# MAC-5754 CML – Sintaxe e Semântica

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# **Sintaxe**

```
/* Expressões necessariamente "atômicas": */
primary_expression
        : IDENTIFIER
         literal
         '{' '}'
         '{' expression_list '}'
          '(' expression ')'
          array_access
         IDENTIFIER '(' ')'
        | IDENTIFIER '(' expression_list ')'
literal
        : INT LITERAL
         REAL LITERAL
          BOOL_LITERAL
          CHAR LITERAL
          STRING_LITERAL
```

```
/* Expressões "não-atômicas" */
expression_list
        : expression
          expression_list ',' expression
expression
        : logical_or_expression
        | IDENTIFIER '=' expression
        | array_access '=' expression
array_access
        : IDENTIFIER '[' expression ']'
        | array_access '[' expression ']'
```

```
logical_or_expression
        : logical_and_expression
        logical_or_expression OR_OP logical_and_expression
logical_and_expression
        : relational_expression
          logical_and_expression AND_OP relational_expression
relational_expression
          additive_expression
          additive_expression '<' additive_expression
          additive_expression '>' additive_expression
          additive_expression LE_OP additive_expression
          additive_expression GE_OP additive_expression
          additive_expression EQ_OP additive_expression
          additive_expression NE_OP additive_expression
```

```
additive_expression
        : multiplicative_expression
          additive_expression '+' multiplicative_expression
          additive_expression '-' multiplicative_expression
multiplicative_expression
        : unary_minus_expression
        multiplicative_expression '*' unary_minus_expression
        multiplicative_expression '/' unary_minus_expression
unary_minus_expression
        : neg_expression
        '-' neg_expression
neg_expression
        : primary_expression
         '!' neg_expression
```

#### Sintaxe - Comandos

```
command
        : compound_command
          expression_command
          selection command
         iteration_command
          jump_command
          SKIP ':'
compound_command
        : '{' declaration_or_command_list '}'
declaration_or_command_list
        : declaration or command
          declaration_or_command_list declaration_or_command
declaration_or_command
        : declaration
          command
```

#### Sintaxe - Comandos

```
expression_command
        expression ';'
selection_command
        : IF '(' expression ')' compound_command ELSE compound_command
        | IF '(' expression ')' compound_command
iteration command
        : WHILE '(' expression ')' command
jump_command
        : RETURN ':'
        RETURN expression ';'
```

# Sintaxe - Declarações

```
declaration
    : type_specifier IDENTIFIER ';'
    | type_specifier IDENTIFIER '=' expression ';'
;

parameter_declaration_list
    : parameter_declaration
    | parameter_declaration_list ',' parameter_declaration
;

parameter_declaration
    : type_specifier IDENTIFIER
    ;
```

# Sintaxe - Declarações

```
type_specifier
    : VOID
    | CHAR
    | INT
    | REAL
    | BOOL
    | STRING
    | DATASET
    | MODEL
    | type_specifier '[' ']'
;
```

# Sintaxe - Definições

# Sintaxe - Programa

# Semântica

- $Int := \{..., -2, -1, 0, 1, 2, ...\} = \mathbb{Z}$
- $\blacksquare$  Real :=  $\mathbb{R}$
- Bool := {true, false}
- Char := {0,..., 127}, onde os números de 0 a 127 são interpretados conforme o padrão ASCII, e.g., 43 representa '+', 49 representa '1', etc. Para mais detalhes, ver tabela no início do relatório.

Definimos, para cada conjunto X, o conjunto X[] dos vetores de elementos de X pondo  $X[] := \bigcup_{n=0}^{+\infty} X^n$  rotulo(X) denota o conjunto  $\{rotulo(x) : x \in X\}$  de elementos de X rotulados pela palavra rotulo. Por exemplo,

$$int(Int) = {\dots, int(-2), int(-1), int(0), int(1), int(2), \dots}.$$

Isto serve para que possamos tomar a união, por exemplo,  $int(Int) \cup char(Char)$ , que é

$$\{\ldots, int(-2), int(-1), int(0), int(1), int(2), \ldots,$$

$$char(0), char(1), \ldots, char(127)$$
,

e saber de que conjunto cada elemento provém. Se simplesmente tomássemos a união sem rótulos,  $Int \cup Char = Int$ , não saberíamos, por exemplo, se  $5 \in Int \cup Char$  provém do conjunto Int ou do conjunto Char, isto é, não saberíamos o tipo do valor 5.

