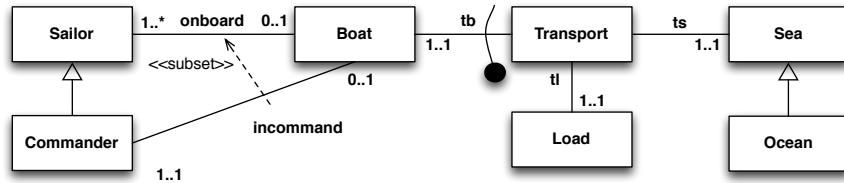


Exercise 1. Express the following UML class diagram in *FOL*.

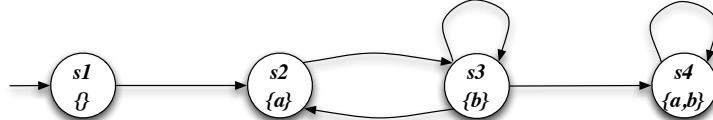


Exercise 2. Consider the above UML class diagram and the following (partial) instantiation.

<i>Sailors</i>	<i>Commander</i>	<i>Boat</i>	<i>onboard</i>	<i>incommand</i>	<i>Sea</i>	<i>Ocean</i>
Dustin Lubber Rusty	Alice Beth	Constitution Enterprise	Dustin Lubber Rusty	Constitution Constitution Constitution	Mediterranean	Atlantic

1. Check whether the above instantiation, once completed, is correct, and explain why it is or it is not.
2. Express in *FOL* the following queries and evaluate them over the completed instantiation:
 - (a) Return the commanders that command a boat with at least a sailor who is not a commander.
 - (b) Return the boats that have on board only commanders
 - (c) Return the boats that have on board all the sailors who are not commanders.

Exercise 3. Model check the Mu-Calculus formula $\nu X.\mu Y.((a \wedge \langle \text{next} \rangle X) \vee (\neg b \wedge \langle \text{next} \rangle Y))$ and the CTL formula $EF(a \wedge EX(AGb))$ (showing its translation in Mu-Calculus) against the following transition system:



Exercise 4. Check whether the following Hoare triple is correct, using as *invariant* $i + j \leq 10$.

{ $i=0$ AND $j=10$ } while($i < 10$) do { $j=j-1$; $i := i+1$ } { $j=0$ }

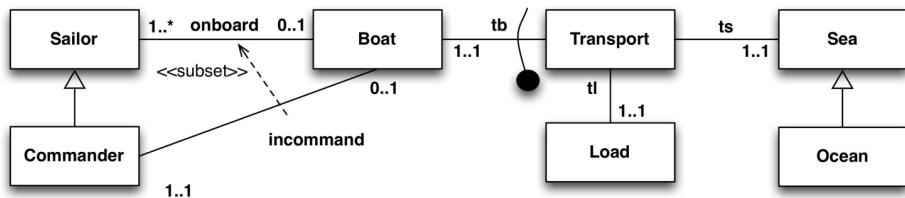
Exercise 5. Given the following boolean conjunctive queries:

```

q1() :- edge(r,b), edge(b,g), edge(g,r).
q2() :- edge(i,f), edge(f,v), edge(v,i), edge(i,a), edge(a,v), edge(a,s), edge(s,i).
  
```

check whether $q1$ is contained into $q2$, explaining the technique used and, in case of containment, showing the homomorphism between the canonical databases.

Exercise 1. Express the following UML class diagram in FOL.



$S(x), C(x), B(x), T(x), L(x), SEA(x), OCEAN(x)$

$ONBOARD(x, y)$

$INCOMMAND(x, y)$

$TB(x, y)$

$TL(x, y)$

$TS(x, y)$

$\forall x, y. ONBOARD(x, y) \supset S(x) \wedge B(y)$

$\forall x. S(x) \supset 0 \leq \#\{y | ONBOARD(x, y)\} \leq 1$

$\forall y. B(y) \supset \exists x. ONBOARD(x, y)$

$\forall x, y. INCOMMAND(x, y) \supset C(x) \wedge B(y)$

$\forall x. C(x) \supset 0 \leq \#\{y | INCOMMAND(x, y)\} \leq 1$

$\forall y. B(y) \supset 1 \leq \#\{x | INCOMMAND(x, y)\} \leq 1$

$\forall x, y. INCOMMAND(x, y) \supset ONBOARD(x, y)$

$\forall x, y. TB(x, y) \supset B(x) \wedge T(y)$

$\forall y. T(y) \supset 1 \leq \#\{x | TB(x, y)\} \leq 1$

$\forall x, x', y. TB(x, y) \wedge TB(x', y) \supset x = x'$

$\forall x, y. TL(x, y) \supset T(x) \wedge L(y)$

$\forall x. T(x) \supset 1 \leq \#\{y | TL(x, y)\} \leq 1$

$\forall x, y. TS(x, y) \supset T(x) \wedge SEA(y)$

$\forall x. T(x) \supset 1 \leq \#\{y | TS(x, y)\} \leq 1$

$\forall x. C(x) \supset S(x)$

$\forall x. OCEAN(x) \supset SEA(x)$

Exercise 2. Consider the above UML class diagram and the following (partial) instantiation.

