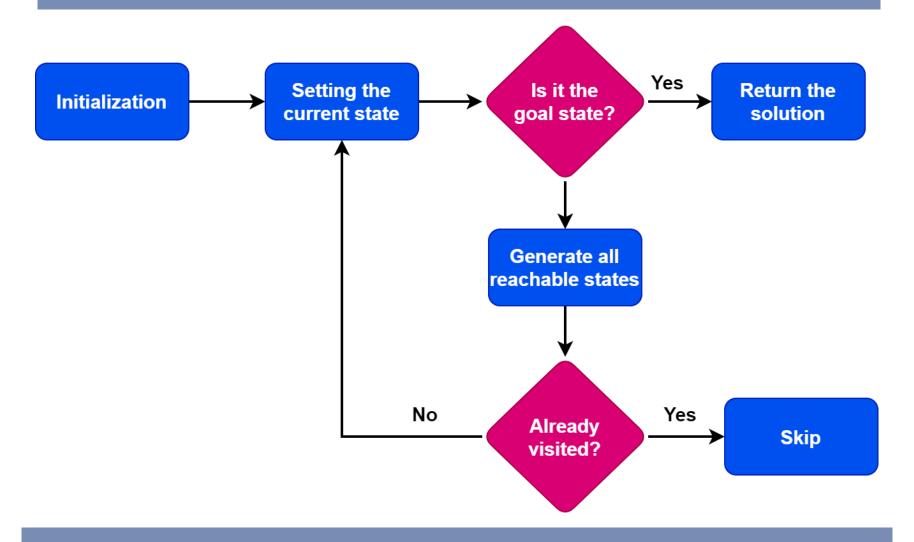




AI Planning Hatem A. Rashwan

State-Space Search

State-Space Search Overview



Overview

- Heuristic Search Strategies
- The A* Algorithm (for Tree Search)
- Graph Search with A*
- (Good) Heuristics

Informed vs uniformed search

| Parameters | Informed Search | Uninformed Search |
|-----------------|---|---|
| Known as | It is also known as Heuristic Search. | It is also known as Blind Search. |
| Using Knowledge | It uses knowledge for the searching process. | It doesn't use knowledge for the searching process. |
| Performance | It finds a solution more quickly. | It finds solution slow as compared to an informed search. |
| Completion | It may or may not be complete. | It is always complete. |
| Cost Factor | Cost is low. | Cost is high. |
| Time | It consumes less time because of quick searching. | It consumes moderate time because of slow searching. |

Heuristic Functions

- heuristic function h: state space $\rightarrow \mathbb{R}$
- h(n) = estimated cost of the cheapest path from node n to a goal node
- if n is a goal node then h(n) must be 0
- heuristic function encodes problem-specific knowledge in a problem-independent way

Best Frist Search

- An instance of the general tree search (or graph) search algorithm
 - strategy: select next node based on an <u>evaluation</u> function f: state space → ℝ
 - select node with lowest value f(n)
- Implementation: selectFrom(fringe, strategy)
 - <u>priority queue</u>: maintains fringe in ascending order of f-values

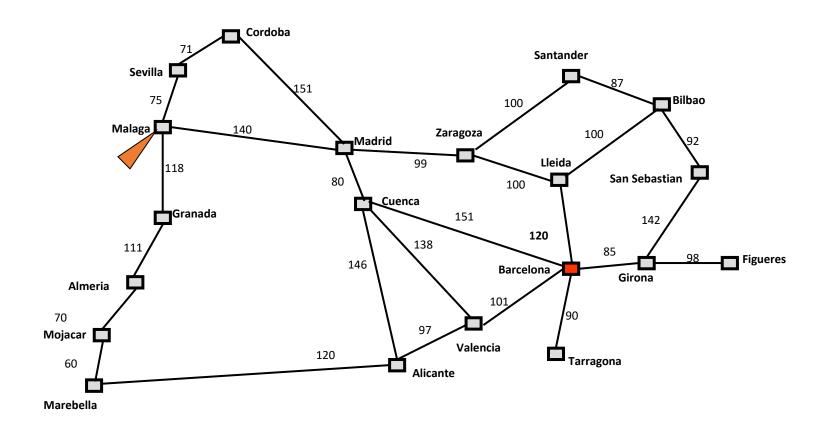
Greedy Best-First Search

 use heuristic function as evaluation function:

$$f(n) = h(n)$$

- always expands the node that is closest to the goal node
- eats the largest chunk out of the remaining distance, hence, "greedy"

