

Introduction to Artificial Intelligence



COMP307

Planning and Scheduling 2: Static Scheduling

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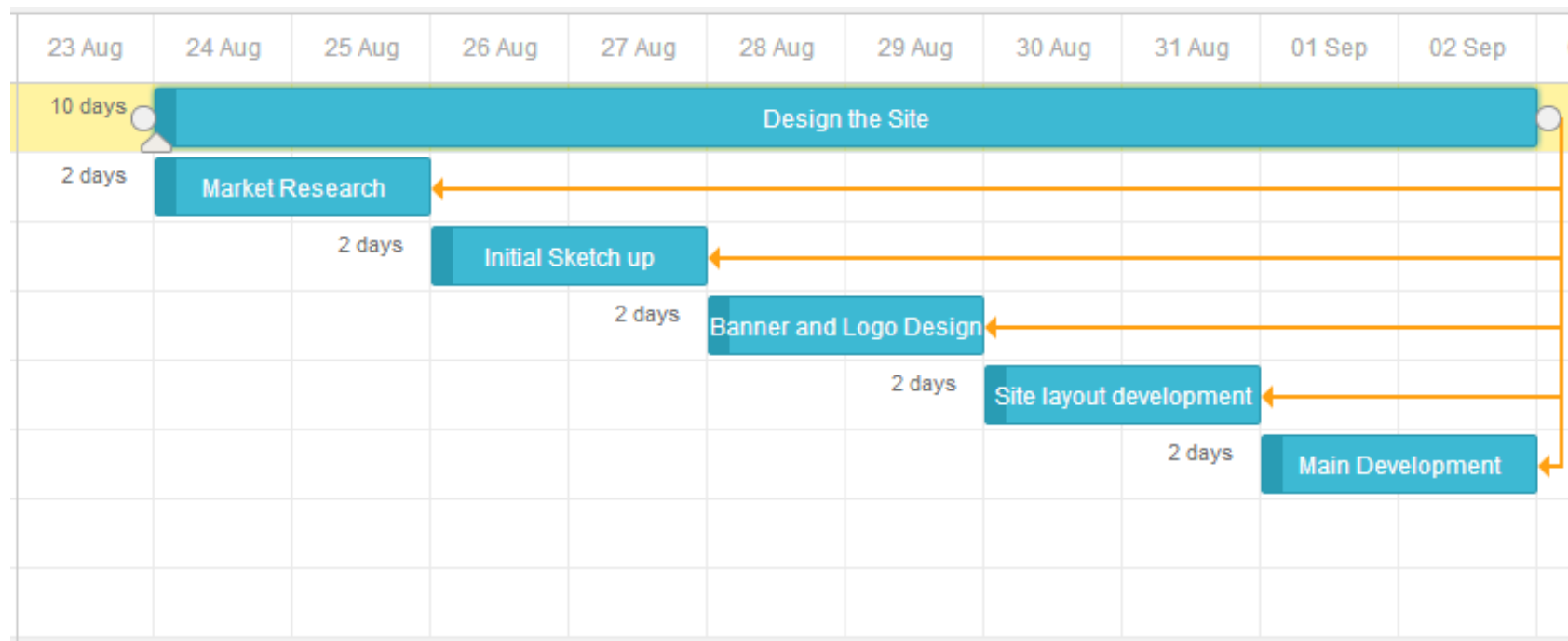
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Outline

- Why Scheduling?
- What is scheduling?
- Job Shop Scheduling
- Solution as Schedule
- Job Shop Scheduling Algorithms

Why Scheduling

- Planning does not consider time
 - States have **no time stamp**
 - Actions have **no starting time**
 - Actions have **no duration**
- Scheduling **deals with time** in planning



Why Scheduling

- Real-world applications need scheduling
 - Cloud computing/Resource allocation
 - Project management
 - Exam/course timetabling
 - Manufacturing
 - ...



