

## 1 Lecture 18: Orthogonal polynomials

We now introduce orthogonal polynomials (OPs). These are **fundamental** for computational mathematics, with applications in

1. Function approximation
2. Quadrature (calculating integrals)
3. Solving differential equations
4. Spectral analysis of Schrödinger operators

We will investigate the properties of *general* OPs, in this lecture:

1. Definition of orthogonal polynomials
2. Three-term recurrence relationships
3. Function approximation with orthogonal polynomials
4. Construction of orthogonal polynomials via GramSchmidt process

### 1.1 Definition of orthogonal polynomials

Let  $p_0(x), p_1(x), p_2(x), \dots$  be a sequence of polynomials such that  $p_n(x)$  is exactly of degree  $n$ , that is,

$$p_n(x) = k_n x^n + O(x^{n-1})$$

where  $k_n \neq 0$ .

Let  $w(x)$  be a continuous weight function on a (possibly infinite) interval  $(a, b)$ : that is  $w(x) \geq 0$  for all  $a < x < b$ . This induces an inner product

$$\langle f, g \rangle := \int_a^b f(x)g(x)w(x)dx$$

We say that  $\{p_0, p_1, \dots\}$  are *orthogonal with respect to the weight  $w$*  if

$$\langle p_n, p_m \rangle = 0 \quad \text{for} \quad n \neq m.$$

Because  $w$  is continuous, we have

$$\|p_n\|^2 = \langle p_n, p_n \rangle > 0.$$

Orthogonal polynomials are not unique: we can multiply each  $p_n$  by a different nonzero constant  $\tilde{p}_n(x) = c_n p_n(x)$ , and  $\tilde{p}_n$  will be orthogonal w.r.t.  $w$ . However, if we specify  $k_n$ , this is sufficient to uniquely define them:

**Proposition (Uniqueness of OPs I)** Given a non-zero  $k_n$ , there is a unique polynomial  $p_n$  orthogonal w.r.t.  $w$  to all lower degree polynomials.

**Proof** Suppose  $r_n(x) = k_n x^n + O(x^{n-1})$  is another OP w.r.t.  $w$ . We want to show  $p_n - r_n$  is zero. But this is a polynomial of degree  $< n$ , hence

$$p_n(x) - r_n(x) = \sum_{k=0}^{n-1} c_k p_k(x)$$

But we have for  $k \leq n-1$

$$\langle p_k, p_k \rangle c_k = \langle p_n - r_n, p_k \rangle = \langle p_n, p_k \rangle - \langle r_n, p_k \rangle = 0 - 0 = 0$$

which shows all  $c_k$  are zero.

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**Corollary (Uniqueness of OPs I)** If  $q_n$  and  $p_n$  are orthogonal w.r.t.  $w$  to all lower degree polynomials, then  $q_n(x) = C p_n(x)$  for some constant  $C$ .

### 1.1.1 Monic orthogonal polynomials

If  $k_n = 1$ , that is,

$$p_n(x) = x^n + O(x^{n-1})$$

then we refer to the orthogonal polynomials as monic.

Monic OPs are unique as we have specified  $k_n$ .

### 1.1.2 Orthonormal polynomials

If  $\|p_n\| = 1$ , then we refer to the orthogonal polynomials as orthonormal w.r.t.  $w$ . We will usually use  $q_n$  when they are orthonormal. Note it's not unique: we can multiply by  $\pm 1$  without changing the norm.

**Remark** The classical OPs are neither monic nor orthonormal (apart from one case). Many people make the mistake of using orthonormal polynomials for computations. But there is a good reason to use classical OPs: their properties result in rational formulae, whereas orthonormal polynomials require square roots. This makes a performance difference.

## 1.2 Function approximation with orthogonal polynomials

A basic usage of orthogonal polynomials is for polynomial approximation. Suppose  $f(x)$  is a degree  $n-1$  polynomial. Since  $\{p_0(x), \dots, p_{n-1}(x)\}$  span all degree  $n-1$  polynomials, we know that

$$f(x) = \sum_{k=0}^{n-1} f_k p_k(x)$$

where

$$f_k = \frac{\langle f, p_k \rangle}{\langle p_k, p_k \rangle}$$

Sometimes, we want to incorporate the weight into the approximation. This is typically one of two forms, depending on the application:

$$f(x) = w(x) \sum_{k=0}^{\infty} f_k p_k(x)$$

or

$$f(x) = \sqrt{w(x)} \sum_{k=0}^{\infty} f_k p_k(x)$$

The  $w(x)p_k(x)$  or  $\sqrt{w(x)}p_k(x)$  are called weighted polynomials.

