

1)

- a) Run a for (i = 0; start <= end; i++)
 If x = A[i], increment the count
 If x < A[i], change start to mid + 1
 Else if x > A[i], change end to end - 1
 If x = -1, the x isn't in the array
- b) for (i = 0; start <= end; i++) // n
 if (x == A[i]) { Output: count++ } // 1
 if (x < A[i]) { Output: start = mid + 1 } // 1
 else { Output: end = mid - 1 } // 1
 if (x == -1) { Output: return -1 } // 1
- c) $O(n) = \log(n)$

2)

- a) Take the stack of coins and split the amount in equal half (if even), split into 3 equal groups (if odd).
 Whichever side is heavier, discard the lighter side and split the heavier side into equal halves.
 Continue this process until you only have 2 coins left.
 Whichever the heavier one is the outlier coin.
- b) $O(n) = n \cdot \log(n)$

3)

- a) $O(n)$
- b) Using Binary Search on a sorted Linked List would give us the same time complexity as using a Linear Search. This because even if the Linked List is sorted, it isn't necessarily contiguous in memory, so it will take roughly the same time to perform a Binary Search as it would Linear.

4)

- a) Check if x is either 0 or 1. If so, return x.
 Run a while (start <= end)
 If mid^2 is equal to x, the number is a perfect square
 If mid^2 is less than x, change the start to mid + 1 and the answer to mid
 Else if mid^2 is greater than x, change the end to mid - 1
- b) if (x == 0 || x == 1) { Output: return x } // 1
 while (start <= end) // n
 mid = (start + end) / 2 // 1
 if (mid * mid == x) { Output: return mid } // 1
 if (mid * mid > x) // n
 start = mid + 1 // 1
 ans = mid // 1

else

// n

end = mid + 1

// 1

c) $O(n) = \log n$