# Computer Engineering & Informatics Department University of Patras

**Scientific Calculation** 

Academic Year 2015-2016 Laboratory Exercise

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### Question 1

### (1) System Characterization

<u>Processor:</u> 2nd generation Intel® Core<sup>TM</sup> i3-2348M Processor

OS: Windows 7 (6.1) Professional 64-bit

*Hard drive:* 500 GB (5,400 rpm) σειριακό ATA

<u>System memory:</u> 4,096 (1x) MB, μέγιστη επεκτασιμότητα: 8,192 MB, τεχνολογία: DDR3 RAM (1.333 MHz)

### 3 levels of cache:

Level 1 cache size:

- 2 x 32 Kbytes, 8way set associative, 64- byte line size (L1 D-cache)
- 2 x 32 Kbytes, 8way set associative, 64- byte line size (L1 I- cache)

Level 2 cache size

• 2 x 256 Kbytes, 8way set associative, 64- byte line size

Level 3 cache size

• 3 Mbytes, 12 way set associative, 64- byte line size

### (II) OS & MATLAB:

- •Windows 7 (6.1) Professional 64-bit
- •*MATLAB R2013a*

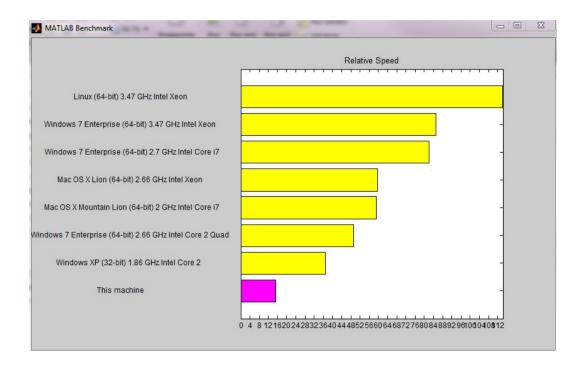
### (III) Metroprogram and command result bench

```
Matlab Benchmark
   -----
Number of times each test is run
  I. Matrix calculation
Creation, transp., deformation of a 1200x1200 matrix (sec): 0.084263
1250x1250 normal distributed random matrix ^1000___ (sec): 0.15168
Sorting of 1,100,000 random values (sec): 0.098116
550x550 cross-product matrix (b = a' * a) (sec): 0.015296
Linear regression over a 700x700 matrix (c = a \ b') (sec): 0.032516
                     _____
                    Trimmed mean (2 extremes eliminated): 0.071632
  II. Matrix functions
                                    (sec): 0.031067
FFT over 900,000 random values_
Eigenvalues of a 220x220 random matrix (sec): 0.090773
Determinant of a 750x750 random matrix (sec): 0.061002
Cholesky decomposition of a 1000x1000 matrix_____ (sec): 0.087038
Inverse of a 500x500 random matrix (sec): 0.13686
                    Trimmed mean (2 extremes eliminated): 0.079604
  III. Programmation
  _____
225,000 Fibonacci numbers calculation (vector calc) (sec): 0.18627
Creation of a 1500x1500 Hilbert matrix (matrix calc) (sec): 0.15504
Grand common divisors of 35,000 pairs (recursion) ___ (sec): 0.10264
Creation of a 220x220 Toeplitz matrix (loops) _____ (sec): 0.0019962
Escoufier's method on a 22x22 matrix (mixed) _____ (sec): 0.54551
                    ______
                    Trimmed mean (2 extremes eliminated): 0.14798
Total time for all 15 tests
Overall mean (sum of I, II and III trimmed means/3) (sec): 0.099739
                    --- End of test ---
```

Using the official matlab platform bench, we have the following results:

```
>> bench
ans =

0.6160  0.3289  0.6830  1.2787  1.8909  0.9542
```



