

Motivation

- Problem:
 - select x from Y
 where z = 'k';
- Plan:

select z = 'k'

↑

scan Y
- Cost?

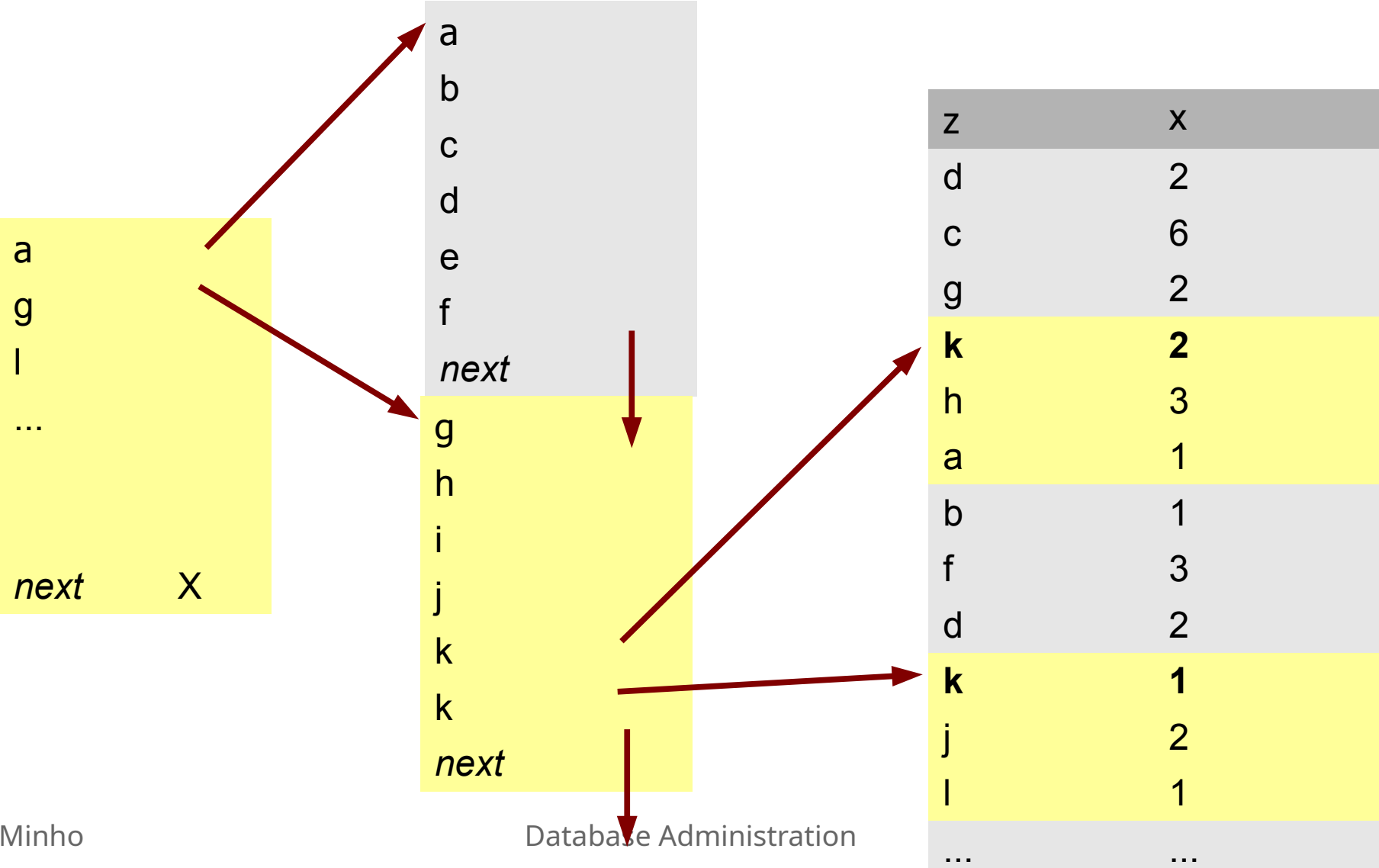
z	x
d	2
c	6
g	2
k	2
h	3
a	1
b	1
f	3
d	2
k	1
j	2
l	1
...	...

