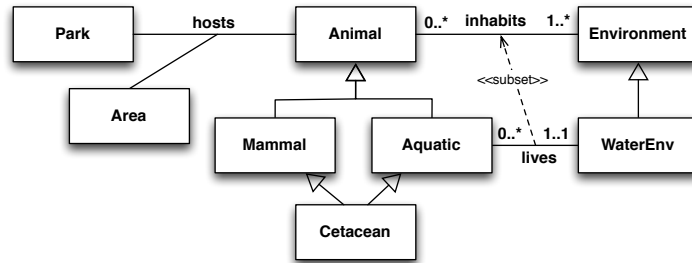
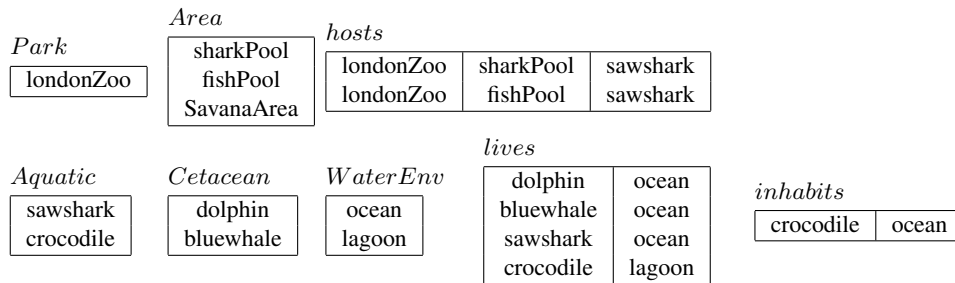


Exercise 1. Express the following UML class diagram in *FOL*.

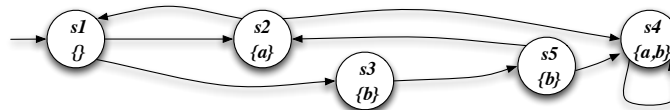


Exercise 2. Consider the above UML class diagram and the following (partial) instantiation.

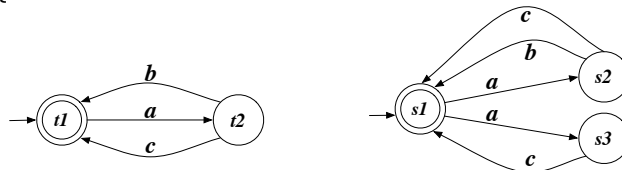


1. Check whether the above instantiation, once completed, is correct, and explain why it is or it is not.
2. Express in FOL and evaluate the following queries:
 - (a) Return animals that inhabit at least two environments.
 - (b) Return parks that they host only aquatic animals.
 - (c) Check if there are parks that host all Cetacean.

Exercise 3. Model check the Mu-Calculus formula $\nu X. \mu Y. ((a \wedge [next]X) \vee (b \wedge [next]Y))$ and the CTL formula $AF(a \supset EXEGb)$ (showing its translation in Mu-Calculus) against the following transition system:



Exercise 4. Consider the following two transition systems:

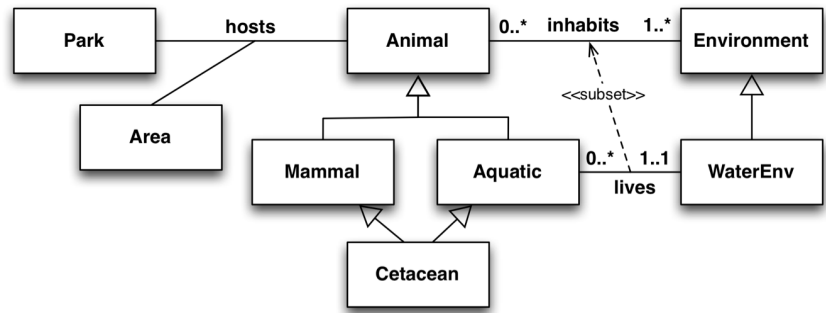


Write the definition of bisimilarity and compute the bisimilarity relation for the two transition systems.

