






**Part 1 - Artificial Intelligence**  
(Time to complete the test: 2:00 hours)

Consider a simplified Pac-Man world consisting in a 3x3 grid, as shown below:

	A	B	C
1			
2			
3			

Besides Pac-Man, the grid contains: power pellets (●); obstacles (■); ghosts.

Consider the following formalization of the scenario:

*Non-Fluents:*

- $Connected(c_1, c_2)$ , denoting that cells  $c_1$  and  $c_2$  are connected: two cells are connected if they share an *entire* side (not just a corner); for instance, B1 is connected to C1 but not to C2; notice also that  $Connected$  is symmetrical, i.e.,  $c_1$  is connected to  $c_2$  if and only if  $c_2$  is connected to  $c_1$ ; cells are not connected to themselves.
- $ObstacleIn(c)$ , denoting that cell  $c$  contains an obstacle.

*Fluents:*

- $In(c)$ , denoting that Pac-Man is in cell  $c$ .
- $Pellet(c)$ , denoting that cell  $c$  contains a power pellet.
- $Super()$ , denoting that Pac-Man has superpowers.
- $Alive()$ , denoting that Pac Man is alive.
- $GhostIn(c)$ , denoting that cell  $c$  contains a ghost (notice that ghosts are static).

*Actions:*

- $move(c)$ , which allows Pac-Man to move to cell  $c$ . The action can be done only if Pac-Man is alive, in a cell  $c'$  connected to  $c$ , and  $c$  contains no obstacles (it may contain a ghost, though). The effect is that: 1. if there are no ghosts in  $c$  or Pac-Man has superpowers then Pac-Man is in cell  $c$ ; 2. any power pellet in  $c$  disappears; 3. if there is a power pellet in  $c$  then Pac-Man has superpowers (once obtained, superpowers last forever); 4. if there is a ghost in  $c$  and Pac-Man does not have superpowers, then Pac-Man is no longer alive and disappears from the grid.
- $eatGhost()$ , which allows Pac-Man to eat a ghost. The action can be done only if Pac-Man has superpowers and is in a cell containing a ghost. The effect is that the ghost disappears from the cell.

*Initial situation:*

Pac-Man is alive. Cells, connections, Pac-Man's initial position, power pellets, and ghosts are as in the figure above.

**Exercise 1.**

1. Formalize the above scenario as a Basic Action Theory.
2. Using regression, check whether the action sequence  $\varrho = move(A2); move(A3)$ ; leads to a situation where Pac-Man has superpowers.

**Exercise 2.**

Considering goal  $\gamma = In(C3)$ ,

1. Formalize the above scenario as a PDDL domain file and a PDDL problem file.
2. Draw the corresponding entire transition system;
3. Solve planning for achieving  $\gamma$ , using forward depth-first search (uninformed), reporting the steps of the forward search computation, and returning the resulting plan. **The search must be exhaustive**, i.e., all states must be visited.

### Exercise 3.

1. Formalize the following sentences in First-Order Logic:

- $\phi_1$ : Every beverage is liquid
- $\phi_2$ : Not all liquids are beverages
- $\phi_3$ : Some liquids are dangerous
- $\phi_5$ : There exists exactly one dangerous beverage

2. Given the following formulas:

- $\phi_1 = (\forall x.P(x) \supset Z(x)) \wedge (\forall x.Q(x) \supset Z(x))$
- $\phi_2 = \forall x.(Z(x) \wedge P(x)) \supset \neg Q(x)$
- $\phi_3 = \neg \exists x.P(x) \wedge Q(x)$

use the Tableaux method to check whether  $\{\phi_1, \phi_2\} \models \phi_3$

## Exercise 1.

1. Formalize the above scenario as a Basic Action Theory.
2. Using regression, check whether the action sequence  $q = \text{move}(A2); \text{move}(A3);$  leads to a situation where Pac-Man has super-powers.