### 1.3 Jacobi operators and three-term recurrences for general orthogonal polynomials

#### 1.3.1 Three-term recurrence relationships

A central theme: if you know the Jacobi operator / three-term recurrence, you know the polynomials. This is the **best** way to evaluate expansions in orthogonal polynomials: even for cases where we have explicit formulae (e.g. Chebyshev polynomials  $T_n(x) = \cos n \arccos x$ ), using the recurrence avoids evaluating trigonometric functions.

Every family of orthogonal polynomials has a three-term recurrence relationship:

**Theorem (three-term recurrence)** Suppose  $\{p_n(x)\}$  are a family of orthogonal polynomials w.r.t. a weight  $w(x)$ . Then there exists constants  $a_n$ ,  $b_n \neq 0$  and  $c_n$  such that

$$\begin{aligned} x p_0(x) &= a_0 p_0(x) + b_0 p_1(x) \\ x p_n(x) &= c_n p_{n-1}(x) + a_n p_n(x) + b_n p_{n+1}(x) \end{aligned}$$

**Proof** The first part follows since  $p_0(x)$  and  $p_1(x)$  span all degree 1 polynomials.

The second part follows essentially because multiplication by  $x$  is "self-adjoint", that is,

$$\langle x f, g \rangle = \int_a^b x f(x) g(x) w(x) dx = \langle f, x g \rangle$$

Therefore, if  $f_m$  is a degree  $m < n - 1$  polynomial, we have

$$\langle x p_n, f_m \rangle = \langle p_n, x f_m \rangle = 0.$$

In particular, if we write

$$x p_n(x) = \sum_{k=0}^{n+1} C_k p_k(x)$$

then

$$C_k = \frac{\langle xp_n, p_k \rangle}{\|p_k\|^2} = 0$$

if  $k < n - 1$ .

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Monic polynomials clearly have  $b_n = 1$ . Orthonormal polynomials have an even simpler form:

**Theorem (orthonormal three-term recurrence)** Suppose  $\{q_n(x)\}$  are a family of orthonormal polynomials w.r.t. a weight  $w(x)$ . Then there exists constants  $a_n$  and  $b_n$  such that

$$\begin{aligned} xq_0(x) &= a_0q_0(x) + b_0q_1(x) \\ xq_n(x) &= b_{n-1}q_{n-1}(x) + a_nq_n(x) + b_nq_{n+1}(x) \end{aligned}$$

**Proof** Follows again by self-adjointness of multiplication by  $x$ :

$$c_n = \langle xq_n, q_{n-1} \rangle = \langle q_n, xq_{n-1} \rangle = \langle xq_{n-1}, q_n \rangle = b_{n-1}$$

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**Corollary (symmetric three-term recurrence implies orthonormality)** Suppose  $\{p_n(x)\}$  are a family of orthogonal polynomials w.r.t. a weight  $w(x)$  such that

$$\begin{aligned} xp_0(x) &= a_0p_0(x) + b_0p_1(x) \\ xp_n(x) &= b_{n-1}p_{n-1}(x) + a_np_n(x) + b_np_{n+1}(x) \end{aligned}$$

with  $b_n \neq 0$ . Suppose further that  $\|p_0\| = 1$ . Then  $p_n$  must be orthonormal.

**Proof** This follows from

$$b_n = \frac{\langle xp_n, p_{n+1} \rangle}{\|p_{n+1}\|^2} = \frac{\langle xp_{n+1}, p_n \rangle}{\|p_{n+1}\|^2} = b_n \frac{\|p_n\|^2}{\|p_{n+1}\|^2}$$

which implies

$$\|p_{n+1}\|^2 = \|p_n\|^2 = \cdots = \|p_0\|^2 = 1$$

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**Remark** We can scale  $w(x)$  by a constant without changing the orthogonality properties, hence we can make  $\|p_0\| = 1$  by changing the weight.

**Remark** This is beyond the scope of this course, but satisfying a three-term recurrence like this such that coefficients are bounded with  $p_0(x) = 1$  is sufficient to show that there exists a  $w(x)$  (or more accurately, a Borel measure) such that  $p_n(x)$  are orthogonal w.r.t.  $w$ . The relationship between the coefficients  $a_n, b_n$  and the  $w(x)$  is an object of study in spectral theory, particularly when the coefficients are periodic, quasi-periodic or random.