- String = Char[]
   Observe que String é simplesmente o conjunto dos vetores de Char's. Usaremos os rótulos string e array para distingui-los.
- Dataset =  $\{0, ..., n\} \times \{1, ..., m\} \rightarrow int(Int) \cup real(Real) \cup string(String)$
- lacktriangledown Model = Dataset ightarrow Dataset
- Array = Location[]
- Input :=  $\bigcup_{n=0}^{+\infty} Dataset^n$
- Output :=  $\bigcup_{n=0}^{+\infty} Dataset^n$

- Function =  $\bigcup_{n=0}^{+\infty} ((Location^n \times Store \times Input \times Output) \times Location \rightarrow Store \times Input \times Output)$
- StorableValue :=  $int(Int) \cup real(Real) \cup bool(Bool) \cup char(Char) \cup string(String) \cup dataset(Dataset) \cup model(Model) \cup array(Array)$
- ExpressibleValue := StorableValue
- DenotableValue := var(Location) ∪ fun(Function)
- Location  $:= \mathbb{N} := \{0, 1, 2, \ldots\}$

- Environment := Identifier  $\rightarrow$  DenotableValue $\cup \{unbound\}^1$
- ReturnFlag := {returnFlag0, returnFlag1}
  A flag indica se houve return no comando.
- Store := Location  $\rightarrow$  StorableValue  $\cup$  {unused}  $\cup$  {undefined}<sup>2</sup>
- ullet  $\Sigma := State := Environment imes Store imes Input imes Output$

<sup>&</sup>lt;sup>1</sup>O valor *unbound* indica que o identificador não foi associado a uma posição de memória ou função, isto é, não foi declarado.

<sup>&</sup>lt;sup>2</sup>O valor *unused* indica que a localização não está sendo utilizada por nenhuma variável. Já o valor *undefined* indica que a localização está sendo utilizada por uma variável que não foi inicializada.

Obs.: acrescentamos o valor especial *error* em todos os domínios, para indicar erro no programa e supomos que erros são propagados pelas funções semânticas.

# Frunções Semânticas

```
E: \texttt{Expression} \to (\Sigma \to \textit{Store} \times \textit{Input} \times \textit{Output} \times \textit{ExpressibleValue})
C: \texttt{Command} \to (\Sigma \to \textit{Store} \times \textit{Input} \times \textit{Output} \times \textit{ReturnFlag})
Dec: \texttt{Declaration} \to (\Sigma \to \Sigma)
Def: \texttt{FunctionDefinition} \to (\Sigma \to \textit{Environment})
P: \texttt{Program} \to (\textit{Input} \to \textit{Output} \times \textit{Int})
P_1: \texttt{Program} \to (\Sigma \to \Sigma)
P_2: \texttt{Program} \to (\Sigma \to \Sigma)
P_3: \texttt{Program} \to (\Sigma \to \Sigma)
```

