

# Lenguajes Formales y de Programación

Práctica1: Trabajo teórico-práctico

Sección: A

Nombre: Registro académico:

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1. Creación de la expresión regular que describa el patrón de cada token.

## 1.1. Identificador:

Expresión regular: ([L])+.([0-9])\*

1.2. Número:

Expresión regular: [0-9]+

1.3. Decimal:

Expresión regular: [0-9]+((.).(0-9)+)\*

1.4. Puntuación:

Expresión regular: ((.)|(,)|(;)|(:))+

1.5. Operador:

Expresión regular: ((+)|(-)|(\*)|(/)|(%))+

1.6. Agrupación:

Expresión regular:  $((()|())|([)|(])|(\{)|(\}))+$ 

2. Gramática regular de cada token.

2.1. Identificador:

Diagrama autómata finito no determinista Método de Thomson: Identificador

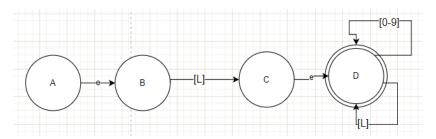


Tabla de transición

FT	е	L	[0-9]
Α	{B}=s0	∂(s0,L)=C	{ (s0, [0-9]) } ={}
С	{D}=s1	∂{ (s1,L) } =D	∂(s2, [0-9]) =D

Optimizar - Tabla de transición

FT	е	L	[0-9]
Α	{B}=s0	∂(s0,L)=s1	{ (s0, [0-9]) } ={}
С	{D}=s1	∂{ (s1,L) } =s1	∂(s2, [0-9]) =s1

Definición formal AFD:  $A=(Q,\Sigma,\partial,A,F)$ 

2.s0

3. 
$$\Sigma = \{[a-z], [A-Z], [0-9]\}$$

$$4.F = {s1}$$

5. Función de transición

$$\partial(S0, [L]) = S1$$
  $\partial(S1, [0-9]) = S1$   $\partial(S1, [L]) = S1$ 

#### 2.2. Número:

Diagrama autómata finito no determinista Método de Thomson: Número

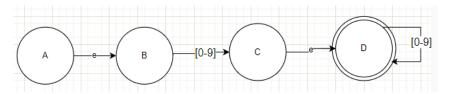


Tabla de transición

FT	е	[0-9]
Α	{B}=s0	$\partial \{ (s0, [0-9]) \} = C$
С	{D}=s1	∂(s1, [0-9]) =D

Optimizar - Tabla de transición

FT	е	[0-9]
Α	{B}=s0	$\partial$ { (s0, [0-9]) } =s1
С	{D}=s1	$\partial$ (s1, [0-9]) =s1

Definición formal AFD:  $A=(Q,\Sigma,\partial,A,F)$ 

2.s0

3. 
$$\Sigma = \{[0-9]\}$$

 $4.F={s1}$ 

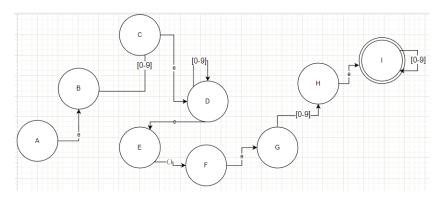
5. Función de transición

$$\partial(S0, [0-9]) = S1$$

$$\partial(S1, [0-9]) = S1$$

#### 2.3. Decimal:

Diagrama autómata finito no determinista Método de Thomson: **Decimal** 



## Tabla de transición

FT	е	[0-9]	
Α	{B}=s0	∂(s0, [0-9])=C	{ (s0, [.]) } ={}
С	{D}=s1	$\partial \{ (s1, [0-9]) \} = D$	{ (s1, [.]) } ={}
D	{E}=s1	$\partial \{ (s1, [0-9]) \} = \{ \}$	∂{ (s1, [.]) } =F
F	{G}=s2	∂{ (s2, [0-9]) } =H	{ (s2, [.]) } ={}
Н	{I}=s3	∂{ (s3, [0-9]) } =H	{ (s3, [.]) } ={}