Index

- Makes it easy to find pages containing interesting data
- Smaller than data
 - Fits in memory?
- Efficient look-up:
 - Identity (=)
 - Ranges
 - LIKE
 - ...

z	x
d	2
c	6
g	2
k	2
h	3
a	1
b	1
f	3
d	2
k	1
j	2
l	1
...	...

B-Tree



B-Tree

- Insert:
 - If free entry not available, split leaf
 - Recursively insert new leaf in upper layer
 - Tree grows towards the root!
 - Add entry to leaf
- Delete:
 - Remove entry from leaf
 - If enough space available, collapse leafs
 - Recursively delete leaf in upper layer

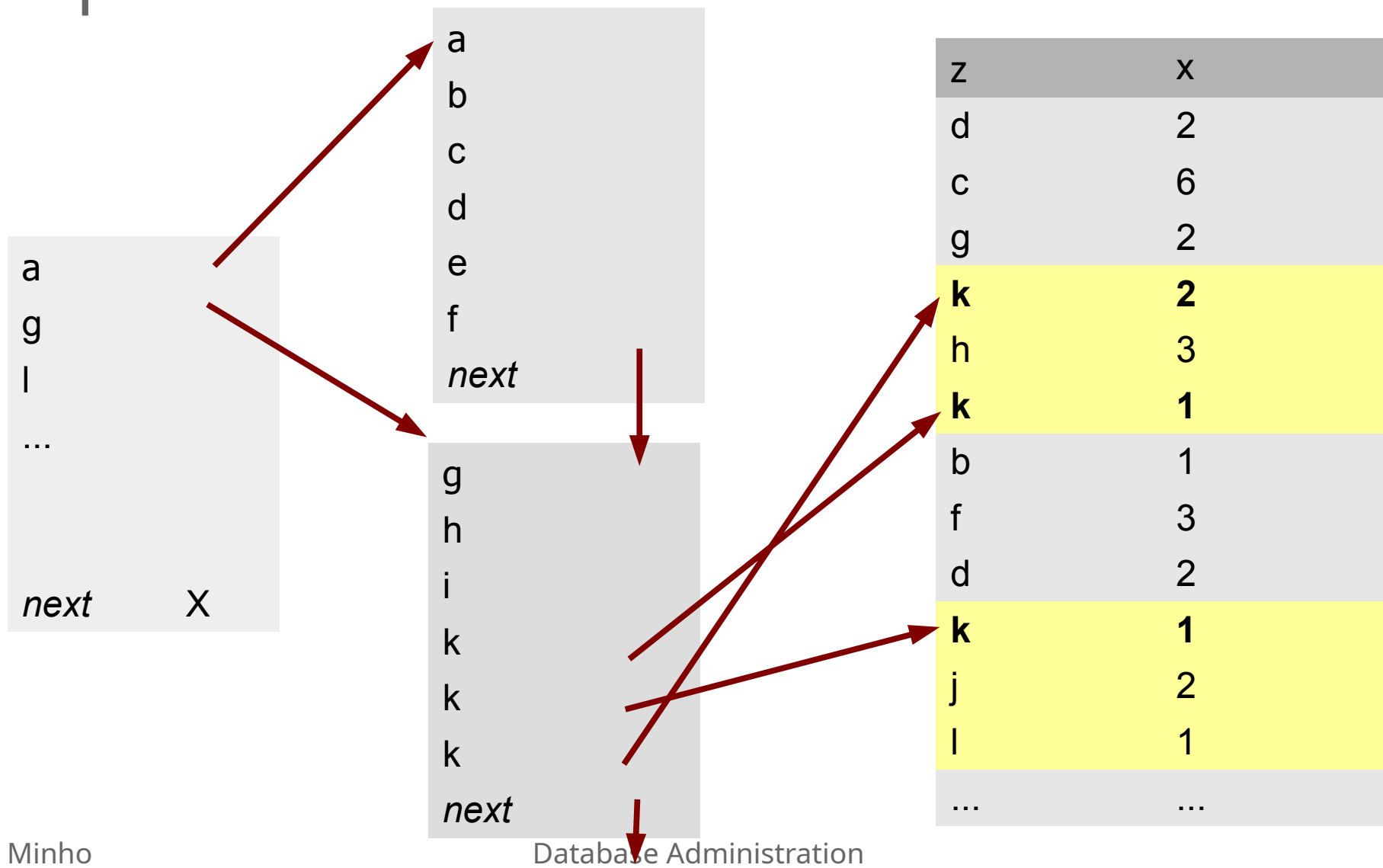
B-Tree

- Desirable characteristics:
 - Balanced
 - $\log(n)$ depth
 - Fit for block I/O
- Supports:
 - Identity look up
 - Range queries / Ordered scan
 - Updates

Composite indexes

- Index on (X,Y):
 - Answers equality and interval on X
 - Answers equality on X and interval on Y
- Index on expression, e.g. $X+Y$
 - Answers equality and interval on $X+Y$

Dispersion



Clustered indexes

- Problem:
 - $\#blocks \gg (\#records / \text{records per block})$
 - Read each block multiple times
- A clustered index:
 - Records are (roughly) sorted according to the index
 - No sorting within a block is needed
 - Free space may be kept for insertions

Multi-criteria filtering

- Problem:

- select x from Y
 where z = 'k'
 and x = 4;

Index on Z

g
h
i
k
k
k
<i>next</i>

Z	X
d	2
c	6
g	4
h	2
h	3
k	4
b	1
f	3
d	2
k	1
j	2
l	1
...	...

