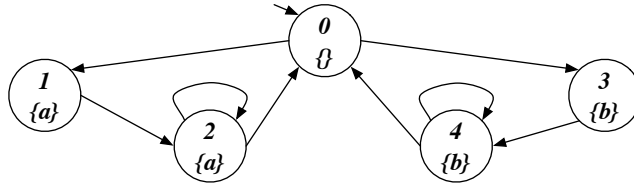
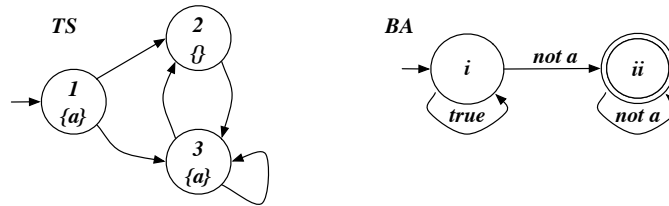


Part 1. Consider the following transition system:

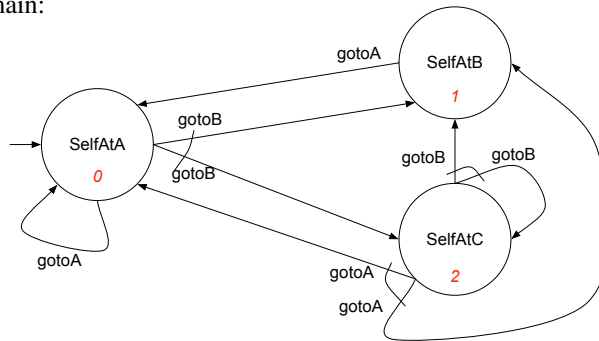


- **Exercise 1.1:** Model check the Mu-Calculus formula: $\nu X. \mu Y. ((a \wedge [next]X) \vee (b \wedge \langle next \rangle Y))$
- **Exercise 1.2:** Model check the CTL formula $AF(EG(a \supset AXEX\neg a))$, by translating it in Mu-Calculus.

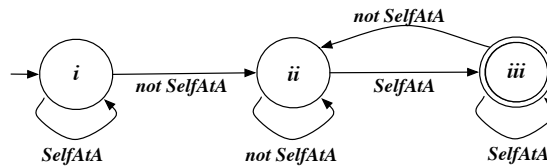
Part 2 Consider the transition system TS below. Model check the LTL formula $\Box \Diamond a$, by considering that the Büchi automaton BA for $\neg \Box \Diamond a$ (i.e., $\Diamond \Box \neg a$) is the one below:



Part 3 Consider the following domain:



- **Exercise 2.1:** Synthesize a strategy (a plan) for realizing the LTLf formula $\Diamond(\neg SelfAtA \wedge \Diamond(SelfAtA \wedge \bullet false))$, by considering that the corresponding DFA is the one below:



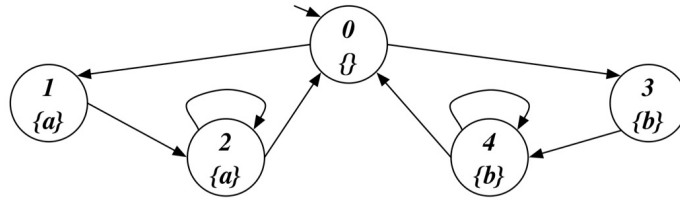
Part 4 (optional) Consider the following program:

```
while (x<10) do x := x + 1
```

Compute its *execution* and *final state*, starting from an *initial state* where $x = 9$, by using:

1. *Evaluation Semantics;*
2. *Transition Semantics.*

Part 1. Consider the following transition system:



- **Exercise 1.1:** Model check the Mu-Calculus formula: $\nu X. \mu Y. ((a \wedge [next]X) \vee (b \wedge \langle next \rangle Y))$
- **Exercise 1.2:** Model check the CTL formula $AF(EG(a \supset AXEX\neg a))$, by translating it in Mu-Calculus.

$$\varphi = \nu X. \mu Y. ((a \wedge [NEXT]X) \vee (b \wedge \langle NEXT \rangle Y))$$

$$[X_0] = \{0, 1, 2, 3, 4\}$$

$$[X_1] = [\mu Y. ((a \wedge [NEXT]X_0) \vee (b \wedge \langle NEXT \rangle Y))]$$

$$[Y_0] = \emptyset$$

$$[Y_1] = ([a] \wedge PREA(NEXT, X_0)) \cup ([b] \wedge FREE(NEXT, Y_0)) =$$

$$= (\{1, 2\} \cap \{0, 1, 2, 3, 4\}) \cup (\{3, 4\} \cap \emptyset) = \{1, 2\}$$