Computer Type	LU	FFT	ODE	Sparse	2-D	3-D
Linux (64-bit) 3.47 GHz Intel Xeon	0.0626	0.0661	0.1617	0.1161	0.2055	0.0911
Windows 7 Enterprise (64-bit) 3.47 GHz Intel Xeon	0.0701	0.0719	0.1094	0.1297	0.2784	0.7044
Windows 7 Enterprise (64-bit) 2.7 GHz Intel Core i7	0.0857	0.0801	0.0958	0.1319	0.2975	0.7003
Mac OS X Lion (64-bit) 2.66 GHz Intel Xeon	0.0547	0.1278	0.2008	0.1877	0.6670	0.6299
Mac OS X Mountain Lion (64-bit) 2 GHz Intel Core i7	0.0725	0.1292	0.1881	0.1587	0.7203	0.6476
Windows 7 Enterprise (64-bit) 2.66 GHz Intel Core 2 Quad	0.1239	0.2333	0.1561	0.2822	0.4819	0.7390
Windows XP (32-bit) 1.86 GHz Intel Core 2	0.3406	0.3178	0.1883	0.3542	0.5775	0.3601
This machine	0.6160	0.3289	0.6830	1.2787	1.8909	0.9542
Place the cursor near a co this data to compare different v see the help for the bend	ersions of MATLAB, or	to download an	updated timing	data file,		

### Question -Timing of Functions

- (I) Basic Operations & Functions of MATLAB
- [L,U]=lu(A): Table L returns a lower triangular array and o U an upper upper triangle so LU = PA, where P is a matrix matrix.
- [L,U,P]=lu(A): Table P returns the matrix P
- qr: It represents a register as the product of a real or complex ortho-regular register (Q) and an upper triangular (R)
- **svd:** Analysis of M at particular prices. Generating a table in the form M = USV \* where U is a ortho-mononial matrix. A rectangular array with non-negative values in the diagonal (its diagonal elements are M's) and V \* are real ortho-mononial matrix.
- det: Calculates the table definition
- rank: Calculates an estimate of the number of linearly independent lines or columns of a complete registry
- polyval (p,x) The x statement can be either an element, or a vector, or a register. The result y is the same size as x and contains the polynomial values at x. The method you use to calculate the value of the polynomial is the Horner method

Intrinsic functions are the functions we do not have access to in their code, while External, the ones we have access to. In addition, svd, qr, lu and det are built in functions, while rank and polyval are exogenous.

We run the following code:

```
clear all;
clc;
%Erwthma 2-(ii)- (a)

k=1;
for i=7:12
    n=2^i;
    A=rand(n);
    b=rand(n,1);

%xronometrisi gia thn lu
tic;
    [L, U, P]=lu(A);
time_lu(k,1)=toc;
gflopi(k,1)=2*n^2/time_lu(k,1)/10^9;

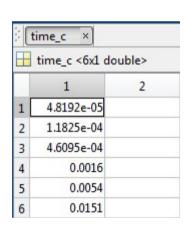
%xronometrisi gia thn A*b
```

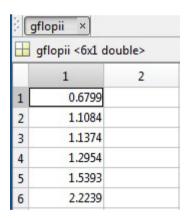
```
tic;
    c=A*b;
    time c(k,1)=toc;
    gflopii(k,1)=2*n^2/time\ c(k,1)/10^9;
    k=k+1;
  end
subplot(1,2,1)
i=2.^{7:12};
plot(i,time_lu(:,1), 'g.-', i, time_c(:,1), 'b.--');
legend('lu(A)', 'c');
xlabel('n');
ylabel('t sec');
title('Xρηση tic-toc');
subplot(1,2,2)
i=2.^{7:12};
plot(i,gflopi(:,1), 'g.-', i, gflopii(:,1), 'b.--');
legend('lu(A)', 'c');
xlabel('n');
ylabel('gflop');
title('Επίδοση gflop');
```

### **Execution times**

time_lu × time_lu <6x1 double>			
	1	2	
1	0.0057		
2	0.0062		
3	0.0070		
4	0.0416		
5	0.2786		
6	2.1254		

⊞ gf	gflopi <6x1 double>				
	1	2			
1	0.0057				
2	0.0212				
3	0.0749				
4	0.0504				
5	0.0301				
6	0.0158				

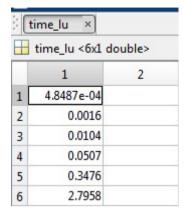


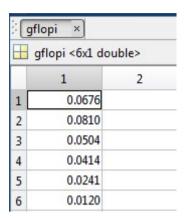


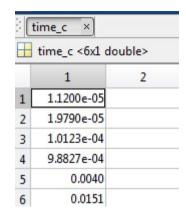
(u)- ( $\beta$ ) Execution times for **tic**, **toc**  $\sigma\tau\alpha$  [L, U, P]=lu(A)  $\kappa\alpha\iota$  c=A\*b for 20 times and gflop performance. We run the following code:

