

# Complexity

---

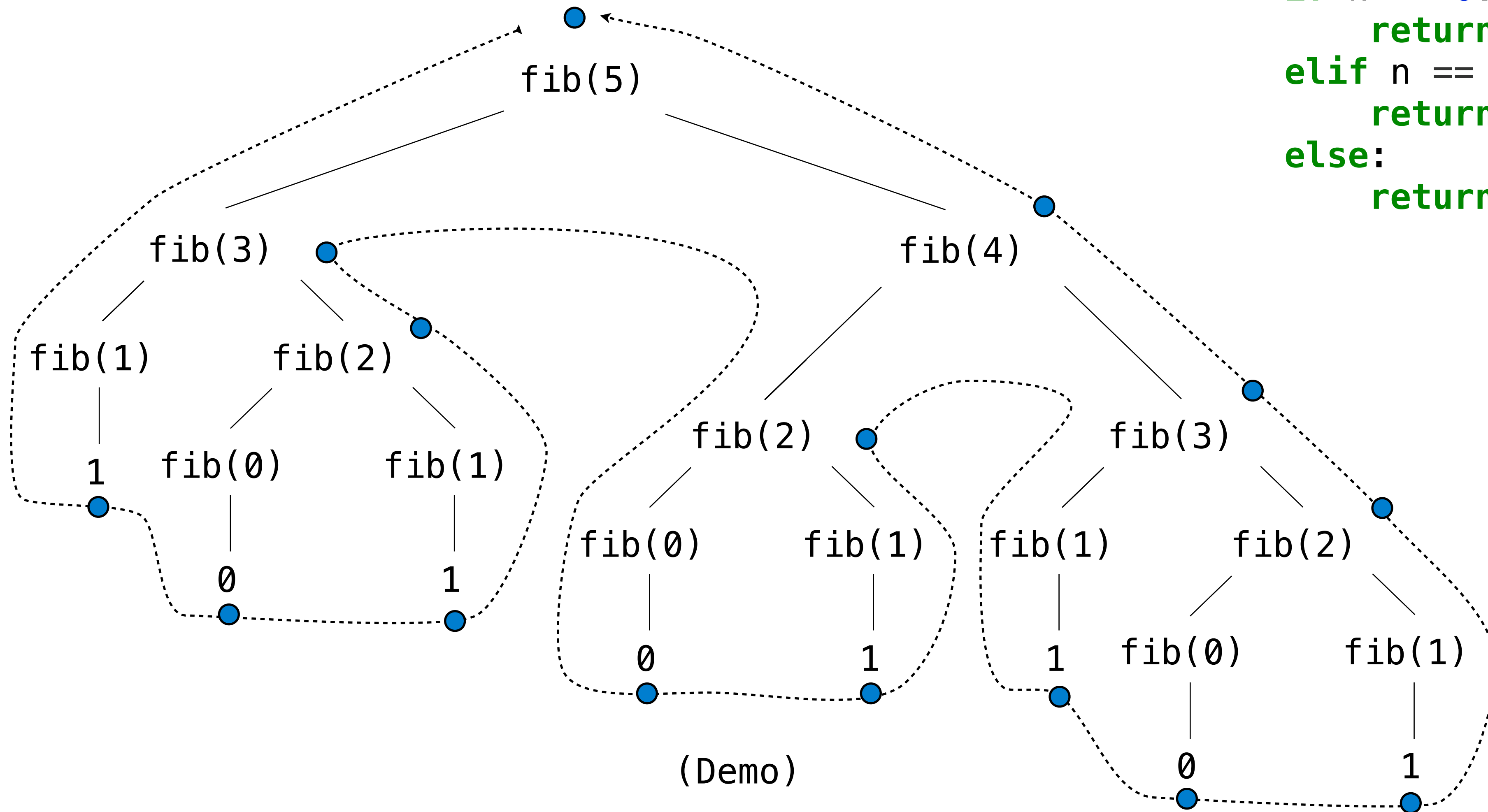
# Announcements

# Measuring Efficiency

# Recursive Computation of the Fibonacci Sequence

Our first example of tree recursion:

```
def fib(n):  
    if n == 0:  
        return 0  
    elif n == 1:  
        return 1  
    else:  
        return fib(n-2) + fib(n-1)
```



