REVERSE ENGINEERING

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- Aims at hardening the process of reverse engineering
 - Increases level of experience required
 - Increases cost (time, money)
 - Imposes the need for specific tools, techniques and procedures
- Applications (some):
 - License protected software: to prevent the generation of arbitrary licenses or subversion of the program code
 - Proprietary software: prevent the recovery of a design pattern or algorithm (IP protection)
 - Malware: to prevent recovery of the actions, prevent detection, Social Engineer users

Static vs Dynamic

- Static obfuscation transforms code before execution
 - Maybe before compilation, or during compilation
 - Countering static analysis
 - An obfuscated program is complex to analyze but it's always the same

- Dynamic obfuscation transforms code during execution
 - Countering Dynamic Analysis
 - The obfuscated program may change it's behavior

Main Categories (Balachandran, TIFS 2013)

- Layout Obfuscation
- Design Obfuscation
- Data Obfuscation
- Control Obfuscation

Also: Content Type Obfuscation

- Dissimulate one file type as another file type or as raw data
 - Exploring how the file is processed
 - Exploring how users interact with it
 - Exploring how researchers and automatic tools process a file
- Purposes (some):
 - Marketing, branding and usability
 - Exploit users through social engineering
 - Increase the cost required for a reverse engineering task
 - Carry a malicious payload while escaping manual analysis
 - Carry a malicious payload bypassing automatic filtering

Marketing, Branding and Usability

- Aims to make a filetype more usable, or to make the brand present to the user
 - Benning and common usage
- Approach: file has one specific type, but uses another file extension
 - Environment has a configuration stating how to handle such file extension
 - Explores the fact that an Environment uses fixed string to know how to open file
- Impact: File explorers will present a content based on the file extension, not based on the content

Marketing, Branding and Usability

- For a PPTX file
 - File reports a zip file and magic is PK
 - DOCX and XLSX are similar

```
$ unzip -1 8\ -\ Obfuscation.pptx
Archive: 8 - Obfuscation.pptx
 Length
              Date
                      Time
                              Name
    5179 1980-01-01 00:00
                              ppt/presentation.xml
          1980-01-01 00:00
                              customXml/item1.xml
          1980-01-01 00:00
                              customXml/itemProps1.xml
          1980-01-01 00:00
                              customXml/item2.xml
          1980-01-01 00:00
                              customXml/itemProps2.xml
          1980-01-01 00:00
                              customXml/item3.xml
          1980-01-01 00:00
                              customXml/itemProps3.xml
          1980-01-01 00:00
                              ppt/slideMasters/slideMaster1.xml
          1980-01-01 00:00
                              ppt/slides/slide1.xml
                              ppt/slides/slide2.xml
          1980-01-01 00:00
          1980-01-01 00:00
                              ppt/slides/slide3.xml
          1980-01-01 00:00
                              ppt/slides/slide4.xml
    4719 1980-01-01 00:00
                              ppt/slides/slide5.xml
```

Explore users through social engineering

- Aims to confuse users about the purpose of a file
 - Malicious and common in phishing campaigns and malware
- Approach: file has a filename and presentation that confuses users
 - Mail client or explorer presents a safe file with known extension
 - But... icon is stored in the file metadata, and file has two extensions (file.txt.exe)
- Impact: User thinks that a file is not malicious (e.g, it's a word document), while in reality, it executes a malicious code

Explore users through social engineering







2018-19.xlsx

- Windows hides extension of known file types
 - Sample.pptx becomes only Sample
- Executable files may have an embedded icon
 - Freely defined by the developer
 - Explorer will show that icon
- A file named Sample.pptx.exe will be shown as Sample.pptx
 - Users recognize the extension and may think the file is safe
- In a RE task, a file may have bogus extensions

Increase the cost required for a reverse engineering task

 Aims to disguise/manipulate files so that a RE task skips the file, or processes the file incorrectly

Approaches:

- Hides content in file without extension, without headers or with modified headers
- Mangles content to make it less human friendly
- Polyglots
- Impact: Reversing or Forensics Analyst will not process the file, or will not process the file with the correct approach/tools
 - May prevent the researcher from recovering the original file

Magic Headers

- Besides extensions, most files can be recognized by a magic value in the file start/end
 - Manipulating headers can lead to incorrect detection and maybe processing
- Some magic values:
 - Office Documents: DØ CF 11 EØ
 - ELF: 7F E L F
 - JPG: FF D8
 - PNG: 89 P N G 0D 0A 1A 0A
 - Java CLASS: CA FE BA BE

Magic Headers

 Headers are important to maintain compatibility with third party software

- Headers may be irrelevant for custom software
 - Software has the filetype hard coded

Magic Headers

PyInstaller allows converting Python code to an executable

68 6C 69 62 72 02 00 00 00

It actually packs the pyc files into a containerContainer is extracted on runtime and compiled python code is executed

l)...socket.

hashlibr....

