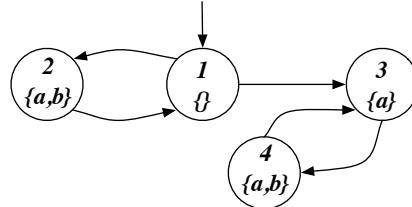
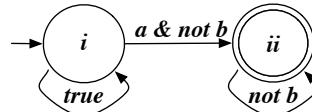


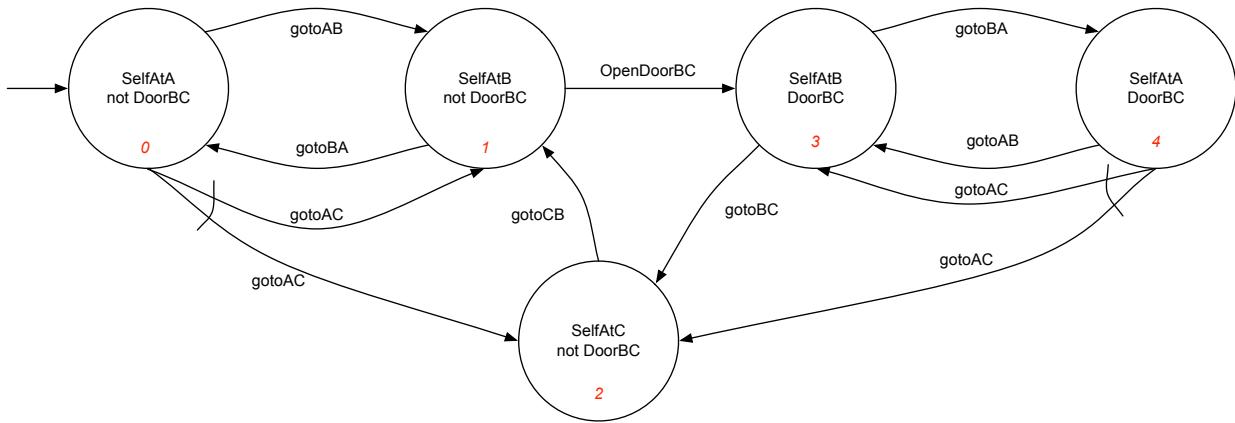
**Part 1.** Consider the following transition system:



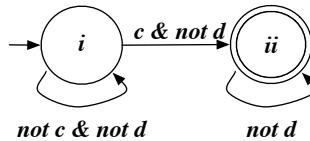
- **Exercise 1.1:** Model check the CTL formula  $EG(b \supset EX AFa)$ , by translating it in Mu-Calculus.
- **Exercise 1.2:** Model check the LTL formula  $\square(a \supset \diamond b)$ , by considering that the Büchi automaton for  $\neg(\square(a \supset \diamond b))$  is the one below:



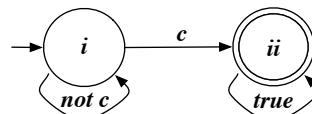
**Part 2** Consider the following domain:



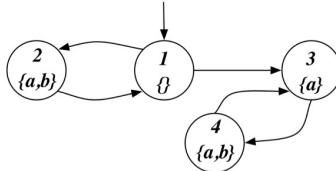
- **Exercise 2.1:** Model check the LTLf formula  $\square(\neg DoorBC) \supset \square(\neg SelfAtC)$ , by considering that the DFA for  $\neg(\square(\neg DoorBC) \supset \square(\neg SelfAtC))$  is the one below:



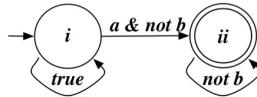
- **Exercise 2.2:** Synthesize a strategy (a plan) for realizing the LTLf formula  $\diamond(SelfAtC)$ , by considering that the DFA for  $\diamond(SelfAtC)$  is the one below:



Part 1. Consider the following transition system:



- Exercise 1.1: Model check the CTL formula  $EG(b \supset EX AF a)$ , by translating it in Mu-Calculus.
- Exercise 1.2: Model check the LTL formula  $\square(a \supset \diamond b)$ , by considering that the Büchi automaton for  $\neg(\square(a \supset \diamond b))$  is the one below:



