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## 1 Template

```
#include <bits/stdc++.h>
using namespace std;
using 11 =
                       long long;
#define vll
                       vector<11>
#define vvll
                       vector <vll>>
#define pll
                       pair<ll, ll>
                       vector <pll>
#define vpll
#define vvpll
                       vector < vpll >
#define endl '\n'
#define all(xs)
                       xs.begin(), xs.end()
#define found(x, xs) (xs.find(x) != xs.end())
```

## 2 Search

## 2.1 Ternary Search

 $O(\log n)$ 

Function f(x) is unimodal on an interval [l, r]. Unimodal means: the function strictly increases first, reaches a maximum, and then strictly decreases OR the function strictly decreases first, reaches a minimum and then strictly decreases

```
double ternary_search(double 1, double r) {
       double eps = 1e-9; // error limit
2
       while (r - 1 > eps) {
           double m1 = 1 + (r-1) / 3;
           double m2 = r - (r-1) / 3;
            double f1 = f(m1);
            double f2 = f(m2);
            if(f1 < f2)
10
                1 = m1;
            else
               r = m2;
       }
14
15
       return f(1);
   }
17
```

# 3 Algebra

#### 3.1 All divisors

 $O(\sqrt{n})$ 

## 3.2 Primality test

 $O(\sqrt{n})$ 

```
bool isPrime(ll n)
{
    if(n!=2 && n % 2==0)
        return false;

    for(ll d=3; d*d <= n; d+=2)
    {
        if(n % d==0)
            return false;
    }

    return n >= 2;
}
```

## 3.3 Binary exponentiation

 $O(\log n)$ 

#### 3.4 Greatest common divisor

 $O(\log \min(a, b))$ 

#### 3.4.1 Least common multiple

#### 3.4.2 Extended Euclides Algorithm

```
11 gcd(ll a, ll b, ll& x, ll& y) {
       if (b == 0) {
2
           x = 1;
           y = 0;
           return a;
       }
6
       ll x1, y1;
       11 d = gcd(b, a \% b, x1, y1);
       x = y1;
       y = x1 - y1 * (a / b);
10
       return d;
11
12
```

## 3.5 Linear Diophantine Equations

 $O(\log \min(a, b))$ 

#### 3.5.1 Any solution

```
bool find_any_solution(ll a, ll b, ll c, ll &x0, ll &y0, ll
    &g) {
       g = gcd(abs(a), abs(b), x0, y0);
       if (c % g) {
            return false;
       }
}
```

```
x0 *= c / g;

y0 *= c / g;

if (a < 0) x0 = -x0;

if (b < 0) y0 = -y0;

return true;

}
```

## 3.6 Integer Factorization

## 3.6.1 Pollard's Rho

 $O(\sqrt[4]{n}\log n)$ 

```
/**
       @param a first multiplier
       @param b second multiplier
       @param mod
       @return a * b mod n (without overflow)
       @brief Multiplies two numbers >= 10^18
       Time Complexity: O(log b)
   ll mult(ll a, ll b, ll mod) {
10
       11 result = 0;
       while (b) {
           if (b & 1)
               result = (result + a) % mod;
           a = (a + a) \% mod;
           b >>= 1;
16
       return result;
17
18
19
   /**
20
    * Oparam x first multiplier
      @param c second multiplier
       @param mod
       Oreturn f(x) = x^2 + c \mod (mod)
       Obrief Polynomial function chosen for pollard's rho
       Time Complexity: 0(1)
26
   11 f(11 x, 11 c, 11 mod) {
       return (mult(x, x, mod) + c) % mod;
29
   }
30
31
   /**
32
      @param n number that we want to find a factor p
       @param x0 number where we will start
       Oparam c constant in polynomial function
       Oreturn fac
```

```
Obrief Pollard's Rho algorithm (works only for composite
37
         numbers)
       if (g==n) try other starting values
38
       Time Complexity: O(n^{(1/4)} \log n)
39
   */
   ll rho(ll n, ll x0=2, ll c=1) {
41
       11 x = x0;
42
       11 y = x0;
43
       11 g = 1;
44
       while (g == 1) {
45
           x = f(x, c, n);
           y = f(y, c, n);
47
           y = f(y, c, n);
48
           g = gcd(abs(x - y), n);
49
50
       return g;
51
```