## Funções Semânticas Auxiliares

```
value: Literal \rightarrow int(Int) \cup real(Real) \cup bool(Bool) \cup char(Char) \cup string(String) \\ emptyEnv: Environment \\ extendEnv: Environment \times Identifier \times DenotableValue \rightarrow Environment \\ applyEnv: Environment \times Identifier \rightarrow DenotableValue \cup \{unbound\} \\ emptySto: Store \\ updateSto: Store \times Location \times (StorableValue \cup \{undefined, unused\}) \rightarrow Store \\ applySto: Store \times Location \rightarrow StorableValue \cup \{undefined, unused\} \\ allocate: Store \rightarrow Store \times Location \\ deallocate: Store \times Location \rightarrow Store
```

```
\begin{split} E[[\{\}]] & (\textit{env}, \textit{sto}, \textit{in}, \textit{out}) = (\textit{sto}, \textit{in}, \textit{out}, \textit{array}(\emptyset)) \\ E[[\{\textit{exp}\}]] & (\textit{env}, \textit{sto}, \textit{in}, \textit{out}) = (\textit{sto}_3, \textit{in}_1, \textit{out}_1, \textit{array}(\textit{loc})) \\ & \quad \text{where } (\textit{sto}_1, \textit{in}_1, \textit{out}_1, \textit{val}) = E[[\textit{exp}]](\textit{env}, \textit{sto}, \textit{in}, \textit{out}) \\ & \quad \text{and } (\textit{sto}_2, \textit{loc}) = \textit{allocate sto}_1 \\ & \quad \text{and } \textit{sto}_3 = \textit{updateSto}(\textit{sto}_2, \textit{loc}, \textit{val}) \\ E[[\{\textit{exp\_list}, \textit{exp}\}]] & (\textit{env}, \textit{sto}, \textit{in}, \textit{out}) = (\textit{sto}_f, \textit{in}_f, \textit{out}_f, \textit{array}(\textit{extendArray}(\textit{arr}, \textit{loc}))) \\ & \quad \text{where } (\textit{sto}_1, \textit{in}_1, \textit{out}_1, \textit{array}(\textit{arr})) = E[[\{\textit{exp\_list}\}]] & (\textit{env}, \textit{sto}, \textit{in}, \textit{out}) \\ & \quad \text{and } (\textit{sto}_2, \textit{in}_f, \textit{out}_f, \textit{val}) = E[[\textit{exp}]] & (\textit{env}, \textit{sto}_1, \textit{in}_1, \textit{out}_1) \\ & \quad \text{and } (\textit{sto}_3, \textit{loc}) = \textit{allocate sto}_2 \\ & \quad \text{where } \textit{sto}_f = \textit{updateSto}(\textit{sto}_3, \textit{loc}, \textit{val}) \end{split}
```

```
E[[exp_1 + exp_2]](env, sto, in, out) = (sto_f, in_f, out_f, val_f)
   where (sto_1, in_1, out_1, int(val_1)) = E[[exp_1]] (env, sto, in, out)
      and (sto_f, in_f, out_f, int(val_2)) = E[[exp_2]] (env, sto_1, in_1, out_1)
      and val_f = int(val_1 + val_2)
   where (sto_1, in_1, out_1, int(val_1)) = E[[exp_1]] (env, sto, in, out)
      and (sto_f, in_f, out_f, real(val_2)) = E[[exp_2]] (env, sto_1, in_1, out_1)
      and val_f = real(val_1 + val_2)
   where (sto_1, in_1, out_1, real(val_1)) = E[[exp_1]] (env, sto, in, out)
      and (sto_f, in_f, out_f, int(val_2)) = E[[exp_2]] (env, sto_1, in_1, out_1)
      and val_f = real(val_1 + val_2)
   where (sto_1, in_1, out_1, real(val_1)) = E[[exp_1]] (env, sto, in, out)
      and (sto_f, in_f, out_f, real(val_2)) = E[[exp_2]] (env, sto_1, in_1, out_1)
      and val_{\epsilon} = real(val_1 + val_2)
```

```
E[[\exp_1 - \exp_2]](env, sto, in, out) = (sto_f, in_f, out_f, val_f)
   where (sto_1, in_1, out_1, int(val_1)) = E[[exp_1]] (env, sto, in, out)
      and (sto_f, in_f, out_f, int(val_2)) = E[[exp_2]] (env, sto_1, in_1, out_1)
      and val_f = int(val_1 - val_2)
   where (sto_1, in_1, out_1, int(val_1)) = E[[exp_1]] (env, sto, in, out)
      and (sto_f, in_f, out_f, real(val_2)) = E[[exp_2]] (env, sto_1, in_1, out_1)
      and val_f = real(val_1 - val_2)
   where (sto_1, in_1, out_1, real(val_1)) = E[[exp_1]] (env, sto, in, out)
      and (sto_f, in_f, out_f, int(val_2)) = E[[exp_2]] (env, sto_1, in_1, out_1)
      and val_f = real(val_1 - val_2)
   where (sto_1, in_1, out_1, real(val_1)) = E[[exp_1]] (env, sto, in, out)
      and (sto_f, in_f, out_f, real(val_2)) = E[[exp_2]] (env, sto_1, in_1, out_1)
      and val_f = real(val_1 - val_2)
```

```
E[[exp_1 * exp_2]](env, sto, in, out) = (sto_f, in_f, out_f, val_f)
   where (sto_1, in_1, out_1, int(val_1)) = E[[exp_1]] (env, sto, in, out)
      and (sto_f, in_f, out_f, int(val_2)) = E[[exp_2]] (env, sto_1, in_1, out_1)
      and val_f = int(val_1 * val_2)
   where (sto_1, in_1, out_1, int(val_1)) = E[[exp_1]] (env, sto, in, out)
      and (sto_f, in_f, out_f, real(val_2)) = E[[exp_2]] (env, sto_1, in_1, out_1)
      and val_f = real(val_1 * val_2)