1)  $EG(b \supset EX AF \alpha)$

$\alpha$   
 $\beta$   
 $\gamma$   
 $\delta$

$$[\alpha] = [\text{AF } \alpha] = [\mu \exists. \alpha \vee \text{NEXT} \exists]$$

$$[\bar{z}_0] = \emptyset$$

$$[\bar{z}_1] = [\alpha] \cup \text{PREA}(\text{NEXT}, \bar{z}_0) =$$

$$= \{2, 3, 4\} \cup \emptyset = \{2, 3, 4\}$$

$$[\bar{z}_2] = [\alpha] \cup \text{PREA}(\text{NEXT}, \bar{z}_1) =$$

$$= \{2, 3, 4\} \cup \{1, 3, 4\} = \{1, 2, 3, 4\}$$

$$[\bar{z}_3] = [\alpha] = \{1, 2, 3, 4\}$$

$$[\beta] = [EX \alpha] = [\text{NEXT } \alpha] = \text{PREE}(\text{NEXT}, \alpha) = \{1, 2, 3, 4\} = [\beta]$$

$$[\gamma] = [b \supset \beta] = [b] \cup [\beta] = \{1, 3\} \cup \{1, 2, 3, 4\} = \{1, 2, 3, 4\} = [\gamma]$$

$$[\delta] = [EG \gamma] = [\mu \exists. \gamma \wedge \text{NEXT} \exists]$$

$$[\bar{z}_0] = \{1, 2, 3, 4\}$$

$$[\bar{z}_1] = [\gamma] \cap \text{PREE}(\text{NEXT}, \bar{z}_0) =$$

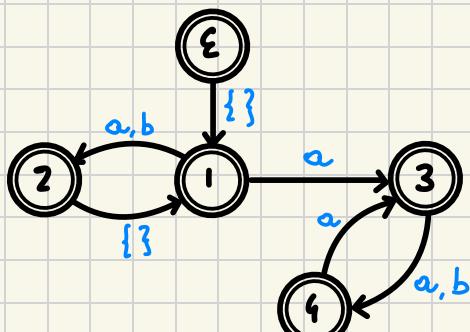
$$= \{1, 2, 3, 4\} \cap \{1, 2, 3, 4\} = \{1, 2, 3, 4\}$$

$$[\bar{z}_0] = [\bar{z}_1] = [\delta] = \{1, 2, 3, 4\}$$

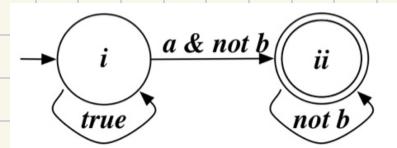
$$\gamma_s, \in \delta ? \rightarrow s, \in [\delta] = \{1, 2, 3, 4\} ? \quad \text{YES!}$$

2)

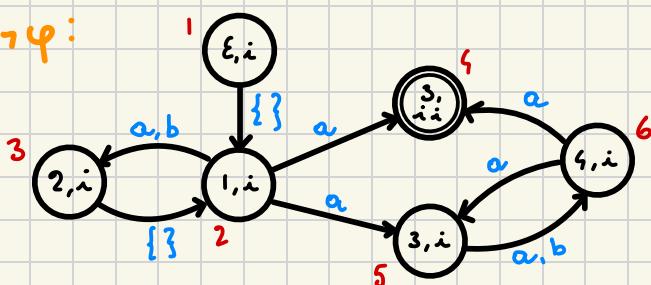
Ar:



$A_{7\varphi} :$



$A_3 \cap A_7 \varphi :$



$$\psi = \exists X. \mu Y (F \wedge \text{NEXT}(X) \vee \text{NEXT}(Y))$$

$$[X_0] = \{1, 2, 3, 4, 5, 6\}$$

$$[x_i] = [\mu y. (F \wedge \text{NEXT} x_i \vee \text{NEXT} y)]$$

$$[Y_0] = \phi$$

$$[Y_1] = [F] \cap \text{PREE}(\text{NEXT}, X_0) \cup \text{PREE}(\text{NEXT}, Y_0) =$$

$$= \{4\} \cap \{1, 2, 3, 5, 6\} \cup \emptyset = \emptyset$$

$$[Y_0] = [Y_+] = [X_+] = \emptyset$$

$$[X_2] = [MY. (F \wedge \text{NEXT} > X, v \wedge \text{NEXT} > Y)]$$

$$[Y_0] = \phi$$

$$[Y_i] = [F] \cap \text{PREE}(\text{NEXT}, X_i) \cup \text{PREE}(\text{NEXT}, Y_0) =$$