Headers are omitted from pyc files. If header is added, extracted file executes as a standard pyc file 0123456789ABC 0123456789ABC 00000000 00000000 0000000D 0000001A d.1.m.Z...e. Z.d.d.l.m.Z. 00000082 .Z.e.e...d.k Reconstructed e...d...Z.e. Extracted ...Z.e.e.d.d 0000012B .md5) . . . check)...md5)...ch .z.0.0.0.0iQ. 00000145 eck output.. 00000152 ..).z.0.0.0.0s4v3 th3 0000016C w0rlds....Inv 00000179 alid..littles h3 w0rlds... Invalid..litt

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hell) ... socke

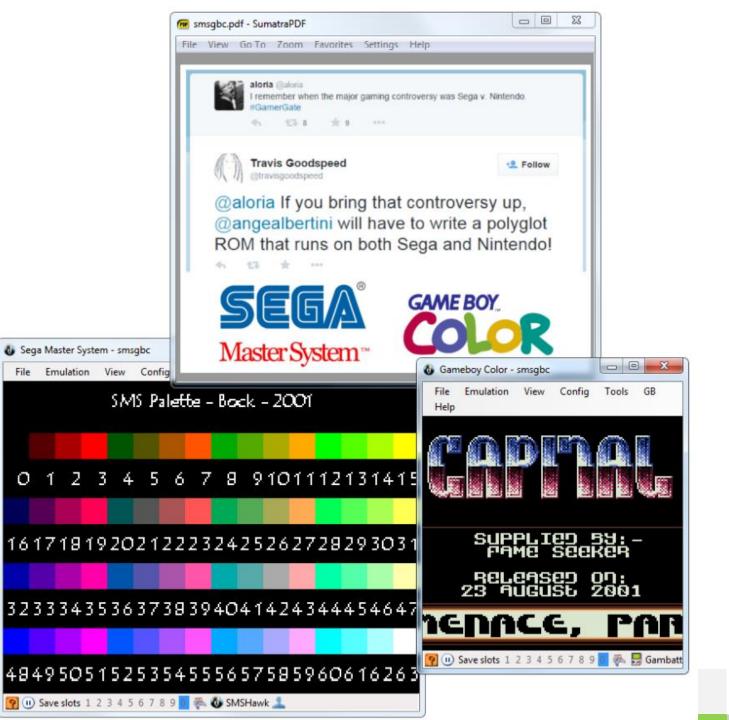
Added header

Polyglots

A file that has different types simultaneously, which may bypass filters and avoid security counter-measures.

pocorgtfo19.pdf (alchemistowl.org)

Technical Note: This file, pocorgtfo19.pdf, is valid as a PDF document, a ZIP archive, and a HTML page. It is also available as a Windows PE executable, a PNG image and an MP4 video, all of which have the same MD5 as this PDF



Types

- Simple Polyglot file: file <u>has</u> different types, accessed depending on how it is handled
- **Schizophrenic file**: is one that <u>is interpreted</u> differently depending on the parser. One parser may crash or fail to process it, while other may return a valid file.
- Chimera file: file has some data that is interpreted as different types

Use in Malware

https://nvd.nist.gov/vuln/detail/CVE-2009-1862

...allows remote attackers to **execute arbitrary code** or cause a **denial of service** (memory corruption) via (1) a **crafted Flash application in a .pdf file** or (2) a crafted .swf file, related to authplay.dll, as exploited in the wild in July 2009.