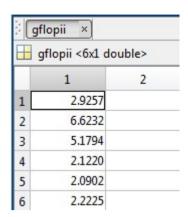
```
clear all;
clc;
%Erwthma 2-(ii)- (b)
k=1;
  for i=7:12
     n=2^i;
      A=rand(n);
     b=rand(n,1);
     %xronometrisi gia thn lu
     tic;
     for 1=1:20
        [L, U, P]=lu(A);
     end
     time lu(k,1)=toc/20;
     gflopi(k,1)=2*n^2/time_lu(k,1)/10^9;
     %xronometrisi gia thn A*b
     tic;
     for 1=1:20
        c=A*b;
     end
     time c(k,1)=toc/20;
     gflopii(k,1)=2*n^2/time c(k,1)/10^9;
    k=k+1;
  end
subplot(1,2,1)
i=2.^[7:12];
plot(i,time_lu(:,1), 'g.-', i, time_c(:,1), 'b.--');
legend('lu(A)', 'c');
xlabel('n');
ylabel('t sec');
title('Χρηση tic-toc');
subplot(1,2,2)
i=2.^{7:12};
plot(i,gflopi(:,1), 'g.-', i,gflopii(:,1), 'b.--');
legend('lu(A)', 'c');
xlabel('n');
ylabel('gflop');
title('Επίδοση gflop');
```

#### **Execution times**









(u) Using the timeit function in [L, U, P] = lu(A) and c = A \* b and performance gflop:

We are running the following code

```
clear all
clc
%Erwthma 2-(g)
k=1;
  for i=7:12
     n=2^i;
     A=rand(n);
     b=rand(n,1);
     f = (a()lu(A);
     time lu(k,1)=timeit(f,3);
     gflopi(k,1)=2*n^2/time lu(k,1)/10^9;
     f = (a)(A*b;
     time c(k,1)=timeit(f,1);
     gflopii(k,1)=2*n^2/time_c(k,1)/10^9;
   k=k+1;
  end
subplot(1,2,1)
i=2.^{7:12};
plot(i,time lu(:,1), 'g.-', i, time c(:,1), 'b.--');
legend('lu(A)', 'c');
xlabel('n');
ylabel('t sec');
```

```
title('Χρηση timeit')

subplot(1,2,2)
i=2.^{[7:12]};
plot(i, gflopi(:,1), 'g.-', i, gflopii(:,1), 'b.--');
legend('lu(A)', 'c');
xlabel('n');
ylabel(' gflop');
title('Επίδοση gflop');
```

### **Execution times**

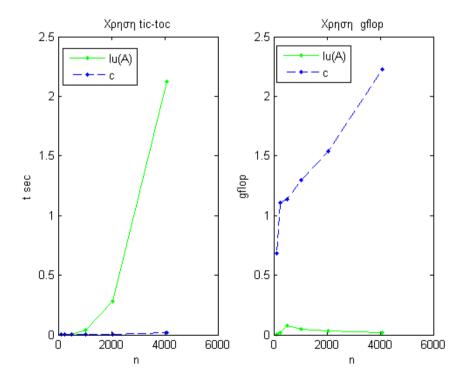
time_lu × time_lu <6x1 double>			
	1	2	
1	2.9004e-04		
2	0.0017		
3	0.0080		
4	0.0444		
5	0.3302		
6	2.8016		

⊞ gi	flopi <6x1 dou	ble>
	1	2
1	0.1130	
2	0.0763	
3	0.0659	
4	0.0472	
5	0.0254	
6	0.0120	

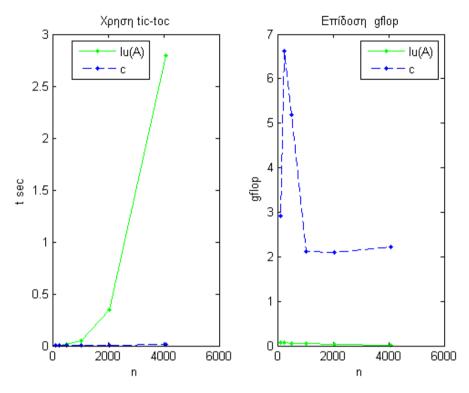
t	ime_c ×	
	time_c <6x1 do	uble>
	1	2
1	1.0833e-05	
2	2.2852e-05	
3	7.6324e-05	
4	9.0531e-04	
5	0.0035	
6	0.0172	