Index on X

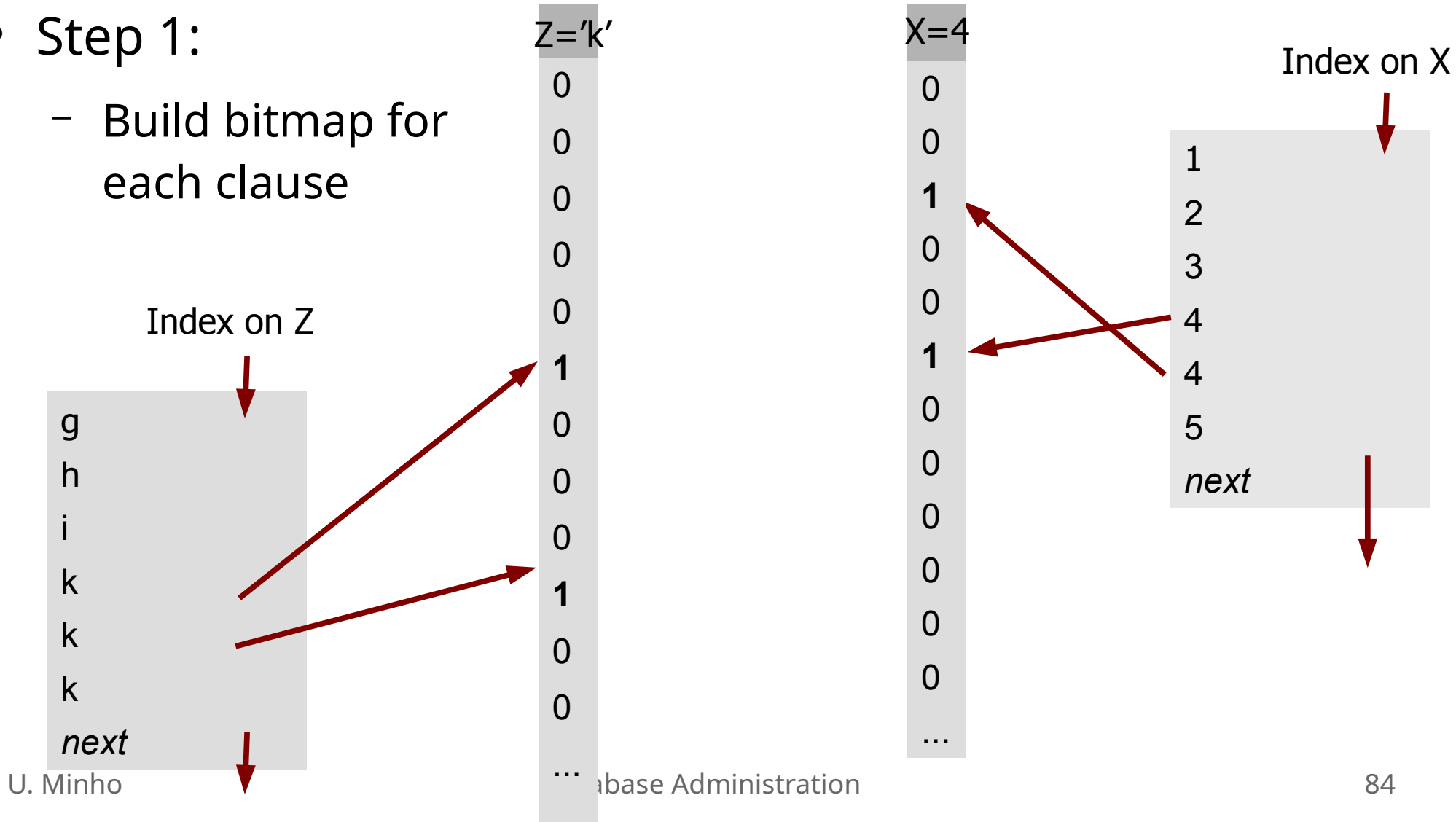
1
2
3
4
4
5
<i>next</i>

Multi-criteria filtering

- Sequential scan
- Either X or Z plus scan
- Composite index on (Z,X):
 - Often, many columns and combinations
- Typical examples:
 - Search with a combination of features
 - Hotel booking, real estate, on-line shopping, ...
- What if “or” instead of “and”?
 - (T2 or T3) and in Braga

