

# Ontinental 3

## LIN



- History and Expectation
- > Technical Features
- > The ISO/OSI reference model and LIN
- > Frames
- Message Frames
- Communication concept of LIN
- Command and Extended Frames
- Sleep mode and wake-up signal
- Error and Exception handling





### **History and Expectation**



## **Growth expectations**

- -3 ... 10 LIN nodes per vehicle
- 1.2 billion LIN nodes per year world wide



First specification Jul. 1994 draft released

The Ultimate Driving Machine



Time

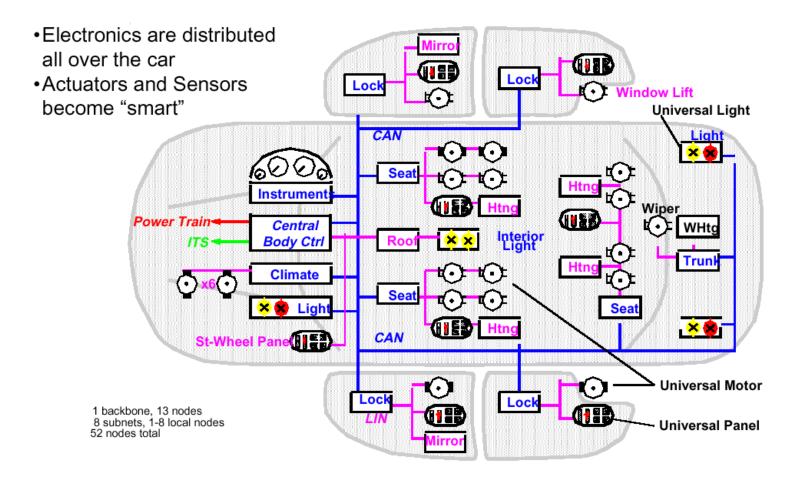
2001 - first application under development start in car production

1998 - LIN Consortium is initiated as a workgroup (5 car manufactures + 1 semiconductor supplier + 1 tool provider)





# Typical example for LIN in automobiles (1)







# Typical example for LIN in automobiles (2)

#### Roof:

(high amount of wiring)
Rain Sensor, Light
Sensor, Light Control,

Sun Roof ...

(Rain Sensor needs to be interrogated every 10-20ms)

#### •Door:

Mirror, Central ECU, Mirror Switch, Window Lift, Seat Control Switch,

Door Lock, etc.

### Climate:

many Small Motors Control Panel

#### **Steering Wheel:**

(very many controls are going to be positioned on the steering wheel)

Cruise Control, Wiper,

Turning Light, ...

Optional: Climate Control,

Radio, Telephone, etc.

#### **Seat:**

many Seat Position Motors, Occupancy Sensor, Control Panel

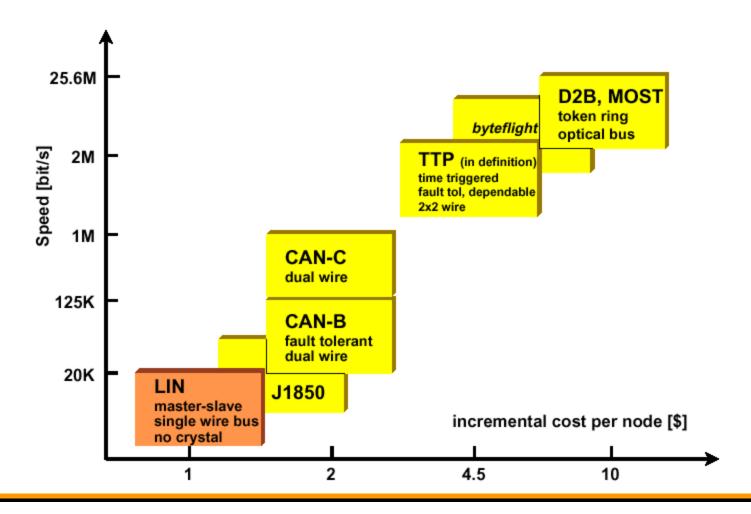
### **Engine:**

Sensors Small Motors





# **Automotive Bus Systems**







- History and Expectation
- > Technical Features
- The ISO/OSI reference model and LIN
- > Frames
- Message Frames
- Communication concept of LIN
- Command and Extended Frames
- Sleep mode and wake-up signal
- Error and Exception handling





## **Technical Features(1)**

- Sub-Bus as a extension to CAN to provide connection to local network clusters
- Low cost single-wire implementation (cheaper than CAN but does not have the same reliability level as CAN)
- Low cost silicon implementation based on common UART/SCI interface hardware (almost any microcontroller has necessary hardware on chip)
- Single Master / Multiple Slave concept

(no Arbitration is necessary)