## 1.4 Jacobi operators and multiplication by $x$

We can rewrite the three-term recurrence as

$$x \begin{pmatrix} p_0(x) \\ p_1(x) \\ p_2(x) \\ \vdots \end{pmatrix} = J \begin{pmatrix} p_0(x) \\ p_1(x) \\ p_2(x) \\ \vdots \end{pmatrix}$$

where  $J$  is a Jacobi operator, an infinite-dimensional tridiagonal matrix:

$$J = \begin{pmatrix} a_0 & b_0 & & & \\ c_1 & a_1 & b_1 & & \\ & c_2 & a_2 & b_2 & \\ & & c_3 & a_3 & \ddots \\ & & & \ddots & \ddots \end{pmatrix}$$

When the polynomials are monic, we have 1 on the superdiagonal. When we have an orthonormal basis, then  $J$  is symmetric:

$$J = \begin{pmatrix} a_0 & b_0 & & & \\ b_0 & a_1 & b_1 & & \\ & b_1 & a_2 & b_2 & \\ & & b_2 & a_3 & \ddots \\ & & & \ddots & \ddots \end{pmatrix}$$

Given a polynomial expansion,  $J$  tells us how to multiply by  $x$  in coefficient space, that is, if

$$f(x) = \sum_{k=0}^{\infty} f_k p_k(x) = (p_0(x), p_1(x), \dots) \begin{pmatrix} f_0 \\ f_1 \\ f_2 \\ \vdots \end{pmatrix}$$

then (provided the relevant sums converge absolutely and uniformly)

$$xf(x) = x(p_0(x), p_1(x), \dots) \begin{pmatrix} f_0 \\ f_1 \\ f_2 \\ \vdots \end{pmatrix} = \left( J \begin{pmatrix} p_0(x) \\ p_1(x) \\ p_2(x) \\ \vdots \end{pmatrix} \right)^{\top} \begin{pmatrix} f_0 \\ f_1 \\ f_2 \\ \vdots \end{pmatrix} = (p_0(x), p_1(x), \dots) X \begin{pmatrix} f_0 \\ f_1 \\ f_2 \\ \vdots \end{pmatrix}$$

where  $X := J^{\top}$ .

### 1.4.1 Evaluating polynomials

We can use the three-term recurrence to construct the polynomials. One way to express this is in the language of linear algebra. Suppose we are given  $p_0(x) = k_0$  (where  $k_0 = 1$  is pretty much always the case in practice). This can be written in matrix form as

$$(1, 0, 0, 0, 0, \dots) \begin{pmatrix} p_0(x) \\ p_1(x) \\ p_2(x) \\ \vdots \end{pmatrix} = k_0$$

We can combine this with the Jacobi operator to get

$$\underbrace{\begin{pmatrix} 1 & & & & \\ a_0 - x & b_0 & & & \\ c_1 & a_1 - x & b_1 & & \\ & c_2 & a_2 - x & b_2 & \\ & & c_3 & a_3 - x & b_3 \\ & & & \ddots & \ddots & \ddots \end{pmatrix}}_{L_x} \begin{pmatrix} p_0(x) \\ p_1(x) \\ p_2(x) \\ \vdots \end{pmatrix} = \begin{pmatrix} k_0 \\ 0 \\ 0 \\ \vdots \end{pmatrix}$$

For  $x$  fixed, this is a lower triangular system, that is, the polynomials equal

$$k_0 L_x^{-1} \mathbf{e}_0$$

This can be solved via forward recurrence:

$$\begin{aligned} p_0(x) &= k_0 \\ p_1(x) &= \frac{(x - a_0)p_0(x)}{b_0} \\ p_2(x) &= \frac{(x - a_1)p_0(x) - c_1 p_0(x)}{b_1} \\ p_3(x) &= \frac{(x - a_2)p_1(x) - c_2 p_1(x)}{b_2} \\ &\vdots \end{aligned}$$