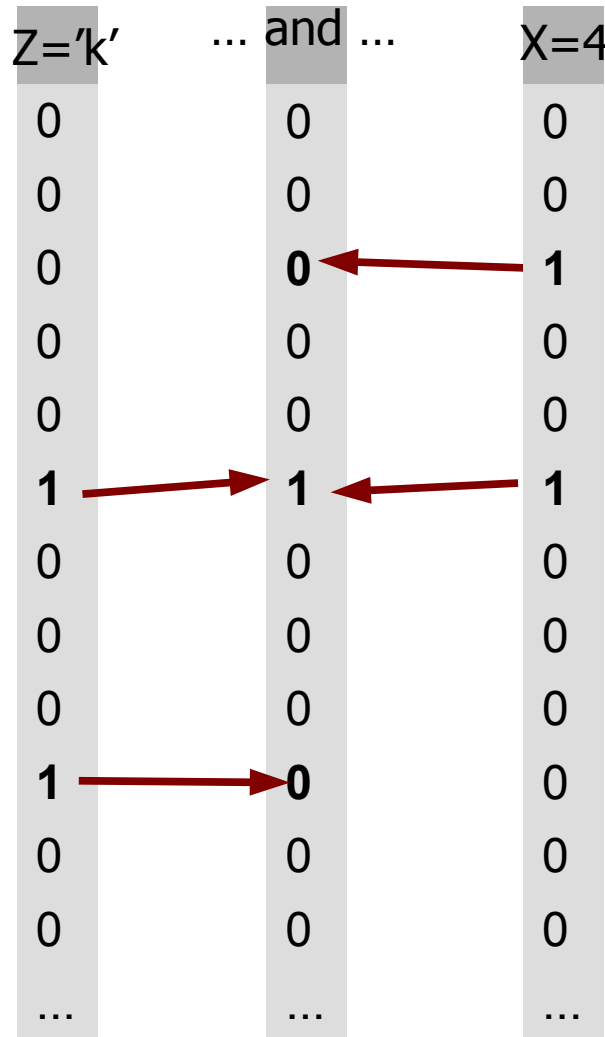
Bitmap indexes

- Step 1:
 - Build bitmap for each clause



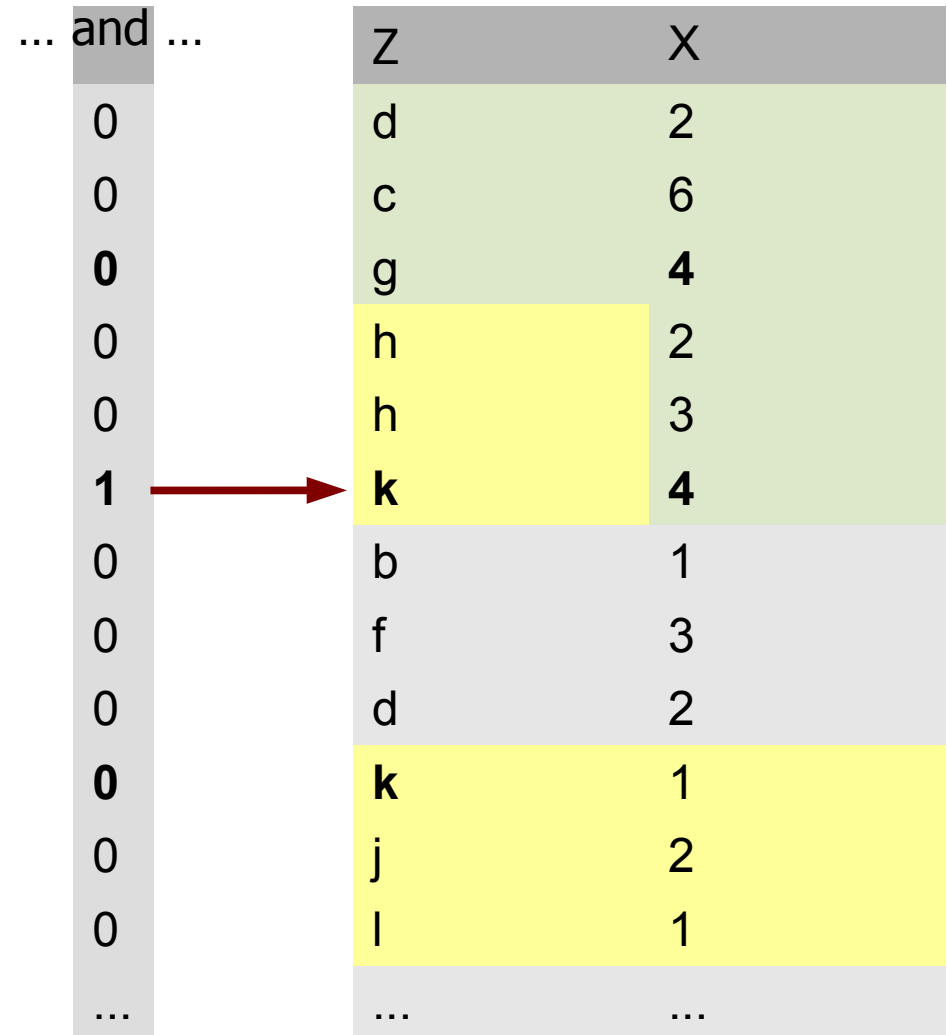
Bitmap indexes

- Step 2:
 - Combine with logical operators



Bitmap indexes

- Step 3:
 - Traverse bitmap and read table



Motivation

- Assumptions:
 - Several TB of data
 - ~50%, $y=1$
 - ~50%, $y=2$
 - a few, $y=3$

z	y
d	1
c	2
g	1
k	2
h	3
a	1
b	1
f	2
d	2
k	1
j	2
l	1
...	...

Motivation

- Problem:
 - select count(*) from X
where y = 1;
- Possible plans:

count(*)
↑
select y = 1
↑
scan X

count(*)
↑
index scan X
(y = 1)

- Cost?

z	y
d	1
c	2
g	1
k	2
h	3
a	1
b	1
f	2
d	2
k	1
j	2
l	1
...	...