What is Scheduling

- Scheduling v.s. Planning

- Planning is to find a **sequence of actions** to achieve the goal state from the initial state
- In Scheduling, each action has a **starting and finishing time**
- $s' = \text{RESULT}(s,a)$ is reached **after the action is finished** (it takes time from s to s')

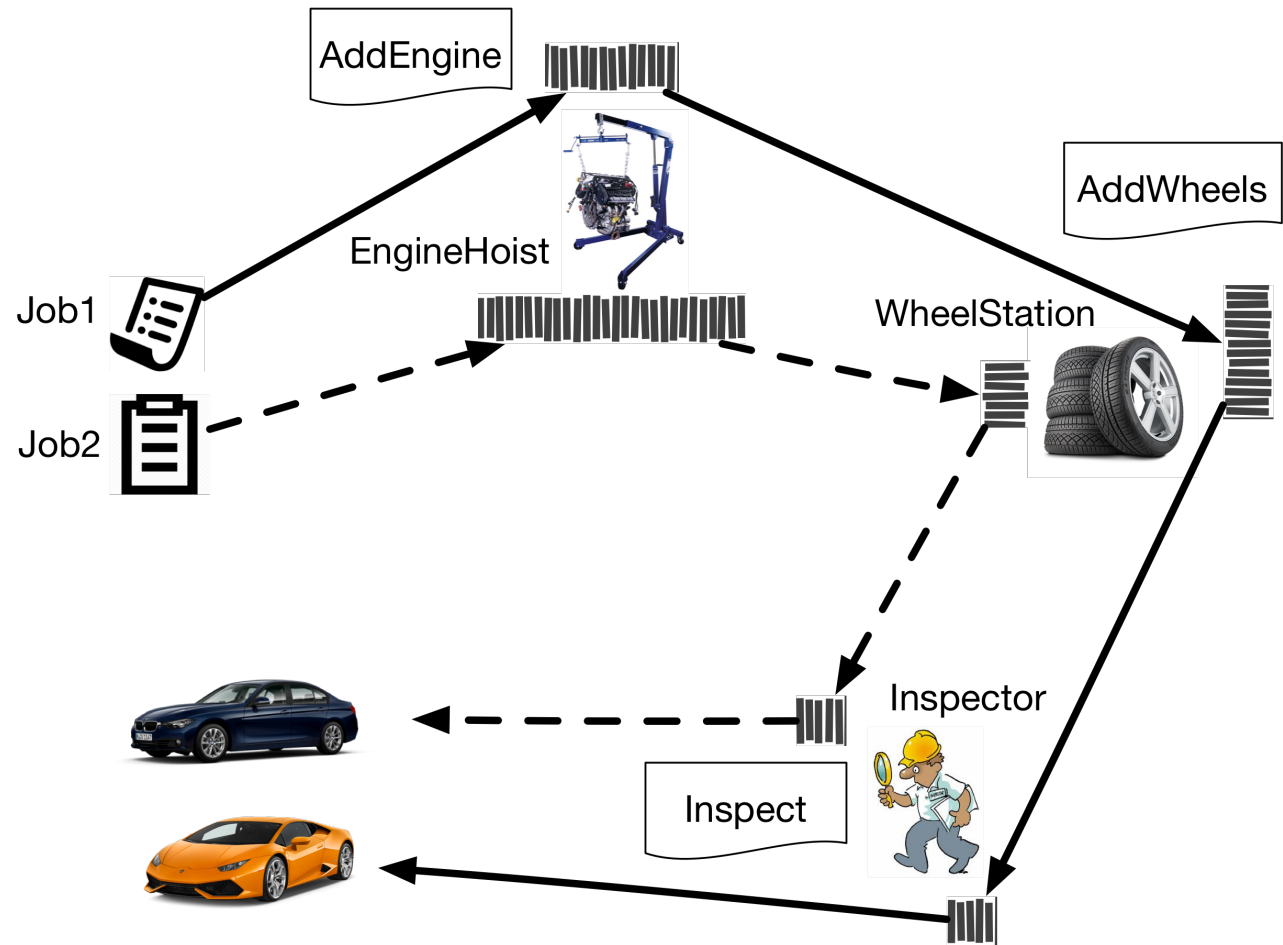
- An example: **Job Shop Scheduling**

Job Shop Scheduling

- A set of **jobs** to be processed
- Each job has a sequence of **operations**, following the **ordering constraint**
 - An operation cannot start until its preceding operations have completed
 - The first operation can start at any time
- Each operation has a **processing time**, and occupies a **machine**
 - (**Resource constraint**) Each machine can process only one operation at a time
- **Objective**: process the operations with the machines to minimise the total duration (makespan) of the plan