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Strategies

- Stacks: Data is appended to the file
- Cavities: Uses blank (non used space) in the file
- Parasites: Uses comments or metadata fields that allow content to be written
- Zippers: mutual comments

Empty Space

- Files sometimes allow empty or unused space
 - Before, in the middle or after actual content (appended)
 - Most common in Block formats (ISO and ROM dumps, TAR archives)
 - NAND dumps, ROM dumps, ISOs are directly mapped to sectors
 - Some formats allow arbitrary bytes before file start (e.g. PDF)
 - PDFs are processed from the end
- "Empty space" can be abused to inject crafted content

bash-pdf.pdf

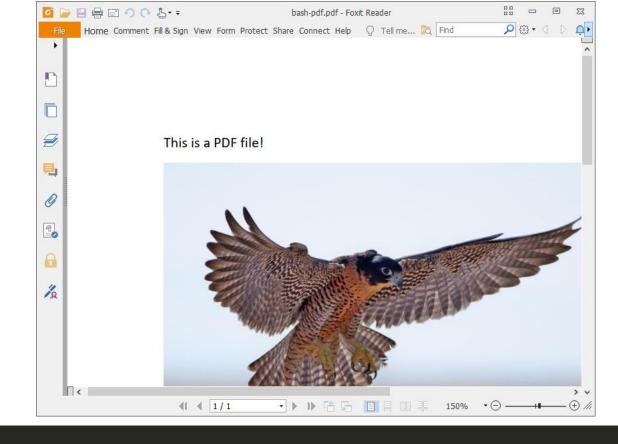
#!/bin/bash

%PDF-1.7

stream

echo "Hello World"; exit

```
$ file bash-pdf.pdf
bash-pdf.pdf: POSIX shell script executable (binary data)
$ ./bash-pdf.pdf
Hello World
```



```
4 %μμμμ
5 1 0 obj
6 <//Type/Catalog/Pages 2 0 R/Lang(en-US) /StructTreeRoot 11 0 R/MarkInfo<</Marked true>>/Metadata 23 0 R/ViewerPreferences 24 0 R>>
7 endobj
8 2 0 obj
9 </Type/Pages/Count 1/Kids[ 3 0 R] >>
10 endobj
11 3 0 obj
12 </Type/Page/Page/Parent 2 0 R/Resources<</Font<</Fi>15 0 R>>/ExtGState<</GS7 7 0 R/GS8 8 0 R>>/XObject<</Image9 9 0 R>>/ProcSet[/PDF/Text/ImageB/ImageC/ImageI]
>>/MediaBox[ 0 0 612 792] /Contents 4 0 R/Group<</Type/Group/S/Transparency/CS/DeviceRGB>>/Tabs/S/StructParents 0>>
13 endobj
14 4 0 obj
15 <//Filter/FlateDecode/Length 245>>
```

Why?

- PDF is a collection of objects
 - Objects are dictionaries of properties with a named type
 - Called "CosObjects" or Carousel Object System
 - Simply added to file. New revisions will create new objects that are appended
 - A PDF can have unused object
 - Objects can contain executable code (the code is not executed by the pdf reader!)
 - Objects can contain anything!
 - Well.... There is the LAUNCH action, and Javascript is a valid object type...

A simple object

```
1 0 obj
<</length 100>>
stream
...100 bytes...
endstream
endobj
```

Two objects

```
1 0 obj
<</le>
stream
...100 bytes...
endstream
Endobj
2 0 obj
<</le>
stream
...100 bytes...
endstream
endobj
```

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Two objects and something else that is not parsed

```
1 0 obj
<</le>
stream
...100 bytes...
endstream
Endobj
I should not be here, but who cares. And I could be anywhere
2 0 obj
<</le>
stream
..100 bytes...
endstream
endobj
```

The XREF Table

- At the end of each PDF there is table with the offset of every object
 - Reader skips to the end of the file, reads the table and parses the objects
 - That's one reason why it ignores garbage between objects

- XREF table also defines where the file magic (%PDF-1.5\n\n) is
 - There may be some bytes before the magic
 - Actually, 1024 random bytes are allowed

```
xref
0 26
0000000011_65535 f
0000000017 00000 n
0000000166 90000 n
0000000222 00000 n
0000000511 00000 n
0000000830 00000 n
0000000998 00000 n
0000001237 00000 n
0000001290 00000 n
0000001343 00000 n
0000055720 00000 n
0000000012 65535 f
0000000013 65535 f
0000000014 65535 f
0000000015 65535 f
0000000016 65535 f
0000000017 65535 f
0000000018 65535 f
0000000019 65535 f
00000000020 65535 f
0000000000 65535 f
0000056466 00000 n
0000056683 00000 n
0000083140 00000 n
0000086318 00000 n
0000086363 00000 n
trailer
<</Size 26/Root 1 0 R/Info 10 0 R/ID[<85F88F67066D2E4AAB78E636585E887B><85F88F67066D2E4AAB78E636585E887B>] >>
```