Real-World Problem: Touring in Spain

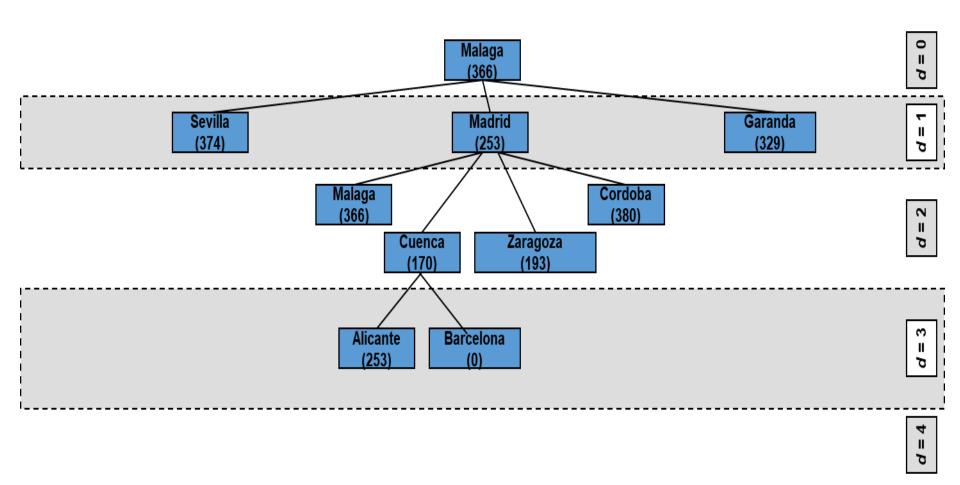


Real-World Problem: Touring in Spain

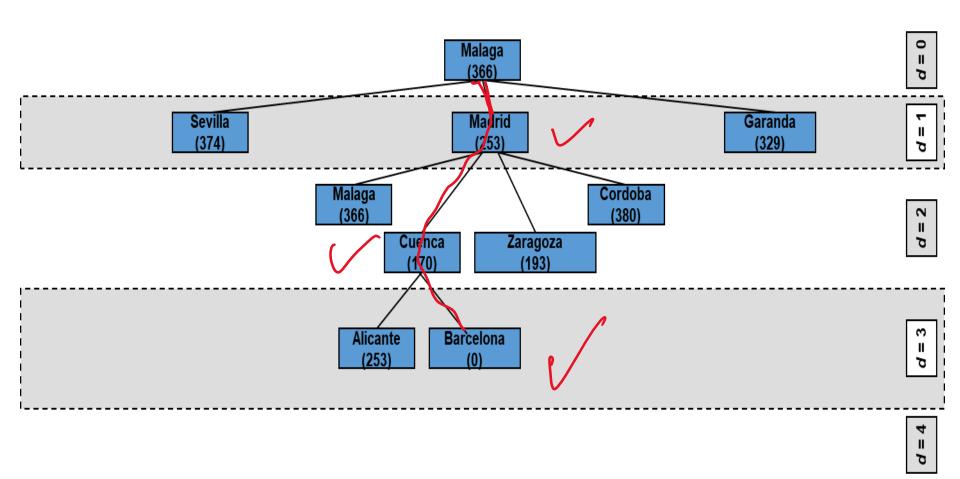
• $h_{SLD}(n)$ = straight-line distance to Barcelona

| Malaga | 366 | Girona | 101 | |
|-----------|-----|---------|-----|--|
| Barcelona | 0 | Bilbo | 226 | |
| Valencia | 87 | Almeria | 340 | |
| Alicante | 198 | Madrid | 253 | |
| Figueres | 130 | | | |
| Zaragoza | 170 | | | |
| Tarragona | 80 | | | |

Greedy Best-First Search: Touring Romania



Greedy Best-First Search: Touring Romania



Overview

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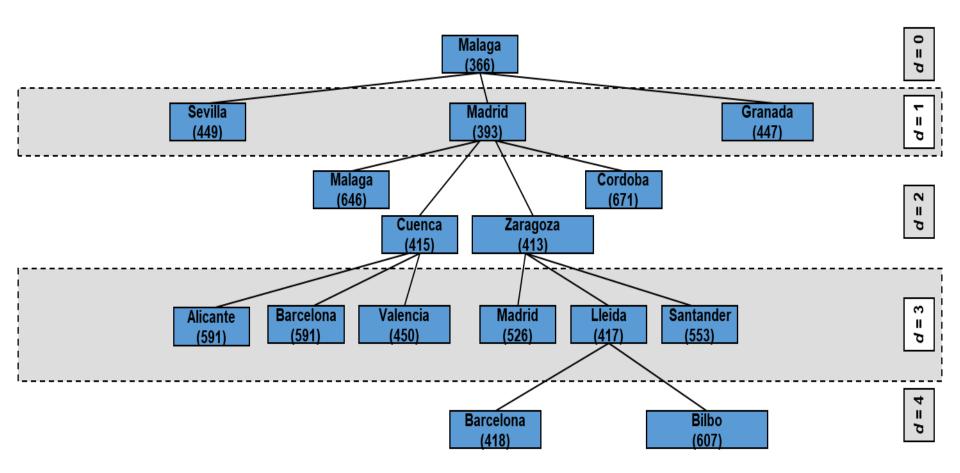
A* Search