# Memoization

# Memoization

---

**Idea:** Remember the results that have been computed before

```
def memo(f):
```

```
    cache = {}
```

Keys are arguments that  
map to return values

```
    def memoized(n):
```

```
        if n not in cache:
```

```
            cache[n] = f(n)
```

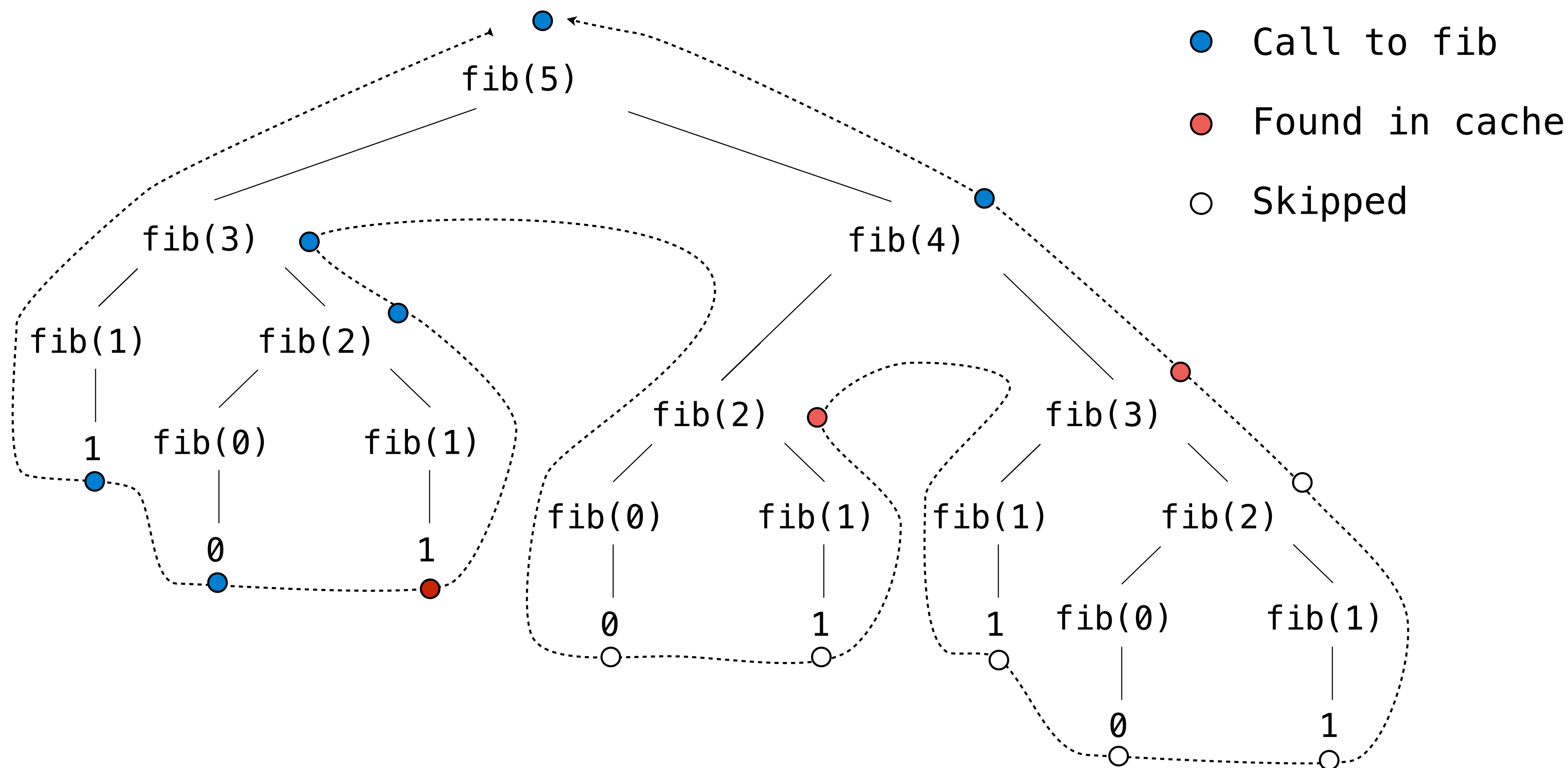
```
        return cache[n]
```

```
    return memoized
```