   where (sto_1, in_1, out_1, real(val_1)) = E[[exp_1]] (env, sto, in, out)
      and (sto_f, in_f, out_f, int(val_2)) = E[[exp_2]] (env, sto_1, in_1, out_1)
      and val_f = real(val_1 * val_2)
   where (sto_1, in_1, out_1, real(val_1)) = E[[exp_1]] (env, sto, in, out)
      and (sto_f, in_f, out_f, real(val_2)) = E[[exp_2]] (env, sto_1, in_1, out_1)
      and val_{\epsilon} = real(val_1 * val_2)
```

```
E[[\exp_1/\exp_2]](env, sto, in, out) = ans
   where (sto_1, in_1, out_1, int(val_1)) = E[[exp_1]] (env, sto, in, out)
      and (sto_f, in_f, out_f, int(val_2)) = E[[exp_2]] (env, sto_1, in_1, out_1)
      and if val_2 = 0 then ans = error else ans = (sto_f, in_f, out_f, int(val_1/val_2))
   where (sto_1, in_1, out_1, int(val_1)) = E[[exp_1]] (env, sto, in, out)
      and (sto_f, in_f, out_f, real(val_2)) = E[[exp_2]] (env, sto_1, in_1, out_1)
      and if val_2 = 0 then ans = error else ans = (sto_f, in_f, out_f, real(val_1/val_2))
   where (sto_1, in_1, out_1, real(val_1)) = E[[exp_1]] (env, sto, in, out)
      and (sto_f, in_f, out_f, int(val_2)) = E[[exp_2]] (env, sto_1, in_1, out_1)
      and if val_2 = 0 then ans = error else ans = (sto_f, in_f, out_f, real(val_1/val_2))
   where (sto_1, in_1, out_1, real(val_1)) = E[[exp_1]] (env, sto, in, out)
      and (sto_f, in_f, out_f, real(val_2)) = E[[exp_2]] (env, sto_1, in_1, out_1)
      and if val_2 = 0 then ans = error else ans = (sto_f, in_f, out_f, real(val_1/val_2))
```

```
\begin{split} E[[-\text{exp}]](\textit{env}, \textit{sto}, \textit{in}, \textit{out}) &= (\textit{sto}_f, \textit{in}_f, \textit{out}_f, \textit{labeledVal}) \\ \text{where } (\textit{sto}_f, \textit{in}_f, \textit{out}_f, \textit{int}(\textit{val})) &= E[[\text{exp}]](\textit{env}, \textit{sto}, \textit{in}, \textit{out}) \\ \text{where } \textit{labeledVal} &= \textit{int}(-\textit{val}) \\ \text{where } (\textit{sto}_f, \textit{in}_f, \textit{out}_f, \textit{real}(\textit{val})) &= E[[\text{exp}]](\textit{env}, \textit{sto}, \textit{in}, \textit{out}) \\ \text{where } \textit{labeledVal} &= \textit{real}(-\textit{val}) \end{split}
```

```
\begin{split} E[[\exp_1 \ \mid \ \mid \ \exp_2]](\mathit{env}, \mathit{sto}, \mathit{in}, \mathit{out}) &= (\mathit{sto}_f, \mathit{in}_f, \mathit{out}_f, \mathit{val}_f) \\ \text{where } (\mathit{sto}_1, \mathit{in}_1, \mathit{out}_1, \mathit{bool}(\mathit{val}_1)) &= E[[\exp_1]] \ (\mathit{env}, \mathit{sto}, \mathit{in}, \mathit{out}) \\ \text{and } (\mathit{sto}_f, \mathit{in}_f, \mathit{out}_f, \mathit{bool}(\mathit{val}_2)) &= E[[\exp_2]] \ (\mathit{env}, \mathit{sto}_1, \mathit{in}_1, \mathit{out}_1) \\ \text{and } \mathit{val}_f &= \mathit{bool}(\mathit{val}_1 \lor \mathit{val}_2) \end{split}
```

```
\begin{split} E[[\exp_1 \ \&\& \ \exp_2]] \ (env, sto, in, out) &= (sto_f, in_f, out_f, val_f) \\ \text{where } (sto_1, in_1, out_1, bool(val_1)) &= E[[\exp_-1]] \ (env, sto, in, out) \\ \text{and } (sto_f, in_f, out_f, bool(val_2)) &= E[[\exp_-2]] \ (env, sto_1, in_1, out_1) \\ \text{and } val_f &= bool(val_1 \wedge val_2) \end{split}
```

```
E[[exp_1 = exp_2]] (env, sto, in, out) =
   if labeledVal_1 = labeledVal_2 then (sto_2, in_2, out_2, bool(true))
else (sto2, in2, out2, bool(false))
       where (sto_1, in_1, out_1, labeledVal_1) = E[[exp_1]] (env, sto, in, out)
       and (sto_2, in_2, out_2, labeledVal_2) = E[[exp_2]](env, sto_1, in_1, out_1)
E[[exp_1 != exp_2]] (env, sto, in, out) =
   if labeledVal_1 \neq labeledVal_2 then (sto_2, in_2, out_2, bool(true))
else (sto<sub>2</sub>, in<sub>2</sub>, out<sub>2</sub>, bool(false))
       where (sto_1, in_1, out_1, labeledVal_1) = E[[exp_1]] (env, sto, in, out)
       and (sto_2, in_2, out_2, labeledVal_2) = E[[exp_2]](env, sto_1, in_1, out_1)
```

```
E[[exp_1 < exp_2]] (env, sto, in, out) =
   if labeledVal_1 < labeledVal_2 then (sto_2, in_2, out_2, bool(true))
else (sto2, in2, out2, bool(false))
      where (sto_1, in_1, out_1, labeledVal_1) = E[[exp_1]] (env, sto, in, out)
      and (sto_2, in_2, out_2, labeledVal_2) = E[[exp_2]](env, sto_1, in_1, out_1)
E[[exp_1 \le exp_2]] (env, sto, in, out) =
   if |abe| = dVal_1 < |abe| = dVal_2 then (sto_2, in_2, out_2, bool(true))
else (sto2, in2, out2, bool(false))
      where (sto_1, in_1, out_1, labeledVal_1) = E[[exp_1]] (env, sto, in, out)
      and (sto_2, in_2, out_2, labeledVal_2) = E[[exp_2]](env, sto_1, in_1, out_1)
```

