

Advanced Logic Homework 6

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1. Soundness and completeness

This exercise is about the soundness lemma for the normal modal logic K_σ . Consider the following tableau in K_σ , which contains only one branch which we call b :

$$1: \quad \neg(\Box p \supset \Box\Box p), 0$$

a) The branch is (open and) complete. Every rule that can be applied has been applied, including the symmetric (σ) restriction.

b) Branch b is *open*. Consider the interpretation $\mathcal{I} = \langle W, R, v \rangle$ with:

- $W = \{w_0, w_1\}$
- $R = \{\langle w_0, w_1 \rangle, \langle w_1, w_0 \rangle\}$
- $v_{w_0}(p) = 1$ and $v_{w_1}(p) = 0$

Moreover, consider partial function f from the natural numbers to W such that $f(1) = w_0$ and $f(0) = f(2) = w_1$.

Question: Does f show that \mathcal{I} is *faithful* to b ?

If so, show this step by step, considering each line from b . If not, provide a counterexample and show that it is one.

$$\begin{aligned} f(1) &= w_0 \\ f(0) &= f(2) = w_1 \end{aligned}$$

→ I have no idea how to interpret the partial functions

$$2) \quad \forall x \Box Px \wedge \Diamond Qa \vdash_{VK_p} \Diamond (Pb \wedge Qa)$$

$$\forall x \square P_x \wedge \Diamond Q_a, 0$$

$$\neg \Diamond (P_b \wedge Q_a), 0$$

$$\square \neg (P_b \wedge Q_a), 0$$

$$\forall x \square p_x, \text{ o}$$

$\diamond Q_a, 0$

oro

$$\neg(P_b \wedge Q_a), c$$

$$\neg P_6, 0$$

$$\neg Q_c, 0$$

or

Or 1

Qa.

Qa,

$$\neg (P_b \wedge Q_c), 1$$

$$\neg (P_b \wedge Q_a), 1$$

$\rightarrow P_{b,1}$

$$-Q_a,$$

$P_6, 1$

λ_a, τ

$$-e_{b,0}$$

$\square P_b, 0$

$P_{b,0}$

$$\neg \mathcal{E}_{b,c}$$

$\square P_{b,0}$

$P_b, 0$

There are two open and complete branches. → to the best of my knowledge. IDK what to do with that → E₆

There are two open and complete branches.