🔢 gflopii <6x1 double>			
	1	2	
1	3.0247		
2	5.7358		
3	6.8693		
4	2.3165		
5	2.3809		
6	1.9455		

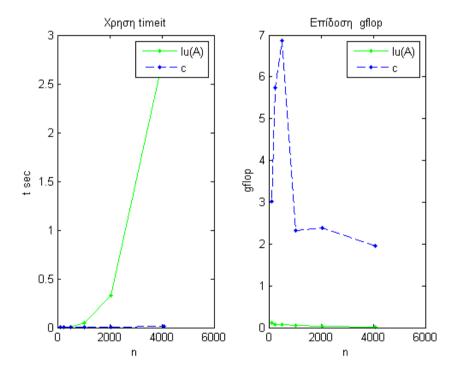
- (iii) Visualization for each of the above numerical results:
- (a) Use tic, toc functions to perform each operation once and gflop performance



## (b) Using tic, toc averaged functions for greater reliability and performance gflop



(c) Use timeit function and gflop performance



### Commenting on Results

- (1) Act lu requires considerably more operations to find the L and U registers from A \* b.
- (2) The lu act exponentially expands as the number of register elements increases, as opposed to A \* b which increases almost linearly at a low rate of change regardless of the value of n.
- (3) The timeit function takes so many measurements to be considered more reliable. In the above examples, the times from the two previous methods (tic, toc and averages) and timeit have a slight deviation.

### Question 3- Evaluating endogenous functions and m-functions

- (a) Using the timeit function, we make timing of mtimes for the above types of registers, size n = 2 ^ [7:12]
- (i) Random registers
- (ii) Registers having random data and being triangular
- (iii) Registers having random data and being upper triangular
- (iv) Registers with random data and is over Hessenberg

We run the following code:

```
clear all;
clc;
%Erwthma 3-(a)- (i)
k=1:
  for i=7:12
     n=2^i;
     %(i)random matrices
     Ai=rand(n);
     Bi=rand(n):
     f = (a()) mtimes(Ai, Bi);
     ti(k,1)=timeit(f,1);
     gflopi(k,1)=2*n^2/ti(k,1)/10^9;
     %(ii) tridiagwnio matrices
     upo diag=rand(n-1,1);
     kuria diag=rand(n,1);
     uper diag=rand(n-1,1);
     Aii=[diag(upo diag,-1)+diag(kuria diag)+diag(uper diag,1)];
     upo diag=rand(n-1,1);
     kuria diag=rand(n,1);
     uper diag=rand(n-1,1);
```

```
Bii=[diag(upo diag,-1)+diag(kuria diag)+diag(uper diag,1)];
   f = (a_i) mtimes(Aii, Bii);
   tii(k,1)=timeit(f,1);
   gflopii(k,1)=2*n^2/tii(k,1)/10^9;
   %(iii) upper triangular matrices
   Aiii=triu(rand(n));
   Biii=triu(rand(n));
   f = (a)() mtimes(Aiii, Biii);
   tiii(k,1)=timeit(f,1);
   gflopiii(k,1)=2*n^2/tiii(k,1)/10^9;
   %(iv) matrices
   Aiv = hess(rand(n));
   Biv= hess(rand(n));
   f = (a()) mtimes(Aiv, Biv);
   tiv(k,1)=timeit(f,1);
   gflopiv(k,1)=2*n^2/tiv(k,1)/10^9;
  k=k+1;
end;
n=2.^{7:12};
subplot(1,2,1)
plot(n, ti(:,1), 'y.-', n,tii(:,1),'r.--', n, tiii(:,1),'b.-', n, tiv(:,1), 'g.-');
legend('i', 'ii', 'iii', 'iv');
xlabel('n');
ylabel('t sec');
title('Question 3 (a)- mtimes');
subplot(1,2,2)
plot(n, gflopi(:,1), 'y.-', n, gflopii(:,1), 'r.--', n, gflopiii(:,1), 'b.-', n, gflopiv(:,1), 'g.-');
legend('i', 'ii', 'iii', 'iv');
xlabel('n');
ylabel('Gflop/s');
title('Question 3 (a)- mtimes');
```

