots, N - 1$), a_k ($k = 2, \dots, N - 1$) of orders r'_k ($k = 1, \dots, N - 1$), upper generators g_i ($i = 1, \dots, N - 1$), h_j ($j = 2, \dots, N$), b_k ($k = 2, \dots, N - 1$) of orders r''_k ($k = 1, \dots, N - 1$) and diagonal entries d_k ($k = 1, \dots, N$). Let us define the numbers ρ_k via recursive relations $\rho_N = 0$, $\rho_{k-1} = \min\{1 + \rho_k, r'_{k-1}\}$, $k = N, \dots, 2$, $\rho_0 = 0$ and the numbers $\rho'_k = r''_k + \rho_k$, $k = 1, \dots, N$.

The matrix R admits the factorization

$$R = QS,$$

where Q is a unitary matrix with lower generators $(p_Q)_i$ ($i = 2, \dots, N$), $(q_Q)_j$ ($j = 1, \dots, N - 1$), $(a_Q)_k$ ($k = 2, \dots, N - 1$) of orders ρ_k ($k = 1, \dots, N - 1$), upper generators $(g_Q)_i$ ($i = 1, \dots, N - 1$), $(h_Q)_j$ ($j = 2, \dots, N$), $(b_Q)_k$ ($k = 2, \dots, N - 1$) of orders ρ_k ($k = 1, \dots, N - 1$) also and diagonal entries $(d_Q)_k$ ($k = 1, \dots, N$) and S is an upper triangular matrix with upper generators $(g_S)_i$ ($i = 1, \dots, N - 1$), $(h_S)_j$ ($j = 2, \dots, N$), $(b_S)_k$ ($k = 2, \dots, N - 1$) of orders ρ'_k ($k = 1, \dots, N - 1$) and diagonal entries $(d_S)_k$ ($k = 1, \dots, N$).

The generators and the diagonal entries of the matrices Q and S are determined using the following algorithm.

1. Using the algorithm from Theorem 3.1 compute generators and diagonal entries of the upper triangular matrix S and of the unitary block triangular matrices V and U such that $R = VUS$.

2. Compute generators and diagonal entries of the matrix $Q = VU$ using generators and diagonal entries of the matrices V, U as follows.

2.1. Compute

$$z_1 = (q_V)_1(g_U)_1,$$

$$(q_Q)_1 = (q_V)_1(d_U)_1, \quad \alpha_1 = (a_V)_2 z_1, \quad (3.1)$$

$$(d_Q)_1 = (d_V)_1(d_U)_1, \quad \beta_1 = z_1, \quad (3.2)$$

$$(g_Q)_1 = (d_V)_1(g_U)_1, \quad \gamma_1 = z_1(b_U)_2. \quad (3.3)$$

Set $(a_V)_N = 0_{0 \times \rho_{N-1}}$, $(b_V)_N = 0_{\rho_{N-1} \times 0}$.

2.2. For $i = 2, \dots, N-1$ perform the following. Set

$$(p_Q)_i = (p_V)_i, \quad (a_Q)_i = (a_V)_i, \quad (b_Q)_i = (b_U)_i, \quad (h_Q)_i = (h_U)_i.$$

Compute

$$z_i = (q_V)_i(g_U)_i,$$

$$(q_Q)_i = (q_V)_i(d_U)_i + \alpha_{i-1}(h_U)_i, \quad \alpha_i = (a_V)_{i+1}[z_i + \alpha_{i-1}(b_U)_i], \quad (3.4)$$

$$(d_Q)_i = (d_V)_i(d_U)_i + (p_V)_i\beta_{i-1}(h_U)_i, \quad \beta_i = z_i + (a_V)_i\beta_{i-1}(b_U)_i, \quad (3.5)$$

$$(g_Q)_i = (d_V)_i(g_U)_i + (q_V)_i\gamma_{i-1}, \quad \gamma_i = [z_i + (a_V)_i\gamma_{i-1}](b_U)_{i+1}. \quad (3.6)$$

2.3. Set $(p_Q)_N = (p_V)_N$, $(h_Q)_N = (h_U)_N$. Compute

$$(d_Q)_N = (d_V)_N(d_U)_N + (p_V)_N\beta_{N-1}(h_U)_N. \quad (3.7)$$

Proof. We should justify the second stage of the algorithm. Let $Q = \{Q_{ij}\}_{i,j=1}^N$, $V = \{V_{ij}\}_{i,j=1}^N$, $U = \{U_{ij}\}_{i,j=1}^N$. For $N \geq i > j \geq 1$ since U is an upper triangular matrix and $(p_V)_i$ ($i = 2, \dots, N$), $(q_V)_j$ ($j = 1, \dots, N-1$), $(a_V)_k$ ($k = 2, \dots, N-1$) are lower generators of the matrix V we have

$$Q_{ij} = \sum_{k=1}^j V_{ik}U_{kj} = \sum_{k=1}^j (p_V)_i(a_V)_{ik}^\times(q_V)_kU_{kj}.$$