### 1) PRECONDITIONS AXIOMS

$$\text{POSS}(\text{MOVE}(c), s) \equiv \exists c'. \text{IN}(c', s) \wedge \text{CONNECTED}(c', c) \wedge \text{ALIVE}(s) \\ \wedge \neg \text{OBSTACLEIN}(c, s)$$

$$\text{POSS}(\text{EATGHOST}(c), s) \equiv \exists c. \text{IN}(c, s) \wedge \text{GHOSTIN}(c, s) \wedge \text{SUPER}(s)$$

### SUCCESSOR STATE AXIOMS

#### EFFECT AXIOMS

$$\omega = \text{MOVE}(c) \supset (\neg \text{GHOSTIN}(c, s) \supset \text{IN}(c, \text{DO}(a, s))) \wedge (\text{SUPER}(s) \supset \text{IN}(c, \text{DO}(a, s))) \wedge \\ \neg \text{PELLET}(c, \text{DO}(a, s)) \wedge (\text{PELLET}(c, s) \supset \text{SUPER}(\text{DO}(a, s))) \wedge \\ (\text{GHOSTIN}(c, s) \wedge \neg \text{SUPER}(s) \supset \neg \text{ALIVE}(\text{DO}(a, s))) \wedge \\ (\forall x. \text{IN}(x, s) \supset \neg \text{IN}(x, \text{DO}(a, s))))$$

$$\omega = \text{EATGHOST}(c) \supset E_c.(\text{GHOSTIN}(c, s) \supset \neg \text{GHOSTIN}(c, \text{DO}(a, s)))$$

#### NORMALIZE

$$(\omega = \text{MOVE}(c) \wedge \neg \text{GHOSTIN}(c, s)) \supset \text{IN}(c, \text{DO}(a, s))$$

$$(\omega = \text{MOVE}(c) \wedge \text{SUPER}(s)) \supset \text{IN}(c, \text{DO}(a, s))$$

$$\omega = \text{MOVE}(c) \supset \neg \text{PELLET}(c, \text{DO}(a, s))$$

$$\exists c. (\omega = \text{MOVE}(c) \wedge \text{PELLET}(c, s)) \supset \text{SUPER}(\text{DO}(a, s))$$

$$\exists c. (\omega = \text{MOVE}(c) \wedge \text{GHOSTIN}(c, s) \wedge \neg \text{SUPER}(s)) \supset \neg \text{ALIVE}(\text{DO}(a, s))$$

$$\exists c. (\omega = \text{MOVE}(c) \wedge \text{GHOSTIN}(c, s) \wedge \neg \text{SUPER}(s) \wedge \text{IN}(x, s)) \supset \neg \text{IN}(x, \text{DO}(a, s))$$

$$(\omega = \text{EATGHOST}(c) \wedge \text{GHOSTIN}(c, s)) \supset \neg \text{GHOSTIN}(c, \text{DO}(a, s))$$

#### EXPLANATION CLOSURE

$$\text{IN}(c, \text{DO}(a, s)) \equiv (\omega = \text{MOVE}(c) \wedge \neg \text{GHOSTIN}(c, s)) \vee (\omega = \text{MOVE}(c) \wedge \text{SUPER}(s)) \\ \vee (\text{IN}(c, s) \wedge \neg (\exists c'. (\omega = \text{MOVE}(c') \wedge \text{GHOSTIN}(c', s) \\ \wedge \neg \text{SUPER}(s))))$$

$$\text{PELLET}(c, \text{DO}(a, s)) \equiv (\text{PELLET}(c, s) \wedge \neg \omega = \text{MOVE}(c))$$

$$\text{SUPER}(\text{DO}(a, s)) \equiv \exists c. (\omega = \text{MOVE}(c) \wedge \text{PELLET}(c, s)) \vee \text{SUPER}(s)$$

$$\text{ALIVE}(\text{DO}(a, s)) \equiv (\text{ALIVE}(s) \wedge \neg (\exists c. a = \text{MOVE}(c) \wedge \text{GHOSTIN}(c, s) \wedge \neg \text{SUPER}(s)))$$

$$\text{GHOSTIN}(c, \text{DO}(a, s)) \equiv (\text{GHOSTIN}(c, s) \wedge \neg a = \text{EATGHOST}(c))$$

