VADCOPs: Leveraging STARKs for Tailored Proof Generation

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Section 1

Introduction

Execution Trace, Transition Function and Arithmetization

Rules of the transition function:

- a starts at 1 and b at 7.
- The sum of a,b,c equals 2.d.
- The next value of c is d.
- Each two states, b is zero.
- ...

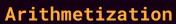
state	а	b	С	d
0	a ₀	b _ø	c ₀	d _ø
1	a ₁	b ₁	с ₁	d ₁
2	a ₂	b ₂	C ₂	d_2
3	a ₃	b ₃	c ₃	d ₃
4	a ₄	b ₄	C ₄	d ₄
				•••

Execution Trace

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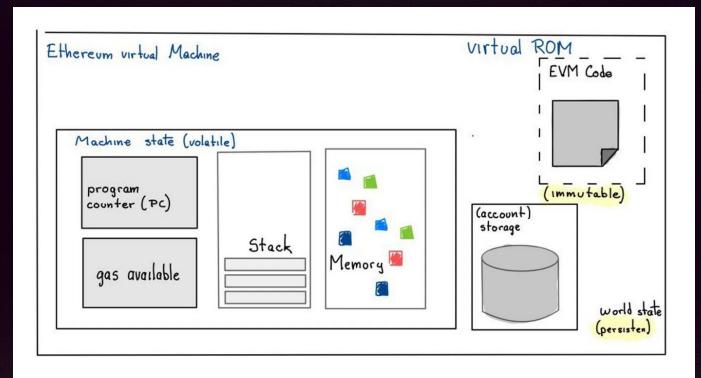


L1 * (a - 1) === 0 L1 * (b - 7) === 0
a + b + c === 2*d
c' === d (∑L2i) * b === 0

state	а	b	С	d
0	1	7	0	10
1	2	0	11	13
2	1	11	0	12
3	9	0	15	24
4	7	15	1	23

Execution Trace

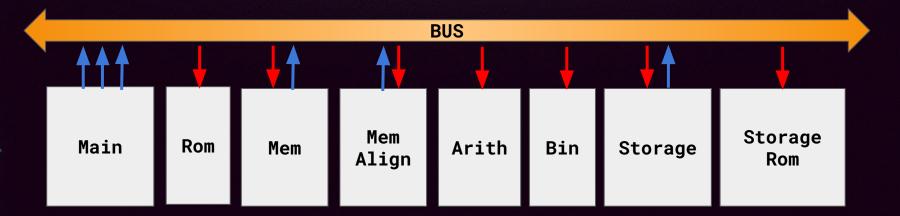
The EVM Arithmetization



The EVM is Too Complex to Arithmetize as a Whole



Unlocking Modularity: The BUS



We didn't know at that time, but we unlocked the BUS:

- The main EVM component (transaction processor) could now "assume" operations, throwing them to the bus.
- Specialized components could now "prove" the operations they get, removing them from the bus.

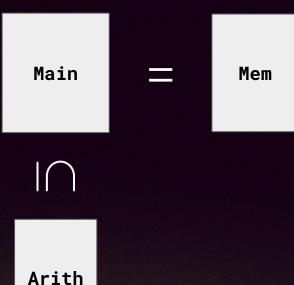
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• Permutations.

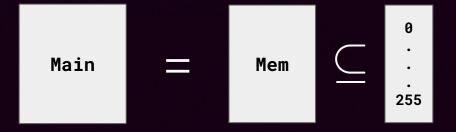
There is a variety of manners that one can throw to the BUS:

- Permutations.
- Lookups.



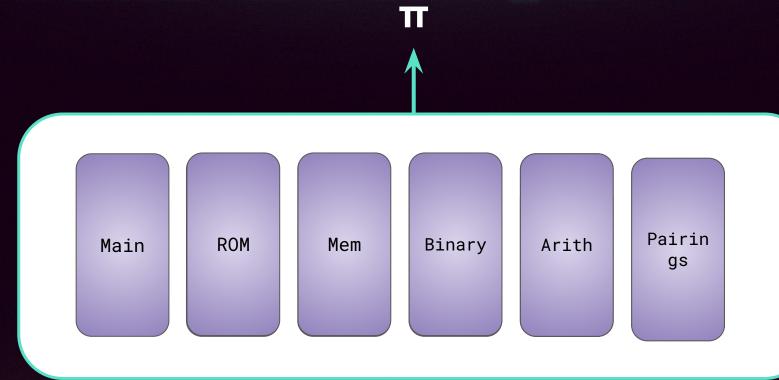
There is a variety of manners that one can throw to the BUS:

- Permutations.
- Lookups.
- Range Checks.
- Connections.
- ...