Motivation

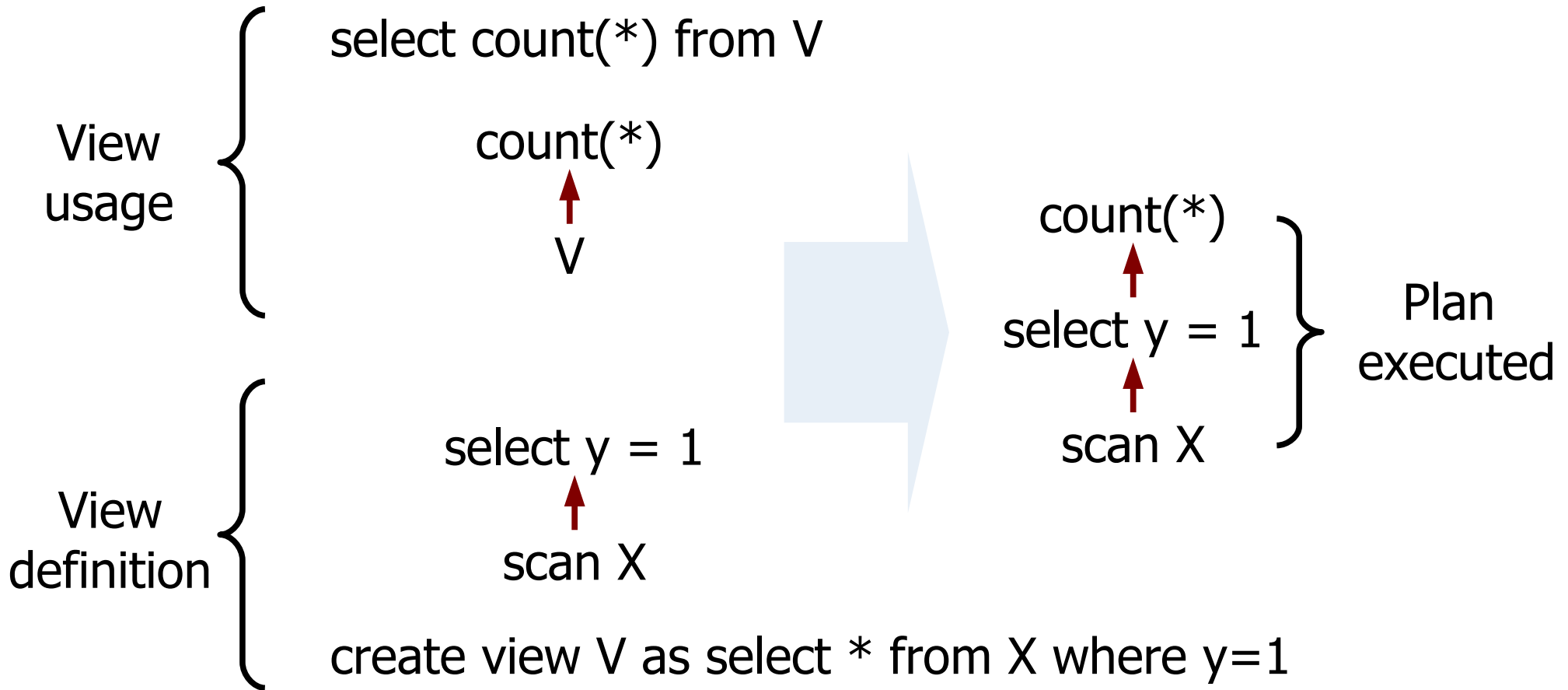
- Keep results cached
- Update when needed