Same behavior as f,  
if f is a pure function

(Demo)

# Memoized Tree Recursion



# Exponentiation



# Exponentiation

---

**Goal:** one more multiplication lets us double the problem size

```
def exp(b, n):  
    if n == 0:  
        return 1  
    else:  
        return b * exp(b, n-1)
```

```
def exp_fast(b, n):  
    if n == 0:  
        return 1  
    elif n % 2 == 0:  
        return square(exp_fast(b, n//2))  
    else:  
        return b * exp_fast(b, n-1)
```

```
def square(x):  
    return x * x
```

$$b^n = \begin{cases} 1 & \text{if } n = 0 \\ b \cdot b^{n-1} & \text{otherwise} \end{cases}$$

$$b^n = \begin{cases} 1 & \text{if } n = 0 \\ (b^{\frac{1}{2}n})^2 & \text{if } n \text{ is even} \\ b \cdot b^{n-1} & \text{if } n \text{ is odd} \end{cases}$$

(Demo)

# Exponentiation

---

**Goal:** one more multiplication lets us double the problem size

```
def exp(b, n):  
    if n == 0:  
        return 1  
    else:  
        return b * exp(b, n-1)
```

```
def exp_fast(b, n):  
    if n == 0:  
        return 1  
    elif n % 2 == 0:  
        return square(exp_fast(b, n//2))  
    else:  
        return b * exp_fast(b, n-1)
```

```
def square(x):  
    return x * x
```

Linear time:

- Doubling the input **doubles** the time
- 1024x the input takes 1024x as much time

Logarithmic time:

- Doubling the input **increases** the time by a constant C
- 1024x the input increases the time by only 10 times C

# Orders of Growth

## Quadratic Time

---

Functions that process all pairs of values in a sequence of length  $n$  take quadratic time

```
def overlap(a, b):  
    count = 0  
    for item in a:  
        for other in b:  
            if item == other:  
                count += 1  
    return count  
  
overlap([3, 5, 7, 6], [4, 5, 6, 5])
```