H ti <6x1 double>		
	1	
1	3.4351e-04	
2	0.0025	
3	0.0218	
4	0.3898	
5	1.1398	
6	14.8938	

	1
1	3.2512e-04
2	0.0026
3	0.0553
4	0.4039
5	1.0837
6	13.7299
7	

tiii <6x1 double			
	1		
1	3.3893e-04		
2	0.0026		
3	0.0603		
4	0.4087		
5	1.6275		
6	15.8906		

	tiv ×						
	tiv <6x1 double>						
	1						
1	3.4116e-04						
2	0.0028						
3	0.0529						
4	0.3907						
5	1.0435						
6	14.4361						
7							

(b) I implement a myTridMult function to perform the multiplication between triangular registers which only takes into account the non-zero elements while ignoring the operations with zeros. By using timeit again, we make the timing of the function for the above triangular registers:

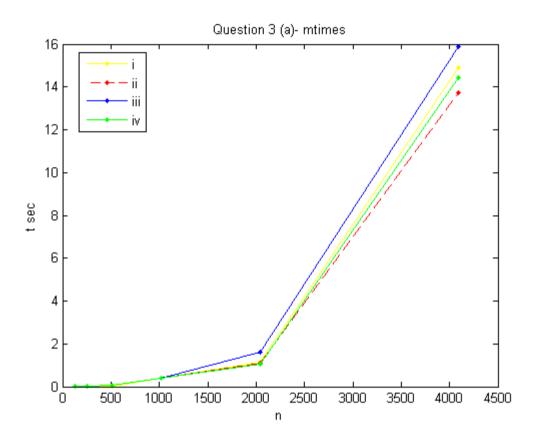
We run the following code:

```
\begin{aligned} & \text{function } [c] = \text{myTridMult}(A,B,n) \\ & \text{c=zeros}(n,n); \\ & \text{for } i = 1:n \\ & \text{for } j = 1:n \\ & \text{for } k = 1:n \\ & \text{if } A(i,k) \sim = 0 \parallel B(k,j) \sim = 0 \\ & \text{c}(i,j) = \text{c}(i,j) + A(i,k) * B(k,j); \\ & \text{end} \end{aligned}
```

```
clear all;
clc;
%Erwthma 3-(b)
k=1:
  for i=7:12
     n=2^i:
     %(ii) tridiagwnio matrices
     upo diag=rand(n-1,1);
     kuria diag=rand(n,1);
     uper diag=rand(n-1,1);
     Aii=[diag(upo diag,-1)+diag(kuria diag)+diag(uper diag,1)];
     upo diag=rand(n-1,1);
     kuria diag=rand(n,1);
     uper diag=rand(n-1,1);
     Bii=[diag(upo diag,-1)+diag(kuria_diag)+diag(uper_diag,1)];
     f = (a)() myTridMult(Aii, Bii,n);
     t myTridMult(k,1)=timeit(f,1);
     f = (a()) mtimes(Aii, Bii);
     t mtimes(k,1)=timeit(f,1);
         gflop(k,1)=2*n^2/t(k,1)/10^9;
    k=k+1;
  end;
  n=2.^{7:12};
  plot(n,t myTridMult(:,1),'r.--',n,t mtimes(:,1),'y.-');
  legend('t myTridMult','t mtimes');
  xlabel('n');
  ylabel('t sec');
  title('Question 3 (b)');
```

The Code while it ran did not end. Possible cause of this is to call the myTridMult function in which many forces are executed and for big n the algorithm drops to a large Loop.

# (c) <u>Visualization of results:</u>



## (d) Commentary on Results

From the times that emerged and from the graphical representation of these according to the size of the registers, we conclude that most of the time for the execution of the operation requires the triangular registers, then the Hessenberg registers, followed by the random registers and finally the upper triangles. The more peculiar features they have, the longer they require to carry out their actions.

We note more specifically that the Hessenberg registers are close to those of the random ones and the upper triangles with the triangles show large deviations as n grows. The preferred embodiment for the different sizes of the problem is that which makes the shortest time, that is, the fastest, even for a large number of n.