Exercise 5. Compute the certain answers to the CQ $q(x) \leftarrow Employee(x), Manages(x, y)$ over the incomplete database (naive tables), by explaining and exploiting the connection with conjunctive query containment:

<i>Employee</i>	<i>Manages</i>														
<table><tr><th><i>name</i></th></tr><tr><td>Smith</td></tr><tr><td><i>null</i>₁</td></tr><tr><td>Brown</td></tr></table>	<i>name</i>	Smith	<i>null</i> ₁	Brown	<table><tr><th><i>mgr</i></th><th><i>mgd</i></th></tr><tr><td>Green</td><td>Smith</td></tr><tr><td>Smith</td><td><i>null</i>₁</td></tr><tr><td><i>null</i>₁</td><td>Brown</td></tr><tr><td>Brown</td><td><i>null</i>₂</td></tr></table>	<i>mgr</i>	<i>mgd</i>	Green	Smith	Smith	<i>null</i> ₁	<i>null</i> ₁	Brown	Brown	<i>null</i> ₂
<i>name</i>															
Smith															
<i>null</i> ₁															
Brown															
<i>mgr</i>	<i>mgd</i>														
Green	Smith														
Smith	<i>null</i> ₁														
<i>null</i> ₁	Brown														
Brown	<i>null</i> ₂														

Exercise 1. Express the following UML class diagram in *FOL*.



$P(x)$, $AREA(x)$, $A(x)$, $M(x)$, $AQ(x)$, $LET(x)$, $ENV(x)$, $WENV(x)$
 $HOSTS(x, y, z)$
 $INHAB(x, y)$
 $LIVES(x, y)$

$\forall x, y, z. HOSTS(x, y, z) \supset P(x) \wedge AREA(y) \wedge A(z)$

$\forall x, y. INHAB(x, y) \supset A(x) \wedge ENV(y)$

$\forall x. A(x) \supset \exists y. INHAB(x, y)$

$\forall y. ENV(y) \supset 0 \leq \# \{x \mid INHAB(x, y)\}$

$\forall x, y. LIVES(x, y) \supset AQ(x) \wedge WENV(y)$

$\forall x. AQ(x) \supset 1 \leq \# \{y \mid LIVES(x, y)\} \leq 1$

$\forall y. WENV(y) \supset 0 \leq \# \{x \mid LIVES(x, y)\}$

$\forall x, y. LIVES(x, y) \supset INHAB(x, y)$

$\forall x. M(x) \supset A(x)$

$\forall x. AQ(x) \supset A(x)$

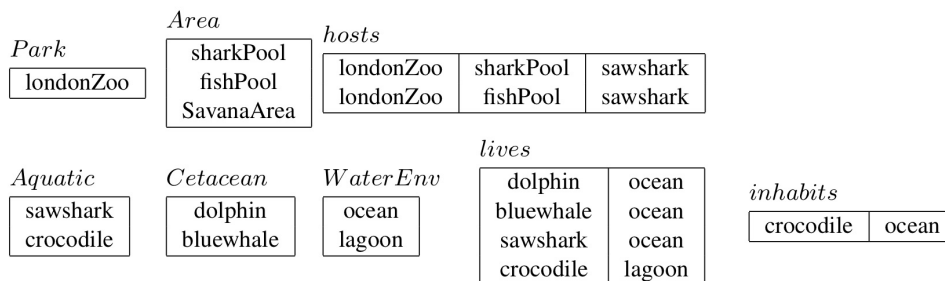
$\forall x. M(x) \supset \neg AQ(x)$

$\forall x. A(x) \supset M(x) \vee AQ(x)$

$\forall x. CET(x) \supset M(x) \wedge AQ(x)$

$\forall x. WENV(x) \supset ENV(x)$

Exercise 2. Consider the above UML class diagram and the following (partial) instantiation.