Sailors	Commander	Boat	onboard	incommand	Sea	Ocean
Dustin Lubber Rusty	Alice Beth	Constitution Enterprise	Dustin Lubber Rusty	Constitution Constitution Constitution	Alice Beth	Constitution Enterprise

1. Check whether the above instantiation, once completed, is correct, and explain why it is or it is not.
2. Express in FOL the following queries and evaluate them over the completed instantiation:
 - (a) Return the commanders that command a boat with at least a sailor who is not a commander.
 - (b) Return the boats that have on board only commanders
 - (c) Return the boats that have on board all the sailors who are not commanders.

$$1) \text{ SAILOR} = \{ \text{DUST, LUB, RUST, ALI, BET} \}$$

$$\text{SEA} = \{ \text{MEDIT, ATLAN} \}$$

$$\text{ONBOARD} = \{ \dots \dots (\text{ALI}, \text{COST}), (\text{BET}, \text{ENT}) \}$$

FOR ALL RELATIONS, ONCE COMPLETED, ALL TUPLES AND CARDINALS ARE OK ✓

$$2) a. \exists b, s. C(x) \wedge \text{INCOMMAND}(x, b) \wedge \text{ONBOARD}(s, b) \wedge \neg C(s)$$

$$\{ \text{ALICE} \}$$

$$b. B(x) \wedge \forall c. (\text{ONBOARD}(c, x) \supset C(c))$$

$$\{ \text{ENT} \}$$

$$c. B(x) \wedge \forall s. (\neg C(s) \supset \text{ONBOARD}(s, x))$$

$$\{ \}$$

Exercise 3. Model check the Mu-Calculus formula $\nu X. \mu Y. ((a \wedge \langle \text{next} \rangle X) \vee (\neg b \wedge \langle \text{next} \rangle Y))$ and the CTL formula $EF(a \wedge EX(AGb))$ (showing its translation in Mu-Calculus) against the following transition system:



$$1) \nu X. \mu Y. ((a \wedge \langle \text{next} \rangle X) \vee (\neg b \wedge \langle \text{next} \rangle Y))$$

$$[x_0] = \{1, 2, 3, 4\}$$

$$[x_1] = [\mu Y. (a \wedge \langle \text{next} \rangle X_0) \vee (\neg b \wedge \langle \text{next} \rangle Y)]$$

$$[y_0] = \emptyset$$

$$\begin{aligned} [y_1] &= ([\alpha] \cap \text{PREE}(\text{next}, x_0)) \cup ([\neg b] \cap \text{PREE}(\text{next}, y_0)) = \\ &= (\{2, 4\} \cap \{1, 2, 3, 4\}) \cup (\{1, 2\} \cap \emptyset) = \{2, 4\} \end{aligned}$$

$$\begin{aligned} [y_2] &= ([\alpha] \cap \text{PREE}(\text{next}, x_0)) \cup ([\neg b] \cap \text{PREE}(\text{next}, y_1)) = \\ &= (\{2, 4\} \cap \{1, 2, 3, 4\}) \cup (\{1, 2\} \cap \{1, 3, 4\}) = \{1, 2, 4\} \end{aligned}$$

$$\begin{aligned} [y_3] &= ([\alpha] \cap \text{PREE}(\text{next}, x_0)) \cup ([\neg b] \cap \text{PREE}(\text{next}, y_2)) = \\ &= (\{2, 4\} \cap \{1, 2, 3, 4\}) \cup (\{1, 2\} \cap \{1, 3, 4\}) = \{1, 2, 4\} \end{aligned}$$

$$[y_4] = [y_3] = [x_1] = \{1, 2, 4\}$$

$$[x_2] = [\mu Y. (a \wedge \langle \text{next} \rangle X_1) \vee (\neg b \wedge \langle \text{next} \rangle Y)]$$

$$[y_0] = \emptyset$$

$$\begin{aligned} [y_1] &= ([\alpha] \cap \text{PREE}(\text{next}, x_1)) \cup ([\neg b] \cap \text{PREE}(\text{next}, y_0)) = \\ &= (\{2, 4\} \cap \{1, 3, 4\}) \cup (\{1, 2\} \cap \emptyset) = \{4\} \end{aligned}$$