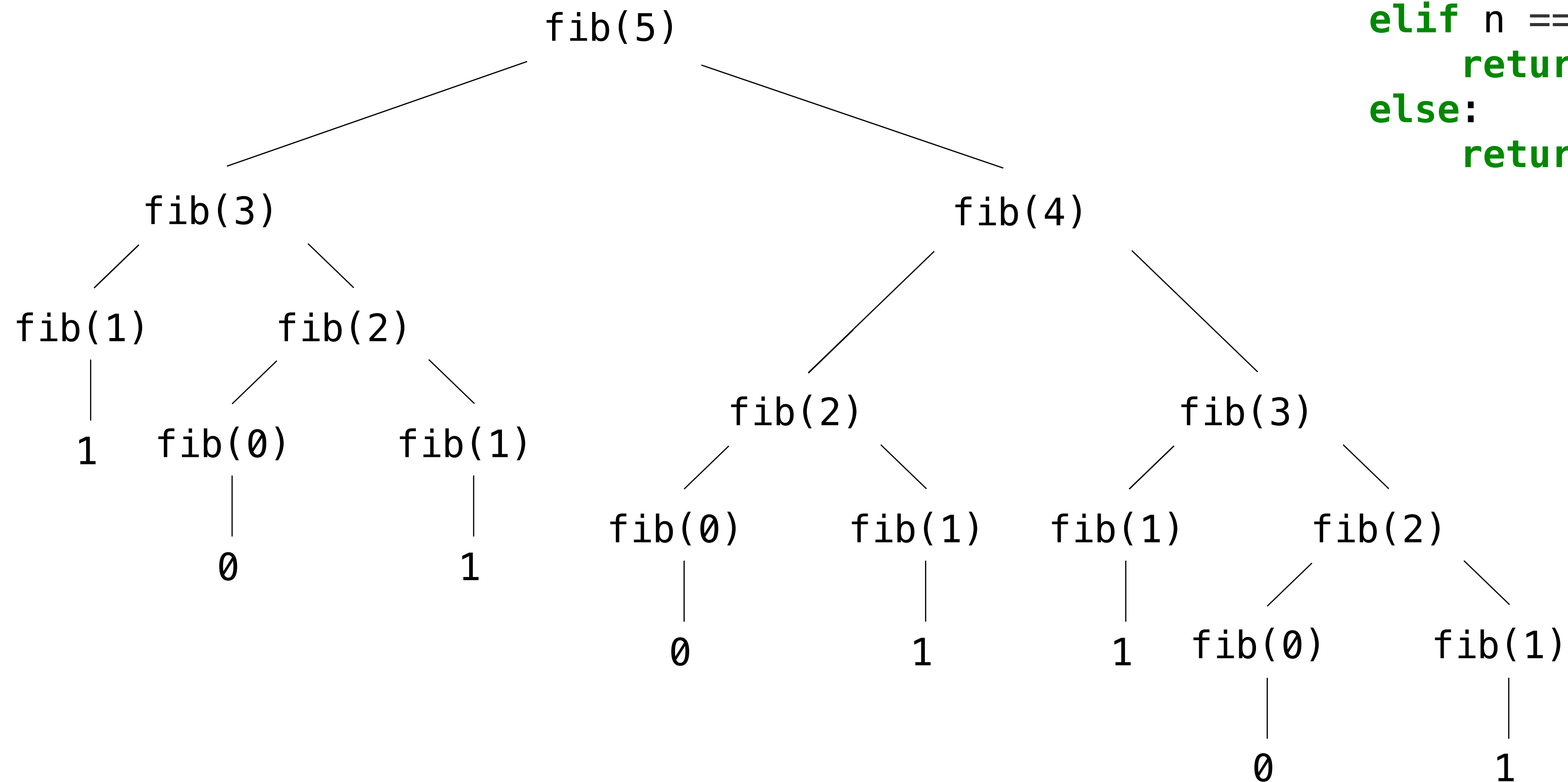
	3	5	7	6
4	0	0	0	0
5	0	1	0	0
6	0	0	0	1
5	0	1	0	0

(Demo)

# Exponential Time

Tree-recursive functions can take exponential time

```
def fib(n):  
    if n == 0:  
        return 0  
    elif n == 1:  
        return 1  
    else:  
        return fib(n-2) + fib(n-1)
```



## Common Orders of Growth

Time for  $n+n$

Time for input  $n+1$

Time for input  $n$

**Exponential growth.** E.g., recursive `fib`

Incrementing  $n$  multiplies *time* by a constant

$$a \cdot b^{n+1} = (a \cdot b^n) \cdot b$$

**Quadratic growth.** E.g., `overlap`

Incrementing  $n$  increases *time* by  $n$  times a constant

$$a \cdot (n+1)^2 = (a \cdot n^2) + a \cdot (2n+1)$$

**Linear growth.** E.g., slow `exp`

Incrementing  $n$  increases *time* by a constant

$$a \cdot (n+1) = (a \cdot n) + a$$

**Logarithmic growth.** E.g., `exp_fast`

Doubling  $n$  only increments *time* by a constant

$$a \cdot \ln(2 \cdot n) = (a \cdot \ln n) + a \cdot \ln 2$$

**Constant growth.** Increasing  $n$  doesn't affect time

# Order of Growth Notation

## Big Theta and Big O Notation for Orders of Growth

---

**Exponential growth.** E.g., recursive `fib`

Incrementing  $n$  multiplies *time* by a constant

$\Theta(b^n)$

$O(b^n)$

**Quadratic growth.** E.g., `overlap`

Incrementing  $n$  increases *time* by  $n$  times a constant

$\Theta(n^2)$

$O(n^2)$

**Linear growth.** E.g., slow `exp`

Incrementing  $n$  increases *time* by a constant

$\Theta(n)$

$O(n)$

**Logarithmic growth.** E.g., `exp_fast`

Doubling  $n$  only increments *time* by a constant

$\Theta(\log n)$

$O(\log n)$

**Constant growth.** Increasing  $n$  doesn't affect time

$\Theta(1)$

$O(1)$



Space

# Space and Environments

---

Which environment frames do we need to keep during evaluation?