 $E[[exp_1 > exp_2]]$  (env, sto, in, out) =

```
if labeledVal_1 > labeledVal_2 then (sto_2, in_2, out_2, bool(true))
else (sto_2, in_2, out_2, bool(false))
      where (sto_1, in_1, out_1, labeledVal_1) = E[[exp_1]] (env, sto, in, out)
      and (sto_2, in_2, out_2, labeledVal_2) = E[[exp_2]] (env, sto_1, in_1, out_1)
E[[exp_1 >= exp_2]] (env, sto, in, out) =
   if |abe| = dVal_1 > |abe| = dVal_2 then (sto_2, in_2, out_2, bool(true))
 else (sto2, in2, out2, bool(false))
       where (sto_1, in_1, out_1, labeledVal_1) = E[[exp_1]] (env, sto, in, out)
       and (sto_2, in_2, out_2, labeledVal_2) = E[[exp_2]](env, sto_1, in_1, out_1)
```

```
\begin{split} &E[[!\exp]] \; (\textit{env}, \textit{sto}, \textit{in}, \textit{out}) = (\textit{sto}_f, \textit{in}_f, \textit{out}_f, \textit{bool}(\neg \textit{val})) \\ & \quad \text{where} \; (\textit{sto}_f, \textit{in}_f, \textit{out}_f, \textit{bool}(\textit{val})) = E[[\exp]] \; (\textit{env}, \textit{sto}, \textit{in}, \textit{out}) \end{split}
```