Using the equality (2.1) we obtain

$$Q_{ij} = (p_V)_i(a_V)_{ij}^\times(q_Q)_j, \quad 1 \leq j < i \leq N$$

where

$$(q_Q)_j = \sum_{k=1}^j (a_V)_{j+1,k}^\times(q_V)_kU_{kj}, \quad j = 1, \dots, N-1. \quad (3.8)$$

This implies that the matrix Q has the lower generators $(p_Q)_i = (p_V)_i$ ($i = 2, \dots, N$), $(a_Q)_k = (a_V)_k$ ($k = 2, \dots, N-1$) and $(q_Q)_j$ ($j = 1, \dots, N-1$) defined in (3.8). This in particular means that the orders ρ_k ($k = 1, \dots, N-1$) of these generators are the same as for the matrix V . Now we must check that the generators $(q_Q)_j$ satisfy the relations (3.1), (3.4). Indeed for $j = 1$ we have

$$(q_Q)_1 = (a_V)_{2,1}^\times(q_V)_1U_{11} = (q_V)_1(d_U)_1$$

and for $j = 2, \dots, N-1$ using $U_{jj} = (d_U)_j$ and the fact that $(g_U)_i$ ($i = 1, \dots, N-1$), $(h_U)_j$ ($j = 2, \dots, N$), $(b_U)_k$ ($k = 2, \dots, N-1$) are the upper generators of the matrix U we get

$$(q_Q)_j = \sum_{k=1}^{j-1} (a_V)_{j+1,k}^\times (q_V)_k (g_U)_k (b_U)_{kj}^\times (h_U)_j + (a_V)_{j+1,j}^\times (q_V)_j (d_U)_j = \alpha_{j-1} (h_U)_j + (q_V)_j (d_U)_j,$$

where

$$\alpha_{j-1} = \sum_{k=1}^{j-1} (a_V)_{j+1,k}^\times (q_V)_k (g_U)_k (b_U)_{kj}^\times.$$

We have

$$\alpha_1 = (a_V)_{3,1}^\times (q_V)_1 (g_U)_1 (b_U)_{2,1}^\times = (a_V)_2 (q_V)_1 (g_U)_1$$

and using the relations (2.1), (2.2) we obtain

$$\begin{aligned} \alpha_j &= \sum_{k=1}^j (a_V)_{j+2,k}^\times (q_V)_k (g_U)_k (b_U)_{k,j+1}^\times \\ &= (a_V)_{j+2,j}^\times (q_V)_j (g_U)_j (b_U)_{j,j+1}^\times + (a_V)_{j+1,j+1}^\times \left(\sum_{k=1}^{j-1} (a_V)_{j+1,k}^\times (q_V)_k (g_U)_k (b_U)_{kj}^\times \right) (b_U)_j \\ &= (a_V)_{j+1} (q_V)_j (g_U)_j + (a_V)_{j+1} \alpha_{j-1} (b_U)_j \end{aligned}$$

which completes the proof of (3.1), (3.4).

For diagonal entries of the matrix Q we have

$$(d_Q)_1 = Q_{11} = V_{11} U_{11} = (d_V)_1 (d_U)_1$$

and for $i = 2, \dots, N$

$$Q_{ii} = \sum_{k=1}^i V_{ik} U_{ki} = V_{ii} U_{ii} + \sum_{k=1}^{i-1} V_{ik} U_{ki} = (d_V)_i (d_U)_i + (p_V)_i \beta_{i-1} (h_U)_i,$$

where

$$\beta_{i-1} = \sum_{k=1}^{i-1} (a_V)_{ik}^\times (q_V)_k (g_U)_k (b_U)_{ki}^\times$$

We have $\beta_1 = (q_V)_1 (g_U)_1$ and using the relations (2.1), (2.2) we obtain

$$\begin{aligned} \beta_i &= \sum_{k=1}^i (a_V)_{i+1,k}^\times (q_V)_k (g_U)_k (b_U)_{k,i+1}^\times \\ &= (a_V)_{i+1,i}^\times (q_V)_i (g_U)_i (b_U)_{i,i+1}^\times + (a_V)_i \left(\sum_{k=1}^{i-1} (a_V)_{ik}^\times (q_V)_k (g_U)_k (b_U)_{ki}^\times \right) (b_U)_i \\ &= (q_V)_i (g_U)_i + (a_V)_i \beta_{i-1} (b_U)_i \end{aligned}$$

which completes the proof of (3.2), (3.5), (3.7).