We can use this to evaluate functions as well:

$$f(x) = (p_0(x), p_1(x), \dots) \begin{pmatrix} f_0 \\ f_1 \\ \vdots \end{pmatrix} = k_0 \mathbf{e}_0^\top L_x^{-\top} \begin{pmatrix} f_0 \\ f_1 \\ \vdots \end{pmatrix}$$

when  $f$  is a degree  $n-1$  polynomial, this becomes a problem of inverting an upper triangular matrix, that is, we want to solve the  $n \times n$  system

$$\underbrace{\begin{pmatrix} 1 & a_0 - x & c_1 & & & \\ & b_0 & a_1 - x & c_2 & & \\ & & b_1 & a_2 - x & \ddots & \\ & & & b_2 & \ddots & c_{n-2} \\ & & & & \ddots & a_{n-2} - x \\ & & & & & b_{n-2} \end{pmatrix}}_{L_x^\top} \begin{pmatrix} \gamma_0 \\ \vdots \\ \gamma_{n-1} \end{pmatrix}$$

so that  $f(x)/k_0 = \gamma_0$ . We can solve this via back-substitution:

$$\begin{aligned}\gamma_{n-1} &= \frac{f_{n-1}}{b_{n-2}} \\ \gamma_{n-2} &= \frac{f_{n-2} - (a_{n-2} - x)\gamma_{n-1}}{b_{n-3}} \\ \gamma_{n-3} &= \frac{f_{n-3} - (a_{n-3} - x)\gamma_{n-2} - c_{n-2}\gamma_{n-1}}{b_{n-4}} \\ &\vdots \\ \gamma_1 &= \frac{f_1 - (a_1 - x)\gamma_2 - c_2\gamma_3}{b_0} \\ \gamma_0 &= f_0 - (a_0 - x)\gamma_1 - c_1\gamma_2\end{aligned}$$

We give examples of these algorithms applied to Chebyshev polynomials in the next lecture.

## 1.5 GramSchmidt algorithm

In general we don't have nice formulae for OPs but we can always construct them via the GramSchmidt process:

**Proposition (GramSchmidt)** Define

$$\begin{aligned}p_0(x) &= 1 \\ q_0(x) &= \frac{1}{\|p_0\|} \\ p_{n+1}(x) &= xq_n(x) - \sum_{k=0}^n \langle xq_n, q_k \rangle q_k(x) \\ q_{n+1}(x) &= \frac{p_{n+1}(x)}{\|p_{n+1}\|}\end{aligned}$$

Then  $q_0(x), q_1(x), \dots$  are orthonormal w.r.t.  $w$ .

**Proof** By linearity we have

$$\langle p_{n+1}, q_j \rangle = \left\langle xq_n - \sum_{k=0}^n \langle xq_n, q_k \rangle q_k, q_j \right\rangle = \langle xq_n, q_j \rangle - \langle xq_n, q_j \rangle \langle q_j, q_j \rangle = 0$$

Thus  $p_{n+1}$  is orthogonal to all lower degree polynomials. So is  $q_{n+1}$ , since it is a constant multiple of  $p_{n+1}$ .