At any moment there is a set of active environments

Values and frames in active environments consume memory

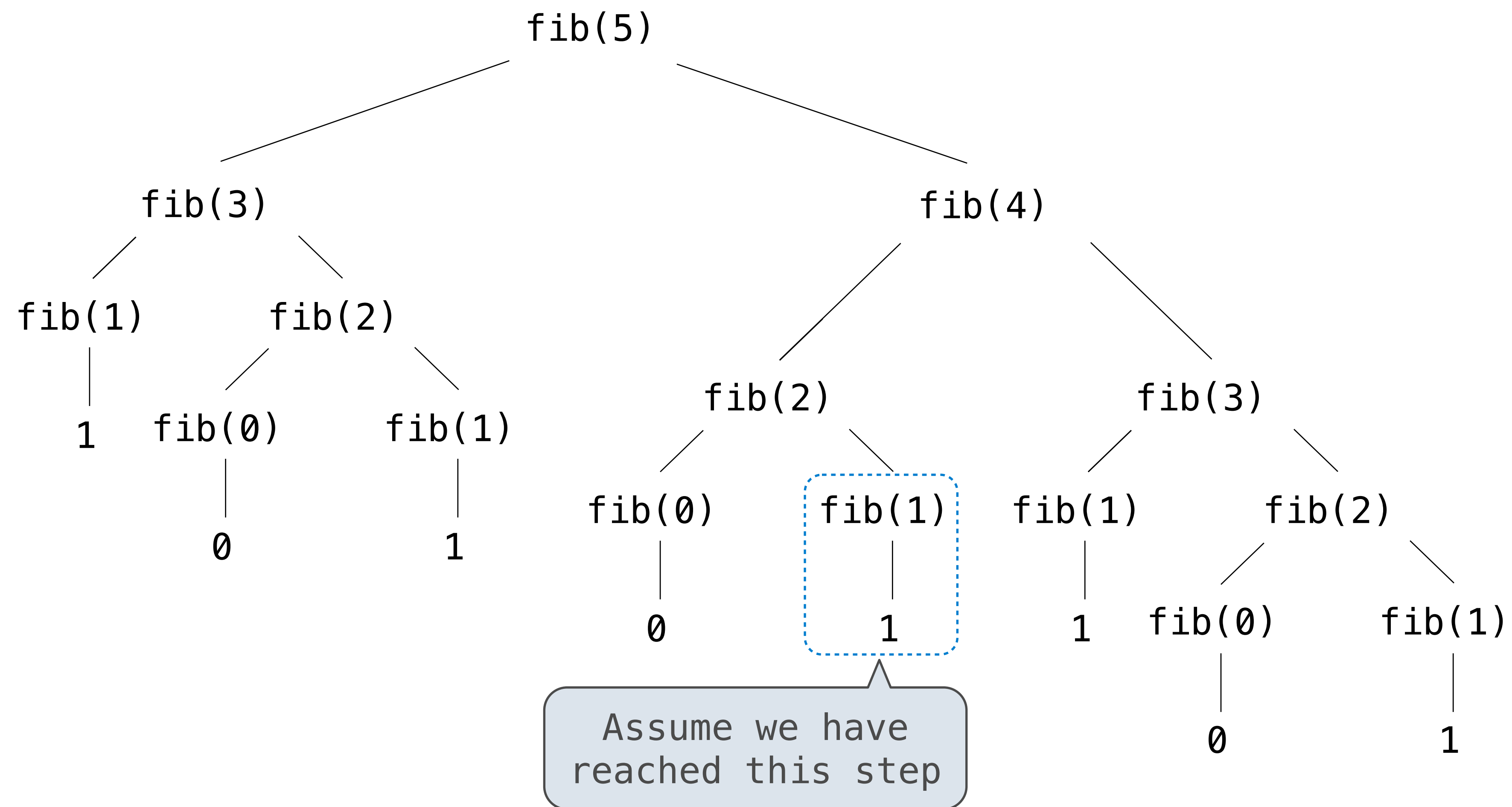
Memory that is used for other values and frames can be recycled

## **Active environments:**

- Environments for any function calls currently being evaluated
- Parent environments of functions named in active environments

(Demo)

# Fibonacci Space Consumption



# Fibonacci Space Consumption