## 4 Graphs

## 4.1 DFS

```
O(n+m)
```

```
void dfs(ll at, ll n ,vpll adj[], bool visited[]) {
   if(visited[at])
      return;

visited[at] = true;

vpll neighbours = adj[at];
for(auto nex: neighbours)
   dfs(nex.first, n, adj, visited);
}
```

## 4.2 BFS

```
O(n+m)
```

```
void bfs(ll s, ll n, vll adj[]) {
   bool visited[n] = {0};
   visited[s] = true;

queue<ll> q;
   q.push(s);
   while (!q.empty())
   {
      vll neighbours = adj[q.front()];
}
```

```
for(auto nex: neighbours) {
                 if(!visited[nex]) {
                      visited[nex]=true;
12
                      q.push(nex);
13
                 }
14
15
            }
            cout << q.front() << '\n';</pre>
16
17
            q.pop();
        }
18
   }
19
```

#### 4.2.1 Shortest path on unweighted graph

O(n+m)

```
vll solve(ll s, ll n, vll adj[]) {
       bool visited[n] = {0};
2
       visited[s] = true;
       queue <11> q;
       q.push(s);
       vll prev(n, -1);
       while (!q.empty())
            vll neighbours = adj[q.front()];
10
            for(auto nex: neighbours) {
11
                if(!visited[nex]) {
                     visited[nex]=true;
13
                     q.push(nex);
14
                     prev[nex] = q.front();
15
                }
16
            }
17
            q.pop();
18
19
20
       return prev;
21
   }
22
23
   vll reconstructPath(ll s, ll e, vll prev) {
24
       vll path;
25
       for(ll i=e; i!=-1; i=prev[i])
26
            path.push_back(i);
27
28
       reverse(path.begin(), path.end());
29
30
       if (path [0] == s)
31
            return path;
32
        else {
33
            vll place;
```

## 4.3 Flood Fill

```
O(n+m)
```

```
int dir_y[] = {};
  int dir_x[] = {};
   int ff(int i, int j, char c1, char c2) {
       if ((i < 0) || (i >= n)) return 0;
       if ((j < 0) || (j >= m)) return 0;
6
       if (grid[i][j] != c1) return 0;
       int ans = 1;
       grid[i][j] = c2;
10
11
       for (int d = 0; d < 8; ++d)</pre>
           ans += floodfill(i+dir_y[d], j+dir_x[d], c1, c2);
13
14
15
       return ans;
  }
```

## 4.4 Topological Sort (Directed Acyclic Graph)

#### 4.4.1 DFS Variation

O(n+m)

```
void dfs(ll at, ll n ,vpll adj[], bool visited[], vll &ts) {
   if(visited[at])
      return;

visited[at] = true;

vpll neighbours = adj[at];
for(auto nex: neighbours)
   dfs(nex.first, n, adj, visited);
ts.push_back(at); // Only change
}
```

#### 4.4.2 Kahn's Algorithm

```
priority_queue<11, v11, greater<11>> pq;
   for(11 at=0; at<n; at++)</pre>
                                     // Push all sources of
       connected components in graph
       if(in_degree[at] == 0)
           pq.push(at);
5
   while(!pq.empty()) {
6
       11 at = pq.top(); pq.pop();
       vll neighbors = adj[at];
       for(auto nex: neighbors) {
           in_degree[nex]--;
10
           if(in_degree[nex]>0) continue;
11
           pq.push(nex);
       }
  }
14
```

## 4.5 Bipartite Graph Check (Undirected Graph)

O(n+m)

```
bool isBipartite(ll s, ll n, vll adj[]) {
       queue <11> q;
       q.push(s);
       vll color(n, -1); color[s]=0;
       bool flag = true;
       while (!q.empty())
6
       {
            vll neighbours = adj[q.front()];
            for(auto nex: neighbours) {
10
                if(color[nex] == -1) {
                    color[nex] = 1-(color[q.front()]);
11
12
                    q.push(nex);
                }
                else if(color[nex] == color[q.front()]) {
14
                    flag = false;
15
                    break;
                }
            }
18
            q.pop();
19
20
21
22
       return flag;
   }
```