Optimizar - Tabla de transición

FT	е	[0-9]	•
Α	{B}=s0	∂(s0, [0-9])=s1	{ (s0, [.]) } ={}
С	{D}=s1	$\partial \{ (s1, [0-9]) \} = s1$	{ (s1, [.]) } ={}
D	{E}=s1	$\partial \{ (s1, [0-9]) \} = \{ \}$	$\partial \{ (s1, [.]) \} = s2$
F	{G}=s2	$\partial \{ (s2, [0-9]) \} = s3$	{ (s2, [.]) } ={}
Н	{I}=s3	$\partial \{ (s3, [0-9]) \} = s3$	{ (s3, [.]) } ={}

Definición formal AFD:  $A=(Q,\Sigma,\partial,A,F)$ 

- 1.Q={s0, s1,s2,s3}
- 2.s0
- 3.  $\Sigma = \{[0-9],(.)\}$
- $4.F={s3}$
- 5. Función de transición

$$\partial(S0, [0-9]) = S1$$

$$\partial(S1, (.)) = S2$$

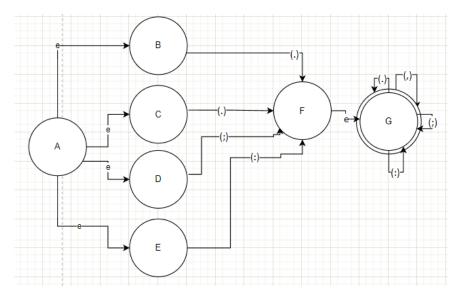
$$\partial(S1, [0-9]) = S1$$

$$\partial(S2, [0-9]) = S3$$

$$\delta(S3, [0-9]) = S3$$

#### 2.4. Puntuación:

Diagrama autómata finito no determinista Método de Thomson: Puntuación



## Tabla de transición

FT	е	•	,	;	:
Α	{A,B,C,D,E}=s0	$\partial \{ (s0, [.]) \} = F$	∂{ (s0, [,]) } =F	∂{ (s0, [;]) } =F	∂{ (s0, [:]) } =F
F	{G}=s1	∂{ (s1, [.]) } =G			∂{ (s1, [:]) } =G

# Optimizar - Tabla de transición

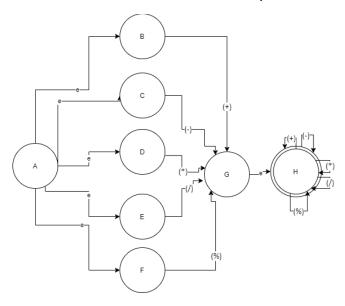
FT	е	•	,	;	:
Α	{A,B,C,D,E}=s0	∂{ (s0, [.]) }	∂{ (s0, [,]) }	∂{ (s0, [;]) }	∂{ (s0, [:]) }
		=s1	=s1	=s1	=s1
F	{G}=s1	∂{ (s1, [.]) }	∂{ (s1, [,]) }	∂{ (s1, [;]) }	∂{ (s1, [:]) }
		=s1	=s1	=s1	=s1

Definición formal AFD:  $A=(Q,\Sigma,\partial,A,F)$ 

- 1.Q={s0, s1}
- 2.s0
- 3.  $\Sigma = \{(.), (,), (;), (:)\}$
- $4.F={s1}$
- 5. Función de transición
- $\partial(SO, (.)) = S1$   $\partial(SO, (,)) = S1$   $\partial(SO, (;)) = S1$   $\partial(SO, (:)) = S1$
- $\delta(\texttt{S1,(.)}) = \texttt{S1} \quad \delta(\texttt{S1,(,)}) = \texttt{S1} \quad \delta(\texttt{S1,(;)}) = \texttt{S1} \quad \delta(\texttt{S1,(:)}) = \texttt{S1}$