# **Technical Features(2)**

Self synchronization without quartz or ceramics resonators in slave node

(significant cost reduction of hardware platform)

Speed up to 20kbit/s

limited by the EMI of single wire transmission. Recommended Bit Rates:

Slow: 2400 bit/sec

Medium: 9600bit/sec

Fast: 19200 bit/sec

- Guarantee of latency times for signal transmission
- Hot plug-in / plug-out





- History and Expectation
- > Technical Features
- The ISO/OSI reference model and LIN
- Frames
- Message Frames
- Communication concept of LIN
- Command and Extended Frames
- Sleep mode and wake-up signal
- Error and Exception handling





### The ISO/OSI reference model and LIN (1)

7 Application Layer

AII

**6** Presentation Layer

**People** 

5 Session Layer

Seem

4 Transport Layer

To

3 Network Layer

Need

2 Data Link Layer

**D**ata

1 Physical Layer

**Processing** 



7	Application Layer	Applications, operating system
6	Presentation Layer	Conversion of data formats
5	Session Layer	Task synchronization, buffers, connection setup and monitoring, access rights control
4	Transport Layer	Address conversion, routing, segmentation
3	Network Layer	Setup of logical connection, transport protocol
2	Data Link Layer	Transmission security, frame setup, error management
1	Physical Layer	Electrical / mechanical characteristics: Transmission medium, wiring, connectors, encoding, signals





### The ISO/OSI reference model and LIN (3)

# 2 Data Link Layer

### LLC- Logical Link Control

is concerning with Message Filtering and Recovery Management



is supervised by a management entity called Fault Confinement

Acceptance Filtering, Recovery Management, Time Base Synchronization, Message Validation



Data Encapsulation/Decapsulation, Error Detection, Error Signaling, Serialization/Deserialization

### **Physical Layer**



Bit Timing, Bit synchronization, Line Driver/Receiver





- History and Expectation
- > Technical Features
- > The ISO/OSI reference model and LIN
- Frames
- Message Frames
- Communication concept of LIN
- Command and Extended Frames
- Sleep mode and wake-up signal
- Error and Exception handling





## Frames (1)

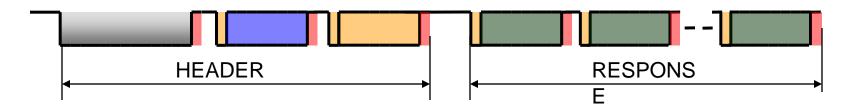
- > Frame: "Envelope" for transmission data
- > 3 different frame types :
  - → Message Frames: used for regular data transmission
  - → Command Frames: used for software updates, network configuration, and diagnostic purposes
  - → Extended Frames: are reserved to allow the embedding of user-defined message formats and future LIN formats into the current LIN protocol without violating the current LIN specification





# Frames (2)

> Frame: Formats



> Byte Fields Formats: transmission with LSB first





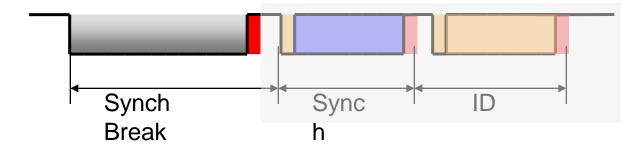


- > History and Expectation
- > Technical Features
- > The ISO/OSI reference model and LIN
- > Frames
- Message Frames
- Communication concept of LIN
- Command and Extended Frames
- Sleep mode and wake-up signal
- Error and Exception handling





## **Message Frames: Header(1)**

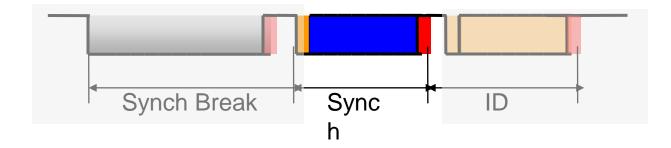


- > Synchronization break field
  - bidentify the beginning of a message frame
  - **⇔consists of 2 parts:** 
    - "dominant" bus value with a minimum duration of  $T_{SYNBRK}$  (13 bits)
    - "recessive" bus value with a minimum duration of  $T_{\text{SYNDEL}}$  (1bit)





# **Message Frames: Header(2)**



### > Synchronization field

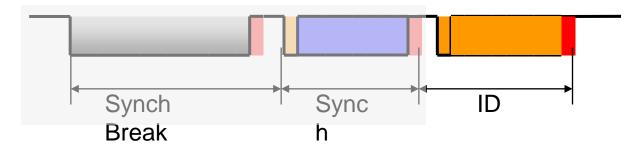
**\$\times\$** contains the information for clock synchronization