JSS: An Example

- A car manufacturing factory
 - Two types of cars
 - Three operations (simplified)
 - Add engine
 - Add wheels
 - Inspect



JSS: Representation

- Jobs({AddEngine1 < AddWheels1 < Inspect1},
• {AddEngine2 < AddWheels2 < Inspect2})
- Machines(EngineHoist, WheelStation, Inspector)
- Operation(AddEngine1, ProcTime: 30, Use: EngineHoist)
- Operation(AddEngine2, ProcTime: 60, Use: EngineHoist)
- Operation(AddWheels1, ProcTime: 30, Use: WheelStation)
- Operation(AddWheels2, ProcTime: 15, Use: WheelStation)
- Operation(Inspect_i, ProcTime: 10, Use: Inspector)

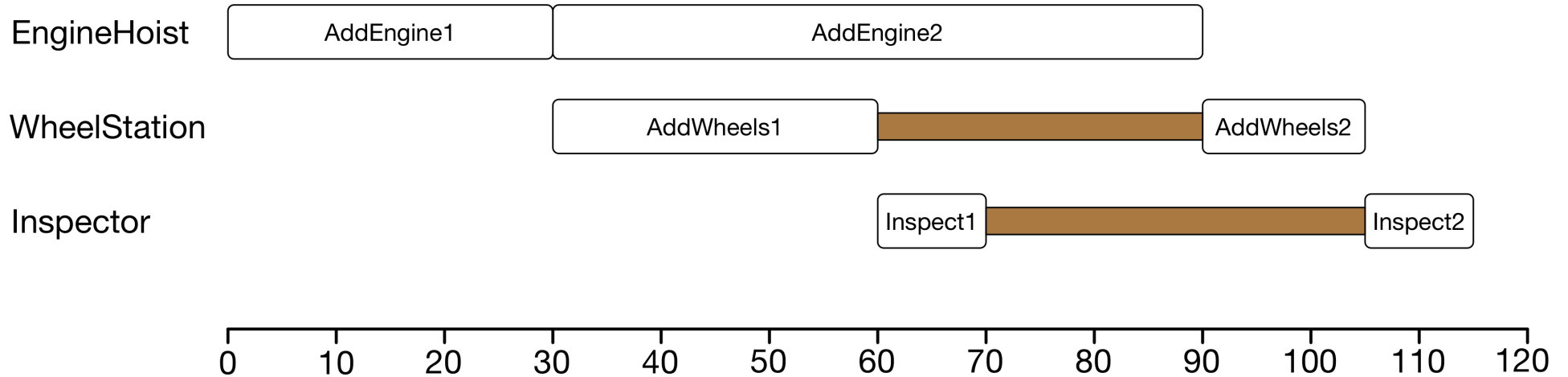
JSS: Table

Job	Operation	Machine	ProcTime
1	AddEngine1	EngineHoist	30
	AddWheels1	WheelStation	30
	Inspect1	Inspector	10
2	AddEngine2	EngineHoist	60
	AddWheels2	WheelStation	15
	Inspect2	Inspector	10

- A solution is a schedule that processes these jobs with the machines (gantt chart)