The proof of the relations (3.3), (3.6) is performed in the same way as the proof of (3.1), (3.4). \square

Corollary 3.3 *Let R be a quasiseparable of order (n_L, n_U) matrix with scalar entries and let $R = QS$ be the factorisation obtained in Theorem 3.2. Then the unitary matrix Q is quasiseparable of order (n_L, n_L) at most and the upper triangular matrix S is upper quasiseparable of order $n_L + n_U$ at most.*

Proof. By Theorem 3.2 the matrix Q has lower and upper generators of the orders ρ_k ($k = 1, \dots, N-1$) defined by the relations

$$\rho_N = 0, \rho_{k-1} = \min\{1 + \rho_k, r'_{k-1}\}, k = N, \dots, 2 \quad (3.9)$$

and by Theorem 3.1 the matrix S has upper generators of orders

$$\rho'_k = r''_k + \rho_k, k = 1, \dots, N-1. \quad (3.10)$$

From the inequalities $r'_k \leq n_L$ ($k = 1, \dots, N-1$) and the relations 3.9 it follows that

$$\rho_k \leq r'_k \leq n_L, k = 1, \dots, N-1 \quad (3.11)$$

and hence the maximal order of generators of the matrix Q is not greater than n_L . Next from (3.10) and (3.11) we conclude that the maximal order of upper generators of the matrix S is not greater than $n_L + n_U$. \square

4 The QR iteration

We consider the QR iteration algorithm for matrices defined via generators. In each iteration step for a given matrix R and for a given real number σ the new iterant R_1 is obtained by the rule

$$\begin{cases} R - \sigma I = QS, \\ R_1 = \sigma I + QR, \end{cases}$$

where Q is a unitary matrix and S is an upper triangular matrix. We show that the matrix R_1 has lower generators with the same order as the lower generators of the matrix Q and hence these orders are not greater than the corresponding generators of the matrix R and obtain an algorithm for computation of these generators and the diagonal entries of the matrix R_1 .

Theorem 4.1 *Let $R = \{R_{ij}\}_{i,j=1}^N$ be a scalar matrix with lower generators p_i ($i = 2, \dots, N$), q_j ($j = 1, \dots, N-1$), a_k ($k = 2, \dots, N-1$) of orders r'_k ($k = 1, \dots, N-1$), upper generators g_i ($i = 1, \dots, N-1$), h_j ($j = 2, \dots, N$), b_k ($k = 2, \dots, N-1$) of orders r''_k ($k = 1, \dots, N-1$) and diagonal entries d_k ($k = 1, \dots, N$) and σ be a real number. Let us define the numbers*

ρ_k via recursive relations $\rho_N = 0$, $\rho_{k-1} = \min\{1 + \rho_k, r'_{k-1}\}$, $k = N, \dots, 2$, $\rho_0 = 0$. Define the matrix R_1 by the rule

$$\begin{cases} R - \sigma I = QS, \\ R_1 = \sigma I + QR, \end{cases}$$

where Q is a unitary matrix and S is an upper triangular matrix.

The matrix R_1 has lower generators of orders ρ_k ($k = 1, \dots, N-1$). These lower generators $p_i^{(1)}$ ($i = 2, \dots, N$), $q_j^{(1)}$ ($j = 1, \dots, N-1$), $a_k^{(1)}$ ($k = 2, \dots, N-1$) and the diagonal entries $d_k^{(1)}$ ($k = 1, \dots, N$) of the matrix R are determined using the following algorithm.

1. Apply to the matrix $R - \sigma I$, which has the same lower and upper generators as the matrix R and the diagonal entries $d_k - \sigma$ ($k = 1, \dots, N$), the algorithm from Theorem 3.2, to compute the lower generators $(p_Q)_i$ ($i = 2, \dots, N$), $(q_Q)_j$ ($j = 1, \dots, N-1$), $(a_Q)_k$ ($k = 2, \dots, N-1$) and the diagonal entries $(d_Q)_k$ ($k = 1, \dots, N$) of the matrix Q and the upper generators $(g_S)_i$ ($i = 1, \dots, N-1$), $(h_S)_j$ ($j = 2, \dots, N$), $(b_S)_k$ ($k = 2, \dots, N-1$) and the diagonal entries $(d_S)_k$ ($k = 1, \dots, N$) of the matrix S .

