1 Syntax

 $code \in Code ::= \frac{code}{rsP}$

```
i \in \text{Integer} ::= \mathbf{0} \mid \mathbf{1} \mid \mathbf{-1} \mid \mathbf{2} \mid \mathbf{-2} \mid \dots
id \in \text{Identifier} ::= \dots
str \in String ::= \dots
p \in PythonCode ::= \dots
sP \in \text{SubProgram} ::= \mathbf{sp}_{id} \mid \dots
rsP \in RawSubProgram ::= `str`
s \in \text{State} ::= \mathbf{s}_{id}
\Gamma \in \text{Environment} ::= \Gamma_{id}
                                                                                                                                                                                   (in code: G)
c \in \text{Construction} ::= rsP \mid s \mid \Gamma \mid sP
cl \in ConstructionList ::= cl_1 , cl_2 \mid c \mid \varepsilon
ap \in ApplyPredicate ::= id(cl_1 | cl_2)
d \in \text{DefinePredicate} ::= id(cl_1 \mid cl_2) \{ p \}
tr \in \text{Transition} ::= \langle c_1, s_1 \rangle \rightarrow \langle c_2, s_2 \rangle \mid \langle c, s_1 \rangle \rightarrow s_2
                                                                                                                                                                                (in code: =>)
uo \in UpOpSemRule ::= uo_1 \ uo_2 \mid ap \mid tr \mid \varepsilon
rel \in \text{Relation} ::= : | <:
ty \in \text{Typing} ::= \Gamma \vdash c_1 \ rel \ c_2
                                                                                                                                                                                  (in code: |-)
ut \in UpTypingRule ::= ut_1 ut_2 \mid ap \mid ty \mid \varepsilon
P \in \text{Program} ::= P_1 P_2 \mid R \mid D \mid \{P\} \mid \text{!popEnv} \mid code
r \in \text{Rule} ::= S \mid O \mid T
rs \in SyntaxRule ::= syntax(id) \{ p \}
ro \in \text{OperationalSemanticsRule} ::= \text{semantics}(id)
                                                                                         }
                                                                                                                                                                            (in code: ---)
rt \in \text{TypingRule} ::= \text{typing}(id)
                                                             }
```

```
State: \hat{s} = \{[E_n, E_{n-1}, \dots, E_0], s\} ({environments, program state: s = \langle st_s, \Gamma_s \rangle \})
Starting configuration (<constrution, state>): < P, \{[\emptyset], <\emptyset, \emptyset > \} >
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                          we expect state as pair of store, gamma
Environment: E = \langle S_E, O_E, T_E, D_E \rangle- elements of environment E
Syntax Environment: S_E: Identifier \rightharpoonup PythonCode
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                             UpOpSemRule