best-first search where

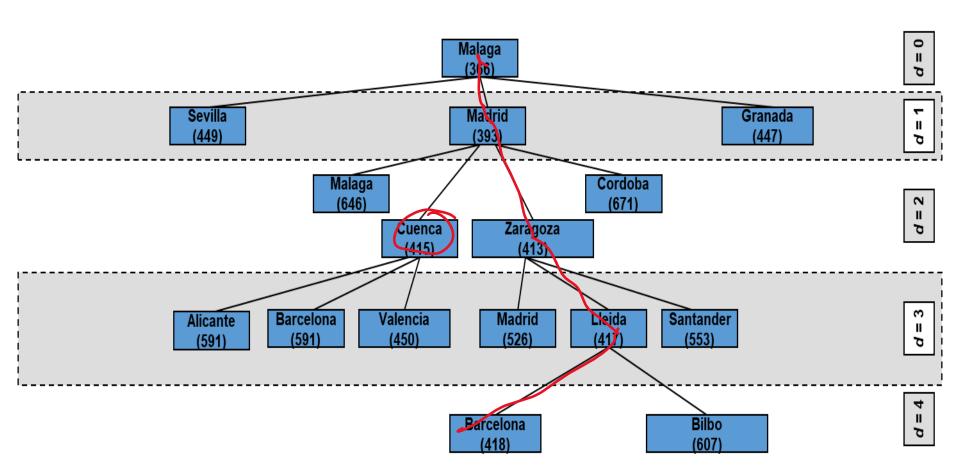
$$f(n) = h(n) + g(n)$$

- h(n) the heuristic function (as before)
- g(n) the cost to reach the node n
- evaluation function:
 - f(n) = estimated cost of the cheapest solution through n
- A* search is optimal if h(n) is admissible

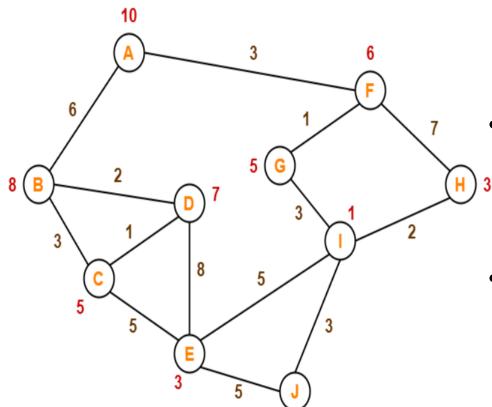
A* (Best-First) Search: Touring Spain



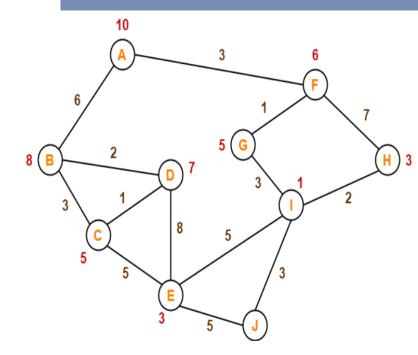
A* (Best-First) Search: Touring Spain



Consider the following graph-



- The numbers written on edges represent the distance between the nodes.
 - The numbers written on nodes represent the heuristic value.
- Find the most cost-effective path to reach from start state
 A to final state J using A*
 Algorithm.



Step-01:

We start with node A.

• Node B and Node F can be reached from node A.

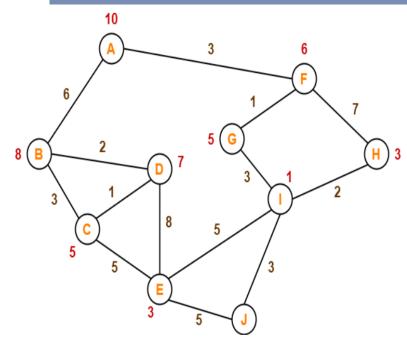
A* Algorithm calculates f(B) and f(F).

•
$$f(B) = 6 + 8 = 14$$

•
$$f(F) = 3 + 6 = 9$$

Since f(F) < f(B), so it decides to go to node F.

Path- $A \rightarrow F$



Step-02:

Node G and Node H can be reached from node F.

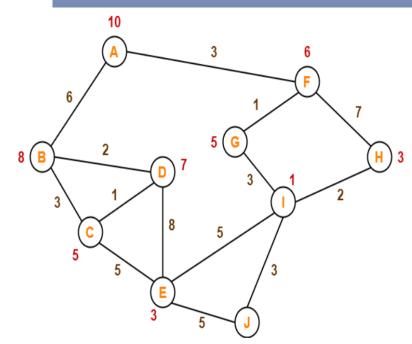
A* Algorithm calculates f(G) and f(H).

$$f(G) = (3+1) + 5 = 9$$

 $f(H) = (3+7) + 3 = 13$

Since f(G) < f(H), so it decides to go to node G.

Path- $A \rightarrow F \rightarrow G$



Step-03:

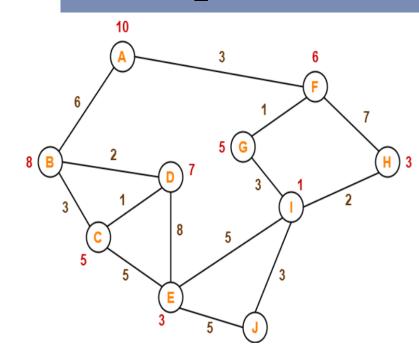
Node I can be reached from node G.

A* Algorithm calculates f(I).

$$f(1) = (3+1+3) + 1 = 8$$

It decides to go to node I.

Path- A \rightarrow F \rightarrow G \rightarrow I



Step-04:

Node E, Node H and Node J can be (H)3 reached from node I.

