A New Local Search Procedure for the Minimization of Tool Switches

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- Flexible Manufacturing Systems (FMS) are known for their great flexibility and adaptability on production planning
 - It allows a wide range of item types to be produced;
 - e.g., the same machinery present on a production line can be quickly adjusted to produce an unrelated new set of items.
 - This characteristic is also important when the production must the replanned.

- A common type of FMS, mostly employed by the automotive and metallurgical industries, uses flexible machines
 - Machines able to perform different types of operations requiring only a few changes on their configuration;
 - Cutting, drilling, etc.

- ► Each item requires a set of such tools to be produced and the magazine is able to store all tools for a single item
 - All available tools cannot be loaded at the same time on the magazine;
- Between the production of two different item types it may be neccessary to switch tools on the magazine;
- ► The tool switching imply the interruption of the production line as the machine should be turned off in order to be configured
 - Thus increasing the cost of production.

- From a predetermined demand for products, it becomes neccessary to determine a production plan such that the production line meets this demand;
- ► The plan is divided into jobs and aims to minimize the downtime of the production machines, in order to maximize productivity and reduce related costs;

- ▶ The production plan consists of two parts:
 - Determine the order in which the jobs are processed; and
 - Oecide when to switch tools and what tools will be switched.
- ► The first one is the Minimization of Tool Switches Problem;
- ► The second one is deterministic polynomial time solvable, using the Keep Tools Needed Soonest algorithm.

Problem Data

- ▶ The set of jobs $T = \{1, ..., n\}$;
- ▶ The of set tools $F = \{1, \dots, m\}$;
- ▶ The subset of tools T_j needed to process job $j \in T$;
- ▶ The magazine capacity *C*.

Formal Definition

- ► The *Minimization of Tool Switches Problem* (MTSP) is defined on a binary sparse matrix *M*:
 - ► The n rows correspond to the tools and the m columns correspond to each job;
 - ▶ Entry $m_{ij} = 1$ if tool i is loaded in the machine during job j processing and $m_{ij} = 0$ otherwise.
- A solution is represented by a permutation π of the elements of T and a tool switching plan, resulting in M_{π} ;
- ▶ The number of inversions in M_{π} is the number of tool switches.

An Example

Jobs	1	2	3	4	5	
Tools	1	1	3	2	1	
	2	3	4	3	4	
	4	5	7	5	6	
	7			6		
Capacity = 4						

A Solution Example

π	1	3	5	2	4	_
Tools	1	1	1	1	0	
	1	0	0	0	1	
	0	1	1	1	1	
	1	1	<u>1</u>	0	0	
	0	0	0	1	1	
	0	0	1	<u>1</u>	1	
	1	1	0	0	0	

- Underlined tools denote they are loaded in advance;
- 8 tool switches (the initial four also counts).

Motivation

- It is NP-hard;
- Its practical application in VLSI design industry
 - Printed circuit board assembly;
 - ▶ The assembly tasks are jobs and components are tools.

Proposed Methods

- ▶ A graph search for generating an initial solution;
- A new local search method;
- ► An Iterated Local Search.

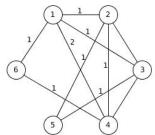
Current State of the Art

- Metaheuristic: Clustering Search + Variable Neighborhood Search + Biased Random Key Genetic Algorithm (2016);
- Exact: Branch & Bound (2009), Integer Linear Programming (2015).

Graph Model

- Nodes represent tools;
- Edges connects tools used in a same job;
- ► No loops or parallel edges;
- Weights are defined by the frequency of tools;
- ► Each job induces a clique.

Jobs	1	2	3	4	5	
Tools	1	1	3	2	1	
	2	3	4	3	4	
	4			5	6	



Graph Search

- ► The main idea is to search for tools that must be kept together in the solution, avoiding switches;
- ▶ Breadth-first search (BFS) in non-decreasing weight of edges order.

Jobs Sequencing

- The BFS returns a list of tools;
- ▶ The jobs sequence is determined in a greedy fashion:
 - ► The list of tools is traversed, and once all tools of a specific job are found, it is sequenced.
- Two tie breakers:
 - Higher number of tools (if no job sequenced yet);
 - Lower number of switches.

Local Search

- Each consecutive pair of 1-blocks in a same row denotes an inversion
 - Minimizing 1-blocks number is a different problem.
- 1-blocks columns are inserted before/after the other 1-blocks;
- The best move is performed.

Iterated Local Search

- Classic version;
- Local search only non deteriorating movements;
- Perturbation: 2-opt of columns
 - Randomly selected, 20% of the columns.
- ▶ All parameters tuned using the irace package.

Methods

- Iterated Local Search (ILS);
- Biased Random Key Genetic Algorithm + VND + Clustering Search (BRKGA + CS).

ILS

- ► Intel Core i5 3.2 GHz processor;
- ▶ 8 GB RAM;
- Ubuntu 15.10 LTS;
- Codes written in C++, compiled with gcc 4.8.4 and the -O3 optimization option.

BRKGA + CS

- Intel Core i7 3.4 GHz processor;
- ▶ 16 GB RAM;

Instances

Two different data sets from the literature were considered, a total of 1670 instances.

- ▶ Yanasse (A, B, C, D, E): 1510 artificial instances;
- ▶ Crama (C_1, C_2, C_3, C_4) : 160 artificial instances.

Average gap from Optimal Solutions (20 runs)

	A	_	~		_
	0.00%				
BRKGA + CS	0.00%	0.00%	0.00%	0.05%	0.00%

^aConsidering only known optimal solutions.

Average Running Times (20 runs)

Method	A	B	C	D	E
ILS					
BRKGA + CS	3.71s	4.05s	9.83s	27.66s	6.54s

Average gap from Best Known Results (20 runs)

Method	C_1	C_2	C_3	C_4
ILS				
BRKGA + CS	0.00%	0.00%	0.00%	0.34%

Average Running Times (20 runs)

Method	C_1	C_2	C_3	C_4
ILS	0.09s	1.11s	168.71s	1063.47s
BRKGA + CS	2.42s	11.58s	123.15s	541.10s

Conclusion

Conclusion

- ► The proposed ILS was able to match a large part of the best results available on the literature and to improve some of them;
- Running time increased on large instances;
- Future work includes:
 - TSP model + solver;
 - New formulations;
 - New lower bounds.

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Questions?