Offsets of object locations

```
<</Size 26/Root 1 0 R/Info 10 0 R/ID[<85F88F67066D2E4AAB78E636585E887B><85F88F67066D2E4AAB78E636585E887B>] /Prev 86664/XRefStm 86363>>
```

startxref

startxref

86664 %%EOF

xref 0 0 trailer

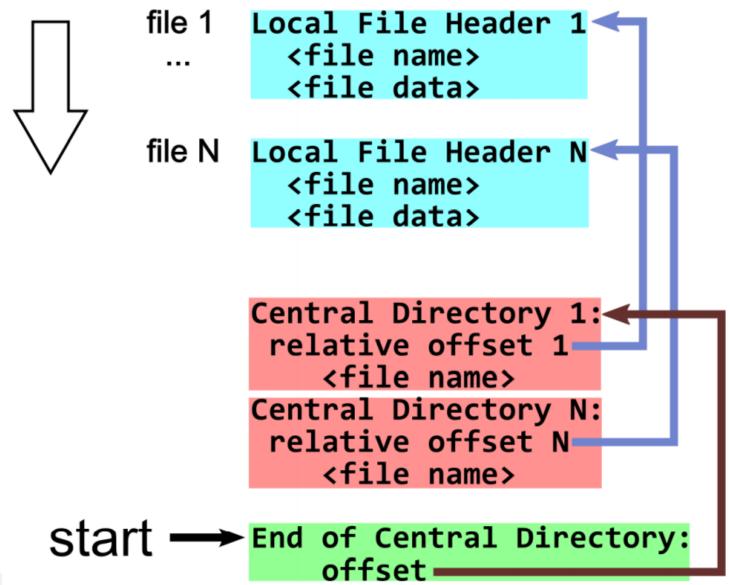
87341

32

34

%%EOF





JPG

JAR (ZIP + CLASS)



AES_{K1}

PNG





FLV

AES,

Practical application

- Malware makes use of polyglots are means to circumvent filters
 - A Packet/Email/Web application firewall will block executables, bit will it block JPGs?
 - If it does, can he do it with a low rate of false positives?

- General process involves download a polyglot and a decoder
 - Polyglot contains malicious code
 - Decode is implemented in a less suspicious manner (e.g. Javascript)
- From a Reversing Perspective: how much effort will we spend analyzing a JPG?
 - Automated tools such as binwalk, TrId and file can help (but are limited)

- Practical application Stegsploit https://stegosploit.info/
- Creates a new way to encode "drive-by" browser exploits and deliver them through image files.
- These payloads are undetectable using current means
- Drive-by browser exploits are steganographically encoded into JPG and PNG images.
- The resultant image file is fused with HTML and Javascript decoder code, turning it into an HTML+Image polyglot.
 - The polyglot looks and feels like an image, but is decoded and triggered in a victim's browser when loaded.



Exercise

- Check https://github.com/corkami/pocs/tree/master/mini repository for some Polyglots
 - Can you detect most types? What safe methodology can be devised?
 - https://github.com/corkami/formats can help with some technical aspects of the file formats
- Use Mitra (https://github.com/corkami/mitra) and create some polyglots with your files.
 - Analyze what was produced

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Exercise

Can you craft your own polyglot?