 Operational Semantics Environment: O_E: Identifier \rightharpoonup
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                       (formally: UpOpSemRule × Transition)
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                  Transition
Typing Environment: T_E: Identifier \rightharpoonup
Predicate Environment: D_E: Identifier \rightarrow (ConstructionList | ConstructionList ) { PythonCode }
Predefined system predicates: PARSE(S, rsP, sP), UNIFY-AND-RUN(( cl_1 \mid cl_2 ), ( cl_3 \mid cl_4 ){ p })
            S \vdash \mathtt{PARSE}(uo, uo')
                                                                                           \frac{SE(uo,\,uo')}{S \vdash PARSE(tr,\,tr')} \underbrace{\begin{array}{c} S \vdash PARSE(tr,\,tr') \\ uo & uo' \\ S \vdash PARSE(\begin{array}{c} -\\ -\\ tr \end{array}, \begin{array}{c} -\\ -\\ tr' \end{array}) \\ P^{-RO} \end{array} \xrightarrow{P^{-RO}} \underbrace{\begin{array}{c} S \vdash PARSE(ut,\,ut') & S \vdash PARSE(ty,\,ty') \\ ut & ut' \\ S \vdash PARSE(\begin{array}{c} -\\ -\\ ty \end{array}, \begin{array}{c} -\\ -\\ ty \end{array})}_{P^{-RT}}
         \frac{S \vdash \mathtt{PARSE}(uo_1,\,uo_1') \qquad S \vdash \mathtt{PARSE}(uo_2,\,uo_2')}{S \vdash \mathtt{PARSE}(uo_1\,\,uo_2,\,uo_1'\,\,uo_2')} \xrightarrow{\mathtt{P-uo12}} \qquad \frac{S \vdash \mathtt{PARSE}(ut_1,\,ut_1') \qquad S \vdash \mathtt{PARSE}(ut_2,\,ut_2')}{S \vdash \mathtt{PARSE}(ut_1\,\,ut_2,\,ut_1'\,\,ut_2')} \xrightarrow{\mathtt{P-ut12}}
         \frac{\text{PARSE}(S,\,c_1,\,c_1')}{S \vdash \text{PARSE}(<\!c_1\,,\,s_1\!> \,\to\, <\!c_2\,,\,s_2\!>,\,<\!c_1'\,,\,s_1\!> \,\to\, <\!c_2'\,,\,s_2\!>)}{S \vdash \text{PARSE}(<\!c_1\,,\,s_1\!> \,\to\, <\!c_2'\,,\,s_1\!> \,\to\, <\!c_2'\,,\,s_2\!>)} \xrightarrow{\text{P-trend}} \frac{\text{PARSE}(S,\,c,\,c')}{S \vdash \text{PARSE}(<\!c_1\,,\,s_1\!> \,\to\, s_2,\,<\!c'\,,\,s_1\!> \,\to\, c'\,,\,s_2)} \xrightarrow{\text{P-trend}} \frac{(s_1,\,s_2)}{(s_1,\,s_2)} \xrightarrow{\text{P-trend}} \frac{(s_2,\,s_2)}{(s_2,\,s_2)} \xrightarrow{\text{P-tre
          \frac{\texttt{PARSE}(S,\,c_1,\,c_1') \qquad \texttt{PARSE}(S,\,c_2,\,c_2')}{S \vdash \texttt{PARSE}(\Gamma \vdash c_1 \;rel \;c_2,\,\Gamma \vdash c_1' \;rel \;c_2')} \xrightarrow{\texttt{P-ty}} \frac{S \vdash \texttt{PARSE}(cl_1,\,cl_1') \qquad S \vdash \texttt{PARSE}(cl_2,\,cl_2')}{S \vdash \texttt{PARSE}(cl_1\;,\,cl_2\;,\,cl_1'\;,\,cl_2')} \xrightarrow{\texttt{P-cl}} \frac{\texttt{PARSE}(S,\,c,\,c')}{S \vdash \texttt{PARSE}(c,\,c')} \xrightarrow{\texttt{P-cl}} \frac{\texttt{PARSE}(S,\,c_2,\,c_2')}{S \vdash \texttt{PARSE}(C,\,c_2')} \xrightarrow{\texttt{P-cl}} \frac{\texttt{PARSE}(S,\,c_2,\,c_2')}{S \vdash \texttt{PARSE}(C,\,c_2')} \xrightarrow{\texttt{P-cl}} \frac{\texttt{PARSE}(S,\,c_2,\,c_2')}{S \vdash \texttt{PARSE}(C,\,c_2')} \xrightarrow{\texttt{P-cl}} \frac{\texttt{PARSE}(S,\,c_2,\,c_2')}{S \vdash \texttt{PARSE}(C,\,c_2')} \xrightarrow{\texttt{P-cl}} \frac{\texttt{P-cl}}{S \vdash \texttt{PARSE}(C,\,c_2')} \xrightarrow{\texttt{P-cl}} \frac{\texttt{PARSE}(S,\,c_2,\,c_2')}{S \vdash \texttt{PARSE}(S,\,c_2')} \xrightarrow{\texttt{P-cl}} \frac{\texttt{P-cl}}{S \vdash \texttt{PARSE}(C,\,c_2')} \xrightarrow{\texttt{P-cl}} \frac{\texttt{P-cl}}{S \vdash \texttt{P-cl}} \xrightarrow{\texttt{P-cl}} \frac{\texttt{P-cl
         \frac{\mathtt{PARSE}(S,\, cl_1,\, cl_1') \quad \mathtt{PARSE}(S,\, cl_2,\, cl_2')}{S \vdash \mathtt{PARSE}(id(\,\, cl_1 \mid cl_2 \,\,),\, id(\,\, cl_1' \mid cl_2' \,\,))} \xrightarrow{\mathtt{P-ap}} \quad \frac{\mathtt{PARSE}(S,\, cl_1,\, cl_1') \quad \mathtt{PARSE}(S,\, cl_2,\, cl_2')}{S \vdash \mathtt{PARSE}(id(\,\, cl_1 \mid cl_2 \,\,)\{\,\, p \,\,\},\, id(\,\, cl_1' \mid cl_2' \,\,)\{\,\, p \,\,\})} \xrightarrow{\mathtt{P-ap}} \quad \frac{\mathtt{P-ap}}{S \vdash \mathtt{PARSE}(id(\,\, cl_1 \mid cl_2 \,\,)\{\, p \,\,\},\, id(\,\, cl_1' \mid cl_2' \,\,)\{\,\, p \,\,\})} \xrightarrow{\mathtt{P-ap}} \quad \frac{\mathtt{P-ap}}{S \vdash \mathtt{PARSE}(id(\,\, cl_1 \mid cl_2 \,\,)\{\, p \,\,\},\, id(\,\, cl_1' \mid cl_2' \,\,)\{\,\, p \,\,\})} \xrightarrow{\mathtt{P-ap}} \quad \frac{\mathtt{P-ap}}{S \vdash \mathtt{PARSE}(id(\,\, cl_1 \mid cl_2 \,\,)\{\, p \,\,\},\, id(\,\, cl_1' \mid cl_2' \,\,)\{\,\, p \,\,\})} \xrightarrow{\mathtt{P-ap}} \quad \frac{\mathtt{P-ap}}{S \vdash \mathtt{PARSE}(id(\,\, cl_1 \mid cl_2 \,\,)\{\, p \,\,\},\, id(\,\, cl_1' \mid cl_2' \,\,)\{\,\, p \,\,\})} \xrightarrow{\mathtt{P-ap}} \quad \frac{\mathtt{P-ap}}{S \vdash \mathtt{PARSE}(id(\,\, cl_1 \mid cl_2 \,\,)\{\, p \,\,\},\, id(\,\, cl_1' \mid cl_2' \,\,)\{\, p \,\,\})} \xrightarrow{\mathtt{P-ap}} \quad \frac{\mathtt{P-ap}}{S \vdash \mathtt{PARSE}(id(\,\, cl_1 \mid cl_2 \,\,)\{\, p \,\,\},\, id(\,\, cl_1' \mid cl_2' \,\,)\{\, p \,\,\})} \xrightarrow{\mathtt{P-ap}} \quad \frac{\mathtt{P-ap}}{S \vdash \mathtt{PARSE}(id(\,\, cl_1 \mid cl_2 \,\,)\{\, p \,\,\},\, id(\,\, cl_1' \mid cl_2' \,\,)\{\, p \,\,\})} \xrightarrow{\mathtt{P-ap}} \quad \frac{\mathtt{P-ap}}{S \vdash \mathtt{PARSE}(id(\,\, cl_1 \mid cl_2 \,\,)\{\, p \,\,\})} \xrightarrow{\mathtt{P-ap}} \quad \frac{\mathtt{P-ap}}{S \vdash \mathtt{PARSE}(id(\,\, cl_1 \mid cl_2 \,\,)\{\, p \,\,\})} \xrightarrow{\mathtt{P-ap}} \quad \frac{\mathtt{P-ap}}{S \vdash \mathtt{PARSE}(id(\,\, cl_1 \mid cl_2 \,\,)\{\, p \,\,\})} \xrightarrow{\mathtt{P-ap}} \quad \frac{\mathtt{P-ap}}{S \vdash \mathtt{PARSE}(id(\,\, cl_1 \mid cl_2 \,\,)\{\, p \,\,\})} \xrightarrow{\mathtt{P-ap}} \quad \frac{\mathtt{P-ap}}{S \vdash \mathtt{PARSE}(id(\,\, cl_1 \mid cl_2 \,\,)\{\, p \,\,\})} \xrightarrow{\mathtt{P-ap}} \quad \frac{\mathtt{P-ap}}{S \vdash \mathtt{PARSE}(id(\,\, cl_1 \mid cl_2 \,\,)\{\, p \,\,\})} \xrightarrow{\mathtt{P-ap}} \quad \frac{\mathtt{P-ap}}{S \vdash \mathtt{PARSE}(id(\,\, cl_1 \mid cl_2 \,\,)\{\, p \,\,\})} \xrightarrow{\mathtt{P-ap}} \quad \frac{\mathtt{P-ap}}{S \vdash \mathtt{PARSE}(id(\,\, cl_1 \mid cl_2 \,\,)\{\, p \,\,\})} \xrightarrow{\mathtt{P-ap}} \quad \frac{\mathtt{P-ap}}{S \vdash \mathtt{PARSE}(id(\,\, cl_1 \mid cl_2 \,\,)\{\, p \,\,\})} \xrightarrow{\mathtt{P-ap}} \quad \frac{\mathtt{P-ap}}{S \vdash \mathtt{PARSE}(id(\,\, cl_1 \mid cl_2 \,\,)\{\, p \,\,\})} \xrightarrow{\mathtt{P-ap}} \quad \frac{\mathtt{P-ap}}{S \vdash \mathtt{PARSE}(id(\,\, cl_1 \mid cl_2 \,\,)\{\, p \,\,\})} \xrightarrow{\mathtt{P-ap}} \quad \frac{\mathtt{P-ap}}{S \vdash \mathtt{PARSE}(id(\,\, cl_1 \mid cl_2 \,\,)\{\, p \,\,\})} \xrightarrow{\mathtt{P-ap}} \quad \frac{\mathtt{P-ap}}{S \vdash \mathtt{PARSE}(id(\,\, cl_1 \mid cl_2 \,\,)\{\, p \,\,\})} \xrightarrow{\mathtt{P-ap}} \quad \frac{\mathtt{P-ap}}{S \vdash \mathtt{P-ap}} \xrightarrow{\mathtt{P-ap}} \quad \frac{\mathtt{P-ap}}{S \vdash \mathtt{P-ap}} \xrightarrow{\mathtt{P-ap}} \quad \frac{\mathtt{P-ap}}{S \vdash \mathtt{P-ap}} \xrightarrow{\mathtt{P-ap}} \quad \frac{\mathtt{P-ap}}{S \vdash \mathtt{P-ap}}
\frac{< P_1, \{[E, \ldots], s\} > \to < P_1', \{[E', \ldots], s'\} >}{< P_1P_2, \{[E, \ldots], s\} > \to < P_1'P_2, \{[E', \ldots], s'\} >} \xrightarrow{\text{SR-P1P2}} \\ \frac{< P_1, \{[E, \ldots], s\} > \to \{[E', \ldots], s'\}}{< P_1P_2, \{[E, \ldots], s\} > \to < P_2, \{[E', \ldots], s'\} >} \xrightarrow{\text{SR-P1P2end}} \\ \frac{< P_1P_2, \{[E, \ldots], s\} > \to < P_2, \{[E', \ldots], s'\} >}{< P_1P_2, \{[E, \ldots], s\} >} \xrightarrow{\text{SR-P1P2end}} \\ \frac{< P_1P_2, \{[E, \ldots], s\} > \to < P_2, \{[E', \ldots], s'\} >}{< P_1P_2, \{[E', \ldots], s\} >} \xrightarrow{\text{SR-P1P2end}} \\ \frac{< P_1P_2, \{[E, \ldots], s\} >}{< P_1P_2, \{[E', \ldots], s'\} >} \xrightarrow{\text{SR-P1P2end}} \\ \frac{< P_1P_2, \{[E, \ldots], s\} >}{< P_1P_2, \{[E', \ldots], s'\} >} \xrightarrow{\text{SR-P1P2end}} \\ \frac{< P_1P_2, \{[E, \ldots], s\} >}{< P_1P_2, \{[E', \ldots], s'\} >} \xrightarrow{\text{SR-P1P2end}} \\ \frac{< P_1P_2, \{[E', \ldots], s\} >}{< P_1P_2, \{[E', \ldots], s'\} >} \xrightarrow{\text{SR-P1P2end}} \\ \frac{< P_1P_2, \{[E', \ldots], s\} >}{< P_1P_2, \{[E', \ldots], s'\} >} \xrightarrow{\text{SR-P1P2end}} \\ \frac{< P_1P_2, \{[E', \ldots], s\} >}{< P_1P_2, \{[E', \ldots], s'\} >} \xrightarrow{\text{SR-P1P2end}} \\ \frac{< P_1P_2, \{[E', \ldots], s\} >}{< P_1P_2, \{[E', \ldots], s'\} >} \xrightarrow{\text{SR-P1P2end}} \\ \frac{< P_1P_2, \{[E', \ldots], s\} >}{< P_1P_2, \{[E', \ldots], s'\} >} \xrightarrow{\text{SR-P1P2end}} \\ \frac{< P_1P_2, \{[E', \ldots], s\} >}{< P_1P_2, \{[E', \ldots], s'\} >} \xrightarrow{\text{SR-P1P2end}} \\ \frac{< P_1P_2, \{[E', \ldots], s\} >}{< P_1P_2, \{[E', \ldots], s'\} >} \xrightarrow{\text{SR-P1P2end}} \\ \frac{< P_1P_2, \{[E', \ldots], s\} >}{< P_1P_2, \{[E', \ldots], s'\} >} \xrightarrow{\text{SR-P1P2end}} \\ \frac{< P_1P_2, \{[E', \ldots], s\} >}{< P_1P_2, \{[E', \ldots], s'\} >} \xrightarrow{\text{SR-P1P2end}} \\ \frac{< P_1P_2, \{[E', \ldots], s\} >}{< P_1P_2, \{[E', \ldots], s'\} >} \xrightarrow{\text{SR-P1P2end}} \\ \frac{< P_1P_2, \{[E', \ldots], s\} >}{< P_1P_2, \{[E', \ldots], s'\} >} \xrightarrow{\text{SR-P1P2end}} \\ \frac{< P_1P_2, \{[E', \ldots], s\} >}{< P_1P_2, \{[E', \ldots], s'\} >} \xrightarrow{\text{SR-P1P2end}} \\ \frac{< P_1P_2, \{[E', \ldots], s\} >}{< P_1P_2, \{[E', \ldots], s'\} >} \xrightarrow{\text{SR-P1P2end}} \\ \frac{< P_1P_2, \{[E', \ldots], s\} >}{< P_1P_2, \{[E', \ldots], s'\} >} \xrightarrow{\text{SR-P1P2end}} \\ \frac{< P_1P_2, \{[E', \ldots], s\} >}{< P_1P_2, \{[E', \ldots], s'\} >} \xrightarrow{\text{SR-P1P2end}} \\ \frac{< P_1P_2, \{[E', \ldots], s\} >}{< P_1P_2, \{[E', \ldots], s'\} >} \xrightarrow{\text{SR-P1P2end}} \\ \frac{< P_1P_2, \{[E', \ldots], s\} >}{< P_1P_2, \{[E', \ldots], s'\} >} \xrightarrow{\text{SR-P1P2end}} 
       E' = \langle S_E[id \mapsto p], O_E, T_E, D_E \rangle
\langle \operatorname{syntax}(id) \{ p \}, \{ [E, \ldots], s \} \rangle \rightarrow \{ [E', \ldots], s \}
SR-RS
         S \vdash \mathtt{PARSE}(\begin{array}{c} uo & uo' \\ \hline tr & tr' \end{array}) \qquad E' = < S_E, O_E[id \mapsto \frac{uo'}{tr'}], T_E, D_E > \\ \\ < \underline{\mathsf{semantics}(id)\{} \qquad \underbrace{}_{SR\text{-RO}} \}, \{[E, \ldots], s\} > \to \{[E', \ldots], s\}
         \frac{S \vdash \mathtt{PARSE}(id(\ cl_1 \ | \ cl_2\ )\{\ p\ \},\ id(\ cl_1' \ | \ cl_2'\ )\{\ p\ \})}{< id(\ cl_1 \ | \ cl_2\ )\{\ p\ \}, \{[E,\ldots],s\} > \to \{[E',\ldots],s\}} \xrightarrow{SR-D} \frac{SR-D}{(\ cl_1 \ | \ cl_2\ )\{\ p\ \}, \{[E,\ldots],s\} > \to \{[E',\ldots],s\}}
         \langle \{P\}, \{[E,\ldots],s\} \rangle \rightarrow \langle P! popEnv, \{[E,E,\ldots],s\} \rangle
```

 $< !popEnv, \{[E_n, E_{n-1}, ...], s\} > \rightarrow \{[E_{n-1}, ...], s\}$ SR-popEnv

```
\frac{\text{PARSE}(S_E, rsP, sP)}{< \operatorname{\mathbf{code}} rsP, \{[E, \ldots], s\} > \to < sP, \{[E, \ldots], s\} >} \xrightarrow{\text{SR-rsP}}
D[id] = (cl_3 \mid cl_4) \{ p \} \qquad \text{UNIFY-AND-RUN}((cl_1 \mid cl_2), (cl_3 \mid cl_4) \{ p \})
OT-sP-ap
                                                             D \vdash id(cl_1 \mid cl_2)