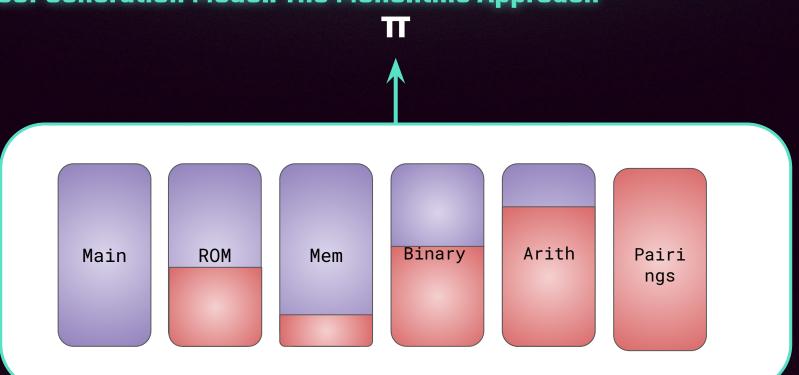




Proof Generation Model: The Monolithic Approach



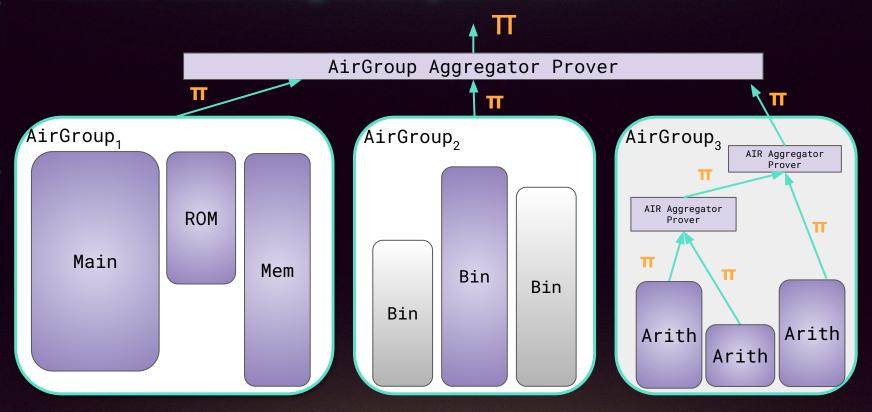
Proof Generation Model: The Monolithic Approach



Section 2

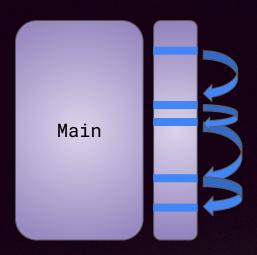
VADCOPs

VADCOPs: Variable Degree Composite Proofs



AirGroup Values and Global Constraints

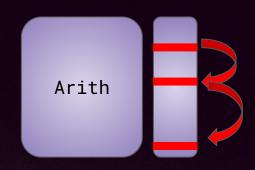
Main Auxiliary Trace



Global Constraint

$$V_{M} = = = V_{A}$$

Arith Auxiliary Trace



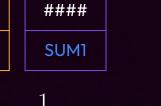
Arith AirGroup Value V_A

Main AirGroup Value V_м

Example: Multi-Domain Lookups

fl	f2	f3
8	1600	88
2	400	22
2	400	22
1	200	11

S1
####
####
####
SUM1



$$S_1' = S_1 \cdot (1 - L_4) + \frac{1}{f_1' + \alpha f_2' + \alpha^2 f_3' + \beta}$$