- Because ZIP is a format ready to include several other files, and PDF is highly flexible, It is easy to build a polyglot
 - Consider the "PDF Objects" all text before the xref and PDF trailer all text after the xref

ZIP Entry Header

Stored PDF Objects

ZIP Central Directory Structure

PDF Trailer (as ZIP comments)

zip -Z store -z polyglot.pdf pdf_objects < pdf_trailer</pre>

Use a hexeditor and check the result

Code Obfuscation

Layout Obfuscation

- Aims at hiding how the <u>source code</u> is structured
 - As source code (or symbols) can present enough information to help reversing a program
- Applied to the source code, and focused on situations where source can be obtained
 - Javascript, HTML, CSS, Java
- Methods:
 - Deleting comments
 - Remove debugging information
 - Renaming classes, methods and variables
 - Removing spaces
 - Stripping a binary

```
#\
                                define C(c
                                                      /**/)#c
                               /*size=3173*/#include<stdio.h>
                            /*crc=b7f9ecff.*/#include<stdlib.h>
                           /*Mile/Adele von Ascham*/#include<time.h>
                           typedef/**/int(I);I/*:3*/d,i,j,a,b,l,u[16],v
                           [18], w[36], x, y, z, k; char*P="\n\40(),",*p,*q,*t[18], m[4];
                          void/**/0(char*q){for(;*q;q++)*q>32?z=111-*q?z=(z+*q)%185,(k?)}
                          k--:(y=z\%37,(x=z/37\%7)?printf(*t,t[x],y?w[y-1]:95):y>14&&y<33?x
                          =y>15, printf(t[15+x], x?2<<y%16:1, x?(1<<y%16)-1:1):puts(t[y%28])))
                          0:z+82:0;void/**/Q(I(p),I*q){for(x=0;x<p;x++)}{q[x]=x;}{for(;--p)}
                         =q[x=rand()\%-p],q[x]=q[p];}char/**/n[999]=C(Average?!nQVQd%R>Rd%
   >1;q[p]=y)y
 R%
              %RNIPRfi#VQ}R;TtuodtsRUd%RUd%RUOSetirwf!RnruterR{RTSniamRtniQ>h.oidts<edulc
ni
                      #V>rebmun<=NIPD-RhtiwRelipmocResaelPRrorre#QNIPRfednfi#V__ELIF__R_
                 nifed#V~-VU0V; }V{R= R][ORrahcRdengisnuRtsnocRcitatsVesle#Vfidne#V53556
Re
             .1RfoRegnarRehtRniRre getniRnaRsiR]NIP[R erehwQQc.tuptuoR>Rtxt.tupniR
           R]NIP[R:egasuV Redulcn i#VfednfiVfednuVenife dVfedfiVQc%Rs%#V);I/**/main(
          f), char**e){if(f){for(i=
                                      time(NULL), p=n,q=n+998, x=18; x; p++){*p>32&&!(}
 Ι(
         *--q=*p>80&&*p<87?P[*p-
                                   81]:*
                                             p)?t [(--x)]=q+1:q;if(f-2||(d=atoi
        (e[1]))<1||65536<d){;0("
                                                      goto 0;}srand(i);Q(16,u);i=0;Q(
       36,w);for(;i<36; i++){w[i]
                                                  [i]<26 ? 97:39; }0(C(ouoo9oBotoo%]#
                                    +=w
      ox^#oy_#ozoou#o{ a#o b#o}c#
                                                  o~d#oo-e
                                                             #00. f#00/g#000h#001i#00
    2j#oo3k#oo4l#o p));for(j
                                                   =8;EOF
                                                                   getchar()); l+=1){a=1+
   rand()%16;for(b =0;b<a||i-
                                                                  (0,e);b++)x=d^d/4^d/8^d/
    32,d= (d/ 2|x<<15)&65535;
                                                          b = 1 < 17; 0(18, v); for(a=0; a<18;
            ){if( (b&(1<<(i=v[a]
                                      )))))*
                                                             m=75+i,0(m),j=i<17&&j<i?i:j;}0(C(
    a++
    !)
                                      /n!));i=
                 ); }0(C(oqovoo970
                                                          0; for(; i < 8; 0(m))m[2] = 35, *m = 56 + u[i], m[1]
    1=
                    +i++;0(C(oA!oro
                                       ogoo9)
                                                        );k=112-j*7;0(C(6o.!Z!Z#5o-!Y!Y#4~!X!X#3}
                                                  T!T#/y!S!S#.x!R!R#-w!Q!Q#ooAv!P!P#+o#!0!0#*t!N!
     !W
        !W
                #2
                      |!V!V#1{!U!U#0z!
                        >s!M!M#oo=r!L!L#oo<q!K!K#
                                                     &pIo@:;= oUm#oo98m##oo9=8m#oo9oUm###oo9;=8m#o
      N#
               00
                    oUm##oo9=oUm#oo98m####
                                                             #o1:^#o2; #o3<o ou#o4=a#o5>b#o6?c#o
               о9
                                                     009]
             7@d#o8A e#o
                            9B
                                  f#o:Cg#o;
                                                     D
                                                            h#o<Ei #o=Fj#o> Gk#o?Hl#oo9os#####
           ));d=0
                                                            ;}
                                                                              for(x=y=0;x<8;++
          x)y =
                                                                               d&(1<<u[x])?
          1<<
                                                                             x:0;return
            y
```

