

DM549/DS820/MM537/DM547

Lecture 12: Partial Orders, Modular Arithmetic

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# Last Time: Transitive Closure

## Definition

For a relation  $R$  on a set  $A$ ,

$$R^* = \bigcup_{i=1}^{\infty} R^i .$$

$\underbrace{\hspace{1.5cm}}_{R \cup R^2 \cup R^3 \cup \dots}$

## Theorem (only proof sketch, Theorem 9.4.2)

The transitive closure of a relation  $R$  is

$$t(R) = R^* .$$

# Last Time: Equivalence Relations

## Definitions:

- Equivalence relation: reflexive, symmetric, and transitive relation.
- Equivalence class of  $a$  with respect to equivalence relation  $R$ :  $[a]_R$ , the set of elements related to  $a$ .

## Theorem (Theorem 9.5.2)

Let  $A$  be a set. There is a one-to-one correspondence between equivalence relations on  $A$  and partitions of  $A$ :

- (1) For any equivalence relation  $R$  on  $A$ ,  $P = \{[a]_R \mid a \in A\}$  is a partition of  $A$ .
- (2) For any partition  $P = \{A_i \mid i \in I\}$  of  $A$ , there exists an equivalence relation  $R$  on  $A$  such that  $\{[a]_R \mid a \in A\} = \{A_i \mid i \in I\}$ .

## Definition (Definition 9.6.1)

A relation  $R$  on a set  $A$  is called a *partial order* (partiel ordning) if it is

- reflexive,
- **anti**symmetric, and
- transitive.

If this is the case,  $(A, R)$  is called a *partially ordered set* (partielt ordnet mængde) or *poset*.

## Remarks:

- Instead of  $R$ , one often uses  $\leq$  or  $\preceq$  for partial orders.
- When using these notations,  $a < b$  ( $a \prec b$ ) can be used to indicate that  $a \leq b$  ( $a \preceq b$ ) and  $a \neq b$ .

# Hasse Diagrams

**Idea:** Special representation of a partially ordered set  $(A, \preceq)$ .

**Specifically:** Like graph representation of relation but

- leave out edges implied by the relation being reflexive and transitive, and
- if  $a \preceq b$  and edge not implied, leave out arrow head but draw  $a$  under  $b$ .
  - ▶ Since partial orders are transitive and antisymmetric, there are no cycles, and this is possible!

# Special Elements of Partially Ordered Sets

## Definition (cf. Section 9.6.4)

Let  $(A, \preceq)$  be a poset. For  $a \in A$ ,  $a$  is called

- a *minimal* element if  $\neg \exists b \in A : b \prec a$ .
- the *least* element if  $\forall b \in A : a \preceq b$ .
- a *maximal* element if  $\neg \exists b \in A : a \prec b$ .
- the *greatest* element if  $\forall b \in A : b \preceq a$ .

## Remarks:

- Every least (greatest) element, is also a minimal (maximal) element, but not necessarily the other way around.
- If  $A$  is non-empty and finite, there always exists a minimal (maximal) element, but not necessarily least (greatest).

# Total Orders

## Definition (Definition 9.6.2)

Let  $(A, \preceq)$  be a poset. We call  $a, b \in A$  *comparable* (sammenlignelige) if  $a \preceq b$  or  $b \preceq a$ .

## Definition (Definition 9.6.3)

Let  $(A, \preceq)$  be a poset. If all  $a, b \in A$  are comparable, we call  $\preceq$  a *total order* (total ordning).

**Different view:** Partial orders can be obtained from total order by removing edges.

# The Lexicographic Order

## Definition (cf. Section 9.6.2)

Let  $(A_1, \preceq_1), (A_2, \preceq_2), \dots, (A_n, \preceq_n)$  be partial orders. Then we can define a *lexicographic order*, a partial order,  $\preceq$  on  $A_1 \times A_2 \times \dots \times A_n$  as follows.

For two different elements  $(a_1, a_2, \dots, a_n), (b_1, b_2, \dots, b_n)$  of  $A_1 \times A_2 \times \dots \times A_n$  that are not equal,  $(a_1, a_2, \dots, a_n) \preceq (b_1, b_2, \dots, b_n)$  holds if and only if

- $a_1 \prec_i b_1$ , or
- there exists an  $i > 0$  such that  $a_1 = b_1, a_2 = b_2, \dots, a_i = b_i$  and  $a_{i+1} \prec_i b_{i+1}$ .