A* Algorithm calculates f(E), f(H) and f(J).

$$\bullet f(E) = (3+1+3+5) + 3 = 15$$

$$\bullet f(H) = (3+1+3+2) + 3 = 12$$

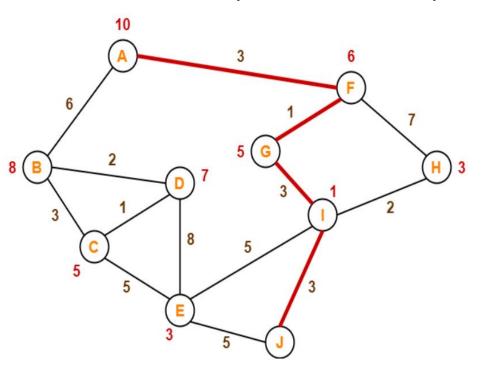
$$\bullet f(J) = (3+1+3+3) + 0 = 10$$

Since f(J) is least, so it decides to go to node J.

Final:.

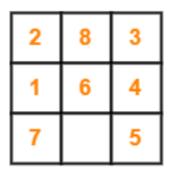
Path- A \rightarrow F \rightarrow G \rightarrow I \rightarrow J

This is the required shortest path from node A to node J.



Important Note-

- A* Algorithm is one of the best path finding algorithms.
- But it does not produce the shortest path always.
- This is because it heavily depends on heuristics.



| 1 | 2 | 3 |
|---|---|---|
| 8 | | 4 |
| 7 | 6 | 5 |

Given an initial state of a 8-puzzle problem and final state to be reached-

Initial State

Final State

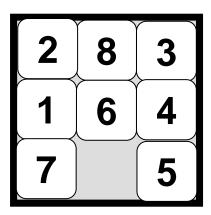
Find the most cost-effective path to reach the final state from initial state using A^* Algorithm. Consider g(n) = Depth of node and h(n) = Number of misplaced tiles.

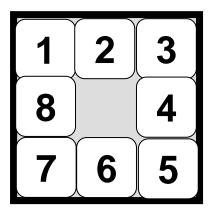
Solution-

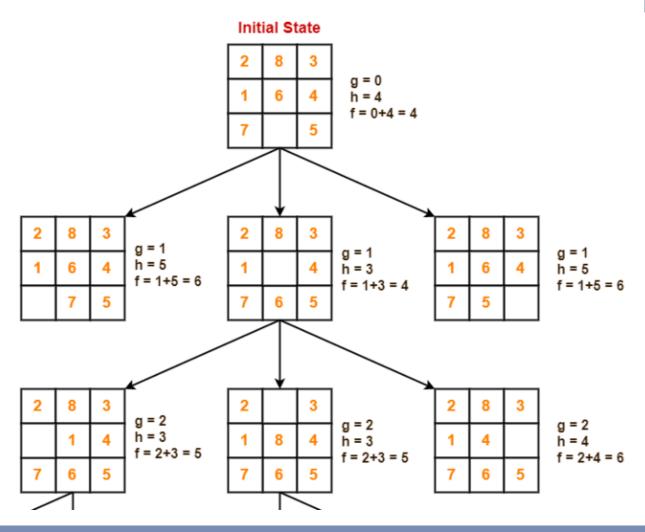
- •A* Algorithm maintains a tree of paths originating at the initial state.
- •It extends those paths one edge at a time.
- •It continues until final state is reached.

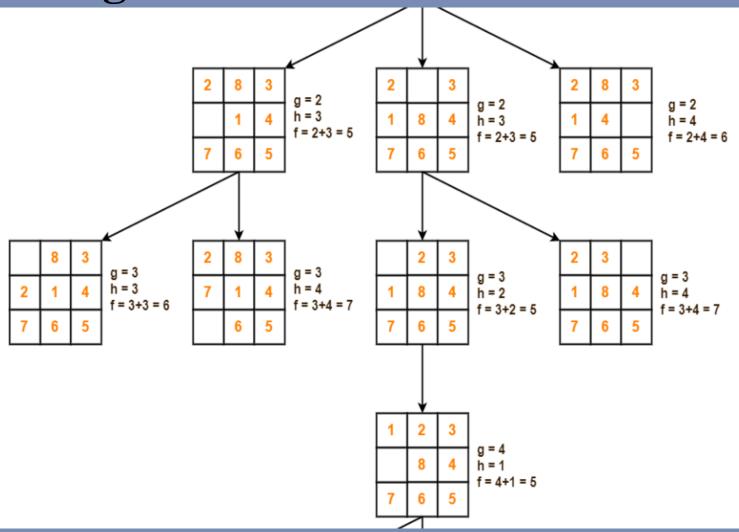
- h_1 : number of misplaced tiles
- h_2 : Manhattan block distance
 - Example:
 - h1 = 44 tiles are misplaced
 - Or h2 = 2+1+1+1 = 5
 - g = depth that initially equals 0

