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# PACO: Laboratorio 4

# Índice:

1. Before starting lab assignment	2
2. Task decomposition analysis for Mergesort	2
2.1. Divide and conquer	2
2.2. Task decomposition analysis with Tareador	2
3. Shared-memory parallelisation with OpenMP tasks	8
3.1. Leaf strategy in OpenMP	8
3.2. Tree strategy in OpenMP	10
3.3. Optimization: Cut-off mechanism	12
4. Shared-memory parallelisation with OpenMP task using dependencies	14

- 1. Before starting lab assignment
- 2. Task decomposition analysis for Mergesort
  - 2.1. Divide and conquer

**For the deliverable**: Take note of the times reported for the sequential execution and use them as reference times to check the scalability of the parallel versions in OpenMP you will develop.

## 2.2. Task decomposition analysis with Tareador

**multisort-tareador.c** is already prepared to insert Tareador instrumentation. Complete the instrumentation to understand the potential parallelism that each one of the two recursive task decomposition strategies (leaf and tree) provide when applied to the sort and merge phases:

- In the leaf strategy you should define a task for the invocations of basicsort and basicmerge once the recursive divide—and—conquer decomposition stops.
- In the tree strategy you should define tasks during the recursive decomposition, i.e. when invoking multisort and merge.

**For the deliverable:** Include both tareador codes (in the .zip file) and show code excepts in the report.

#### Leaf strategy:

Definir las tareas del tareador en el basicmerge y el basicsort.

```
vold merge(long n, T left[n], T right[n], T result[n*2], long start, long length) {
    if (length < MIN_MERGE_SIZE*2L) {
        // Base case
            tareador_start_task("basicmerge");
            basicmerge(n, left, right, result, start, length);
            tareador_end_task("basicmerge");
    } else {
        // Recursive decomposition
            merge(n, left, right, result, start, length/2);
            merge(n, left, right, result, start + length/2, length/2);
    }
}

void multisort(long n, T data[n], T tmp[n]) {
    if (n >= MIN_SORT_SIZE*4L) {
            // Recursive decomposition
            multisort(n/4L, &data[0], &tmp[0]);
            multisort(n/4L, &data[0], &tmp[n/4L]);
            multisort(n/4L, &data[n/4L], &tmp[n/4L]);
            multisort(n/4L, &data[n/4L], &tmp[n/4L]);

            merge(n/4L, &data[0], &data[n/4L], &tmp[0], 0, n/2L);
            merge(n/4L, &data[0], &data[3L*n/4L], &tmp[n/2L], 0, n/2L);

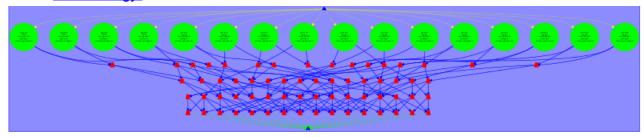
            merge(n/4L, &data[0], &tmp[n/2L], &data[0], 0, n);
} else {
            // Base case
            tareador_start_task("basicsort");
            basicsort(n, data);
            tareador_end_task("basicsort");
}
}
```

#### Tree strategy:

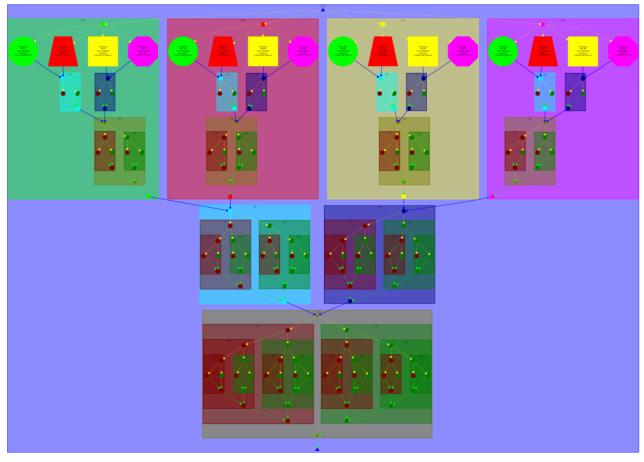
Definir las tareas del tareador en el merge y el multisort.

**For the deliverable:** From the task dependence graphs that are generated for leaf and tree, do you observe any major differences in terms of structure, types, number and granularity of tasks, ...? Save the image of the two graphs in order to include them in the report.

### Leaf strategy:



## Tree strategy:



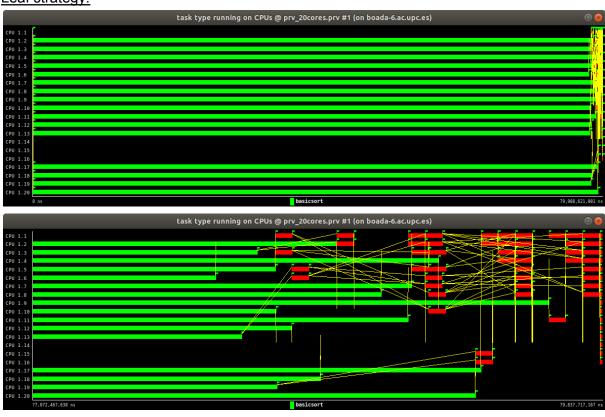
Como podemos observar los dos gráficos son muy diferentes para cada estrategia. La estrategia del árbol tiene muchas mas tareas generadas porque hay muchas mas invocaciones para el basicsort y el mergesort, por lo contrario la estrategia de hoja genera solo tareas en el caso base de las funciones de merge y el multisort, no como el caso recursivo de la estrategia del árbol.