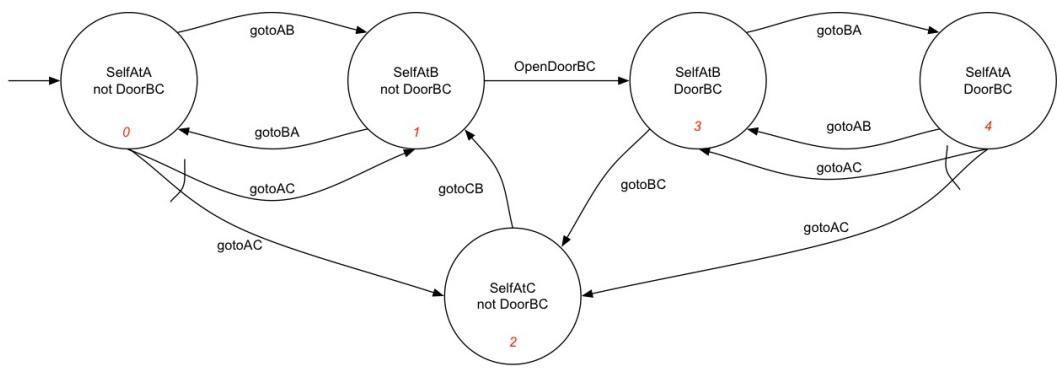
$$= \{ \emptyset \} \cap \emptyset = \emptyset$$

$$[Y_0] = [Y_1] = [x_2] = \emptyset$$

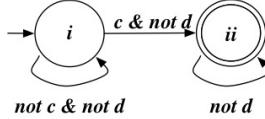
$$[x_1] = [x_2] = \emptyset$$

$\Sigma, \epsilon [\varphi] = \Phi$  ? **no!**

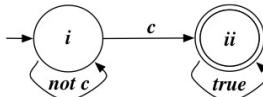
Part 2 Consider the following domain:



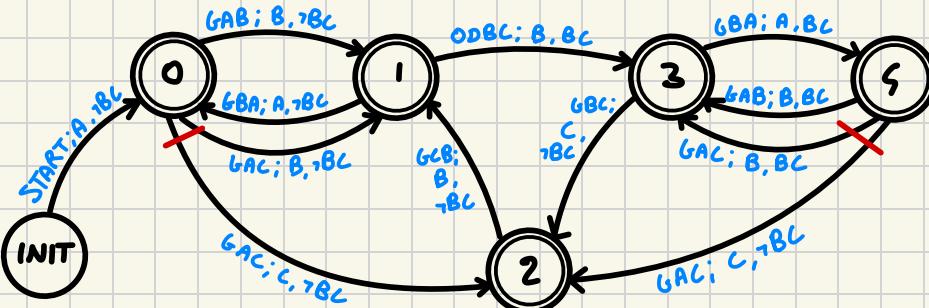
- Exercise 2.1: Model check the LTLf formula  $\square(\neg \text{DoorBC}) \supset \square(\neg \text{SelfAtC})$ , by considering that the DFA for  $\neg(\square(\neg \text{DoorBC}) \supset \square(\neg \text{SelfAtC}))$  is the one below:



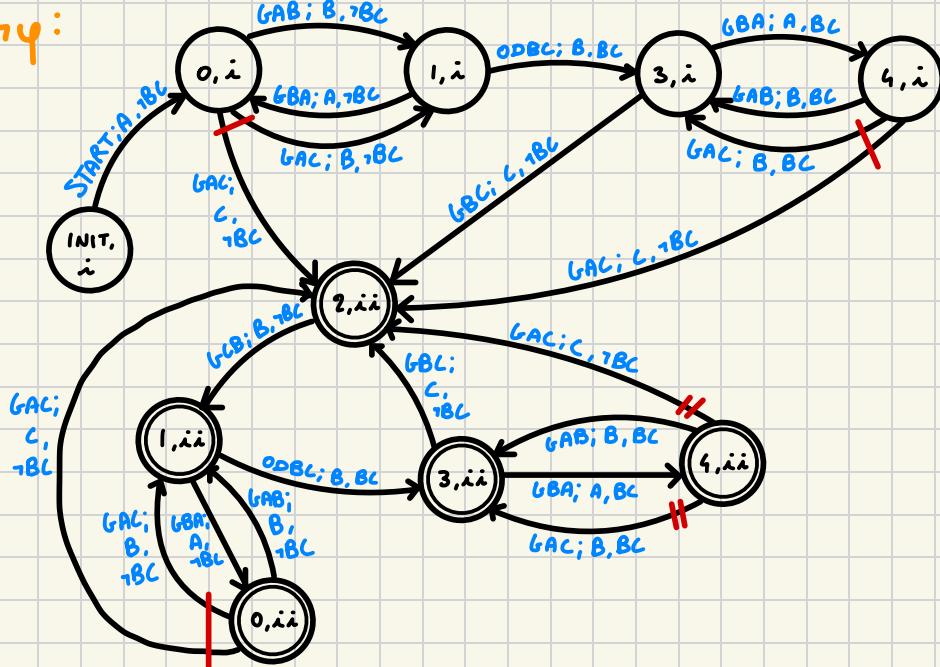
- Exercise 2.2: Synthesize a strategy (a plan) for realizing the LTLf formula  $\diamond(\text{SelfAtC})$ , by considering that the DFA for  $\diamond(\text{SelfAtC})$  is the one below:



i)  $A_D$ :

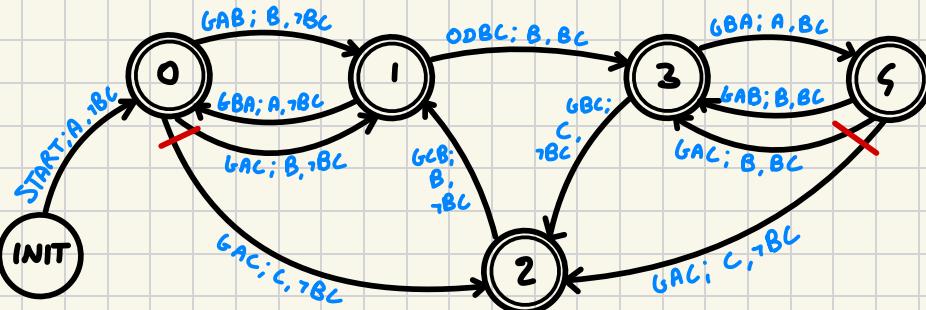


$A_D \times A_{Df}$ :