Code Obfuscation

Design Obfuscation

- Aims at making the design nonobvious, more difficult to recover
 - Usually done by a tool before compilation or during compilation
 - GCC can do this automatically by inlining functions (-O3 -finline)

Methods:

- Merging and splitting methods
- Merging and splitting classes
- Splitting binary code, while inserting dummy instructions
- Splitting loops and conditions, maybe interleaved with dummy code
- Inlining functions
- Dead Code

Code Obfuscation

Code inserted, but never executed. JMP before dummy code effectively only splits code

Design Obfuscation – Breaking Code

```
#include <stdio.h>
    #include <stdlib.h>
    unsigned long long factorial(unsigned long long a) {
        unsigned long long r = 1;
        while(a > 0){
            unsigned long long v = r * a;
            if(v < r)
                printf("ERROR: Overflow\n");
12
                exit(-1);
            r = v:
            a = a - 1;
        return r;
    int main(int argc, char** argv) {
        unsigned long long v = 0;
        if(argc != 2) {
            printf("Need a positive integer argument\n");
            return -1;
        v = atol(argv[1]);
        if(v \leftarrow 0)
            printf("Need a positive integer argument\n");
            return -1;
        printf("Result: %llu\n", factorial(v));
        return 0;
```

```
int main(int argc, char** argv) {
22
         unsigned long long v = 0;
         if(argc != 2) {
23
             printf("Need a positive integer argument\n");
25
             return -1;
27
         asm("jmp label");
         factorial(factorial(argc));
         asm("label:");
29
        v = atol(argv[1]);
31
32
        if(v \leftarrow 0)
33
             printf("Need a positive integer argument\n");
             return -1;
37
         asm("jmp label b");
         factorial(factorial(v * factorial(-v)));
         asm("label b:");
41
         printf("Result: %llu\n", factorial(v));
42
43
         return 0;
44
45
```

return 0;

Code inserted, but never executed.

JMP before dummy code effectively only splits code

Design Obfuscation – Breaking Code

```
#include <stdio.h>
                                                         int main(int argc, char** argv) {
   #include <stdlib.h>
                                                             unsigned long long v = 0;
                                                    22
                                                    23
                                                             if(argc != 2) {
   unsigned long long factorial(unsigned long long a) {
                                                                 printf("Need a positive integer argument\n");
                                                    25
                                                                 return -1;
      unsigned long long r = 1;
      while(a > 0){
                                                             asm("jmp label");
          unsigned long long v = r * a;
          if(v < r){}
                         What about the output binary?
             printf("ERROR:
12
             exit(-1);
          r = v;
          a = a - 1;
                         Compile with gcc -O0 -o factorial-split factorial-split.c
       return r;
   int main(int argc, char**
                         Does it effect static or dynamic analysis?
      unsigned long long v :
      if(argc != 2) {
          printf("Need a pos
                         Check with objdump -d and ghidra
          return -1;
      v = atol(argv[1]);
                         gcc may also inline functions (the opposite) when using -O3
      if(v \leftarrow 0)
          printf("Need a pos
          return -1;
                         or —finline-functions
      printf("Result: %llu\n", factorial(v));
```

Design Obfuscation – Dead Code

- Aims at inserting dummy code (not executed) to confuse the analysis
 - Code may follow some pattern (previous example), or be random
 - Code may lock the analysis tool if recursive disassembly is used
 - Decompilation to Pseudo C will surely be affected

- Dead code can be added after compilation
 - May contain fingerprinting information by making binaries unique

Design Obfuscation – Dead Code

```
unsigned long long factorial(unsigned long long a) {
22
23
        unsigned long long r = 1;
24
        while(a > 0){
            unsigned long long v = r * a;
            if(v < r){
                printf("ERROR: Overflow\n");
                 exit(-1);
             r = v;
            a = a - 1;
            if(v != r) {
                 asm (REP(3,3,3,"nop;"));
         return r;
```

r=v, therefore, if(v!=r) will be always false. Compiler will not easily discard this code.