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Let's make our own family of OPs:

```
using ApproxFun, Plots
x = Fun()
w = exp(x)
ip = (f,g) -> sum(f*g*w)
nrm = f -> sqrt(ip(f,f))
```

```

n = 10
q = Array{Fun}(undef,n)
p = Array{Fun}(undef,n)
p[1] = Fun(1, -1 .. 1 )
q[1] = p[1]/nrm(p[1])

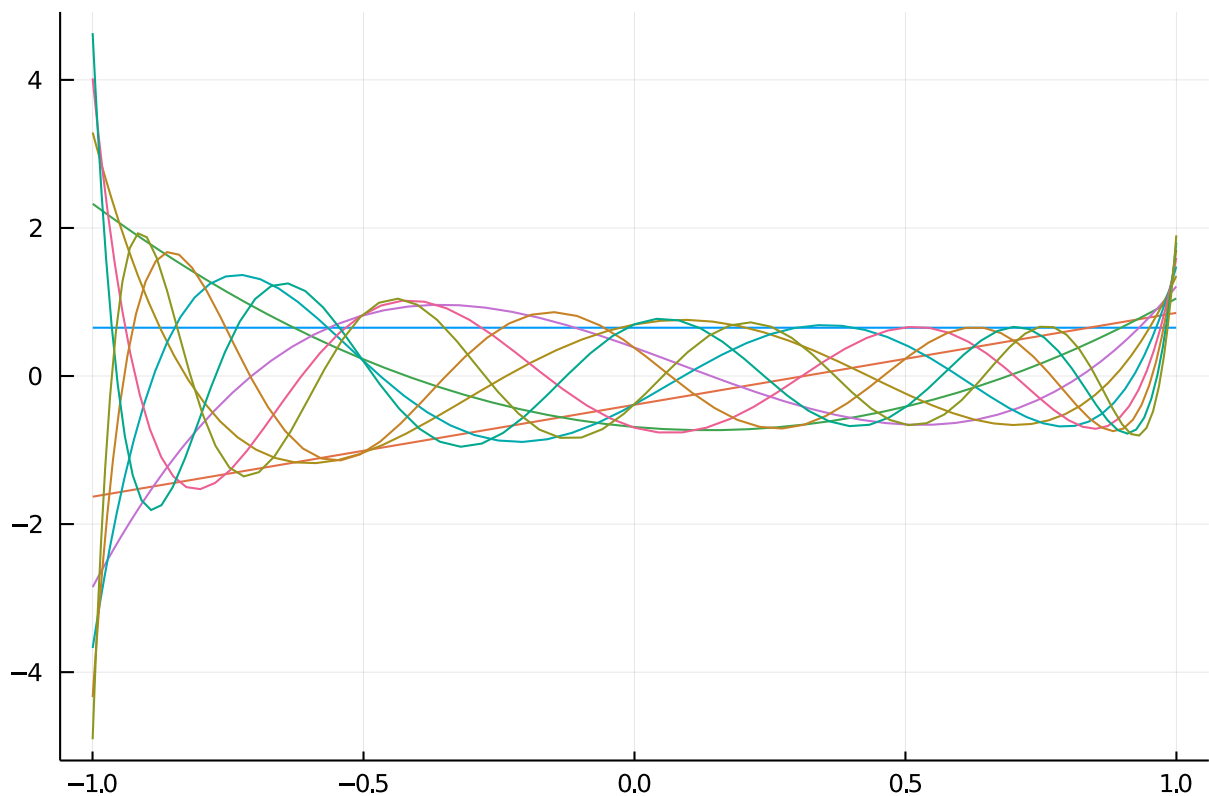
for k=1:n-1
    p[k+1] = x*q[k]
    for j=1:k
        p[k+1] -= ip(p[k+1],q[j])*q[j]
    end
    q[k+1] = p[k+1]/nrm(p[k+1])
end

@show sum(q[2]*q[4]*w)

p = plot(; legend=false)
for k=1:10
    plot!(q[k])
end
p

sum(q[2] * q[4] * w) = 8.167294965333305e-16

```



The three-term recurrence means we can simplify GramSchmidt, and calculate the recurrence coefficients at the same time:

**Proposition (GramSchmidt)** Define



$$\begin{aligned}
p_0(x) &= 1 \\
q_0(x) &= \frac{1}{\|p_0\|} \\
a_n &= \langle xq_n, q_n \rangle \\
b_{n-1} &= \langle xq_n, q_{n-1} \rangle \\
p_{n+1}(x) &= xq_n(x) - a_nq_n(x) - b_{n-1}q_{n-1}(x) \\
q_{n+1}(x) &= \frac{p_{n+1}(x)}{\|p_{n+1}\|}
\end{aligned}$$

Then  $q_0(x), q_1(x), \dots$  are orthonormal w.r.t.  $w$ .

**Remark** This can be made a bit more efficient by using  $\|p_{n+1}\|$  to calculate  $b_n$ .

```

x = Fun()
w = exp(x)
ip = (f,g) -> sum(f*g*w)
nrm = f -> sqrt(ip(f,f))
n = 10
q = Array{Fun}(undef, n)
p = Array{Fun}(undef, n)
a = zeros(n)
b = zeros(n)
p[1] = Fun(1, -1 .. 1)
q[1] = p[1]/nrm(p[1])

p[2] = x*q[1]
a[1] = ip(p[2],q[1])
p[2] -= a[1]*q[1]
q[2] = p[2]/nrm(p[2])

for k=2:n-1
    p[k+1] = x*q[k]
    b[k-1] = ip(p[k+1],q[k-1])
    a[k] = ip(p[k+1],q[k])
    p[k+1] = p[k+1] - a[k]q[k] - b[k-1]q[k-1]
    q[k+1] = p[k+1]/nrm(p[k+1])
end

ip(q[5],q[2]) # shows orthogonality (to numerical accuracy)

1.0755285551056204e-15

```

Here we see a plot of the first 10 polynomials:

```

p = plot(; legend=false)
for k=1:10
    plot!(q[k])
end
p

```