  \frac{T_{E}, O_{E}, D_{E} \vdash \langle c_{1}, s_{1} \rangle \rightarrow \langle c_{2}, s_{2} \rangle}{\langle c_{1}, \{[E, \ldots], s_{1}\} \rangle \rightarrow \langle c_{2}, \{[E, \ldots], s_{2}\} \rangle} \xrightarrow{\text{SR-sPend}} \frac{T_{E}, O_{E}, D_{E} \vdash \langle c, s_{1} \rangle \rightarrow s_{2}}{\langle c, \{[E, \ldots], s_{1}\} \rangle \rightarrow \langle \{[E, \ldots], s_{2}\} \rangle} \xrightarrow{\text{SR-sPend}}
                uo_1 \ uo_2
                                          \in Img(O) T, O, D \vdash uo_1, uo_2
O-sP-uo12end
         \langle c, s_1 \rangle \to s_2
                                T, O, D \vdash \langle c, s_1 \rangle \rightarrow s_2
                   uo_1 \ uo_2
                                                     \in Img(O) T, O, D \vdash uo_1, uo_2 T, D \vdash \Gamma_{s_2} \vdash c_2 : unit
     \langle c_1, s_1 \rangle \to \langle c_2, s_2 \rangle
                                                     T, O, D \vdash \langle c_1, s_1 \rangle \rightarrow \langle c_2, s_2 \rangle
                                 = \in Img(O) T, O, D \vdash uo _{\text{O-sP-uoend}}
         \langle c_1, s_1 \rangle \to s_2
                          T, O, D \vdash \langle c_1, s_1 \rangle \rightarrow s_2
                                                     \in Img(O) T, O, D \vdash uo T, D \vdash \Gamma_{s_2} \vdash c_2 :  unit
     \langle c_1, s_1 \rangle \rightarrow \langle c_2, s_2 \rangle
                                                T, O, D \vdash \langle c_1, s_1 \rangle \rightarrow \langle c_2, s_2 \rangle
                                                                                                                                          \in Img(O) T, D \vdash \Gamma_{s_2} \vdash c_2 : \mathbf{unit}
                                     \in Img(O) O-sPend
                                                                                      c_1, s_1 > \to < c_2, s_2 > 
T, O, D \vdash < c_1, s_1 > \to < c_2, s_2 > 
         \langle c_1, s_1 \rangle \rightarrow s_2

T, O, D \vdash \langle c_1, s_1 \rangle \rightarrow s_2
                           \in Img(T) T, D \vdash ut_1, ut_2 <sub>T-sP-ut12</sub>
        ut_1 ut_2
                                                                                                                   ut \in Img(T)
                                                                                                                                                     T,D \vdash ut
            ty
                                                                                                                                                                                                        T, D \vdash ty
                                                                                                                                    T, D \vdash ty
                                  T, D \vdash ty
```