$$E[[(e)]]$$
 (env, sto, in, out) =  $E[[e]]$  (env, sto, in, out)

```
\begin{split} E\big[ [\mathtt{id\_or\_arr\_access}] \big] & (\textit{env}, \textit{sto}, \textit{in}, \textit{out}) = (\textit{sto}_f, \textit{in}_f, \textit{out}_f, \textit{val}_f) \\ \text{where } (\textit{sto}_f, \textit{in}_f, \textit{out}_f, \textit{val}_f) = \textit{applyArray}(\mathtt{id\_or\_arr\_access}, \textit{sto}, \textit{in}, \textit{out}) \\ \text{and } \textit{applyArray}(\mathtt{id}, \textit{sto}, \textit{in}, \textit{out}) = (\textit{sto}, \textit{in}, \textit{out}, \textit{applySto}(\textit{sto}, \textit{loc})) \\ \text{where } \textit{loc} = \textit{applyEnv}(\textit{env}, \mathtt{id}) \\ \text{and } \textit{applyArray}(\mathtt{id\_or\_arr\_access}[\texttt{exp}], \textit{sto}, \textit{in}, \textit{out}) = (\textit{sto}_2, \textit{in}_2, \textit{out}_2, \textit{arr}_\textit{pos}) \\ \text{where } (\textit{sto}_1, \textit{in}_1, \textit{out}_1, \textit{arr}) = \textit{applyArray}(\mathtt{id\_or\_arr\_access}, \textit{sto}, \textit{in}, \textit{out}) \\ \text{and } (\textit{sto}_2, \textit{in}_2, \textit{out}_2, \textit{pos}) = E[[\texttt{exp}]] & (\textit{env}, \textit{sto}_1, \textit{in}_1, \textit{out}_1) \\ \end{split}
```

```
E[[\operatorname{id\_or\_arr\_access} = \operatorname{exp}]] \ (env, sto, in, out) = (sto_2, in_2, out_2, \operatorname{expVal})  \text{where} \ (sto_1, in_1, out_1, \operatorname{expVal}) = E[[\operatorname{exp}]] \ (env, sto, in, out)  \text{where} \ (sto_2, in_2, out_2) = \operatorname{updateArray}((\operatorname{env}, \operatorname{sto}_1, \operatorname{in}_1, \operatorname{out}_1), \operatorname{id\_or\_arr\_access}, \operatorname{expVal})  \text{where}   \operatorname{updateArray}: \Sigma \times (\operatorname{Identifier} \cup \operatorname{ArrayAccess}) \times \operatorname{ExpressibleValue} \to \operatorname{Store} \times \operatorname{Input} \times \operatorname{Output}   \text{where} \ \operatorname{updateArray}((\operatorname{env}, \operatorname{sto}, \operatorname{in}, \operatorname{out}), \operatorname{id}, \operatorname{val}) = (\operatorname{updateSto}(\operatorname{sto}, \operatorname{loc}, \operatorname{val}), \operatorname{in}, \operatorname{out})   \text{where} \ \operatorname{loc} = \operatorname{applyEnv}(\operatorname{env}, \operatorname{id})   \text{and} \ \operatorname{updateArray}((\operatorname{env}, \operatorname{sto}, \operatorname{in}, \operatorname{out}), \operatorname{id\_or\_arr\_access}[\operatorname{exp}], \operatorname{val}) = (\operatorname{sto}_f, \operatorname{in}_f, \operatorname{out}_f)   \text{where} \ (\operatorname{sto}_1, \operatorname{in}_1, \operatorname{out}_1, \operatorname{array}(\operatorname{arr})) = E[[\operatorname{id\_or\_arr\_access}]] \ (\operatorname{env}, \operatorname{sto}, \operatorname{in}, \operatorname{out})   \text{where} \ (\operatorname{sto}_2, \operatorname{in}_f, \operatorname{out}_f, \operatorname{int}(\operatorname{pos})) = E[[\operatorname{exp}]] \ (\operatorname{env}, \operatorname{sto}_1, \operatorname{in}_1, \operatorname{out}_1)   \text{where} \ \operatorname{sto}_f = \operatorname{updateSto}(\operatorname{sto}_2, \operatorname{arr}_{\operatorname{pos}}, \operatorname{val})
```

