

CS-E5740 Complex Networks, Answers to exercise set X

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Compile with `pdflatex ex_template.tex`

Problem 1

- a) Adjacency matrix is used to represent the edges in a graph, where 1 is an edge 0 is no connectin

$a)$	1	2	3	4	5	6	7	8
1	0	0	0	1	0	0	0	0
2	0	0	0	1	0	0	0	0
3	0	0	0	1	1	1	0	0
4	1	1	1	0	0	1	0	1
5	0	0	1	0	0	1	0	0
6	0	0	1	1	1	0	0	0
7	0	0	0	0	0	0	1	1
8	0	0	0	1	0	0	1	0

- b) Edge density (ρ) is the number of edges compared to a fully connected network

$$b) \rho = \frac{2 \cdot m}{N(N-1)} = \frac{2 \cdot 9}{8 \cdot 7} \approx 0,32$$

- c) Degree (k) is the number of neighbors of a node

c) Node number	Degree
1	1
2	1
3	3
4	5
5	2
6	3
7	1
8	2

- d) Average degree measures the number of edges over to the number of nodes

$$d) \langle k \rangle = \frac{\sum k_i}{N} = \frac{2m}{N} = \frac{2 \cdot 9}{8} = 2.25$$

- e) Diameter (d) of a network is the maximum distance found in it, where each node is checked for the shortest path

$$e) d \Rightarrow 7 \rightarrow 8 \rightarrow 6 \rightarrow 3 \rightarrow 5 = 4$$

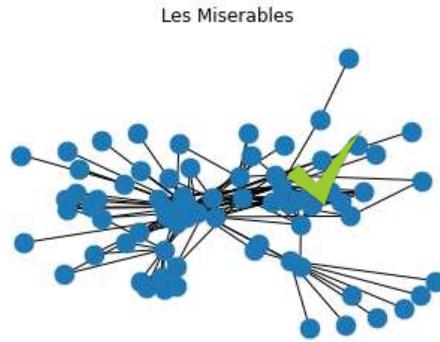
(longest path)
example

f) Clustering coefficient (c) of a node is the number of neighbors connected together

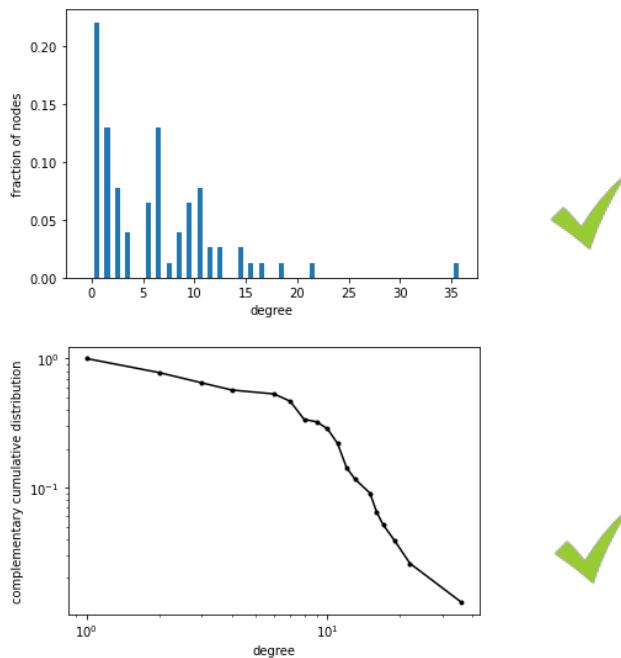
C_1	excluded	$(K_{\text{in}} = 1)$
C_2	excluded	
C_3	$\frac{2 \cdot 2}{4 \cdot 3 \cdot 3 \cdot 2} = 0, \overline{6}$	
C_4	$\frac{2 \cdot 1}{5 \cdot 4} = 0, 1$	
C_5	$\frac{2 \cdot 1}{2 \cdot 1} = 1$	✓
C_6	$\frac{2 \cdot 2}{3 \cdot 2} = 0, \overline{6}$	
C_7	excluded	
C_8	$\frac{2 \cdot 0}{2 \cdot 1} = 0$	

Problem 2

- a) Look at plot below.



- b) D from self-written algorithm: 0.08680792891319207
D from NetworkX: 0.08680792891319207 ✓✓
- c) $\langle l \rangle$ from NetworkX: 2.6411483253588517 ✓✓
- d) $\langle c \rangle$ from NetworkX: 0.5731367499320135 ✓✓
- e) Look plot below



Problem 3

- a) The total is 12 and it's the same in both cases.

a) $V^* = \{1, 2, 3, 4\}$

~~(*)~~ Number of nodes of len 2 of G^*

196

	1	2	3	4
1	✓	✓	✓	X
2	✓	✓	✓	X
3	✓	✓	✓	X
4	X	X	X	X

just used to count

~~(*)~~ A^2 of G^*

$$A^2 = \begin{bmatrix} 1 & 2 & 3 & 4 \\ 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 3 \end{bmatrix}$$

Total 12

\checkmark

b) Look picture

B)

$$\begin{array}{l}
 \left. \begin{array}{l}
 3 \rightarrow 4 \rightarrow 3 \rightarrow 4 \\
 3 \rightarrow 4 \rightarrow 2 \rightarrow 4 \\
 3 \rightarrow 4 \rightarrow 1 \rightarrow 4
 \end{array} \right\} 3 \text{ Walks} \\
 \begin{array}{c}
 A^2 \\
 \begin{array}{c|ccccc}
 & 1 & 2 & 3 & 4 \\
 \hline
 1 & 1 & 1 & 1 & 0 \\
 2 & 1 & 1 & 1 & 0 \\
 3 & 1 & 1 & 1 & 0 \\
 4 & 0 & 0 & 0 & 3
 \end{array}
 \end{array} \\
 \begin{array}{c}
 A^3 = A \cdot A^2 \\
 \Downarrow \\
 (A^3)_{3,4} = 0 \cdot 0 + 0 \cdot 0 + 0 \cdot 0 + 1 \cdot 3 = 3
 \end{array}
 \end{array}$$

~~(A^3)_{3,4} sum any node connected with step 2~~

c) Look picture

Base case: for $m = 1$ using the definition of adjacency matrix the base case is that if there is an edge between the nodes i and j then we write 1, otherwise 0. So $a_{i,j}^{(1)} = A_{i,j}^{(1)}$.

We assume $a_{i,j}^{(m)}$ gives the number of walks of length m , we can express a walk of length $m + 1$ splitting in in length m from v_i to v_k and 1 from v_k to v_j

$$a_{i,j}^{(m+1)} = \sum_{k=1}^{|V|} A_{ik} a_{kj}^{(m)}$$

$\xrightarrow{\text{graph}}$ $\xrightarrow{\text{graph}}$

$$a_{ik}^{(m)} = \sum_{j=1}^{|V|} A_{kj} a_{ij}^{(m)}$$

↓ ↓ ↓ ↓

lem 1 lem m lem 1 lem m



Index of comments

1.1 Total:
6+6+5=17