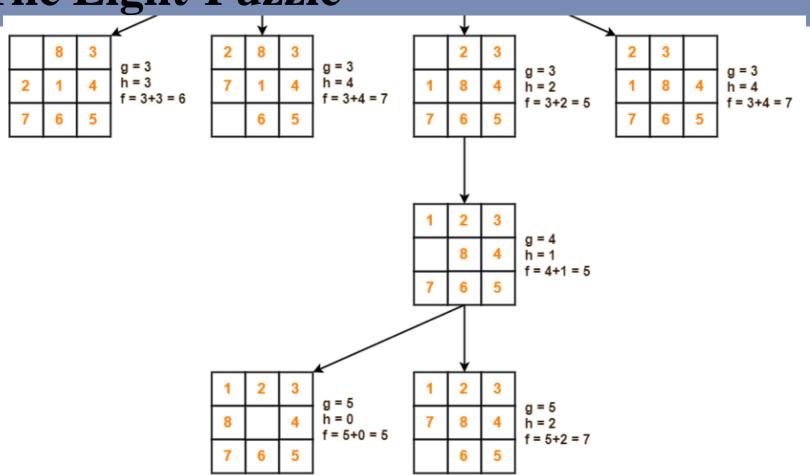
No movement yet











Final State

A* (Best-First) Search: A* search properties

- Admissible: If a solution exists for the given problem, the first solution found by A* is an optimal solution.
- Complete: A* is also complete algorithm, meaning if a solution exists, the answer is bound to be found in a finite amount of time.
- Optimal: A* is optimally efficient for a given heuristic-of the optimal search algorithm that expand search paths from the root node, it can be shown that no other optimal algorithm will expand fewer nodes and find a solution.

A* (Best-First) Search: Advantages/disadvantages

Advantages:

- It is optimal search algorithm in terms of heuristics.
- It is one of the best heuristic search techniques in Al.
- It is used to solve complex search problems.
- There is no other optimal algorithm guaranteed to expand fewer nodes than A*.

Disadvantages:

- This algorithm is complete if the branching factor is finite and every action has fixed cost.
- The performance of A* search is dependent on accuracy of heuristic algorithm used to compute the function h(n).

Overview

- Heuristic Search Strategies
- The A* Algorithm (for Tree Search)
- Graph Search with A*
- (Good) Heuristics

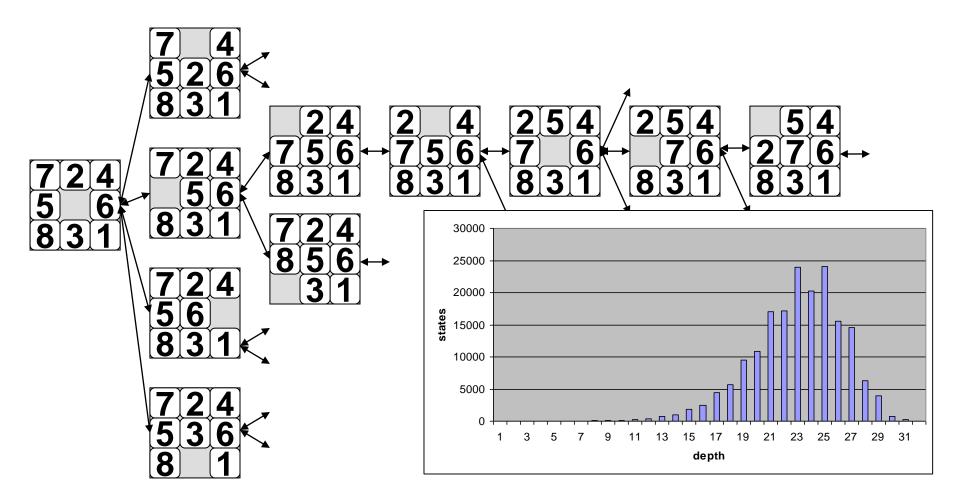
A* Tree/Graph Search Algorithm

```
function aStarTreeSearch(problem, h)
 fringe \leftarrow priorityQueue(new searchNode(problem.initialState))
 allNodes ← hashTable(fringe)
 loop
    if empty(fringe) then return failure
    node \leftarrow selectFrom(fringe)
    if problem.goalTest(node.state) then
       return pathTo(node)
    for successor in expand(problem, node)
       if not allNodes.contains(successor) then
         fringe ← fringe + successor @ f(successor)
         allNodes.add(successor)
```

A* and Exponential Space

- A* has worst case time and space complexity of $O(b^l)$
- exponential growth of the fringe is normal
 - exponential time complexity may be acceptable
 - exponential space complexity will exhaust any computer's resources all too quickly

The Eight-Puzzle Search Space



Overview

- Heuristic Search Strategies
- The A* Algorithm (for Tree Search)
- Properties of A*
- Graph Search with A*
- (Good) Heuristics

Heuristics

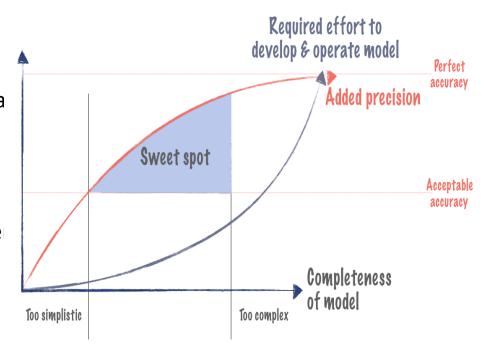
- heuristics are criteria, methods, or principles for deciding which among several alternative courses of action promises to be the most effective in order to achieve some goal
- example: you need to select a course to learn
 - Review the materials of each course
 - Use this information to decide which of the courses that were available you want to take.
 - So you used heuristic function about the course to make this choice.