# Equações Semânticas - Expressões

```
\begin{split} E[[\mathrm{id}()]] & (\mathit{env}, \mathit{sto}, \mathit{in}, \mathit{out}) = (\mathit{sto}_f, \mathit{in}_f, \mathit{out}_f, \mathit{val}_f) \\ & \text{where } \mathit{fun}(f) = \mathit{applyEnv}(\mathit{env}, \mathrm{id}) \\ & \text{and } (\mathit{sto}_f, \mathit{in}_f, \mathit{out}_f) = f((\emptyset, \mathit{sto}', \mathit{in}, \mathit{out}), \mathit{loc}) \\ & \text{where } (\mathit{sto}', \mathit{loc}) = \mathit{allocate} \; \mathit{sto} \\ & \text{and} \; \mathit{val}_f = \mathit{applySto}(\mathit{sto}_f, \mathit{loc}) \end{split}
```

#### Equações Semânticas - Expressões

```
E[[id(exp_list)]] (env, sto, in, out) = (sto<sub>f</sub>, in<sub>f</sub>, out<sub>f</sub>, val<sub>f</sub>)
   where fun(f) = applyEnv(env, id)
   and (sto_f, in_f, out_f) = f(locations(exp_list, sto', in, out), loc)
      where (sto', loc) = allocate sto
      and
locations: ExpressionList \times Store \times Input \times Output \rightarrow Location[] \times Store \times Input \times Output
        locations(exp, sto, in, out) = ((loc), sto_3, in_1, out_1)
          where (sto_1, in_1, out_1, val) = E[[exp]] (env, sto, in, out)
          and (sto_2, loc) = allocate sto
          and sto_3 = updateSto(sto_2, loc, val)
       and locations(exp_list ", "exp, sto, in, out) = (concat(init, last), sto_2, in_2, out_2)
          where (init, sto_1, in_1, out_1) = locations(exp_list, sto, in, out)
          and (last, sto_2, in_2, out_2) = locations(exp, sto_1, in_1, out_1)
```

```
\begin{split} &C[[\{\deg_{\mathsf{C}} \mathsf{cmd}\_\mathsf{list}\}]] \ (env, sto, in, out) = (sto_f, in_f, out_f, retFlag) \\ & \text{ where } (env_f, sto_f, in_f, out_f, retFlag) = sequence(\deg_{\mathsf{C}} \mathsf{cmd}\_\mathsf{list}) \ (env, sto, in, out) \\ & \text{ and } sequence(\deg_{\mathsf{C}} (env, sto, in, out) = (Dec[[\deg_{\mathsf{C}}]] \ (env, sto, in, out), returnFlag0) \\ & \text{ and } sequence(\mathsf{cmd}) \ (env, sto, in, out) = (env, C[[\mathsf{cmd}]] \ (env, sto, in, out)) \\ & \text{ and } sequence(\deg_{\mathsf{C}} \mathsf{cmd}\_\mathsf{list} \ \deg_{\mathsf{C}} \mathsf{cmd}) \ (env, sto, in, out) = \mathbf{if} \ retFlagRec = returnFlag1 \\ & \text{ then } (env_1, sto_1, in_1, out_1, retFlagRec) \\ & \text{ else } (env_2, sto_2, in_2, out_2, retFlagRec') \\ & \text{ where } (env_1, sto_1, in_1, out_1, retFlagRec) = sequence(\deg_{\mathsf{C}} \mathsf{cmd}\_\mathsf{list}) \ (env, sto, in, out) \\ & \text{ and } (env_2, sto_2, in_2, out_2, retFlagRec') = sequence(\deg_{\mathsf{C}} \mathsf{cmd}) \ (env_1, sto_1, in_1, out_1) \end{split}
```

```
C[[\exp;]] (env, sto, in, out) = (env, sto_f, in_f, out_f, returnFlag0)
\mathbf{where} (sto_f, in_f, out_f, val) = E[[\exp]] (env, sto, in, out)
```

```
\begin{split} &C[[\mathsf{IF}(\mathsf{exp})\ \mathsf{comp\_cmd}]]\ (\mathit{env}, \mathit{sto}, \mathit{in}, \mathit{out}) = \mathsf{if}\ b\\ &\quad \mathsf{then}\ C[[\mathsf{comp\_cmd}]]\ (\mathit{env}, \mathit{sto}_e, \mathit{in}_e, \mathit{out}_e)\\ &\quad \mathsf{else}\ (\mathit{sto}_e, \mathit{in}_e, \mathit{out}_e, \mathit{returnFlag}\,0)\\ &\quad \mathsf{where}\ (\mathit{sto}_e, \mathit{in}_e, \mathit{out}_e, \mathit{bool}(\mathit{b})) = E[[\mathsf{exp}]]\ (\mathit{env}, \mathit{sto}, \mathit{in}, \mathit{out}) \end{split} &C[[\mathsf{IF}(\mathsf{exp})\ \mathsf{comp\_cmd}_1\ \mathsf{ELSE}\ \mathsf{comp\_cmd}_2]]\ (\mathit{env}, \mathit{sto}, \mathit{in}, \mathit{out}) = \mathsf{if}\ \mathit{b}\\ &\quad \mathsf{then}\ C[[\mathsf{comp\_cmd}_1]]\ (\mathit{env}, \mathit{sto}_e, \mathit{in}_e, \mathit{out}_e)\\ &\quad \mathsf{else}\ C[[\mathsf{comp\_cmd}_2]]\ (\mathit{env}, \mathit{sto}_e, \mathit{in}_e, \mathit{out}_e)\\ &\quad \mathsf{where}\ (\mathit{sto}_e, \mathit{in}_e, \mathit{out}_e, \mathit{bool}(\mathit{b})) = E[[\mathsf{exp}]]\ (\mathit{env}, \mathit{sto}, \mathit{in}, \mathit{out}) \end{split}
```

```
\begin{split} &C[[\texttt{while} \ (\texttt{exp}) \ \texttt{cmd}]] \ (\texttt{env}, \texttt{sto}, \textit{in}, \textit{out}) = \textit{loop}(\texttt{sto}, \textit{in}, \textit{out}, \textit{returnFlag0}) \\ & \text{where} \ \textit{loop}(\texttt{sto}, \textit{in}, \textit{out}, \textit{retFlag}) = \ \textbf{if} \ \textit{retFlag} = \textit{returnFlag1} \\ & \text{then} \ (\texttt{sto}, \textit{in}, \textit{out}, \textit{returnFlag1}) \\ & \text{else if} \ \textit{b} \\ & \text{then} \ \textit{loop}(C[[\texttt{cmd}]] \ (\textit{env}, \texttt{sto}_e, \textit{in}_e, \textit{out}_e)) \\ & \text{else} \ (\texttt{sto}_e, \textit{in}_e, \textit{out}_e, \textit{returnFlag0}) \\ & \text{where} \ (\texttt{sto}_e, \textit{in}_e, \textit{out}_e, \textit{bool}(\textit{b})) = E[[\texttt{exp}]] \ (\textit{env}, \texttt{sto}, \textit{in}, \textit{out}) \end{split}
```

```
C[[\texttt{RETURN}\,;]] \ (\textit{env}, \textit{sto}, \textit{in}, \textit{out}) = \ (\textit{sto}, \textit{in}, \textit{out}, \textit{returnFlag1}) C[[\texttt{RETURN} \ \texttt{exp}\,;]] \ (\textit{env}, \textit{sto}, \textit{in}, \textit{out}) = \ (\textit{sto}', \textit{in}', \textit{out}', \textit{returnFlag1})  \  \text{where} \ (\textit{sto}_{\textbf{e}}, \textit{in}', \textit{out}', \textit{val}) = E[[\texttt{exp}]] \ (\textit{env}, \textit{sto}, \textit{in}, \textit{out})  \  \text{and} \ \textit{sto}' = \textit{updateSto}(\textit{sto}_{\textbf{e}}, \textit{applyEnv}(\textit{env}, \textit{return}), \textit{val}) C[[\texttt{SKIP};]] \ (\textit{env}, \textit{sto}, \textit{in}, \textit{out}) = \ (\textit{sto}, \textit{in}, \textit{out}, \textit{returnFlag0})
```

## Equações Semânticas - Declarações

```
\begin{aligned} Dec[[\mathsf{type\_spec}\ \mathsf{id}\,;]]\ (env,sto,in,out) &= (env_f,sto_f,in,out)\\ &\quad \mathsf{where}\ env_f = extendEnv(env,\mathsf{id},var(loc))\\ &\quad \mathsf{and}\ (sto_f,loc) = \mathit{allocate}\ sto \end{aligned} Dec[[\mathsf{type\_spec}\ \mathsf{id}\ = \exp\,;]]\ (env,sto,in,out) &= (env_f,sto_f,in_f,out_f)\\ &\quad \mathsf{where}\ (env_f,sto_1,in_1,out_1) &= Dec[[\mathsf{type\_spec}\ \mathsf{id}\,;]]\ (env,sto,in,out)\\ &\quad \mathsf{and}\ (sto_f,in_f,out_f,val) &= E[[\mathsf{id}\ = \exp\,]](env_1,sto_1,in_1,out_1) \end{aligned}
```