# 2.5. Operador:

Diagrama autómata finito no determinista Método de Thomson: **Operador** 



## Tabla de transición

FT	е	+	-	*	/	%
Α	{A,B,C,D,E,F}=s0	∂{ (s0, [+]) } =G		∂{ (s0, [*]) } =G		∂{ (s0, [%]) } =G
G	{H}=s1	∂{ (s1, [+]) } =H		∂{ (s1, [*]) } =H		∂{ (s1, [%]) } =H

## Optimizar - Tabla de transición

FT	е	+	-	*	/	%
Α	{A,B,C,D,E,F}=s0	∂{ (s0, [+]) }	∂{ (s0, [-]) }	∂{ (s0, [*]) }	∂{ (s0, [/]) }	∂{ (s0, [%]) }
		=s1	=s1	=s1	=s1	=s1
G	{H}=s1	∂{ (s1, [+]) }	∂{ (s1, [-]) }	∂{ (s1, [*]) }	∂{ (s1, [/]) }	∂{ (s1, [%]) }
		=s1	=s1	=s1	=s1	=s1

Definición formal AFD:  $A=(Q,\Sigma,\partial,A,F)$ 

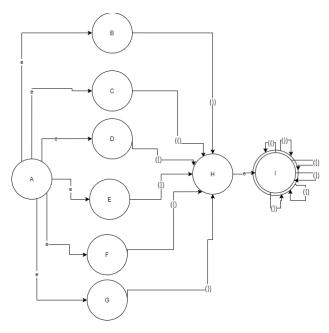
- 1.Q={s0, s1}
- 2.s0
- 3.  $\Sigma = \{(+), (-), (*), (/), (%)\}$
- $4.F={s1}$
- 5. Función de transición

$$\partial(SO, (+)) = S1 \quad \partial(SO, (-)) = S1 \quad \partial(SO, (*)) = S1 \quad \partial(SO, (/)) = S1 \quad \partial(SO, (%)) = S1$$

$$\partial(S1, (+)) = S1 \quad \partial(S1, (-)) = S1 \quad \partial(S1, (*)) = S1 \quad \partial(S1, (/)) = S1 \quad \partial(S1, (%)) = S1$$

# 2.6. Agrupación:

Diagrama autómata finito no determinista Método de Thomson: Agrupación



## Tabla de transición

FT	е	(	)	[	]	{	}
Α	{A,B,C,D,E,F,G}=s0	∂{ (s0, [(]) }	∂{ (s0, [)]) }	∂{ (s0, [ [ ])	∂{ (s0, [ ] ])	∂{ (s0, [{]) }	∂{ (s0, [}]) }
		=H	=H	} =H	} =H	=H	=H
Н	{I}=s1	∂{ (s1, [(]) }	∂{ (s1, [)]) }	∂{ (s1, [ [ ])	∂{ (s1, [ ] ])	∂{ (s1, [{]) }	∂{ (s1, [}]) }
		=I	=I	} =I	} =I	=I	=I

## Optimizar - Tabla de transición

FT	е	(	)	[	]	{	}
Α	{A,B,C,D,E,F,G}=s0	∂{ (s0, [(]) }	∂{ (s0, [)]) }	∂{ (s0, [ [ ])	∂{ (s0, [ ] ])	∂{ (s0, [{]) }	∂{ (s0, [}]) }
		=S1	=S1	} =S1	} =S1	=S1	=S1
Н	{I}=s1	∂{ (s1, [(]) }	∂{ (s1, [)]) }	∂{ (s1, [ [ ])	∂{ (s1, [ ] ])	∂{ (s1, [{]) }	∂{ (s1, [}]) }
		=S1	=S1	} =S1	} =S1	=S1	=S1

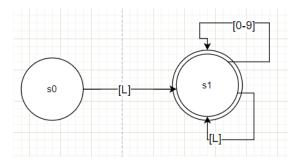
- 1.Q={s0, s1}
- 2.s0
- 3.  $\Sigma = \{((), ()), ([), (]), (\{), (\})\}$
- $4.F={s1}$
- 5. Función de transición

$$\partial(SO, (()) = S1 \quad \partial(SO, ())) = S1 \quad \partial(SO, ([)) = S1 \quad \partial(SO, ([))) = S1 \quad \partial(SO, (\{))) = S1$$