Two different solutions

- Job 1 first: total duration = 115



- Job 2 first: total duration = 130



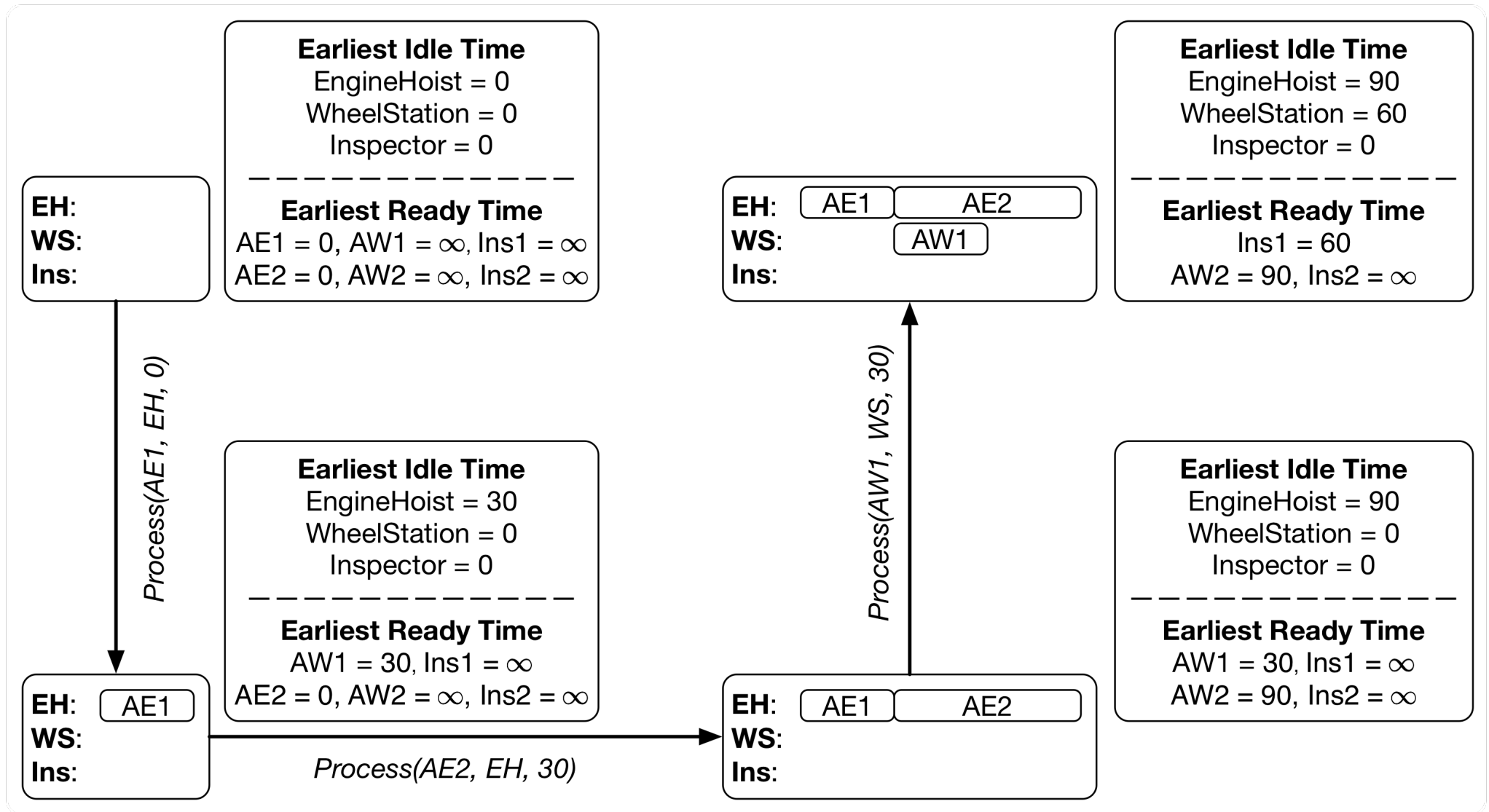
Search for Schedules

- **Initial state**
 - Empty schedule, $t=0$, all operations unprocessed
 - The first operation of each job is **ready** at time 0, all the other operations are not ready
 - All machines are **idle** at time 0
- **Goal state**: all operations processed
- **Actions**: $\text{Process}(o, m, t)$
 - Start processing operation o with machine m at time t
 - Precondition:
 - o unprocessed, and is ready at time t
 - m is idle at t
 - Effect:
 - o processed
 - $\text{next}(o)$ (if exists) is ready at time $t + \text{ProcTime}(o)$, and is idle at $t + \text{ProcTime}(o)$
- How to decide **t** ?