# Equações Semânticas - Definições

```
\begin{split} \textit{Def} & [ [\texttt{type\_spec id() comp\_cmd}]] \ (\textit{env}, \textit{sto}, \textit{in}, \textit{out}) = \textit{env}_1 \\ & \text{where } \textit{env}_1 = \textit{extendEnv}(\textit{env}, \textit{id}, \textit{fun}(\textit{f})) \\ & \text{and } \textit{f}(\emptyset, \textit{sto}, \textit{in}, \textit{out}) = \textit{C}[[\texttt{comp\_cmd}]] \ (\textit{env}_1, \textit{sto}, \textit{in}, \textit{out}) \end{split}
```

#### Equações Semânticas - Definições

```
Def[[type_spec id(param_dec_list) comp_cmd]] (env, sto, in, out) = env1
where env1 = extendEnv(env, id, fun(f))
and f(array, sto, in, out) = C[[comp_cmd]] (env', sto, in, out)
and env' = modifyEnv(env1, param_dec_list, array)
where modifyEnv(env, type_spec id, array) = extendEnv(env, id, array_0)
and modifyEnv(env, param_dec_list, type_spec id, array) = extendEnv(env2, id, last)
where env2 = modifyEnv(env, param_dec_list, init)
where init = init(array)
and last = last(array)
```

#### Equações Semânticas - Programa

```
P_1[[\mathtt{type\_spec}\ \mathtt{id}\,;]]\ (\mathit{env},\mathit{sto},\mathit{in},\mathit{out}) = \mathit{Dec}[[\mathtt{type\_spec}\ \mathtt{id}\,;]]\ (\mathit{env},\mathit{sto},\mathit{in},\mathit{out})
P_1[[\mathtt{type\_spec}\ \mathtt{id}\,=\,\mathtt{exp}\,;]]\ (\mathit{env},\mathit{sto},\mathit{in},\mathit{out}) = \mathit{Dec}[[\mathtt{type\_spec}\ \mathtt{id}\,;]]\ (\mathit{env},\mathit{sto},\mathit{in},\mathit{out})
P_1[[\mathtt{fun\_def}]]\ (\mathit{env},\mathit{sto},\mathit{in},\mathit{out}) = (\mathit{env}_1,\mathit{sto},\mathit{in},\mathit{out})
\mathbf{where}\ \mathit{env}_1 = \mathit{Def}[[\mathtt{fun\_def}]]
P_1[[\mathtt{prog}\ \mathtt{dec\_or\_fun\_def}]]\ (\mathit{env},\mathit{sto},\mathit{in},\mathit{out}) = P_1[[\mathtt{dec\_or\_fun\_def}]]\ (P_1[[\mathtt{prog}]]\ (\mathit{env},\mathit{sto},\mathit{in},\mathit{out}))
```

$$P_{2} = P_{1}$$

## Equações Semânticas - Programa

```
P_{3}[[\mathtt{dec}]] = Dec[[\mathtt{dec}]]
P_{3}[[\mathtt{fun\_def}]] (\mathit{env}, \mathit{sto}, \mathit{in}, \mathit{out}) = (\mathit{env}, \mathit{sto}, \mathit{in}, \mathit{out})
P_{3}[[\mathtt{prog} \ \mathtt{dec\_or\_fun\_def}]] (\mathit{env}, \mathit{sto}, \mathit{in}, \mathit{out}) = P_{3}[[\mathtt{dec\_or\_fun\_def}]] (P_{3}[[\mathtt{prog}]] (\mathit{env}, \mathit{sto}, \mathit{in}, \mathit{out}))
```

## Equações Semânticas - Programa

```
\begin{split} P[[\mathsf{prog}]](\mathit{env}, \mathit{sto}, \mathit{in}, \mathit{out}) &= (\mathit{sto}_f, \mathit{in}_f, \mathit{out}_f) \\ \text{where } (\mathit{env}_1, \mathit{sto}_1, \mathit{in}_1, \mathit{out}_1) &= (P_3[[\mathsf{prog}]] \circ P_2[[\mathsf{prog}]] \circ P_1[[\mathsf{prog}]]) \ (\mathit{env}, \mathit{sto}, \mathit{in}, \mathit{out}) \\ \text{and } (\mathit{sto}_f, \mathit{in}_f, \mathit{out}_f, \mathit{retVal}) &= E[[\mathsf{main}()]] \ (\mathit{env}_f, \mathit{sto}_1, \mathit{in}_1, \mathit{out}_1) \end{split}
```