## 4.6 Cycle Check (Directed Graph)

O(n+m)

```
enum { UNVISITED = -1, VISITED = -2, EXPLORED=-3};
2
   void cycleCheck(ll at, ll n ,vll adj[], int visited[], ll
3
      dfs_parent[]) {
       visited[at] = EXPLORED;
       vll neighbours = adj[at];
       for(auto nex: neighbours) {
           if(visited[nex] == UNVISITED) {
               // Tree edges (part of the DFS spanning tree)
               dfs_parent[nex] = at;
               cycleCheck(nex, n, adj, visited);
           else if(visited[nex] == EXPLORED) {
13
               if(nex == dfs_parent[at]) {
14
                    // Trivial cycle
                    // Do something
16
               }
17
               else {
                    // Non trivial cycle - Back Edge ((u, v)
19
                       such that v is the ancestor of node u but
                        is not part of the DFS tree)
                    // Do something
20
21
           else if(visited[nex] == VISITED) {
               // Forward/Cross edge ((u, v) such that v is a
25
                   descendant but not part of the DFS tree)
               // Do something
26
27
28
30
       visited[at] = VISITED;
31
32
```

## 4.7 Dijkstra

 $O(n\log n + m\log n)$ 

```
void dijkstra(ll s, vll & d, vll & p) {
    d.assign(n, LLONG_MAX);
    p.assign(n, -1);

d[s] = 0;
```

```
priority_queue<pll, vpll, greater<pll>> q;
6
       q.push({0, s});
       while (!q.empty()) {
            11 v = q.top().second;
9
            ll d_v = q.top().first;
10
            q.pop();
11
            if (d_v != d[v])
                continue;
13
14
            for (auto edge : adj[v]) {
15
                11 to = edge.first;
                11 len = edge.second;
17
18
                if (d[v] + len < d[to]) {</pre>
19
                     d[to] = d[v] + len;
20
                     p[to] = v;
21
                     q.push({d[to], to});
22
                }
            }
24
       }
25
   }
26
```

## 5 Math Formulas

## 5.1 Sum of an arithmetic progression

$$S_n = \frac{n}{2}(a_1 + a_n)$$

## 5.2 Permutation with repeated elements

$$P_n = \frac{n!}{n_1! n_2! \dots n_k!}$$

## 5.3 Check if is geometric progression

$$a_i^2 = a_{i-1} a_{i+1}$$

## 5.4 Bitwise equations

$$\begin{aligned} a|b &= a \oplus b + a\&b \\ a \oplus (a\&b) &= (a|b) \oplus b \\ (a\&b) \oplus (a|b) &= a \oplus b \\ \\ a+b &= a|b+a\&b \\ a+b &= a \oplus b + 2(a\&b) \\ \\ a-b &= (a \oplus (a\&b)) - ((a|b) \oplus a) \\ a-b &= ((a|b) \oplus b) - ((a|b) \oplus a) \end{aligned}$$

$$\begin{array}{l} a-b=(a\oplus(a\&b))-(b\oplus(a\&b))\\ a-b=((a|b)\oplus b)-(b\oplus(a\&b)) \end{array}$$

## 5.5 Cube of Binomial

$$(a+b)^3 = a^3 + 3a^2b + 3ab^2 + b^3$$
$$(a-b)^3 = a^3 - 3a^2b + 3ab^2 - b^3$$

## 5.5.1 Sum of Cubes

$$a^3 + b^3 = (a+b)(a^2 - ab + b^2)$$

#### 5.5.2 Difference of Cubes

$$a^3 - b^3 = (a - b)(a^2 + ab + b^2)$$

## 5.6 Binomial expansion

$$\binom{n}{k} = \frac{n!}{k!(n-k)!}$$
$$(a+b)^n = \sum_{k=0}^n \binom{n}{k} a^k b^{n-k}$$

## 6 Facts

## 6.1 XOR

#### 6.1.1 Self-inverse property

To cancel a XOR, you can XOR again the same value because  $a\oplus a=0,$  so  $(value\oplus a)\oplus a=value$ 

#### 6.1.2 Identity element

$$a \oplus 0 = a$$

## 6.1.3 Commutative

$$a\oplus b=b\oplus a$$

### 6.1.4 Associative

$$(a \oplus b) \oplus c = a \oplus (b \oplus c)$$