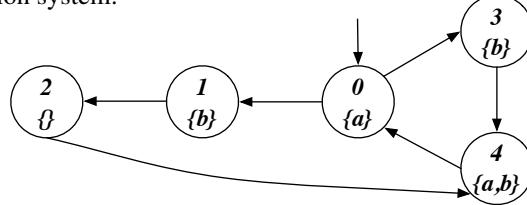
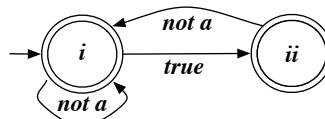


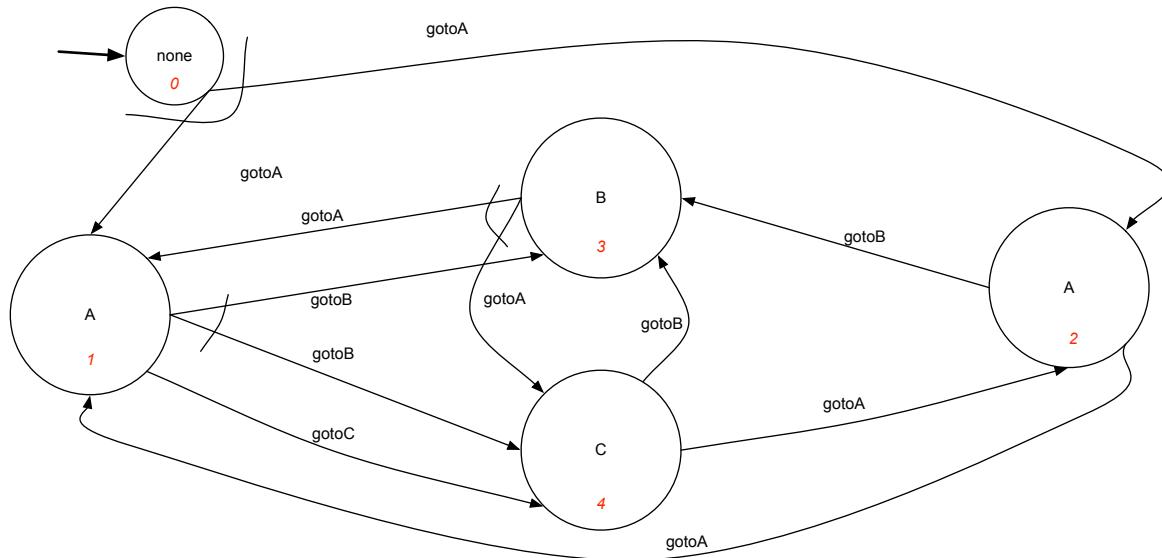
**Part 1.** Consider the following transition system:



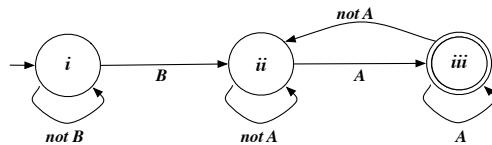
- **Exercise 1.1:** Model check the CTL formula  $AG(EG(AFa \vee EFb))$  by translating it in Mu-Calculus.
- **Exercise 1.2:** Model check the LTL formula  $\diamond(a \wedge \bigcirc a)$ , by considering that the Büchi automaton for  $\neg\diamond(a \wedge \bigcirc a)$  is:



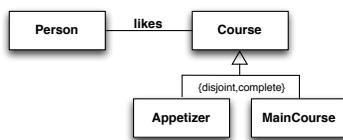
**Part 2.** Consider the following domain:



Synthesize a strategy for realizing the LTLf formula  $\diamond(B \wedge \bigcirc \diamond(A \wedge \bullet \text{false}))$  by considering that the corresponding DFA is:

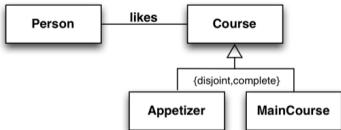


**Part 3.** Consider the following simple UML class diagram, and express in FOL the following boolean queries stating which ones are CQs (do not use abbreviations for cardinalities):



1. Return persons who like an appetizer and a main course.
2. Check if there exists a person who likes two appetizers and a main course.
3. Check if there exists a person who likes exactly one appetizer.
4. Return persons who like all appetizers.
5. Return persons who likes only appetizers.
6. Check if there is a pair of persons such that the first likes all appetizers that the second likes.

**Part 3.** Consider the following simple UML class diagram, and express in FOL the following boolean queries stating which one are CQs (do not use abbreviations for cardinalities):



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1.  $\exists \alpha, m. P(x) \wedge APP(\alpha) \wedge MAIN(m) \wedge LIKES(x, \alpha) \wedge LIKES(x, m)$  ✓
2.  $\exists p, \alpha, \alpha', m. P(p) \wedge APP(\alpha) \wedge APP(\alpha') \wedge MAIN(m) \wedge \alpha \neq \alpha' \wedge LIKES(p, \alpha) \wedge LIKES(p, m) \wedge LIKES(p, \alpha')$  ✗
3.  $\exists p, \alpha. P(p) \wedge APP(\alpha) \wedge LIKES(p, \alpha) \wedge \forall \alpha'. (APP(\alpha') \wedge LIKES(p, \alpha')) \supset \alpha = \alpha'$  ✗
4.  $P(x) \wedge \forall \alpha. (APP(\alpha) \supset LIKES(x, \alpha))$  ✗
5.  $P(x) \wedge \forall \alpha. (LIKES(x, \alpha) \supset APP(\alpha))$  ✗
6.  $\exists p, p'. P(p) \wedge P(p') \wedge \forall \alpha. (APP(\alpha) \wedge LIKES(p', \alpha) \supset LIKES(p, \alpha))$  ✗