Good Heuristics

- good heuristics
 - indicate a way to reduce the number of states that need to be evaluated
 - help obtain solutions within reasonable time constraints
- trade-off:
 - simplicity
 - must provide a simple means of discriminating between choices
 - accuracy
 - no guarantee that they identify the best course of actions
 - but they should do so sufficiently often

Good Heuristics

- trade-off:
 - simplicity
 - must provide a simple mea of discriminating between choices
 - accuracy
 - no guarantee that they ide the best course of actions
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How can we find good heuristics for a given problem?

Heuristics from Simplified Problems How to find a good heuristic function?

- relaxed problem: a problem with fewer restrictions on the actions than the original problem
- The cost of an optimal solution for a relaxed problem is an admissible and consistent heuristic for the original problem.

8-Puzzle Actions

- a tile can move from square A to square B if A is horizontally or vertically adjacent to B and B is blank
- Relaxed conditions:
 - a tile can move from square A to square B if A is adjacent to B (⇒ Manhattan distance)
 - a tile can move from square A to square B if B is blank
 - a tile can move from square A to square B
 (⇒ misplaced tiles)

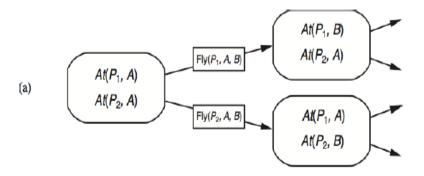
State Space Search

Classical Representations "Stripes"

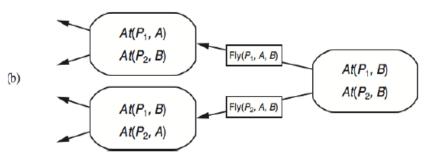
Classical Planning Generate a plan

> Planning Algorithms

Progression: Forward state-space search



Regression: Backward state-space search



Properties of Planning Algorithms

Soundness

- A planning algorithm is sound if all solutions are legal plans
 - All preconditions, goals, and any additional constraints are satisfied

Completeness

- A planning algorithm is complete if a solution can be found whenever one actually exists
- A planning algorithm is *strictly complete* if all solutions are included in the search space

Optimality

 A planning algorithm is optimal if it maximizes a predefined measure of plan quality

Linear Planning

■ Basic Idea – Goal **stack**

- Work on one goal until completely solved before moving on to the next goal
- Search by reducing the difference between the state and the goals
- What means (operators) are available to achieve the desired ends (goal)

Linear Planning

- (Linear Planner) (initial-state, goals)
 - state = initial-state; plan = []; stack = [] Push goals on stack
 - Repeat until stack is empty
 - If top of stack is goal that matches state, then pop stack
 - Else if top of stack is a conjunctive goal g, then
 - Select an ordering for the subgoals of g, and push them on stack
 - Else if top of stack is a simple goal sg, then
 - Choose an operator o whose add-list matches goal sg
 - Replace goal sg with operator o
 - Push the preconditions of o on the stack
 - Else if top of stack is an operator o, then
 - state = apply(o, state)
 - plan = [plan; o]

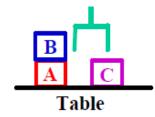
Means-ends Analysis in Linear Planning

(Newell and Simon)

GPS Algorithm (state, goals, plan)

- If $goals \subseteq state$, then return (state, plan) (The symbol \subseteq is used to denote containment of sets.)
- Choose a difference $d \in goals$ between *state* and *goals*
- Choose an operator o to reduce the difference d
- If no applicable operators, then return False
- (state, plan) = GPS (state, preconditions(o), plan)
- If state, then return **GPS** (apply (o, state), goals, [plan, o])

Initial call: GPS (initial-state, initial-goals, [])



Pickup_from_table(b)

Pre: Block(b), Handempty Clear(b), On(b, Table)

Add: Holding(b)

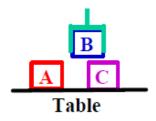
Delete: Handempty, On(b, Table)

Putdown_on_table(b)

Pre: Block(b), Holding(b)

Add: Handempty, On(b, Table)

Delete: Holding(b)



Pickup_from_block(b, c)

Pre: Block(b), Handempty Clear(b), On(b, c), Block(c)

Add: Holding(b), Clear(c)

Delete: Handempty, On(b, c)

Putdown_on_block(b, c)

Pre: Block(b), Holding(b)

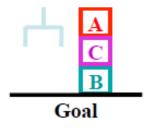
Block(c), Clear(c), $b \neq c$

Add: Handempty, On(b, c)

Delete: Holding(b), Clear(c)

1. Search Stack

On(A, C) On(C, B)



3. Search Stack

On(A, C) On(C, B)

On(A, C)

Put Block(C, B)

Holding(C) Clear(B)

State

Clear(B)
Clear(C)
On(C, A)
On(A, Table)
On(B, Table)
Handempty

C A B

Initial State

State

Clear(B)
Clear(C)
On(C, A)
On(A, Table)
On(B, Table)
Handempty

2. Search Stack

On(A, C) On(C, B)
On(A, C)

