# DM549/DS(K)820/MM537/DM547

Lecture 3: More on Quantifiers

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## More on Quantifiers

#### Remarks:

- We say that the quantifier binds variables x.
- In the above statements, we call *D* the *domain* (domæne) or universe (univers).
- We also say that we *quantify over* (kvantificerer over) D.
- When clear from the context, the domain is sometimes left out.
- Some authors leave out the colon.
- How to memorize?
  - ► for ∀II.
  - ▶ there ∃xists.
  - "!" looks a bit like "1".
- Quantifiers have a *higher* preference (i.e., they are evaluated earlier) than the operators  $\neg$ ,  $\wedge$ ,  $\vee$ ,  $\Rightarrow$ ,  $\Leftrightarrow$ ,  $\oplus$ .

# A Quiz

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- Read: "... x in D with Q(x) ...".
- Notice the difference in how the quantification with restricted domain can be translated into a quantification without restricted domain.

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**Interpretation:** Can pull " $\neg$ " to the right, but, by doing so, we "flip" quantifiers ( $\forall$  changes to  $\exists$  and  $\exists$  changes to  $\forall$ ).

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- By adding another quantifier in front (and binding another variable), one obtains a *proposition* (with no variables!).
- The resulting proposition has two *nested* quantifiers.

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**In general:** We may exchange consecutive quantifiers if and only if they are of the same type!

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### Seriously:

- If you express a mathematical statement in natural language, make sure it is unambiguous!
- When writing down as formal logic, such ambiguities cannot happen.
  - ► (This is why we are learning about this topic!)

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- Or directly: Flip all quantifiers and move the negation all the way inwards.

## More Quantifiers

#### Rule of Thumb

number of quantifier flips (as you read from left to right)

complexity of proposition

### Test 1

Don't forget that the test has been opened today. Good luck!