### **Query 4- Comparison of Embodiments**

In this question we will evaluate both the time, pace and actions of different ways of implementing the same computational problem.

### (1) Theoretical Calculation of Operations

```
We have the act x = (u * u '+ v * v') ^ P * b
```

For the multiplication between u and u 'we will need (2n-1) operations, as well as for the product v \* v', so I have 4n-2. In the end I will have to add the data in parentheses +1. So far I have (4n + 1). After I raise to force p I will have (4n + 1) \* p and finally + n by multiplication with vector b. I have therefore (4n + 1) \* p + n = 4np + p + n = n (4p + 1) + p operations = = 0 (n).

(2) We create the following function rank2\_power returning the result of operation x ,. U, v and b are random vectors

```
function x= rank2_power(u,v,b) x=(u^*u'+v^*v')^{\wedge}10^*b; end
```

We then run the following code:

```
clear all;
clc;
%Question 4

k=1;

for i=7:12
    n=2^i;
    u=rand(n,1);
```

```
v=rand(n,1);
    b=rand(n,1);
  f = (a()) \text{ rank2 power}(u,v,b);
  ta(k,1)=timeit(f,1);
  gflopa(k,1)=2*n^2/ta(k,1)/10^9;
  f = (a)() my rank2 power(u,v,b);
  tb(k,1)=timeit(f,1);
  gflopb(k,1)=2*n^2/tb(k,1)/10^9;
  k=k+1;
end
i=2.^{7:12};
subplot(2,2,1)
plot(i, ta(:,1), 'y.-');
legend('rank2 power');
xlabel('n');
ylabel('t sec');
title('Σύγκριση χρόνου');
i=2.^{7:12};
subplot(2,2,2)
plot(i, gflopa(:,1), 'g.-');
legend('rank2 power');
xlabel('n');
ylabel('gflopi');
title('Σύγκριση gflopi');
i=2.^{7:12};
subplot(2,2,3)
plot(i, tb(:,1), 'y.-');
legend('my rank2 power');
xlabel('n');
ylabel('t sec');
title('Σύγκριση χρόνου');
i=2.^{7:12};
subplot(2,2,4)
plot(i, gflopb(:,1), 'g.-');
legend('my rank2 power');
xlabel('n');
ylabel('gflopii');
title('Σύγκριση gflopii');
```

(3) Transformation into the given expression will be implemented as follows: We will make paralleling from right to left with b after first breaking the force that is raised in the

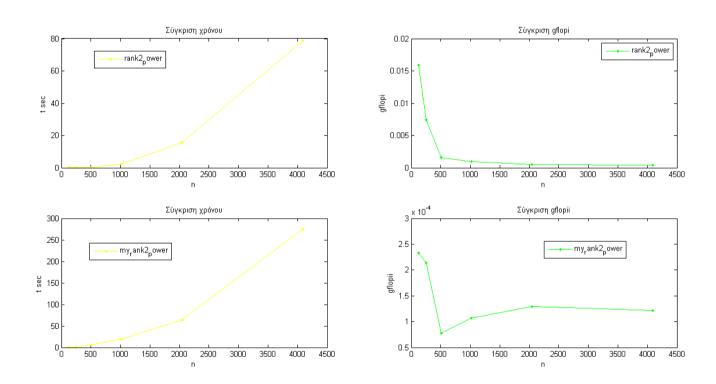
brackets.

```
function x= my_rank2_power(u,v,b)

x=(u*u'+v*v');
x= x*b;

i=1:10
x=(u*u'+v*v')*x;
end
```

- (4) The performance of the two functions for  $n = 2 ^ [7:12]$  in
  - (a) a common chart of execution times
  - (b) execution rates (in Gflop / s) in a common graph



### **Commenting on Results**

We notice that the gflop performance over time is countermetric as the table size increases.

### **Bibliography**

- [1] CLEVE B. MOLER. NUMERICAL METHODS WITH MATLAB. PUBLICATIONS KLIDARITHMOS LTD. 2010.
- [2] GILBERT STRANG. INTRODUCTION TO LINE ALGER. PUBLICATIONS OF PATRAS UNIVERSITY, 2006.
- [3] Efstratios Galopoulos. SCIENTIFIC CALCULATION I. University of Patras, Autumn 2013.