$$\begin{aligned} [y_2] &= ([\alpha] \cap \text{PREE}(\text{next}, x_1)) \cup ([\neg b] \cap \text{PREE}(\text{next}, y_1)) = \\ &= (\{2, 4\} \cap \{1, 3, 4\}) \cup (\{1, 2\} \cap \{4\}) = \{4\} \end{aligned}$$

$$[y_3] = [y_2] = [x_2] = \{4\}$$

$$[x_3] = [\mu Y. (a \wedge \langle \text{next} \rangle X_2) \vee (\neg b \wedge \langle \text{next} \rangle Y)]$$

$$[y_0] = \emptyset$$

$$\begin{aligned} [y_1] &= ([\alpha] \cap \text{PREE}(\text{next}, x_2)) \cup ([\neg b] \cap \text{PREE}(\text{next}, y_0)) = \\ &= (\{2, 4\} \cap \{4\}) \cup (\{1, 2\} \cap \emptyset) = \{4\} \end{aligned}$$

$$[Y_2] = ([\alpha] \cap \text{PREE}(\text{NEXT}, X,)) \cup ([\neg b] \cap \text{PREE}(\text{NEXT}, Y,)) = \\ = (\{2, 4\} \cap \{4\}) \cup (\{1, 2\} \cap \{4\}) = \{4\}$$

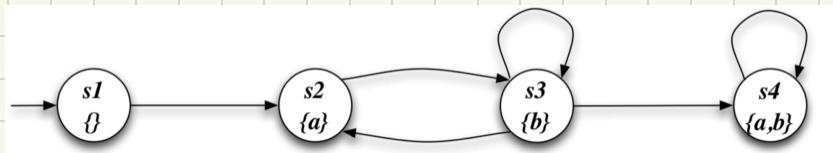
$$[Y_1] = [Y_2] = [X_3] = \{4\}$$

$$[X_2] = [X_3] = \{4\}$$

$$S, \in [\cup X. \mu Y. ((\alpha \wedge \langle \text{NEXT} > X) \vee (\neg b \wedge \langle \text{NEXT} > Y))] = \{4\} ? \quad \text{No!}$$

2) $\text{EF}(\alpha \wedge \text{EX}(\text{AG } b))$

$$\frac{\frac{\alpha}{\beta}}{\delta}$$



$$[\alpha] = [AG b] = [\cup z. b \wedge [NEXT] z]$$

$$[z_0] = \{1, 2, 3, 4\}$$

$$[z_1] = [b] \cap \text{PREA}(\text{NEXT}, z_0) =$$

$$= \{3, 4\} \cap \{1, 2, 3, 4\} = \{3, 4\}$$

$$[z_2] = [b] \cap \text{PREA}(\text{NEXT}, z_1) =$$

$$= \{3, 4\} \cap \{2, 4\} = \{4\}$$

$$[z_3] = [b] \cap \text{PREA}(\text{NEXT}, z_2) =$$

$$= \{3, 4\} \cap \{4\} = \{4\}$$

$$[z_2] = [z_3] = [\alpha] = \{4\}$$

$$[\beta] = [\text{EX } \alpha] = [\langle \text{NEXT} \rangle \alpha] = \text{PREE}(\text{NEXT}, \alpha) = \{3, 4\} = [B]$$

$$[\gamma] = [\alpha \wedge \beta] = [\alpha] \cap [\beta] = \{2, 4\} \cap \{3, 4\} = \{4\} = [\delta]$$

$$[\delta] = [\text{EF } \delta] = [\mu z. \delta \vee \langle \text{NEXT} \rangle z]$$

$$[z_0] = \emptyset$$

$$[z_1] = [\delta] \cup \text{PREE}(\text{NEXT}, z_0) =$$

$$= \{4\} \cup \emptyset = \{4\}$$

$$[z_2] = [\delta] \cup \text{PREE}(\text{NEXT}, z_1) =$$

$$= \{4\} \cup \{3, 4\} = \{3, 4\}$$

$$[z_3] = [\delta] \cup \text{PRE}(\text{NEXT}, z_2) = \\ = \{4\} \cup \{2, 3, 4\} = \{2, 3, 4\}$$

$$[z_4] = [\delta] \cup \text{PRE}(\text{NEXT}, z_3) = \\ = \{4\} \cup \{1, 2, 3, 4\} = \{1, 2, 3, 4\}$$

$$[z_5] = [\delta] \cup \text{PRE}(\text{NEXT}, z_4) = \\ = \{4\} \cup \{1, 2, 3, 4\} = \{1, 2, 3, 4\}$$

$$[z_6] = [z_5] = [d] = \{1, 2, 3, 4\}$$

$$\gamma_{s_i} \in \delta ? \rightarrow s_i \in \delta = \{1, 2, 3, 4\} ? \quad \text{YES!}$$