**⇔consists of pattern: "0x55"** 

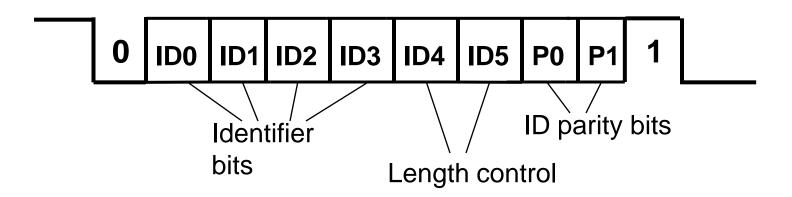




# **Message Frames: Header(3)**



➢ Identifier field: contains the content and length of message







### Message Frames: Header(4)



ID0 - ID5 IDENTIFIER  $\Rightarrow$  64 (26) identifiers divided in four subsets of 16 identifiers with 2, 4, 8 data fields.

ID5	ID4	N <sub>DATA</sub> (number of data fields) [byte]
0	0	2
0	1	2
1	0	4
1	1	8

PO - P1 are the parity check bits of identifier

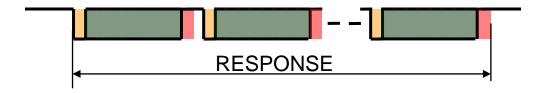
P0 =  $ID0 \oplus ID1 \otimes ID2 \oplus ID4$  (even parity)

P1 =  $\overline{\text{ID1}} \oplus \overline{\text{ID3}} \otimes \overline{\text{ID4}} \oplus \overline{\text{ID5}}$  (odd parity)





### Message Frames: Response Fields



- Contains the data field and the checksum
  - **♦ DATA field consists of BYTE fields**
  - **♦ CHECKSUM** field contains the inverted modulo-256 sum over all data bytes calculated by "add with carry"

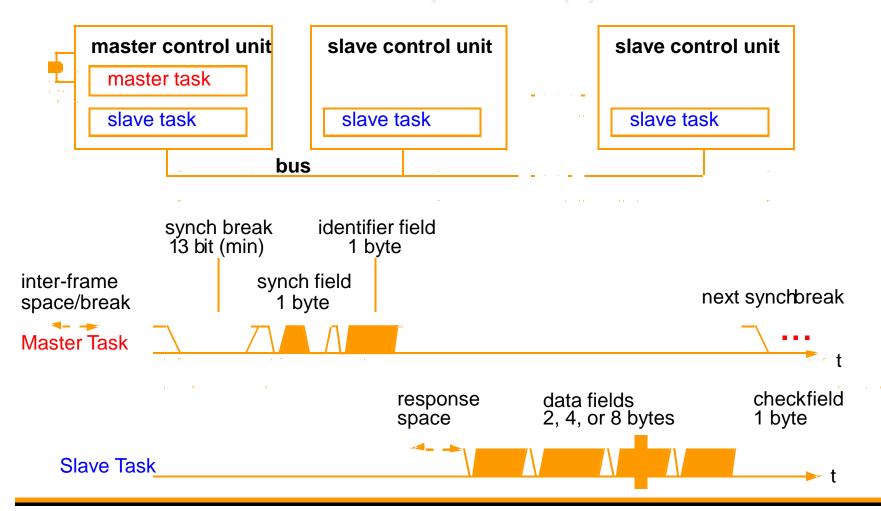


- **≻**History and Expectation
- > Technical Features
- The ISO/OSI reference model and LIN
- > Frames
- Message Frames
- Communication concept of LIN
- Command and Extended Frames
- Sleep mode and wake-up signal
- Error and Exception handling





# **Communication Concept of LIN(1)**







# Communication Concept of LIN(2) Master Task

- Has control the whole Bus and Protocol:
  - \$sends the header of a message
  - monitors data bytes and check sum byte and evaluates them on consistency
  - **∜receives WakeUpBreak from a slave node** when the bus is inactive and they request some action
  - \$\serves as a reference with it's clock base
    (stable clock necessary)





# Communication Concept of LIN(3) Slave Task

➤ Is one of 2-16 members on a bus and receives or transmits data when the appropriate ID is sent by the master:

waits for Synch Break

**synchronize** on Synch Byte

⋄ snoops for ID. According to ID slave determine what to do: receive data, transmit data or do nothing

when transmitting sends 2, 4, or 8 data bytes and the checksum byte





- History and Expectation
- > Technical Features
- > The ISO/OSI reference model and LIN
- > Frames
- Message Frames
- Communication concept of LIN
- Command and Extended Frames
- Sleep mode and wake-up signal
- Error and Exception handling





### **Command and Extended Frames: Command Frames**

- > are used to broadcast general command requests for service purposes from the master to all bus participants:
  - **♦ ID = 0x3C(0x3C)** "Master request frame" to send commands and data from the master to the slave node
  - ⇔ ID = 0x3D(0x7F) "Slave response frame" that triggers one node to send data to the master node
- First data byte of a command frame containing the values from 0x00 to 0x7F are reserved:
  - **♦ SLEEP MODE command used to broadcast the sleep mode to all bus nodes: ID = 0x3C and DATA(0)=0x00**