**For the deliverable:** Identify the points of the code where you should include synchronizations to guarantee the dependences

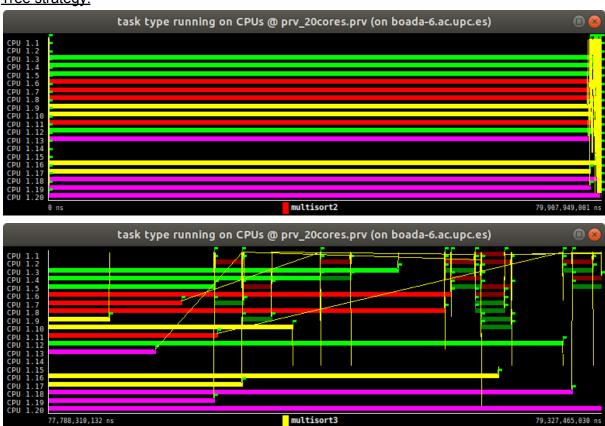
En las dos estrategias, la función merge crea dependencias porque es la función que ordena las dos partes del vector. original, mientras que en la función multisort como no depende de estas partes del vector no crea estas dependencias. Y podemos ver que la función multisort necesita mas tiempo porque hay muchas más tareas creadas.

**For the deliverable:** Capture the necessary Paraver windows in order to support your explanations in the deliverable.

# Leaf strategy:



# Tree strategy:



## 3. Shared-memory parallelisation with OpenMP tasks

#### 3.1. Leaf strategy in OpenMP

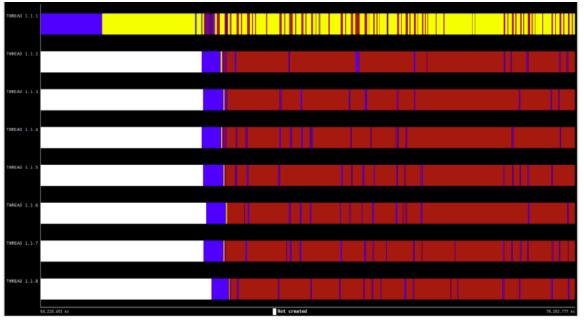
#### Text files:

# speedup-global:

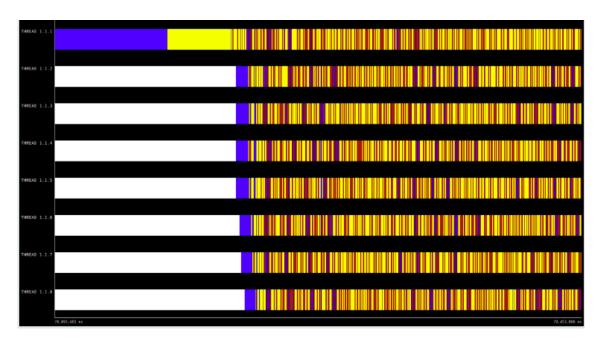
```
paco1208@boada-8:~/lab4$ cat speedup-global-multisort-omp-32768-1024-50-boada-11.txt 1.00340715502555366269
    .98659966499162479061
    .98659966499162479061
    .98659966499162479061
    .98659966499162479061
    .98494983277591973244
    .98659966499162479061
    .98659966499162479061
    .98659966499162479061
10
11
12
13
14
     .98494983277591973244
     .98659966499162479061
     .98659966499162479061
     .98659966499162479061
     .98659966499162479061
15
16
     .98659966499162479061
     .98659966499162479061
17
     .98659966499162479061
18
     .98659966499162479061
     .98659966499162479061
19
      .98659966499162479061
20
paco1208@boada-8:~/lab4$
```

#### speedup partial:

```
paco1208@boada-8:~/lab4$ cat speedup-partial-multisort-omp-32768-1024-50-boada-11.txt
1 1.00830189427418099218
    .98888954552386762238
    .98913109595354516001
    .98950902616324540527
    .98917636442929254116
    .98721702042463897998
    .98906980149653627586
    .98929238381344851418
    .98896401570577317759
10
     .98833617825310176912
11
12
13
14
     .98868090022942009149
     .98923786074591995465
     .98847799375888946648
     .98935615890215112129
15
     .98888011971824273445
16
     .98932294954591591329
17
     .98917391228071513619
18
     .98942805724863597218
19
     .98907866514815151669
20
      .98859234212987288212
paco1208@boada-8:~/lab4$
```



#### 3.2. Tree strategy in OpenMP



Mientras la versión leaf es implementada creando tascas en los casos base, y la versión del árbol solo las crea antes de las llamadas del multisort y merge functions.