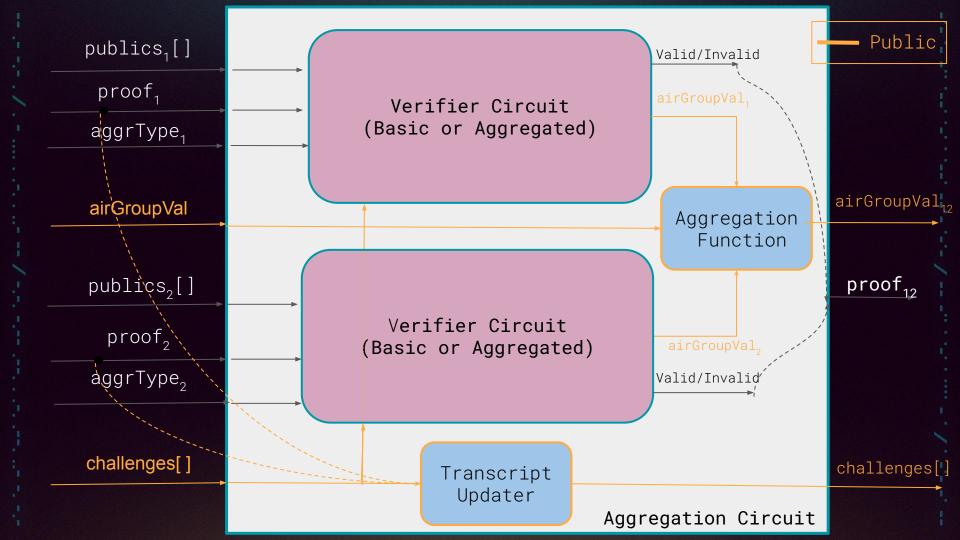
$$L_4 \cdot (S_1 - \text{SUM1}) = 0$$

Global Constraint

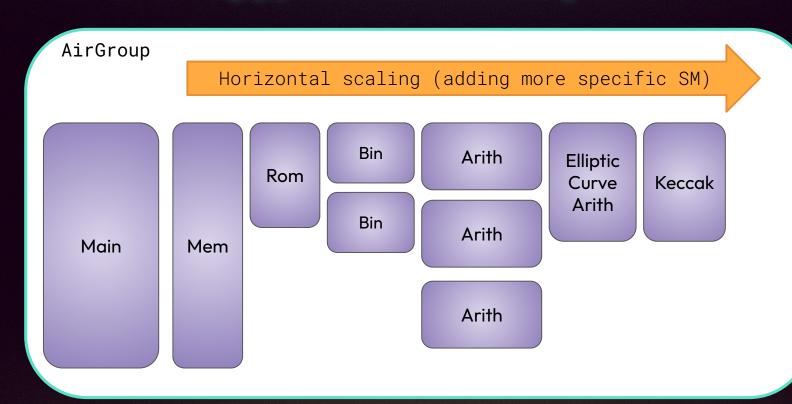
mul	t1	† 2	†3	S
1	1	200	11	##
2	2	400	22	##
0	3	600	33	##
0	4	800	44	##
0	5	1000	55	##
0	6	1200	66	##
0	7	1400	77	##
1	8	1600	88	SU

$$S_2' = S_2 \cdot (1 - L_8) + \frac{\text{mul}}{t_1' + \alpha t_2' + \alpha^2 t_3' + \beta}$$