$$[Y_2] = ([a] \wedge PREA(NEXT, X_0)) \cup ([b] \wedge FREE(NEXT, Y_1)) =$$

$$= (\{1, 2\} \cap \{0, 1, 2, 3, 4\}) \cup (\{3, 4\} \cap \{0, 1, 2\}) = \{1, 2\}$$

$$[Y_1] = [Y_2] = [X_1] = \{1, 2\}$$

$$[X_2] = [\mu Y. ((a \wedge [NEXT]X_1) \vee (b \wedge \langle NEXT \rangle Y))]$$

$$[Y_0] = \emptyset$$

$$[Y_1] = ([a] \wedge PREA(NEXT, X_1)) \cup ([b] \wedge FREE(NEXT, Y_0)) =$$

$$= (\{1, 2\} \cap \{1\}) \cup (\{3, 4\} \cap \emptyset) = \{1\}$$

$$[Y_2] = ([a] \wedge PREA(NEXT, X_1)) \cup ([b] \wedge FREE(NEXT, Y_0)) =$$

$$= (\{1, 2\} \cap \{1\}) \cup (\{3, 4\} \cap \{0\}) = \{1\}$$

$$[Y_1] = [Y_2] = [X_2] = \{1\}$$

$$[X_3] = [\mu Y. ((a \wedge [NEXT]X_2) \vee (b \wedge \langle NEXT \rangle Y))]$$

$$[Y_0] = \emptyset$$

$$[Y_1] = ([a] \wedge PREA(NEXT, X_2)) \cup ([b] \wedge FREE(NEXT, Y_0)) =$$

$$= (\{1, 2\} \cap \emptyset) \cup (\{3, 4\} \cap \emptyset) = \emptyset$$

$$[Y_0] = [Y_1] = [X_3] = \emptyset$$

$$[X_4] = [\mu Y. ((a \wedge [NEXT]X_3) \vee (b \wedge \langle NEXT \rangle Y))]]$$

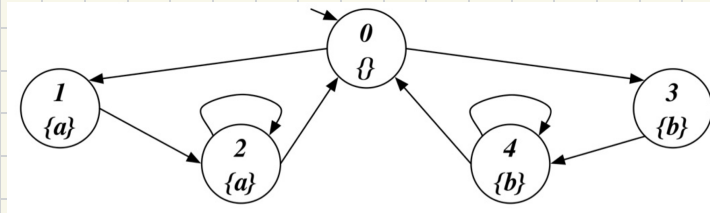
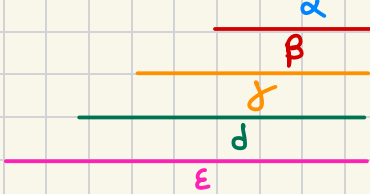
$$[Y_0] = \emptyset$$

$$[Y_1] = ([a] \cap PREA(NEXT, X_3)) \cup ([b] \cap FREE(NEXT, Y_0)) = \\ = (\{1, 2\} \cap \emptyset) \cup (\{3, 4\} \cap \emptyset) = \emptyset$$

$$[Y_0] = [Y_1] = [X_4] = \emptyset$$

$$[X_3] = [X_4] = \emptyset$$

2) $AF(EG(a \supset AXEX \neg a))$



$$[\alpha] = [EX \neg a] = [\langle NEXT \rangle \neg a] = FREE(NEXT, \neg a) = \{0, 2, 3, 4\} = [\alpha]$$

$$[\beta] = [AX \alpha] = [[NEXT] \alpha] = PREA(NEXT, \alpha) = \{1, 2, 3, 4\} = [\beta]$$

$$[\gamma] = [a \supset \beta] = [\neg a] \cup [\beta] = \{0, 3, 4\} \cup \{1, 2, 3, 4\} = \{0, 1, 2, 3, 4\} = [\gamma]$$

$$[\delta] = [EG \gamma] = [\cup Z. \gamma \wedge \langle NEXT \rangle Z]$$

$$[Z_0] = \{0, 1, 2, 3, 4\}$$

$$[Z_1] = [\gamma] \cap FREE(NEXT, Z_0) =$$

$$= \{0, 1, 2, 3, 4\} \cap \{0, 1, 2, 3, 4\} = \{0, 1, 2, 3, 4\} \quad [Z_0] = [Z_1] = [\delta] = \{0, 1, 2, 3, 4\}$$