__asm__... Instruction will insert 333 NOPs (which will not be executed) This is a placeholder that can be used later for post processing

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Exercise

Design Obfuscation – Dead Code

- Consider the example presented in factorial-dead.c
- Create a script that replaces a sequence of NOP (0x90) with another sequence of instructions
- Take care that in X86, instructions are of variable length. The last 15 instructions added to a placeholder should be handled with care, as they can "mask" real instructions after the placeholder
 - For this exercise, leave them with 0x90
- Compile the code and then use objdump and ghidra to reverse the binary. Is this useful? Can it be detected?

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Design Obfuscation – Dead Code

```
2 undefined8 main(int param 1,long param 2)
     undefined8 uVarl:
     long 1Var2;
     if (param 1 == 0x2) {
       1Var2 = atol(*(char **)(param_2 + 0x8));
10
       if (1Var2 == 0x0) {
        puts("Need a positive integer argument");
11
        uVarl = 0xffffffff;
13
       else {
        uVarl = factorial(1Var2);
        printf("Result: %llu\n",uVarl);
         uVar1 = 0x0;
17
18
19
20
      puts("Need a positive integer argument");
       uVarl = 0xfffffffff:
23
     return uVarl:
25
26
```

```
Decompile: main - (factorial-dead-obf)
     undefined4 *local 28;
     int local lc;
     long local 10;
     local 10 = 0x0;
     local 28 = param 2;
     local 1c = param 1;
     if (param 1 == 0x2) {
       puVar4 = *(undefined4 **)(param 2 + 0x2);
       uStack48 = 0x10136a;
       local 10 = atol((char *)puVar4);
       if (local 10 == 0x0) {
        uStack48 = 0x101381;
         puts("Need a positive integer argument");
         pcVar2 = (char *) 0xffffffff;
28
       else {
         if (local lc * local 10 == 0x0) {
           *puVar4 = *param 2;
           if ((POPCOUNT(local lc * local 10 & 0xff) & 0x1U) != 0x0) {
                       /* WARNING: Bad instruction - Truncating control flow here */
             halt baddata();
           puVar4 = (undefined4 *) (ulong) ((int)param 4 - 0x44);
           puVar3 = &uStack48:
           cVar1 = '\x12';
           do {
             puVar4 = puVar4 + -0x1;
             puVar3 = (undefined8 *)((long)puVar3 + -0x4);
             *(undefined4 *)puVar3 = *puVar4;
             cVarl = cVarl + -0x1;
           } while ('\0' < cVarl);</pre>
                       /* WARNING: Bad instruction - Truncating control flow here */
           halt baddata();
47
         uStack48 = 0x10172f;
         factorial(local 10);
         uStack48 = 0x101743;
         printf("Result: %llu\n");
         pcVar2 = (char *) (local_lc * local_10);
         if (pcVar2 + -(local 10 + -0x3) != NULL) {
           pcVar2 = NULL;
```

Data Obfuscation

- Encrypts, or otherwise encodes data contents
 - Contents are decrypted in real time, as the program is executed
 - Static analysis, or fingerprint matching may fail to correctly recover useful information
 - Frequent tactic to evade filters

Why?

- Strings frequently carry semantic information, that may help analysis
- E.g. Str="Please input your AES key": we will know that this a key, and know the algorithm

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Data Obfuscation - how

- Split the string in parts
 - May be combined with two conditions or loops to validate both parts individually
- Erase strings right after use
- Common XOR is frequently found as it requires no dependencies and is fast
 - More recent malware will use RC4 or even AES for this purpose
 - Decryption key can also be encrypted, and some key may be obtained dynamically
 - E.g. from a hardware token as a form of licensing enforcement
- Create a custom encoding based on a complex state machine
 - May use flow information, voiding the decoding of strings if the execution order it changed

Control Obfuscation

- Introduces dummy control structures, with little impact to execution
 - Impact is only from a performance point of view
 - However, analysis tools will interpret the control structures and create complex CFGs
- Makes use of Opaque Predicates: predicates for which the programmer already knows the result.
 - E.g. if (1 > 0)