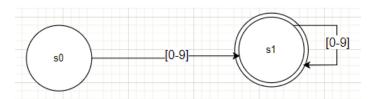
$$\delta(S1, (()\ ) = S1 \quad \delta(S1, ())\ ) = S1 \quad \delta(S1, ([)) = S1 \quad \delta(S1,$$

# 3. AFD de cada token.

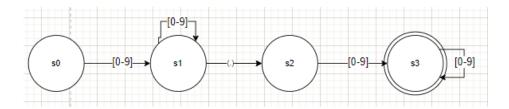
# 3.1. Identificador



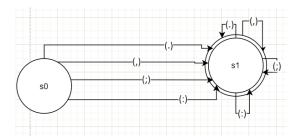
3.2.



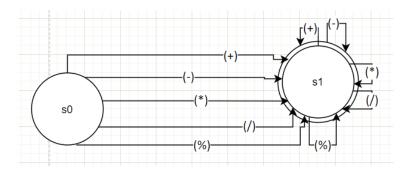
3.3.

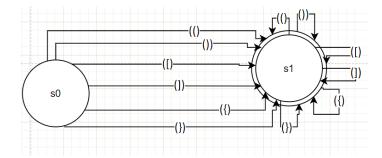


3.4.



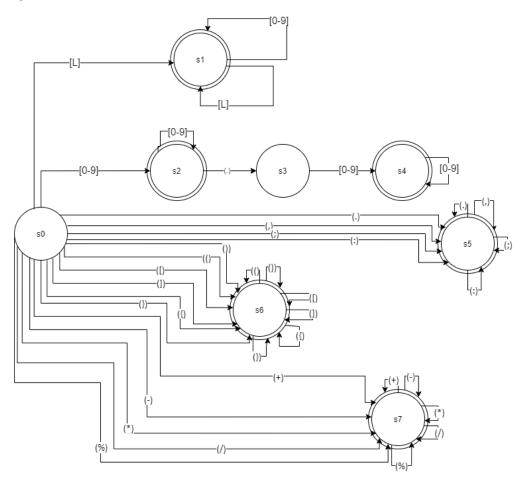
3.5.





# 4. Creación del AFD que acepte todos los tokens

# 4.1. Diagrama de transiciones del AFD



## 4.2. Tabla de transiciones del AFD

FT	s0	s1	s2	s3	s4	s5	s6	s7
L	∂(s0, [L])=s1	∂(s1,						
		[L])=s1						
[0-9]	∂(s0, [0-9])=s2	∂(s1, [0-	∂(s2, [0-	∂(s3, [0-	∂(s4, [0-			
		9])=s1	9])=s2	9])=s4	9])=s4			
(.)	∂(s0, (.))=s5		∂(s2,			∂(s5,		
			(.))=s3			(.))=s5		
(,)	∂(s0, (,))=s5					∂(s5,		
						(,))=s5		
(;)	∂(s0, (;))=s5					∂(s5,		
						(;))=s5		
(:)	∂(s0, (:))=s5					∂(s5,		
						(:))=s5		
(+)	∂(s0, (+))=s7							∂(s7, (+))=s7
(-)	∂(s0, (-))=s7							∂(s7, (-))=s7
(*)	∂(s0, (*))=s7							∂(s7, (*))=s7
(/)	∂(s0, (/))=s7							∂(s7, (/))=s7
(%)	∂(s0, (%))=s7							∂(s7, (%))=s7
(()	∂(s0, (() )=s6						∂(s6, (()	
							)=s6	
())	∂(s0, ()) )=s6						∂(s6, ())	
							)=s6	
([)	∂(s0, ([))=s6						∂(s6, ([))=s6	
(])	∂(s0, (]))=s6						∂(s6, (]))=s6	
({)	∂(s0, ({))=s6						$\partial$ (s6, ({))=s6	
(})	∂(s0, (}))=s6						$\partial$ (s6, (}))=s6	