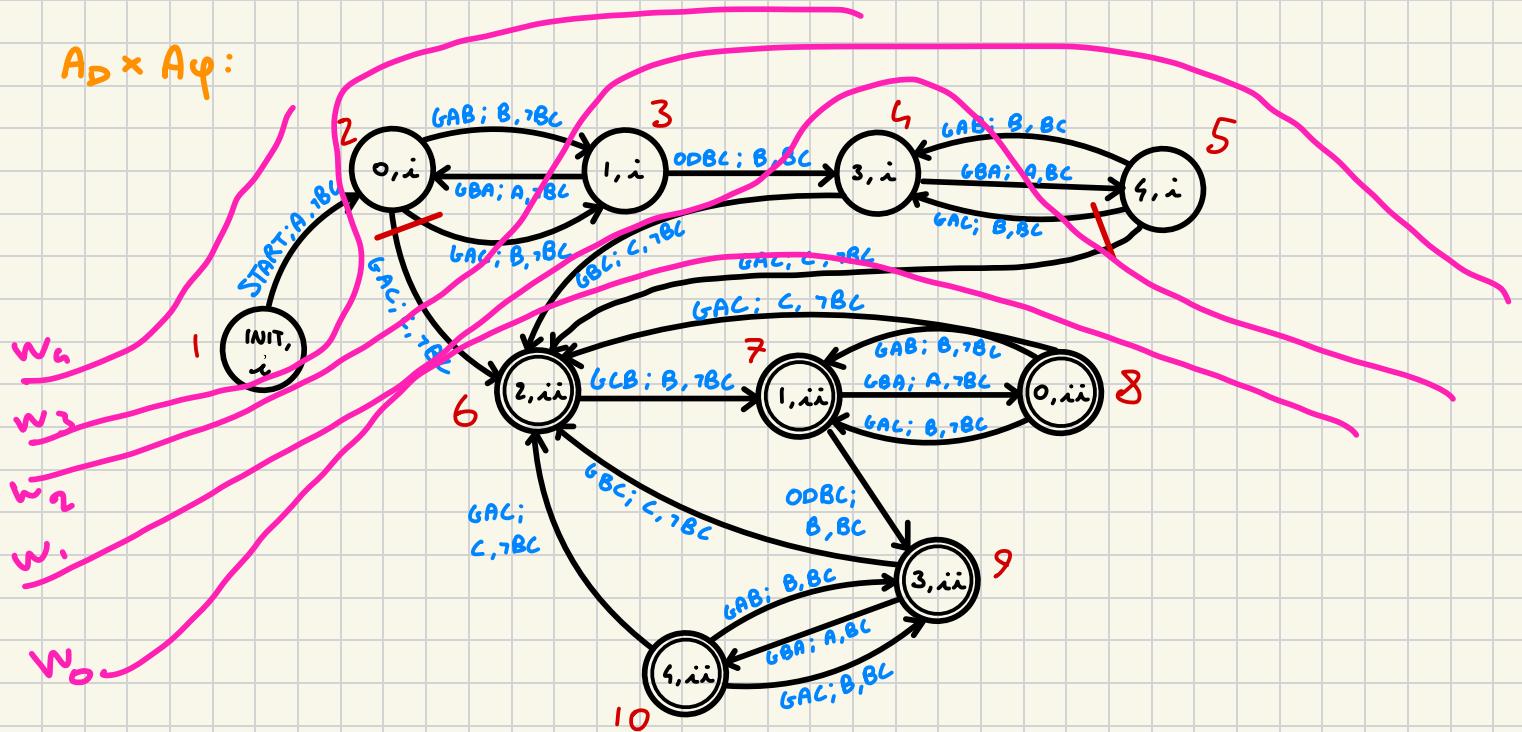


SINCE THERE EXIST A PATH FROM THE INIT TO A FINAL STATE,  $\tau\phi$  IS SATISFIABLE AND SO  $\phi$  IS NOT SAT.

2)  $A_D$ :



$A_D \times A_\varphi$ :



$$w_0 = \{6, 7, 8, 9, 10\}$$

$$w_1 = w_0 \cup \text{PREADV}(w_0) = \{4, 6, 7, 8, 9, 10\}$$

$$w_2 = w_1 \cup \text{PREADV}(w_1) = \{3, 4, 5, 6, 7, 8, 9, 10\}$$

$$w_3 = w_2 \cup \text{PREADV}(w_2) = \{2, 3, 4, 5, 6, 7, 8, 9, 10\}$$

$$w_4 = w_3 \cup \text{PREADV}(w_3) = \{1, 2, 3, 4, 5, 6, 7, 8, 9, 10\}$$

$$w_5 = w_4 \cup \text{PREADV}(w_4) = \{1, 2, 3, 4, 5, 6, 7, 8, 9, 10\}$$

$$w_6 = w_5 = \{1, 2, 3, 4, 5, 6, 7, 8, 9, 10\}$$

$$\begin{aligned} w(1) &= \{\text{START}\} \\ w(2) &= \{GAB, GAC\} \\ w(3) &= \{ODBC\} \\ w(4) &= \{GBC\} \\ w(5) &= \{GAB, GAC\} \\ w(6) &= \text{WIN} \\ w(7) &= \text{WIN} \\ w(8) &= \text{WIN} \\ w(9) &= \text{WIN} \\ w(10) &= \text{WIN} \end{aligned}$$

$$\begin{aligned} w_c(1) &= \text{START} \\ w_c(2) &= GAC \\ w_c(3) &= ODBC \\ w_c(4) &= GBC \\ w_c(5) &= GAC \\ w_c(6) &= \text{WIN} \\ w_c(7) &= \text{WIN} \\ w_c(8) &= \text{WIN} \\ w_c(9) &= \text{WIN} \\ w_c(10) &= \text{WIN} \end{aligned}$$

$$T = (2^x, S, s_0, p, w_c)$$

$$S = \{1, 2, 3, 4, 5, 6, 7, 8, 9, 10\}$$

$$s_0 = \{1\}$$

$$p(s, x) = \delta(s, (w_c(s), x))$$

$$w_c =$$