### **Command and Extended Frames: Extended Frames**

➤ are reserved to allow the embedding of user-defined message formats and future LIN formats into LIN protocol without violating the current LIN specification:

 $\Rightarrow$  ID = 0x3E(0xFE) user defined extended frame

 $\Rightarrow$  ID = 0x3F(0xBF) future LIN extension

>the identifier can be followed by an arbitrary number of LIN bytes fieldc





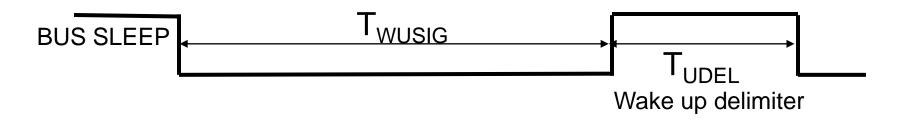
- **≻**History and Expectation
- > Technical Features
- > The ISO/OSI reference model and LIN
- > Frames
- Message Frames
- Communication concept of LIN
- Command and Extended Frames
- Sleep mode and wake-up signal
- Error and Exception handling





## Sleep mode and wake-up signal (1)

- Sleep mode on the bus is determine by:
  - **♦** a sleep command frame
  - ⇔ a time-out (T<sub>TIME\_OUT</sub> = 25000 T<sub>BIT</sub>) on slave nodes in case the sleep command message was corrupted
- ➤ Sleep mode on the bus can be terminated by any node (slave task) by sending a Wake-Up Signal (detected by the master as a valid character 0x80, 0xC0, 0x00):c







# Sleep mode and wake-up signal (2)

### Wake up procedure:

♦ After detecting wake-up signal all nodes run through the Start-up procedure and wait for the master task to send a synch break field followed by a synch field

 $\$  If no synch field is detected before "Time out after wake up" ( $T_{TOBRK} \le 128 \ T_{BIT}$ ) a new wake-up signal is issued by the node requesting first wake-up.

 $\$  This sequence is issued not more than 3 times then the transmission is suspended for a "Time out after three breaks" ( $T_{T3BRK} \ge 5000 T_{BIT}$ )

**♦ The re-transmission of the wake-up signal has to be decided by the application** 





- > History and Expectation
- > Technical Features
- > The ISO/OSI reference model and LIN
- Frames
- Message Frames
- Communication concept of LIN
- Command and Extended Frames
- Sleep mode and wake-up signal
- Error and Exception handling





## **Error and exception handling**

- Possible error types:
  - **♥** Bit Error
  - **♦ Checksum Error**
  - **♦ Identifier Parity Error**
  - **♦ Slave Not Responding Error**





# Error and exception handling(1) Bit Errors

### > Error description

The bit actually appearing on the bus is different that the one transmitted

### Method of detection

Sending unit monitors the bus while transmitting. A BIT\_ERROR has to be detected at that bit time.

### > Fault Confinement

This error is detected by - master task in master

- slave task in slave

while reading back its own transmission





# **Error and exception handling(2)**

### **Checksum Errors**

### > Error description

The inverted modulo-256 sum over all received data bytes does not match with the receive checksum byte

#### Method of detection

The sum of the inverted modulo-256 sum over all received data bytes and the checksum byte does not result in 0xFF

### Fault Confinement

This error is detected by - slave task in master when expecting or reading data from the bus

- slave task in slave while reading

from the bus





# Error and exception handling(3) Identifier-Parity Errors

### > Error description

The parity identifier bits does not match with the correct calculated values

### Method of detection

Typical LIN slave application do not distinguish between an unknown but valid identifier and a corrupted identifier

### Fault Confinement

This error is detected by - master task in master while reading back its own transmission

slave task in slave while reading

from the bus





# Error and exception handling(4) Slave-Not-Responding Errors

### > Error description

The message frame is not fully completed within maximum length

### > Method of detection

A slave task waits the entire message upon transmission of the new header

### > Fault Confinement

This error is detected by - slave task in master when expecting or reading data from the bus

- slave task in slave while reading from the bus only when a slave expects a message from another slave





# **Error and exception handling(5)**

### **Inconsistent-Synch-Field Errors**

> Error description

Synch field is different than the pattern 0x55

Method of detection

Slave task detects edges of Synch field outside the given tolerance

> Fault Confinement

This error is detected only by slave task in slave





# **Error and exception handling(6)**

### Causes for message errors

- Local Disturbance of Ground Potential
- **➤ Local Disturbance of Supply Voltage**
- Global Electric Disturbance of the Bus Signal
- > Unsynchronized Time Base