2. Compute the lower generators and the diagonal entries of the matrix Q as follows.

2.1. Compute

$$z_N = (h_S)_N(p_Q)_N,$$

$$p_N^{(1)} = (d_S)_N(p_Q)_N, \quad \alpha_N = z_N(a_Q)_{N-1}, \quad (4.1)$$

$$d_N^{(1)} = (d_S)_N(d_Q)_N, \quad \beta_N = z_N, \quad (4.2)$$

Set $(a_Q)_1 = 0_{\rho_1 \times 0}$.

2.2. For $i = N-1, \dots, 2$ perform the following. Set

$$q_i^{(1)} = (q_Q)_i, \quad a_i^{(1)} = (a_Q)_i.$$

Compute

$$z_i = (h_S)_i(p_Q)_i,$$

$$p_i^{(1)} = (d_S)_i(p_Q)_i + (g_S)_i \alpha_{i+1}, \quad \alpha_i = [(h_S)_i(p_Q)_i + (b_S)_i \alpha_{i+1}](a_Q)_{i-1}, \quad (4.3)$$

$$d_i^{(1)} = (d_S)_i(d_Q)_i + (g_S)_i \beta_{i+1}(q_Q)_i, \quad \beta_i = z_i + (b_S)_i \beta_{i+1}(a_Q)_i. \quad (4.4)$$

2.3. Set $q_1^{(1)} = (q_Q)_1$. Compute

$$d_1^{(1)} = (d_S)_1(d_Q)_1 + (g_S)_1 \beta_2(q_Q)_1. \quad (4.5)$$

Proof. We should justify the second stage of the algorithm. Let $Q = \{Q_{ij}\}_{i,j=1}^N$, $S = \{S_{ij}\}_{i,j=1}^N$ and $R_1 = \{R_{ij}^{(1)}\}_{i,j=1}^N$. For $N \geq i > j \geq 1$ using the fact S is an upper triangular matrix and $(p_Q)_i$ ($i = 2, \dots, N$), $(q_Q)_j$ ($j = 1, \dots, N-1$), $(a_Q)_k$ ($k = 2, \dots, N-1$) are lower generators of the matrix Q we have

$$R_{ij}^{(1)} = \sum_{k=i}^N S_{ik} Q_{kj} = \sum_{k=i}^N S_{ik} (p_Q)_k (a_Q)_{kj}^\times (q_Q)_j.$$

Using the equality (2.1) we obtain

$$R_{ij}^{(1)} = p_i^{(1)}(a_Q)_{ij}^\times(q_Q)_j, \quad 1 \leq j < i \leq N$$

where

$$p_i^{(1)} = \sum_{k=i}^N S_{ik}(p_Q)_k(a_Q)_{k,i-1}^\times, \quad i = 2, \dots, N. \quad (4.6)$$

This implies that the matrix $R^{(1)}$ has the lower generators $a_k^{(1)} = (a_Q)_k$ ($k = 2, \dots, N-1$), $q_j^{(1)} = (q_Q)_j$ ($j = 1, \dots, N-1$) and $p_i^{(1)}$ ($i = 2, \dots, N$) defined in (4.6). This in particular means that the orders ρ_k ($k = 1, \dots, N-1$) of these generators are the same as for the matrix Q . Now we must check that the generators $p_i^{(1)}$ satisfy the relations (4.1), (4.3). Indeed for $i = N$ we have

$$p_N^{(1)} = S_{NN}(p_Q)_N(a_Q)_{N,N-1}^\times = (d_S)_N(p_Q)_N$$

and for $i = N-1, \dots, 2$ using $S_{jj} = (d_S)_j$ and the fact that $(g_S)_i$ ($i = 1, \dots, N-1$), $(h_S)_j$ ($j = 2, \dots, N$), $(b_S)_k$ ($k = 2, \dots, N-1$) are the upper generators of the matrix S we get

$$p_i^{(1)} = (g_S)_i \sum_{k=i+1}^N (b_S)_{ik}^\times (h_S)_k (p_Q)_k (a_Q)_{k,i-1}^\times + (d_S)_i (p_Q)_i (a_Q)_{i,i-1}^\times = (d_S)_i (p_Q)_i + (g_S)_i \alpha_{i+1},$$

where

$$\alpha_{i+1} = \sum_{k=i+1}^N (b_S)_{ik}^\times (h_S)_k (p_Q)_k (a_Q)_{k,i-1}^\times.$$