## INITIAL SITUATION

ALIVE( $s_0$ )

CONNECTED( $c_1, c_2$ )  $\equiv (c_1 = A_1, c_2 = B_1) \dots$

IN( $c, s_0$ )  $\equiv (c = A_1)$

PELLET( $c, s_0$ )  $\equiv (c = A_3)$

GHOSTIN( $c, s_0$ )  $\equiv (c = c_2)$

OBSTACLES( $c$ )  $\equiv (c = B_2) (c = B_3)$

2)  $q = \text{move}(A_2); \text{move}(A_3);$

$s_1 = \text{DO}(\text{MOVE}(A_2), s_0)$

$s_2 = \text{DO}(\text{MOVE}(A_3), s_1)$

$D_0 \cup D_{UNA} \models R[\text{SUPER}(s_2)]$

$$\begin{aligned} R[\text{SUPER}(s_2)] &= R[\text{SUPER}(\text{DO}(\text{MOVE}(A_3), s_1))] = \\ &= R[\exists c. (\text{MOVE}(A_3) = \text{MOVE}(c) \wedge \text{PELLET}(c, s_1)) \vee \text{SUPER}(s_1)] = \\ &= R[\text{PELLET}(s_1)] \vee R[\text{SUPER}(s_1)] = \end{aligned}$$

$$\begin{aligned} R[\text{PELLET}(s_1)] &= R[\text{PELLET}(\text{DO}(\text{MOVE}(A_2), s_0))] = \\ &= R[\text{PELLET}(c, s_0) \wedge \text{MOVE}(A_2) \neq \text{MOVE}(c)] = \\ &= \text{PELLET}(c, s_0) \wedge \text{MOVE}(A_2) \neq \text{MOVE}(c) \end{aligned}$$

$$\begin{aligned} R[\text{SUPER}(s_1)] &= R[\text{SUPER}(\text{DO}(\text{MOVE}(A_2), s_0))] = \\ &= R[\exists c. (\text{MOVE}(A_2) = \text{MOVE}(c) \wedge \text{PELLET}(c, s_0)) \vee \text{SUPER}(s_0)] = \\ &= R[\text{FALSE}] = \text{FALSE} \end{aligned}$$

$$R[\text{SUPER}(s_2)] = R[\text{PELLET}(s_1)] \vee R[\text{SUPER}(s_1)]$$

TRUE FOR  $c = A_3 \rightarrow p$  LEADS TO A SITUATION WHERE  $\mathcal{B}$  IS SUPER ✓

## Exercise 2.

Considering goal  $\gamma = In(C3)$ ,

1. Formalize the above scenario as a PDDL domain file and a PDDL problem file.
2. Draw the corresponding entire transition system;
3. Solve planning for achieving  $\gamma$ , using forward depth-first search (uninformed), reporting the steps of the forward search computation, and returning the resulting plan. **The search must be exhaustive**, i.e., all states must be visited.

```

1) (DEFINE (DOMAIN PAC-DOM)
  (:REQUIREMENTS :ADL)
  (:TYPES CELL)
  (:PREDICATES
    (CONNECTED ?c ?c2 - CELL)
    (OBSTACLEIN ?c - CELL)
    (IN ?c - CELL)
    (PELLET ?c - CELL)
    (SUPER)
    (ALIVE)
    (GHOSTIN ?c - CELL)
  )
  (:ACTION MOVE
    :PARAMETERS (?c - CELL)
    :PRECONDITIONS (AND
      (EXISTS (?c2 - CELL) (IN ?c2) (CONNECTED ?c ?c2)
        (ALIVE) (NOT (OBSTACLEIN ?c2)))
    )
    :EFFECT (AND
      (WHEN (NOT (GHOSTIN ?c)) (IN ?c))
      (WHEN (SUPER) (IN ?c))
      (NOT (PELLET ?c))
      (WHEN (PELLET ?c) (SUPER))
      (WHEN (AND (GHOSTIN ?c) (NOT (SUPER))))
      (AND (NOT (ALIVE))
        (FORALL (?x - CELL)
          (WHEN (IN ?x) (NOT (IN ?x)))
        )
      )
    )
  )
); END OF MOVE
(:ACTION EAT GHOST
  :PARAMETERS ()
  :PRECONDITION (EXISTS (?c - CELL) (AND (IN ?c) (GHOSTIN ?c) (SUPER)))
  :EFFECT (EXISTS (?c - CELL) (WHEN (GHOSTIN ?c) (NOT (GHOSTIN ?c))))
); END OF EAT GHOST
); END OF DEFINE DOM
  