1. Check whether the above instantiation, once completed, is correct, and explain why it is or it is not.
2. Express in FOL and evaluate the following queries:
 - (a) Return animals that inhabit at least two environments.
 - (b) Return parks that they host only aquatic animals.
 - (c) Check if there are parks that host all Cetacean.

1) $C = \{ \text{DOLPHIN, BLUEW, SAWSHARK, CROCO} \}$ $LIVES = \{ \dots (CROCO, OCEAN) \}$

$$\forall x, y, z. \text{HOSTS}(x, y, z) \supset P(x) \wedge \text{AREA}(y) \wedge A(z)$$

LONDON ZOO IS A PARK
SHARK, FISH ARE AREAS
SAWSHARK IS AN ANIMAL → CARDINALS OK!

$$\forall x, y. \text{INHAB}(x, y) \supset A(x) \wedge \text{ENV}(y)$$

CROCO IS AN ANIMAL → CARDINALS
OCEAN IS AN WENN OK!

$$\forall x, y. \text{LIVES}(x, y) \supset \text{AQ}(x) \wedge \text{WENV}(y)$$

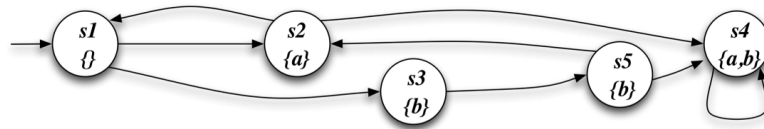
D, B, S, C ARE CETACEANS → CARDINALS
OLEAN, LAGOON ARE WENV OK!

2) a. $\exists e, e'. \text{C}(x) \wedge \text{INHAB}(x, e) \wedge \text{INHAB}(x, e') \wedge e \neq e'$
 $\{ \}$

b. $P(x) \wedge \forall a, c (MOSTS(x, a, c) \supset AQ(c))$
 $\{LONDON \geq 100\}$

c. $\exists p. P(p) \wedge \forall c. (C(c) \supset \exists a. \text{HOSTS}(p, a, c))$
 {FALSE}

Exercise 3. Model check the Mu-Calculus formula $\nu X. \mu Y. ((a \wedge [next]X) \vee (b \wedge [next]Y))$ and the CTL formula $AF(a \supset EX EGb)$ (showing its translation in Mu-Calculus) against the following transition system:



1) $\nu X. \mu Y. ((a \wedge [NEXT]X) \vee (b \wedge [NEXT]Y))$

$$[X_0] = \{1, 2, 3, 4, 5\}$$

$$[X_1] = [\mu Y. ((a \wedge [NEXT]X_0) \vee (b \wedge [NEXT]Y))]$$

$$[Y_0] = \emptyset$$

$$\begin{aligned} [Y_1] &= ([a] \wedge PREA(NEXT, X_0)) \cup ([b] \wedge PREA(NEXT, Y_0)) = \\ &= (\{2, 4\} \cap \{1, 2, 3, 4, 5\}) \cup (\{3, 4, 5\} \cap \emptyset) = \{2, 4\} \end{aligned}$$

$$\begin{aligned} [Y_2] &= ([a] \wedge PREA(NEXT, X_0)) \cup ([b] \wedge PREA(NEXT, Y_1)) = \\ &= (\{2, 4\} \cap \{1, 2, 3, 4, 5\}) \cup (\{3, 4, 5\} \cap \{4, 5\}) = \{2, 4, 5\} \end{aligned}$$

$$\begin{aligned} [Y_3] &= ([a] \wedge PREA(NEXT, X_0)) \cup ([b] \wedge PREA(NEXT, Y_2)) = \\ &= (\{2, 4\} \cap \{1, 2, 3, 4, 5\}) \cup (\{3, 4, 5\} \cap \{2, 3, 4, 5\}) = \{2, 3, 4, 5\} \end{aligned}$$