Deciding Starting Time of Action

- **Non-delay**: start the action as soon as possible
 - Operation **earliest ready time**
 - Machine **earliest idle time**
 - Earliest starting time: the **later** between the above two
- Find operation earliest ready time
 - Initial: 0 for the first operation, and infinity for others
 - When $\text{Process}(o, m, t)$ is scheduled, then the earliest ready time of $\text{next}(o)$ becomes $t + \text{ProcTime}(o)$
- Find machine earliest idle time
 - Initial: 0
 - When $\text{Process}(o, m, t)$ is scheduled, then the earliest idle time of machine m becomes $t + \text{ProcTime}(o)$

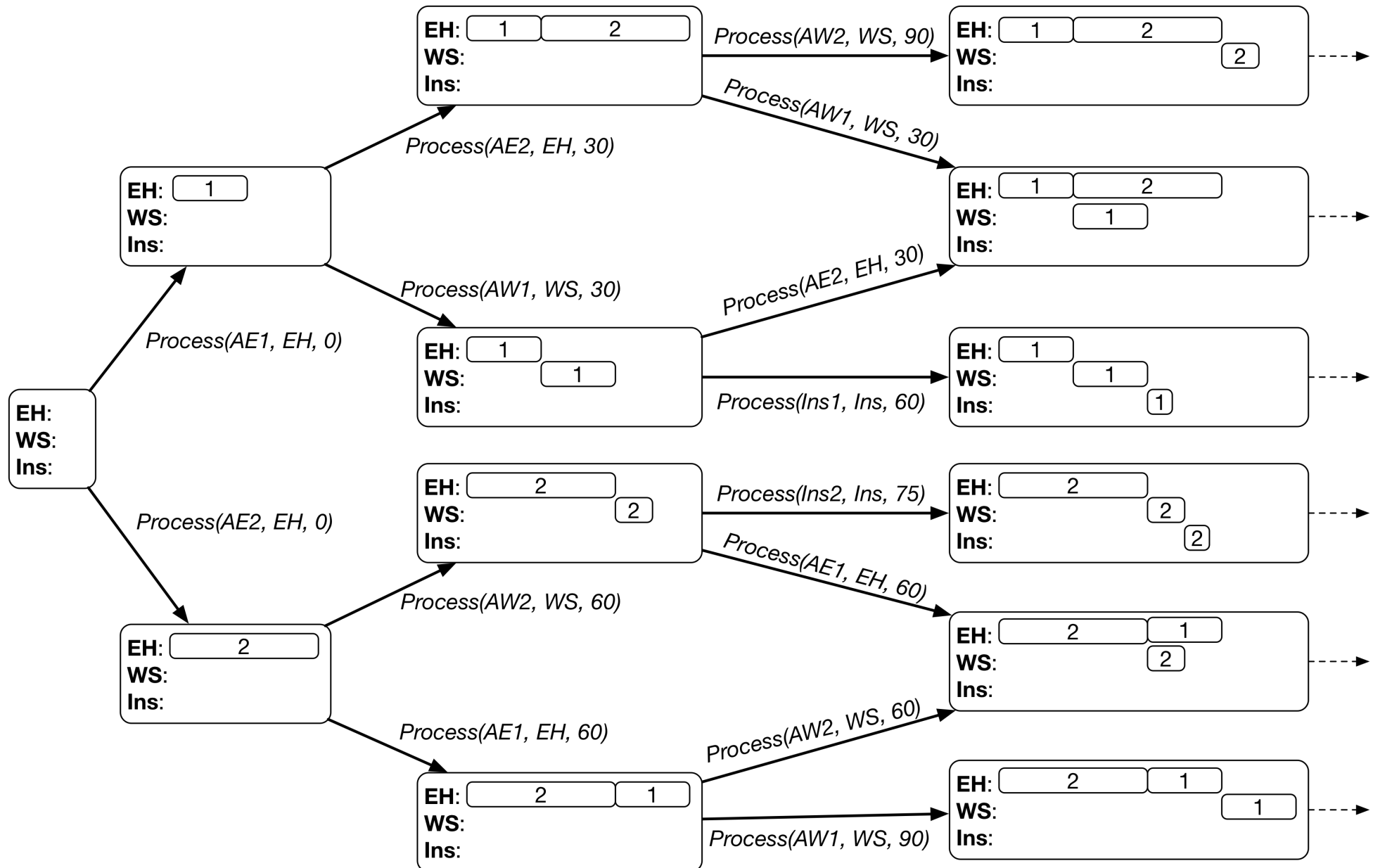
Update Earliest Ready and Idle Time



Forward State-Space Search

- Start from the initial state
 - Empty schedule, $t=0$, all operations unprocessed
 - The first operation of each job is **ready** at time 0, all the other operations are not ready
 - All machines are **idle** at time 0
- Examine all the **applicable** actions $\text{Process}(o,m,t)$
 - Enumerate each unprocessed operation o and its machine m
 - Calculate the earliest starting time t
 - Applicable if $t < \infty$
- **All the leaf nodes are goal states** (all operations processed)
- Each schedule is a path from the root node to a leaf node