On(C, B)

State

Clear(B) Clear(C)

On(C, A)

On(A, Table)

On(B, Table)

Handempty

4. Search Stack

On(A, C) On(C, B)

On(A, C)

Put_Block(C, B)

Holding(C) Clear(B)

Holding(C)

Clear(B)

State

Clear(B)

Clear(C)

On(C, A)

On(A, Table)

On(B, Table)

Handempty

5. Search Stack

On(A, C) On(C, B)

On(A, C)

Put_Block(C, B)

Holding(C) Clear(B)

Holding(C)

State

Clear(B)

Clear(C)

On(C, A)

On(A, Table)

On(B, Table)

Handempty

6. Search Stack

On(A, C) On(C, B)

On(A, C)

Put_Block(C, B)

Holding(C) Clear(B)

Pick Block(C)

State

Clear(B)

Clear(C)

On(C, A)

On(A, Table)

On(B, Table)

Handempty

Handempty Clear(C) On(C, ?b)

7. Search Stack

On(A, C) On(C, B)

On(A, C)

Put_Block(C, B)

Holding(C) Clear(B)

Pick_Block(C)

State

Clear(B)

Clear(C)

On(C,A)

On(A, Table)

On(B, Table)

Handempty

8. Search Stack

On(A, C) On(C, B)

On(A, C)

Put_Block(C, B)

Holding(C) Clear(B)

State

Clear(B)

Clear(C)

On(A, Table)

On(B, Table)

Holding(C)

Clear(A)

[Pick_Block(C)]

9. Search Stack

On(A, C) On(C, B)

On(A, C)

Put Block(C, B)

State

Clear(B)

Clear(C)

On(A, Table)

On(B, Table)

Holding(C)

Clear(A)

10. Search Stack

On(A, C) On(C, B)

On(A, C)

State

Clear(C)

On(A, Table)

On(B, Table)

Clear(A)

Handempty

On(C, B)

[Pick Block(C)]

On(A, C) On(C, B)

11. Search Stack

Put Block(A, C)

Holding(A) Clear(C)

State

Clear(C)

On(A, Table)

On(B, Table)

Clear(A)

Handempty

On(C, B)

12. Search Stack

On(A, C) On(C, B)

[Pick Block(C); Put Block(C, B)]

Put Block(A, C)

Holding(A) Clear(C)

Holding(A)

Clear(C)

[Pick Block(C) Put Block(C, B)]

State

Clear(C)

On(A, Table)

On(B, Table)

Clear(A)

Handempty

On(C, B)

[Pick Block(C) Put Block(C, B)]

13. Search Stack

On(A, C) On(C, B)

Put Block(A, C)

Holding(A) Clear(C)

Holding(A)

State

Clear(C)

On(A, Table)

On(B, Table)

Clear(A)

Handempty

On(C, B)

14. Search Stack

On(A, C) On(C, B)

Put Block(A, C)

Holding(A) Clear(C)

Pick Table(A)

Handempty Clear(A)
On(A, Table)

State

Clear(C)

On(A, Table)

On(B, Table)

Clear(A)

Handempty

On(C, B)

[Pick_Block(C); Put Block(C, B)]

15. Search Stack

On(A, C) On(C, B)

Put Block(A, C)

Holding(A) Clear(C)

Pick_Table(A)

State

Clear(C)

On(A, Table)

On(B, Table)

Clear(A)

Handempty

On(C, B)

[Pick_Block(C); Put_Block(C, B)]

16. Search Stack

On(A, C) On(C, B)

Put_Block(A, C)

Holding(A) Clear(C)

State

Clear(C)

On(B, Table)

Clear(A)

On(C, B)

Holding(A)

[Pick_Block(C); Put Block(C, B)] [Pick_Block(C); Put_Block(C, B); Pick_Table(A)]

17. Search Stack

On(A, C) On(C, B)

Put_Block(A, C)

[Pick_Block(C); Put_Block(C, B); Pick Table(A)]

State

Clear(C) On(B, Table) Clear(A)

On(C, B)

Holding(A)

18. Search Stack

On(A, C) On(C, B)

[Pick_Block(C); Put_Block(C, B); Pick_Table(A); Put_Block(A, C)]

State

On(B, Table) Clear(A)

On(C, B)

Handempty

On(A, C)

19. Search Stack

[Pick_Block(C);

Put_Block(C, B);
Pick Table(A);

Put_Block(A, C)]

State

On(B, Table)

Clear(A)

On(C, B)

Handempty

On(A, C)

Linear Planning

- Sound?
- Optimal?
- Complete?