$$[\epsilon] = [AF \delta] = [\mu Z. \delta \vee [NEXT] Z]$$

$$[Z_0] = \emptyset$$

$$[Z_1] = [\delta] \cup PREA(NEXT, Z_0) =$$

$$= \{0, 1, 2, 3, 4\} \cup \emptyset = \{0, 1, 2, 3, 4\}$$

$$[Z_2] = [\delta] \cup PREA(NEXT, Z_1) =$$

$$= \{0, 1, 2, 3, 4\} \cup \{0, 1, 2, 3, 4\} = \{0, 1, 2, 3, 4\}$$

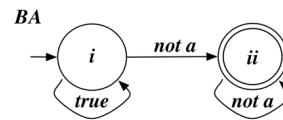
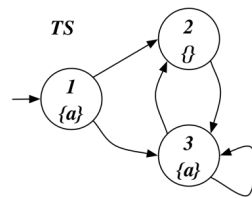
$$[Z_3] = [\delta] \cup PREA(NEXT, Z_2) =$$

$$= \{0, 1, 2, 3, 4\} \cup \{0, 1, 2, 3, 4\} = \{0, 1, 2, 3, 4\}$$

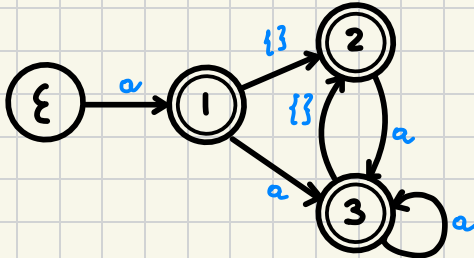
$$[Z_2] = [Z_3] = [\epsilon] = \{0, 1, 2, 3, 4\}$$

$s_0 \in [\epsilon] = ?$ YES!

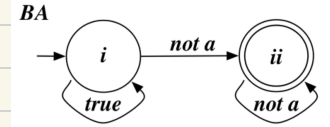
Part 2 Consider the transition system TS below. Model check the LTL formula $\Box \Diamond a$, by considering that the Büchi automaton BA for $\neg \Box \Diamond a$ (i.e., $\Diamond \Box \neg a$) is the one below:



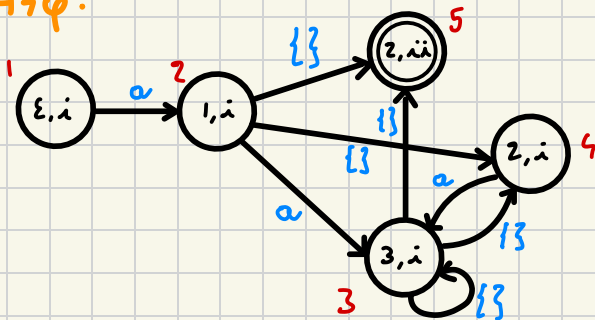
$A \gamma$:



$A \neg \varphi$:



$A \gamma \cap A \neg \varphi$:



$$\varphi = \Box X. \mu Y. (F \wedge \langle \text{NEXT} \rangle X \vee \langle \text{NEXT} \rangle Y)$$

$$[X_0] = \{1, 2, 3, 4, 5\}$$

$$[X_1] = [\mu Y. (F \wedge \langle \text{NEXT} \rangle X_0 \vee \langle \text{NEXT} \rangle Y)]$$

$$[Y_0] = \emptyset$$

$$[Y_1] = [F] \cap \text{PRE}(\text{NEXT}, X_0) \cup \text{PRE}(\text{NEXT}, Y_0) =$$

$$= \{5\} \cap \{1, 2, 3, 4\} \cup \emptyset = \emptyset$$

$$[Y_0] = [Y_1] = [X_1] = \emptyset$$

$$[X_2] = [\mu Y. (F \wedge \langle \text{NEXT} \rangle X_1 \vee \langle \text{NEXT} \rangle Y)]$$

$$[Y_0] = \emptyset$$

$$[Y_1] = [F] \cap \text{PRE}(\text{NEXT}, X_1) \cup \text{PRE}(\text{NEXT}, Y_0) =$$

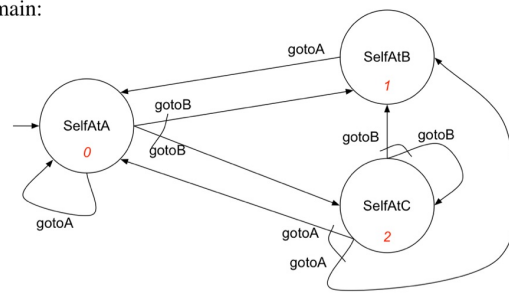
$$= \{5\} \cap \emptyset \cup \emptyset = \emptyset$$

$$[Y_0] = [Y_1] = [X_2] = \emptyset$$

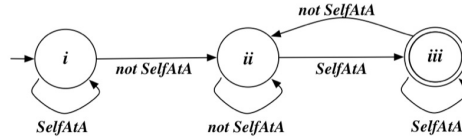
$$[X_1] = [X_2] = \emptyset$$

$S_1 \in [\varphi] = \emptyset$? **no!**

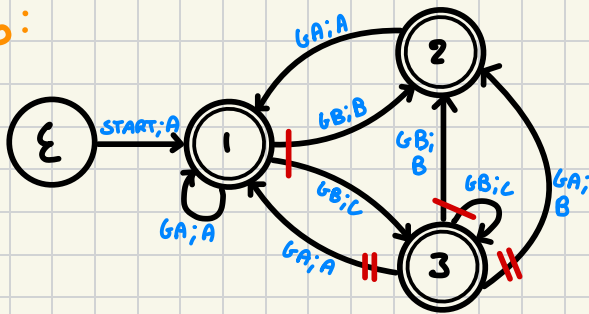
Part 3 Consider the following domain:



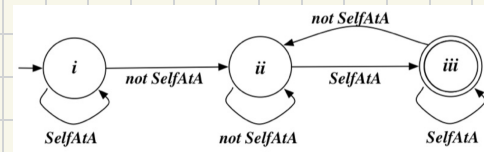
- **Exercise 2.1:** Synthesize a strategy (a plan) for realizing the LTL formula $\Diamond(\neg \text{SelfAtA} \wedge \Diamond(\text{SelfAtA} \wedge \bullet \text{false}))$, by considering that the corresponding DFA is the one below:



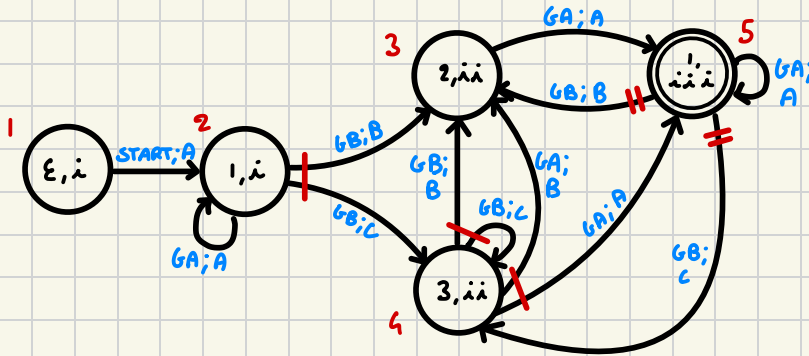
A_D :



A_φ :



$A_D \times A_\varphi$



$$w_0 = \{s\}$$

$$w_1 = w_0 \cup \text{PREADV}(w_0) = \{3, 5\}$$

$$w_2 = w_1 \cup \text{PREADV}(w_1) = \{3, 4, 5\}$$

$$w_3 = w_2 \cup \text{PREADV}(w_2) = \{2, 3, 4, 5\}$$

$$w_4 = w_3 \cup \text{PREADV}(w_3) = \{1, 2, 3, 4, 5\}$$

$$w_5 = w_4 \cup \text{PREADV}(w_4) = \{1, 2, 3, 4, 5\}$$

$$w_4 = w_5$$

$$w(1) = \{\text{START}\}$$

$$w(2) = \{GB\}$$

$$w(3) = \{GA\}$$

$$w(4) = \{GA, GB\}$$

$$w(5) = \text{WIN}$$

$$w_c(1) = \text{START}$$

$$w_c(2) = GB$$

$$w_c(3) = GA$$

$$w_c(4) = GA$$

$$w_c(5) = \text{WIN}$$

$$T = (2^X, S, s_0, p, w_c)$$

$$S = \{1, 2, 3, 4, 5\}$$

$$s_0 = \{1\}$$

$$p(s, x) = \downarrow(s, (w_c(s), x))$$

Part 4 (optional) Consider the following program:

```
while (x < 10) do x := x + 1
```

Compute its *execution* and *final state*, starting from an *initial state* where $x = 9$, by using:

1. *Evaluation Semantics*;
2. *Transition Semantics*.

1) $S_0 = \{x = 9\}$

$$\frac{\text{WHILE } (x < 10) \text{ DO } x = x + 1, S_0 \rightarrow S_F}{\frac{(x = x + 1, S_0) \rightarrow S_1 \quad \wedge \quad (\text{WHILE } (x < 10) \text{ DO } x = x + 1, S_1) \rightarrow S_F}{\text{TRUE}} \quad \text{TRUE}}$$

$$S_1 = S_F = \{x = 10\}$$

2) **FIRST STEP**

$$\frac{(\text{WHILE } (x < 10) \text{ DO } x = x + 1, S_0) \rightarrow (d_1, S_1)}{\frac{x = x + 1, S_0 \rightarrow (\varepsilon, S_1)}{\text{TRUE}} \quad S_1 = \{x = 10\}}$$

SECOND STEP

$$\frac{\varepsilon; \text{WHILE } (x < 10) \text{ DO } x = x + 1, S_1 \rightarrow (d_2, S_2)}{\text{TRUE}}$$

$$S_F = S_2 = S_1 = \{x = 10\}$$