$$\begin{aligned} [Y_4] &= ([a] \wedge PREA(NEXT, X_0)) \cup ([b] \wedge PREA(NEXT, Y_3)) = \\ &= (\{2, 4\} \cap \{1, 2, 3, 4, 5\}) \cup (\{3, 4, 5\} \cap \{1, 3, 4, 5\}) = \{2, 3, 4, 5\} \end{aligned}$$

$$[Y_3] = [Y_4] = [X_1] = \{2, 3, 4, 5\}$$

$$[X_2] = [\mu Y. ((a \wedge [NEXT]X_1) \vee (b \wedge [NEXT]Y))]$$

$$[Y_0] = \emptyset$$

$$\begin{aligned} [Y_1] &= ([a] \wedge PREA(NEXT, X_1)) \cup ([b] \wedge PREA(NEXT, Y_0)) = \\ &= (\{2, 4\} \cap \{1, 3, 4, 5\}) \cup (\{3, 4, 5\} \cap \emptyset) = \{4\} \end{aligned}$$

$$\begin{aligned} [Y_2] &= ([a] \wedge PREA(NEXT, X_1)) \cup ([b] \wedge PREA(NEXT, Y_1)) = \\ &= (\{2, 4\} \cap \{1, 3, 4, 5\}) \cup (\{3, 4, 5\} \cap \{4\}) = \{4\} \end{aligned}$$

$$[Y_1] = [Y_2] = [X_2] = \{4\}$$

$$[X_3] = [\mu Y. ((a \wedge [NEXT]X_2) \vee (b \wedge [NEXT]Y))]$$

$$[Y_0] = \emptyset$$

$$[Y_1] = ([a] \cap \text{PREA}(\text{NEXT}, X_2)) \cup ([b] \cap \text{PREA}(\text{NEXT}, Y_0)) =$$

$$= (\{2, 4\} \cap \{4\}) \cup (\{3, 4, 5\} \cap \emptyset) = \{4\}$$

$$[Y_2] = ([a] \cap \text{PREA}(\text{NEXT}, X_2)) \cup ([b] \cap \text{PREA}(\text{NEXT}, Y_1)) =$$

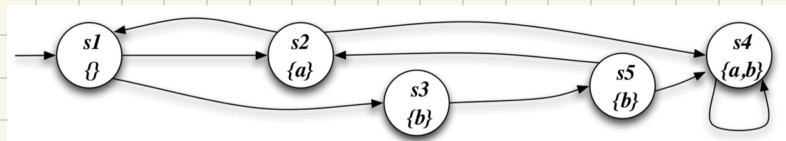
$$= (\{2, 4\} \cap \{4\}) \cup (\{3, 4, 5\} \cap \{4\}) = \{4\}$$

$$[Y_1] = [Y_2] = [X_3] = \{4\}$$

$$[X_2] = [X_3] = \{4\}$$

$$S, \in [\bigvee X. \mu Y. ((a \wedge [\text{NEXT}]X) \vee (b \wedge [\text{NEXT}]Y))] = \{4\} ? \text{ NO!}$$

$$2) \text{ AF } (a \supset \text{EX EG } b)$$



$$[\alpha] = [EG \ b] = [\bigvee Z. b \wedge \langle \text{NEXT} \rangle Z]$$

$$[Z_0] = \{1, 2, 3, 4, 5\}$$

$$[Z_1] = [b] \cap \text{PREE}(\text{NEXT}, Z_0) =$$

$$= \{3, 4, 5\} \cap \{1, 2, 3, 4, 5\} = \{3, 4, 5\}$$

$$[Z_2] = [b] \cap \text{PREE}(\text{NEXT}, Z_0) =$$

$$= \{3, 4, 5\} \cap \{1, 2, 3, 4, 5\} = \{3, 4, 5\} \quad [Z_1] = [Z_2] = [\alpha] = \{3, 4, 5\}$$