`select y, count(*) from X group by y`

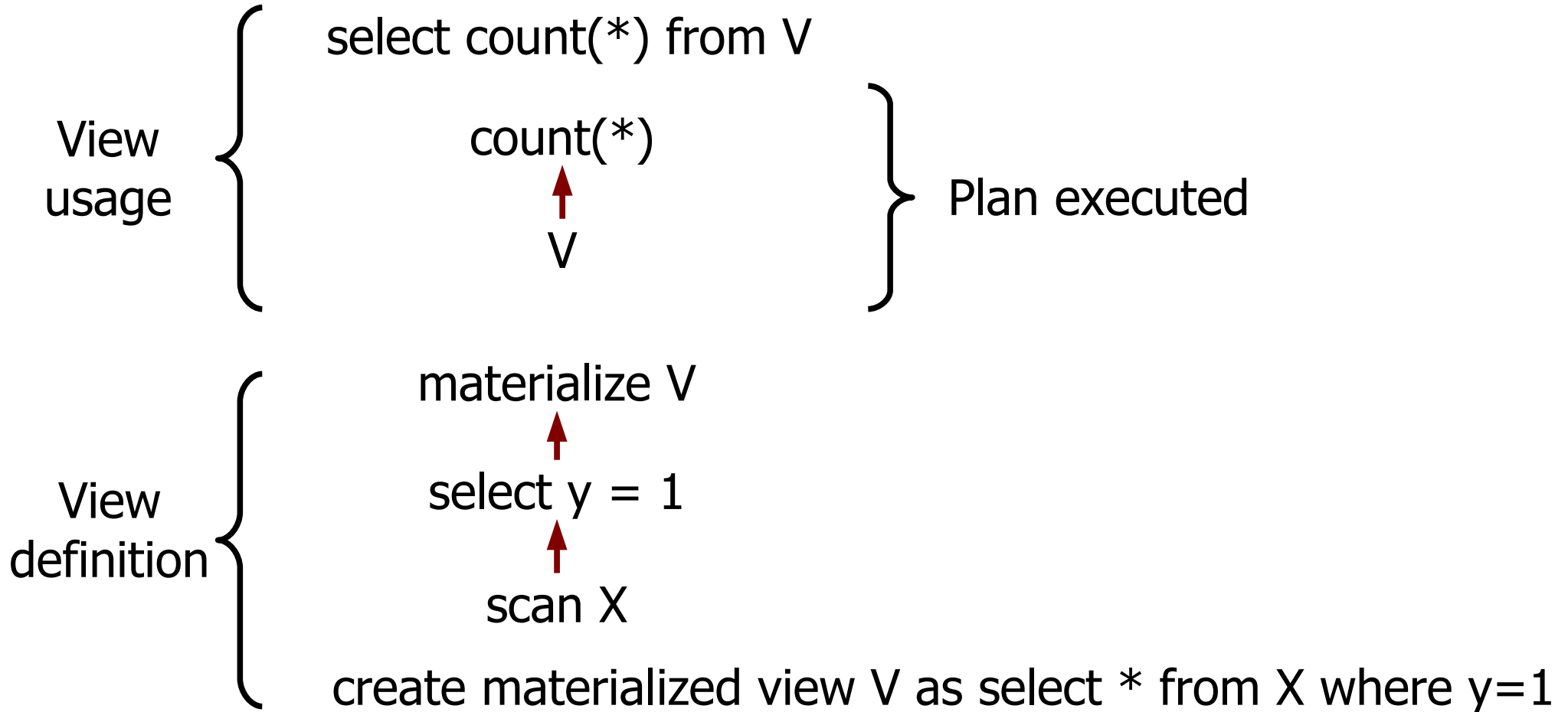
y	count
1	773647263
2	765732332
3	1

X	
z	y
d	1
c	2
g	1
k	2
h	3
a	1
b	1
f	2
d	2
k	1
j	2
l	1
...	...

Views



Materialized views



Maintaining materialized views

- Periodically run the query and update the view
- Update the view when data changes

DIY Materialized Views

- Updating with AFTER triggers:

```
-- Create view
select sum(value) into mv_sum_items from items;

-- Update view
create function upd_sum_items() returns trigger as '
    BEGIN
        update mv_sum_items set sum = sum + new.value - old.value;
        return new;
    END
' language 'plpgsql';

create trigger upd_sum
    after update on items
    for each row execute procedure upd_sum_items();
```

Using materialized views

- Automatically used by the planner:
 - Indexed views in MS SQL Server
- Used explicitly in queries:
 - Materialized views in Oracle
 - DIY materialized views everywhere
 - Developer tip:
 - Using views allows the DBA to select which ones to materialize

Conclusions

- Indexes and mat. views = Redundancy!
- Trade-off between:
 - Complexity of operationsand:
 - Disk space used
 - Usage of main memory
 - Effort when updating
- Usefulness depends on workload mix