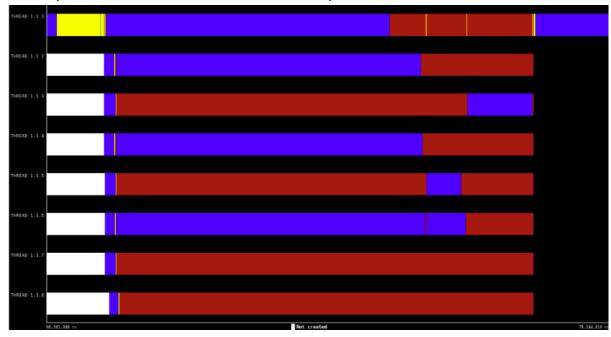
Mientras que en el tiempo de ejecución no hay mucha diferencia, pero la versión árbol es más eficiente.

Podemos observar que las dos versiones tienen problemas de escalabilidad, però la versión árbol mejora considerablemente bien el speedup con la función multisort que puede seguir más las escalabilidad.

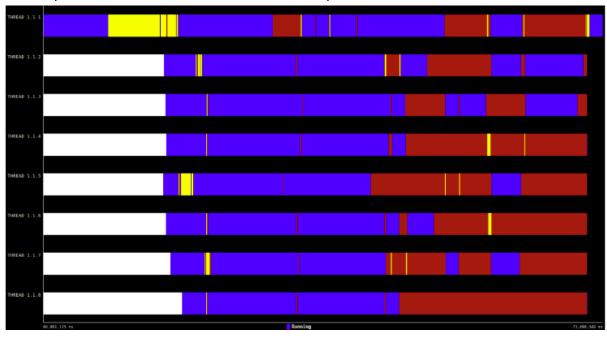
## 3.3. Optimization: Cut-off mechanism

```
GNU nano 6.2
                                                                                    multisort-omp.c
                         g n, 1
                                      left[n], T right[n], T
_MERGE_SIZE*2L) {
                                                                                  result[n*2], long start, long length,int d) {
        merge(lor
      if (length <</pre>
              // Base case
             basicmerge(n, left, right, result, start, length);
             // Recursive decomposition
if(!omp_in_final())
             #pragma omp task final (d>=CUTOFF)
             merge(n, left, right, result, start, length/2,d+1);
#pragma omp task final (d>=CUTOFF)
             merge(n, left, right, result, start + length/2, length/2,d+1);
#pragma omp taskwait
             }else{
             merge(n, left, right, result, start, length/2,d+1);
merge(n, left, right, result, start + length/2, length/2,d+1);
/oid multisort(long n, T data[n], T tmp[n], int d) {
     if (n >= MIN_SORT_SIZE*4L) {
    // Recursive decomposition
    if(!omp_in_final())
            #pragma omp task final (d>=CUTOFF)
multisort(n/4L, &data[0], &tmp[0],d+1);
#pragma omp task final (d>=CUTOFF)
multisort(n/4L, &data[n/4L], &tmp[n/4L],d+1);
#pragma omp task final (d>=CUTOFF)
multisort(n/4L, &data[n/2L], &tmp[n/2L],d+1);
#pragma omp task final (d>=CUTOFF)
multisort(n/4L, &data[3L*n/4L], &tmp[3L*n/4L],d+1);
#pragma omp taskwait
             #pragma omp task final (d>=CUTOFF)
             #pragma omp task final (d>=CUTOFF)
             merge(n/4L, &data[0], &data[n/4L], &tmp[0], 0, n/2L,d+1);
#pragma omp task final (d>=CUTOFF)
             merge(n/4L, &data[n/2L], &data[3L*n/4L], &tmp[n/2L], 0, n/2L,d+1);
             #pragma omp taskwait
             #pragma omp task final (d>=CUTOFF)
             merge(n/2L, &tmp[0], &tmp[n/2L], &data[0], 0, n,d+1);
#pragma omp taskwait
             else{
            else{
    multisort(n/4L, &data[0], &tmp[0],d+1);
    multisort(n/4L, &data[n/4L], &tmp[n/4L],d+1);
    multisort(n/4L, &data[n/2L], &tmp[n/2L],d+1);
    multisort(n/4L, &data[3L*n/4L], &tmp[3L*n/4L],d+1);
    merge(n/4L, &data[0], &data[n/4L], &tmp[0], 0, n/2L,d+1);
    merge(n/2L, &tmp[0], &tmp[n/2L], &data[0], 0, n,d+1);
    ]
```

sbatch -p exectuion submit-extrae.sh multisort-omp 8 0



sbatch -p exectuion submit-extrae.sh multisort-omp 8 1



4. Shared-memory parallelisation with OpenMP task using dependencies For the deliverable: Include the tables of modelfactors, scalability plots and Paraver windows to support the comparison.

#### OMP\_NUM\_THREADS=8 CUTOFF=16 ./multisort-omp -n 32768 -s 1024 -m 1024