Pickup_from_Table (?b)

Pre: (Armempty)

(clear ?b)

(on-table ?b)

Add: (holding ?b)

Delete: (Armempty)

(on-table ?b)

(clear ?b)

Putdown_on_Table(?b)

Pre: (holding ?b)

Add: (Armempty)

(on-table ?b)

Delete: (holding ?b)

The Sussman Anomaly

С

Α

В

A

В

C

Pickup_fom_block (?a, ?b)

Pre: (Armempty)

(clear ?a) (on ?a ?b)

Add: (holding ?a) (clear ?b)

Delete: (Armempty)

(on ?a ?b) (clear ?a)

Putdown_on_block (?a, ?b)

Pre: (holding ?a) (clear ?b)

Add: (Armempty)

(on ?a ?b)

Delete: (holding ?a)

(clear ?b)

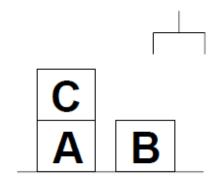
The Sussman Anomaly

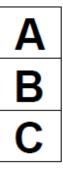
Linear Solution:

- (on B, C)
 - *Pickup* (*B*)
 - Stack (B, C)
- (on A, B)
 - *Unstack* (*B*, *C*)
 - Putdown (B)
 - *Unstack* (*C*, *A*)
 - Putdown (C)
 - Stack (A, B)
- (on B, C)
 - *Unstack* (*A*, *B*)
 - Putdown (A)
 - *Pickup* (*B*)
 - Stack (**B**, **C**)
- (on A, B)
 - Pickup(A)
 - *Stack* (*A*,*B*)

Linear Solution:

- (on A, B)
 - *Unstack* (*C*, *A*)
 - Putdown (C)
 - Stack (A, B)
- (on B, C)
 - *Unstack* (*A*, *B*)
 - Putdown (A)
 - *Pickup* (*B*)
 - Stack (**B**, **C**)
- (on A, B)
 - *Pickup* (*A*)
 - *Stack* (*A*,*B*)





Linear Planning – Goal Stack

✓ Advantages

- Reduced search space, since goals are solved one at a time,
 and not all possible goal orderings are considered
- Advantageous if goals are (mainly) independent
- Linear planning is sound

Disadvantages

- Linear planning may produce *suboptimal* solutions (based on the number of operators in the plan)
- Planner's efficiency is sensitive to goal orderings
 - Control knowledge for the "right" ordering
 - Random restarts
 - Iterative deepening
- Completeness?

Incompleteness of Linear Planning

Initial state:

Goal statement:

(at obj1 locA)
(at obj2 locA)
(at ROCKET locA)
(has-fuel ROCKET)

(and (at obj1 locB) (at obj2 locB))

INCOMPLETE EXAMPLE?

| Goal | Plan |
|----------------|---|
| (at obj1 locB) | (LOAD-ROCKET obj1 locA) (MOVE-ROCKET) (UNLOAD-ROCKET obj1 locB) |
| (at obj2 locB) | failure |

(OPERATOR LOAD-

ROCKET

:preconds

?roc ROCKET

?obj OBJECT

?loc LOCATION

(and (at ?obj ?loc)

(at ?roc ?loc))

:effects

add (inside ?obj ?roc)

del (at ?obj ?loc))

(OPERATOR UNLOAD-

ROCKET

:preconds

?roc ROCKET

?obj OBJECT

?loc LOCATION

(and (inside ?obj ?roc)

(at ?roc ?loc))

:effects

add (at ?obj ?loc)

del (inside ?obj ?roc))

(OPERATOR MOVE-

ROCKET

:preconds

?roc ROCKET

?from-l LOCATION

?to-1 LOCATION

(and (at ?roc ?from-l)

(has-fuel ?roc))

:effects

add (at ?roc ?to-l)

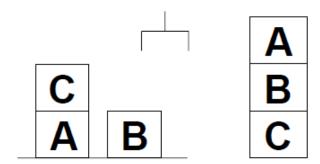
del (at ?roc ?from-l)

del (has-fuel ?roc))

The Sussman Anomaly-Nonlinear planning

NonLinear Solution:

- (on A B)
 - Unstack (C, A)
 - Putdown (C)
- (on **B C**)
 - *Pickup (B)*
 - Stack (B, C)
- (on A B)
 - Pickup (A)
 - Stack (A, B)



- Nonlinear planning is used to a goal set instead of a goal stack,
 and
- It is included in the search space of all possible subgoal orderings.
- It handles the goal interactions by interleaving method.

The Sussman Anomaly- Nonlinear planning

✓ Advantage

Non-linear planning may be an optimal solution with respect to plan length (depending on search strategy used).

Disadvantages

- It takes larger search space, since all possible goal orderings are taken into consideration.
- Complex algorithm to understand.

Why is Planning Hard?

Planning involves a complex search:

- Alternative operators to achieve a goal
- Multiple goals that interact
- Solution optimality, quality
- Planning efficiency, soundness, completeness

State representation

- The frame problem
- The "choice" of predicates
 - On-table (x), On (x, table), On-table-A, On-table-B,...

Action representation

– Many alternative definitions

Reduce to "needed" definition

Conditional effects

Uncertainty

- Quantification
- Functions
 - Generation planning algorithm(S)

Summary

- ☐ **Planning:** selecting one sequence of actions (operators) that transform (apply to) an initial state to a final state where the goal statement is true.
- ☐ Means-ends analysis: identify and reduce, as soon as possible, differences between state and goals.
- ☐ Linear planning: backward chaining with means-ends analysis using a stack of goals potentially efficient, possibly suboptimal, incomplete; GPS, STRIPS.
- □ Nonlinear planning with means-ends analysis: backward chaining using a set of goals; reason about *when* "to reduce the differences;" **Prodigy4.0**.

End



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