```

```

(DEFINE (PROBLEM PAC-PROB) (: DOMAIN PAC-DOM)
  (: OBJECTS  A, A2 A3 B, B2 B3 C, C2 C3 - CELL)
  (: INIT  (ALIVE) (CONNECTED A, B,) ..... (IN A1)
            (OBSTACLEIN B2) (OBSTACLEIN B3) (PELLET A3) (GHOSTIN A1)
  )
  (: GOAL  (IN C3)
  )
); END OF DEF PROB

```

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2. Given the following formulas:

- $\phi_1 = (\forall x. P(x) \supset Z(x)) \wedge (\forall x. Q(x) \supset Z(x))$
- $\phi_2 = \forall x. (Z(x) \wedge P(x)) \supset \neg Q(x)$
- $\phi_3 = \neg \exists x. P(x) \wedge Q(x)$

use the Tableaux method to check whether  $\{\phi_1, \phi_2\} \models \phi_3$

1)  $\phi_1: \forall x. \text{BEVERAGE}(x) \supset \text{LIQUID}(x)$

$\phi_2: \exists x. \text{LIQUID}(x) \wedge \neg \text{BEVERAGE}(x)$

$\phi_3: \exists x. \text{LIQUID}(x) \wedge \text{DANGEROUS}(x)$

$\phi_4: \exists x. (\text{BEVERAGE}(x) \wedge \text{DANGEROUS}(x)) \wedge \forall y. (\text{BEVERAGE}(y) \wedge \text{DANGEROUS}(y) \supset y=x)$

2)

$$(\forall x. P(x) \supset Z(x)) \wedge (\forall x. Q(x) \supset Z(x)) \quad 1$$

$$\forall x. (Z(x) \wedge P(x)) \supset \neg Q(x) \quad 2$$

$$\exists x. P(x) \wedge Q(x) \quad 3$$

$$\mid \delta \text{ on } 3$$

$$P(a) \wedge Q(a) \quad 4$$

$$\mid \alpha \text{ on } 4$$

$$P(a) \quad 5$$

$$Q(a) \quad 6$$

$$\mid \gamma \text{ on } 2$$

$$\neg (Z(a) \wedge P(a)) \vee \neg Q(a) \quad 7$$

$$\beta \text{ on } 7$$

$$\neg Z(a) \vee \neg P(a) \quad 8$$

$$\neg Q(a) \quad 9$$

$$\beta \text{ on } 8$$

$$\neg Z(a) \quad 9$$

$$\neg P(a) \quad 10$$

$$\mid \gamma \text{ on } 1$$

$$\mid$$

$$(P(a) \supset Z(a)) \wedge (Q(a) \supset Z(a)) \quad 11 \quad \times$$

$$\mid \alpha \text{ on } 11$$

$$\neg P(a) \vee Z(a) \quad 12$$

$$\neg Q(a) \vee Z(a) \quad 13$$

$$\beta \text{ on } 12$$

$$\neg P(a)$$

$$Z(a)$$

$\times$

$$(\phi_1, \phi_2) \not\models \neg \phi_3 \longrightarrow \{\phi_1, \phi_2\} \models \phi_3$$