## How to think about this:

- Think of how words are ordered in a dictionary.
- For instance, consider  $n = 4$  and all four partial (in fact, total) orders are  $(\{a, b, c, \dots, z\}, \leq)$  where  $\ell_1 \leq \ell_2$  iff  $\ell_1$  is no later than  $\ell_2$  in the alphabet.
- Then the corresponding lexicographic order can be viewed as the total order in which the corresponding four-letter words would appear in a dictionary.



# A Quiz

Go to [pollev.com/kevs](https://pollev.com/kevs)



# On the Exam

**Date:** 8 January, 2025.

**Duration:**

- DM547, MM537: 3 hours
- DM549, DS820: 4 hours

**Allowed resources:** Must not require the internet.

**Tips:**

- Start by getting an overview of the exam.
- Use paper and pen while taking the exam.
- Justify each answer to yourself (it may help to even *write* down reason).
- Use old exams to practice.

**Q&A session:**

The End of MM537 and DM547!

**Definition:** A branch of Mathematics devoted to the study of integers and their relations (such as divisibility).

**Applications:**

- Cryptology
- Hasing
- Pseudorandom numbers
- many more!

**Beware:** This topic may seem easy at first sight, but it is really one of the harder ones!

# Divisibility

## Definition (Definition 4.1.1)

For  $a, b \in \mathbb{Z}$  with  $a \neq 0$ , we say that  $a$  *divides*  $b$  ( $a$  går op i  $b$ ) if there exists  $c \in \mathbb{Z}$  such that  $ac = b$ . Then we write  $a \mid b$  (and otherwise  $a \nmid b$ ).

We call  $a$  a *factor* (faktor) or *divisor* of  $b$ , and we call  $b$  a *multiple* (multiplum) of  $a$ .

## Theorem (Theorem 4.1.1)

Let  $a, b \in \mathbb{Z}$  with  $a \neq 0$ . Then:

- (i) If  $a \mid b$  and  $a \mid c$  for some  $c \in \mathbb{Z}$ , then  $a \mid (b + c)$ .
- (ii) If  $a \mid b$ , then  $a \mid bc$  for all  $c \in \mathbb{Z}$ .
- (iii) If  $a \mid b$  and  $b \mid c$  for some  $c \in \mathbb{Z}$ , then  $a \mid c$ .

## Corollary (Corollary 4.1.1)

Let  $a, b \in \mathbb{Z}$  with  $a \neq 0$ . Then, if  $a \mid b$  and  $a \mid c$ , then  $a \mid (kb + \ell c)$  for all  $k, \ell \in \mathbb{Z}$ .

# Quotient and Remainder

## Theorem (no proof, Theorem 4.1.3)

Let  $a \in \mathbb{Z}$  and  $d \in \mathbb{Z}^+$ . Then there exist precisely one pair  $q \in \mathbb{Z}$ ,  $r \in \{0, \dots, d-1\}$  such that

$$a = dq + r.$$

## Definition (Definition 4.1.2)

In the theorem above, we call  $d$  the *divisor* (divisor),  $a$  the *dividend* (dividend),  $q$  the *quotient* (quotient), and  $r$  the *remainder* (rest).

We also write

$$a \operatorname{div} d = q \quad \text{and} \quad a \operatorname{mod} d = r,$$

where we say “modulo” for “mod”.

## Definition (Definition 4.1.3)

Let  $a, b \in \mathbb{Z}$  and  $m \in \mathbb{Z}^+$ . Then we have the *congruence* (kongruens)

$$a \equiv b \pmod{m}$$

if and only if  $m$  divides  $a - b$ . We also say that  $a$  and  $b$  are *congruent* (kongruente) modulo  $m$ .

## Theorem (only proof sketch, Theorems 4.1.3 and 4.1.4)

Let  $a, b \in \mathbb{Z}$  and  $m \in \mathbb{Z}^+$ . Then the following statements are equivalent:

- (i)  $a \equiv b \pmod{m}$
- (ii)  $a \bmod m = b \bmod m$
- (iii) There exists  $k \in \mathbb{Z}$  with  $a = b + km$ .

# Adding and Multiplying Congruences

## Theorem (Theorem 4.5.1)

Let  $a, b, c, d \in \mathbb{Z}$  and  $m \in \mathbb{Z}^+$ . If  $a \equiv b \pmod{m}$  and  $c \equiv d \pmod{m}$ , then

$$a + c \equiv b + d \pmod{m} \quad \text{and} \quad a \cdot c \equiv b \cdot d \pmod{m}.$$

### Remark:

- In particular, that means that we can add the same number to both sides of a congruence or multiply them with the same number.
- Question to think about until next lecture: Does the same work for subtraction and (assuming  $c \mid a$  and  $d \mid b$ ) division?