Forward State-Space Search

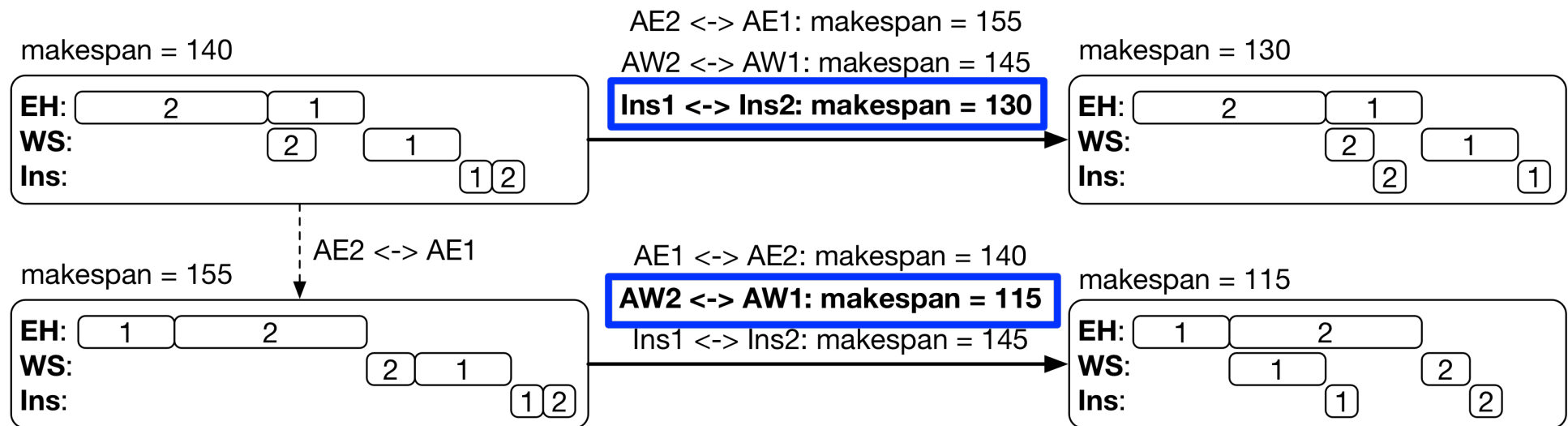


Local Search (Hill Climbing)

- **Step 1.** Random generate a scheduling s ;
- **Step 2.** Examine all the neighbours in the **neighbourhood** of s , and select the best neighbour s' ;
- **Step 3.** If s' is better than s , set $s \leftarrow s'$, and go to Step 2. Otherwise return s .
- How to define the **neighbourhood**?
 - Local/Subtle modification
 - Still feasible
 - An example: swap operations on the same machine
 - **Shift the starting time of all actions**

Local Search (Hill Climbing)

- Step 1. Random generate a scheduling s ;
- Step 2. Examine all the neighbours in the **neighbourhood** of s , and select the best neighbour s' ;
- Step 3. If s' is better than s , set $s \leftarrow s'$, and go to Step 2. Otherwise return s .



- Jump out of local optima: simulated annealing, genetic algorithms, ...

Jobs with Different Sequences

Job	Operation	Machine	ProcTime
1	AddEngine1	EngineHoist	30
	AddWheels1	WheelStation	30
	Inspect1	Inspector	10
2	AddWheels2	WheelStation	15
	AddEngine2	EngineHoist	60
	Inspect2	Inspector	10



Summary

- What is Scheduling? – Temporal planning (actions have starting/finishing time)
- Solution as schedule
- Search for schedules: state-space search, local search, ...
- Suggested Reading: Textbook, Chapter 11
- Next: Dynamic scheduling