I

Exercise 4. Check whether the following Hoare triple is correct, using as invariant $i + j \leq 10$.

$$\{i=0 \text{ AND } j=10\} \quad \text{while } (i < 10) \text{ do } \{j=j-1; i := i+1\} \quad \{j=0\}$$

$$1. P \triangleright I \quad 2. \{g \wedge I\} \triangleright Q \quad 3. \{g \wedge I\} \circ \{I\}$$

$$1. \{i=0 \wedge j=10\} \triangleright i+j \leq 10 \quad \checkmark$$

$$2. \{i \geq 10 \wedge i+j \leq 10\} \triangleright j=0 \quad \checkmark$$

$$3. \{i < 10 \wedge i+j \leq 10\} \quad j=j-1; i=i+1 \quad \{i+j \leq 10\}$$

$$\{i < 10 \wedge i+j \leq 10\} \triangleright \text{wp}(j=j-1; i=i+1)\{i+j \leq 10\} \quad ?$$

$$\begin{aligned} \{i+j \leq 9\} [j/j-1] &= \{i+j \leq 10\} \\ j &= j-1; \\ \{i+j \leq 10\} [i/i+1] &= \{i+j \leq 9\} \\ i &= i+1; \\ \{i+j \leq 10\} & \end{aligned}$$

$$\{i < 10 \wedge i+j \leq 10\} \triangleright \{i+j \leq 10\} \quad ?$$

$\{i+j \leq 10\}$ IS AN INVARIANT

Exercise 5. Given the following boolean conjunctive queries:

$q_1() :- \text{edge}(r, b), \text{edge}(b, g), \text{edge}(g, r).$

$q_2() :- \text{edge}(i, f), \text{edge}(f, v), \text{edge}(v, i), \text{edge}(i, a), \text{edge}(a, v), \text{edge}(a, s), \text{edge}(s, i).$

check whether q_1 is contained into q_2 , explaining the technique used and, in case of containment, showing the homomorphism between the canonical databases.

BUILD CANONICAL INTERPRETATION

$I_{q_1} : \Delta I_{q_1} : \{r, g, b\}$

$\text{EDGE } q_1^{\alpha} : \{(r, b), (b, g), (g, r)\}$

$I_{q_2} : \Delta I_{q_2} : \{a, i, f, v, s\}$

$\text{EDGE } q_2^{\alpha} : \{(i, f), (f, v), (v, i), (i, a), (a, v), (a, s), (s, i)\}$

QUERY ANSWERING

$I_{q_1} \models I_{q_2} ?$

$$\alpha(i) = r$$

$$\alpha(f) = b$$

$$\alpha(v) = g$$

$$\alpha(a) = b$$

$$\alpha(s) = g$$

$I_{q_1}, \alpha \models q_2$ YES!

HOMOMORPHISM

$$h(i) = \alpha(i) = r$$

$$h(f) = \alpha(f) = b$$

$$h(v) = \alpha(v) = g$$

$$h(a) = \alpha(a) = b$$

$$h(s) = \alpha(s) = g$$

$$(i, f) \in \text{EDGE } q_2^{\alpha} \Rightarrow (h(i), h(f)) \in \text{EDGE } q_1^{\alpha}$$

$$(f, v) \in \text{EDGE } q_2^{\alpha} \Rightarrow (h(f), h(v)) \in \text{EDGE } q_1^{\alpha}$$

$$(v, i) \in \text{EDGE } q_2^{\alpha} \Rightarrow (h(v), h(i)) \in \text{EDGE } q_1^{\alpha}$$

$$(i, a) \in \text{EDGE } q_2^{\alpha} \Rightarrow (h(i), h(a)) \in \text{EDGE } q_1^{\alpha}$$

$$(a, v) \in \text{EDGE } q_2^{\alpha} \Rightarrow (h(a), h(v)) \in \text{EDGE } q_1^{\alpha}$$

$$(a, s) \in \text{EDGE } q_2^{\alpha} \Rightarrow (h(a), h(s)) \in \text{EDGE } q_1^{\alpha}$$

$$(s, i) \in \text{EDGE } q_2^{\alpha} \Rightarrow (h(s), h(i)) \in \text{EDGE } q_1^{\alpha}$$