$$[\beta] = [EX \ \alpha] = [\langle \text{NEXT} \rangle \alpha] = \text{PREE}(\text{NEXT}, \alpha) = \{1, 2, 3, 4, 5\} = [\beta]$$

$$[\gamma] = [\alpha \supset \beta] = [\neg \alpha] \cup [\beta] = \{1, 3, 5\} \cup \{1, 2, 3, 4, 5\} = \{1, 2, 3, 4, 5\} = [\gamma]$$

$$[\delta] = [\text{AF } \gamma] = [\mu Z. \gamma \vee [\text{NEXT}] Z]$$

$$[Z_0] = \emptyset$$

$$[Z_1] = [\gamma] \cup \text{PREA}(\text{NEXT}, Z_0) =$$

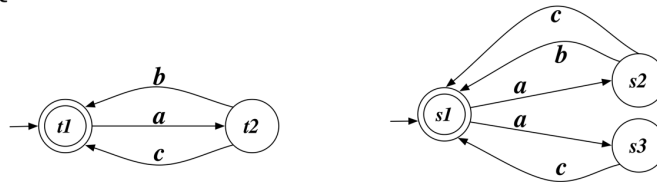
$$= \{1, 2, 3, 4, 5\} \cup \emptyset = \{1, 2, 3, 4, 5\}$$

$$[Z_2] = [\gamma] \cup \text{PREA}(\text{NEXT}, Z_1) =$$

$$= \{1, 2, 3, 4, 5\} \cup \{1, 2, 3, 4, 5\} = \{1, 2, 3, 4, 5\} \quad [Z_1] = [Z_2] = [\delta] = \{1, 2, 3, 4, 5\}$$

$$\top S, \in \delta ? \rightarrow S, \in [\delta] = \{1, 2, 3, 4, 5\} ? \text{ YES!}$$

Exercise 4. Consider the following two transition systems:



Write the definition of bisimilarity and compute the bisimilarity relation for the two transition systems.

TWO TS ARE BISIMILIAR IF THEY HAVE THE SAME BEHAVIOR:

- LOCALLY THEY LOOK INDISTINGUISHABLE
- EVERY ACTION THAT CAN BE DONE ON ONE, CAN ALSO BE DONE ON THE OTHER, AND THEY REACH THE SAME RESULT

$$R_0 = T \times S = \{(t_1, s_1), (t_1, s_2), (t_1, s_3), (t_2, s_1), (t_2, s_2), (t_2, s_3)\}$$

$$R_1 = \{(t_1, s_1), (t_2, s_2), (\cancel{t_1}, s_3)\}$$

$$R_2 = \{(\cancel{t_1}, s_1), (t_2, s_2)\}$$

$$R_3 = \{(\cancel{t_1}, s_2)\}$$

$$R_4 = \{\}$$

$$R_5 = \{\}$$

$$R_4 = R_5 \quad \text{GFP FOUND}$$

$(t_1, s_1) \notin \text{GFP}$ SO T AND S ARE NOT BISIMILIAR

Exercise 5. Compute the certain answers to the CQ $q(x) \leftarrow Employee(x), Manages(x, y)$ over the incomplete database (naive tables), by explaining and exploiting the connection with conjunctive query containment:

Employee

<i>name</i>
Smith
$null_1$
Brown

Manages

<i>mgr</i>	<i>mgd</i>
Green	Smith
Smith	$null_1$
$null_1$	Brown
Brown	$null_2$

$q(x) \leftarrow Employee(x), Manages(x, y)$

THE CERTAIN ANSWER IS $\{SMITH, BROWN\}$. THESE ARE THE ONLY TUPLE THAT DOESN'T CONTAIN NULL VALUES, SINCE THESE CAN TAKE ON ANY VALUE.