Control Obfuscation – Opaque Predicates

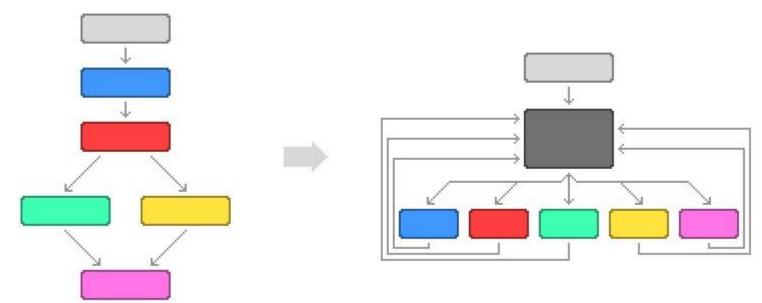
- Introduces dummy control structures, with little impact to execution
 - Impact is only from a performance point of view (additional branch)
 - However, analysis tools will interpret the control structures and create complex CFGs
- Makes use of Opaque Predicates: predicates for which the programmer already knows the result.
 - E.g. if (1 > 0) or v=r; if(v==r)

Control Obfuscation – Opaque Predicates

- Opaque predicates can be more complex
- Manipulate pointers, linked lists, use computation processes
- Result of a predicate can be dynamic, and related to execution state
 - Dynamic analysis may change execution sequence, therefore the predicate result and invalidate the execution
 - Similar to TPMs, where keys are provided at a valid situation
 - Predicate can use dynamic data, received from external services
- Concurrency can be used to create predicates
 - If two threads are executing with some relation, one can update data, that the other uses to construct a
 predicate
 - Timing information can also be used, to further increase the complexity (information not available statically)

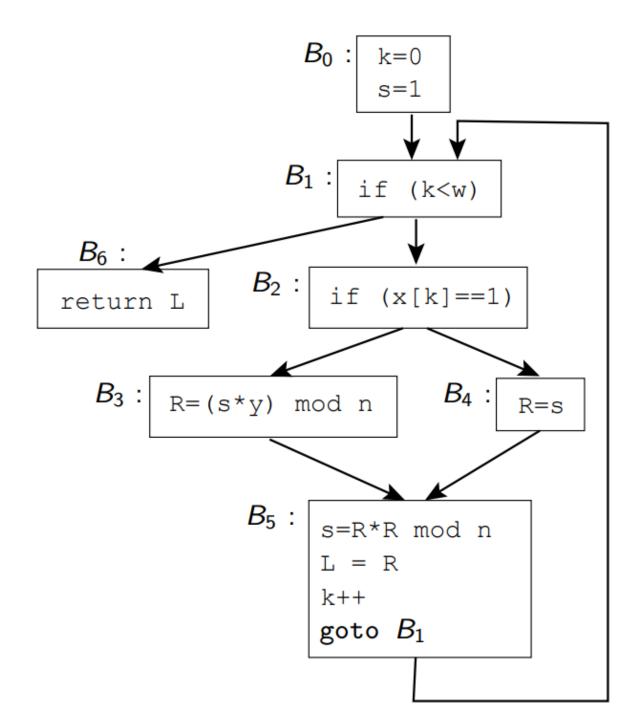
Control Obfuscation – Control Flow Flattening

- Removes control flow structures from program
 - Converts the program to a gigantic Switch, where each condition is a case
 - Program runs on an infinite loop around the switch
- Program becomes ~4 times slower, and 2 times larger



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```
int modexp(int y,int x[],
           int w,int n) {
   int R, L;
   int k = 0;
   int s = 1;
   while (k < w) {
      if (x[k] == 1)
         R = (s*y) \% n;
      else
         R = s;
      s = R*R \% n;
      L = R;
      k++;
   return L;
```



```
int modexp(int y, int x[], int w, int n) {
   int R, L, k, s;
   int next=0;
  for(;;)
      switch(next) {
         case 0 : k=0; s=1; next=1; break;
         case 1 : if (k<w) next=2; else next=6; break;</pre>
         case 2: if (x[k]==1) next=3; else next=4; break;
         case 3 : R=(s*y)%n; next=5; break;
         case 4 : R=s; next=5; break;
         case 5 : s=R*R%n; L=R; k++; next=1; break;
         case 6 : return L;
```

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