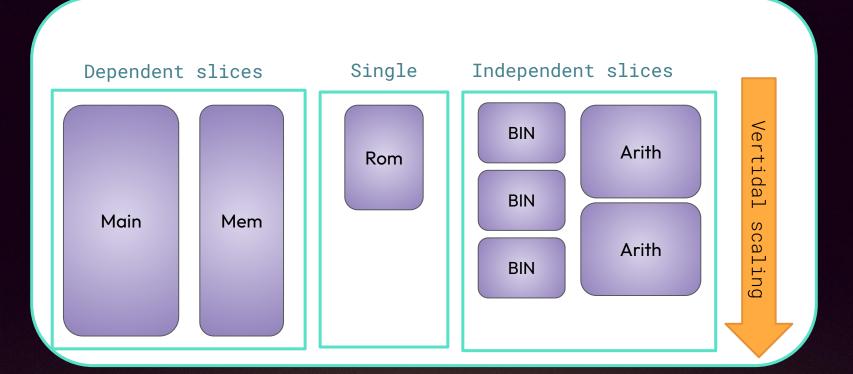
$$L_8 \cdot (S_2 - \text{SUM2}) = 0$$



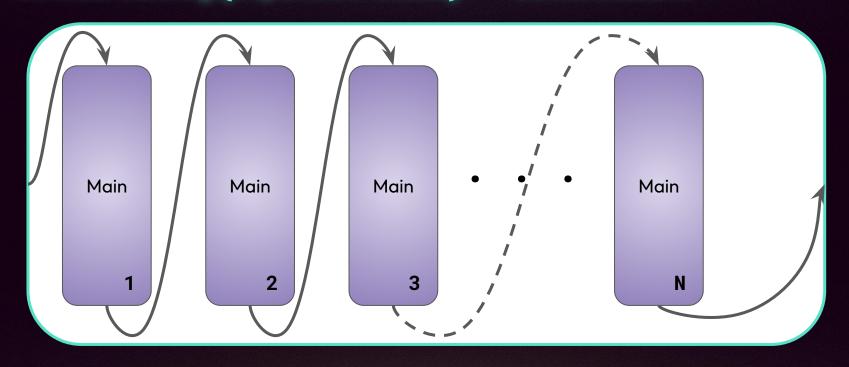
Horizontal Scaling (specific state machines)



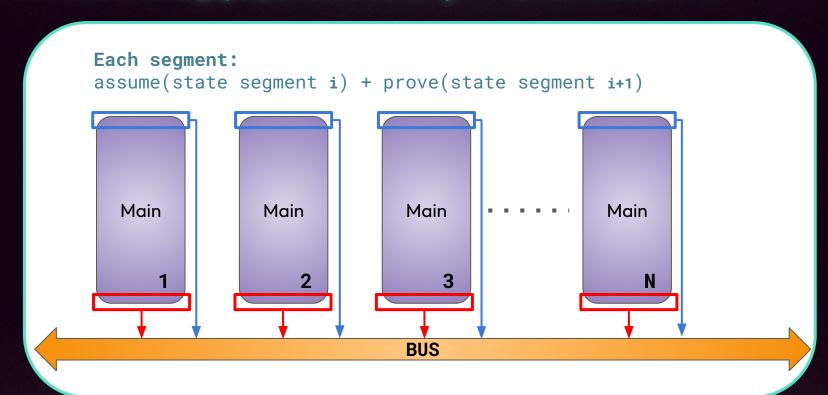
Vertical Scaling (slice trace, types)



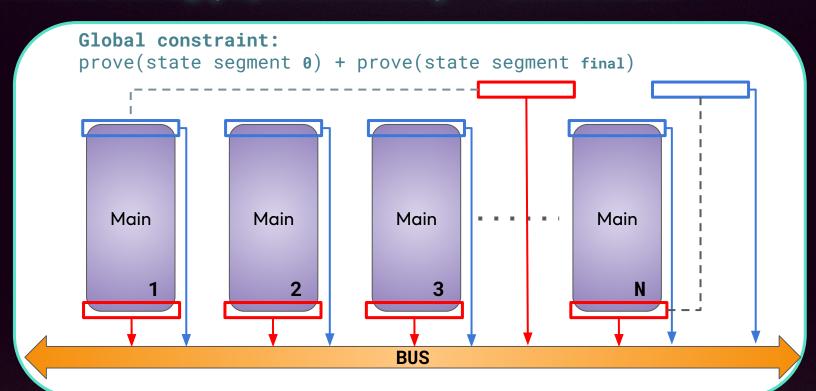
Vertical Scaling (dependent slices) ⇒ Continuations



Vertical Scaling (dependent slices) ⇒ Continuations



Vertical Scaling (dependent slices) ⇒ Continuations



Continuations

- Each segment assumes a prior state and proves the next, ensuring continuity.
- Initial and Final State Control, with global constraints set the initial state (e.g., segment_id = 0, pc = BOOT) and assume the final state, addressing multi-cycle security concerns.
- **Segment vs. Row State.** Each segment has two distinct states, which aren't row-based but rather segment/instance states. To avoid redundant columns, we use **air values** like segment_id and pc.
- Within each segment, constraints are applied to row values and air values, allowing flexible trace continuation beyond row boundaries.
- Unlimited Trace Continuation. With this approach, the trace is not limited by row boundaries.

A glimpse into the future... what's next will change the game

Stay tuned for our next steps

Thank you!

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