We have

$$\alpha_N = (b_S)_{N-1,N}^\times (h_S)_N (p_Q)_N (a_Q)_{N,N-2}^\times = (h_S)_N (p_Q)_N (a_Q)_{N-1}$$

and using the relations (2.1), (2.2) we obtain

$$\begin{aligned} \alpha_i &= \sum_{k=i}^N (b_S)_{i-1,k}^\times (h_S)_k (p_Q)_k (a_Q)_{k,i-2}^\times \\ &= (b_S)_{i-1,i}^\times (h_S)_i (p_Q)_i (a_Q)_{i,i-2}^\times + (b_S)_i \left(\sum_{k=i+1}^N (b_S)_{ik}^\times (h_S)_k (p_Q)_k (a_Q)_{k,i-1}^\times \right) (a_Q)_{i-1} \\ &= [(h_S)_i (p_Q)_i + (b_S)_i \alpha_{i+1}] (a_Q)_{i-1} \end{aligned}$$

which completes the proof of (4.1), (4.3).

For diagonal entries of the matrix S we have

$$d_N^{(1)} = R_{NN}^{(1)} = S_{NN} Q_{NN} = (d_S)_N (d_Q)_N$$

and for $i = N - 1, \dots, 1$

$$R_{ii}^{(1)} = \sum_{k=i}^N S_{ik} Q_{ki} = S_{ii} Q_{ii} + \sum_{k=i+1}^N S_{ik} Q_{ki} = (d_S)_i (d_Q)_i + (g_S)_i \beta_{i+1} (h_S)_i,$$

where

$$\beta_{i+1} = \sum_{k=i+1}^N (b_S)_{ik}^\times (h_S)_k (p_Q)_k (a_Q)_{ki}^\times$$

We have $\beta_1 = (q_V)_1 (g_U)_1$ and using the relations (2.1), (2.2) we obtain

$$\begin{aligned} \beta_i &= \sum_{k=i}^N (b_S)_{i-1,k}^\times (h_S)_k (p_Q)_k (a_Q)_{k,i-1}^\times \\ &= (b_S)_{i-1,i}^\times (h_S)_i (p_Q)_i (a_Q)_{i,i-1}^\times + (b_S)_i \left(\sum_{k=i+1}^N (b_S)_{ik}^\times (h_S)_k (p_Q)_k (a_Q)_{ki}^\times \right) (a_Q)_i = \\ &\quad (h_S)_i (p_Q)_i + (b_S)_i \beta_{i+1} (a_Q)_i \end{aligned}$$

which completes the proof of (4.2), (4.4), (4.5). \square

Corollary 4.2 *Let R be a lower quasiseparable of order n_L matrix with scalar entries and let R_1 be the matrix obtained in Theorem 3.2. Then the unitary matrix Q is quasiseparable of order (n_L, n_L) at most and the upper triangular matrix S is upper quasiseparable of order $n_L + n_U$ at most.*

Proof follows directly from Theorem 4.1 and Corollary 3.3.

Now assume that the matrix R is Hermitian. Then the new iterant R_1 is a Hermitian matrix which is quasiseparable of the same order as the matrix R . This means that for a quasiseparable of a given order Hermitian matrix, the result of QR iteration has the same structure as the original matrix. Moreover an algorithm for computation of this structure is given.

Theorem 4.3 *Let $R = \{R_{ij}\}_{i,j=1}^N$ be a scalar Hermitian quasiseparable of order (n, n) matrix with lower generators p_i ($i = 2, \dots, N$), q_j ($j = 1, \dots, N-1$), a_k ($k = 2, \dots, N-1$) of orders r'_k ($k = 1, \dots, N-1$), upper generators q_i^* ($i = 1, \dots, N-1$), p_j^* ($j = 2, \dots, N$), a_k^* ($k = 2, \dots, N-1$) and diagonal entries d_k ($k = 1, \dots, N$) and σ be a real number. Define the matrix R_1 by the rule*

$$\begin{cases} R - \sigma I = QS, \\ R_1 = \sigma I + QR, \end{cases}$$

where Q is a unitary matrix and S is an upper triangular matrix.

Then R_1 is a Hermitian quasiseparable of order (n, n) at most matrix and generators and diagonal entries of this matrix are obtained using the algorithm from Theorem 4.1.

References

- [1] Y. Eidelman and I. Gohberg, A modification of the Dewilde-van der Veen method for inversion of finite structured matrices. *Linear Algebra and Application* 343-344: 419-450 (2002).