

# Query Processing on Dynamic Networks with Customizable Contraction Hierarchies on Neo4j

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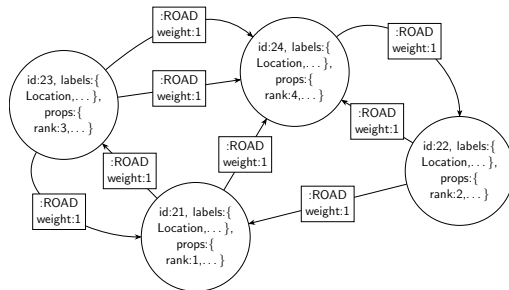
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- 2 Dijkstra
- 3 Contraction Hierarchies
- 4 Vertex Order
- 5 External Memory
- 6 Experiments
- 7 Conclusion

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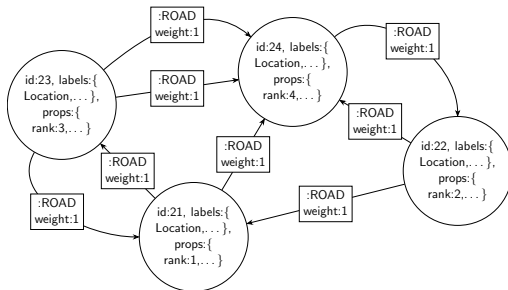
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- Graph Databases  $\Rightarrow$  External memory
- Accelerate Shortest Path Queries in Databases
- Why Customizable Contraction Hierarchies?
  - fast for main memory applications
  - reasonable preprocessing time
  - It is updatable
- Test Data  $\Rightarrow$  Road Networks

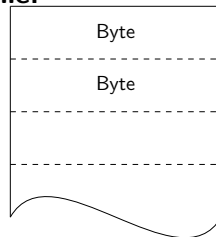
## Property Graph:



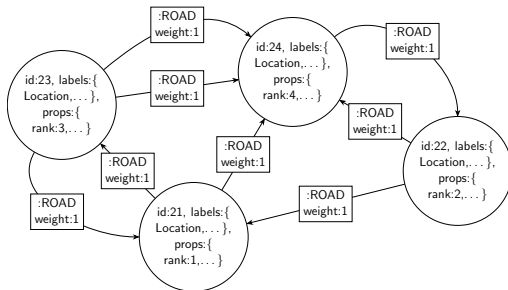
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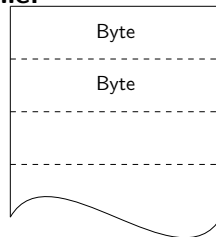
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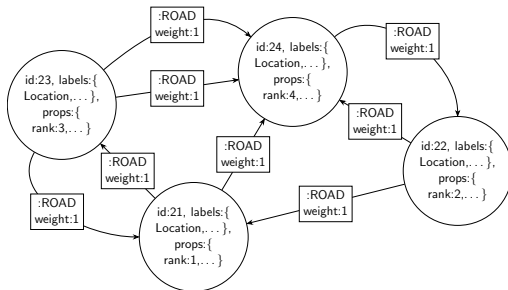


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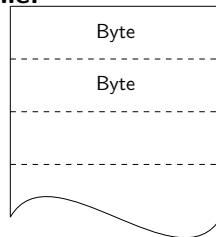


- transformation to data structure with a single dimension

## Property Graph:



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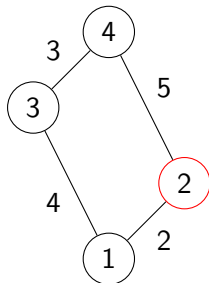
- transformation to data structure with a single dimension
- Databases use HDDs  $\Rightarrow$  slow random access



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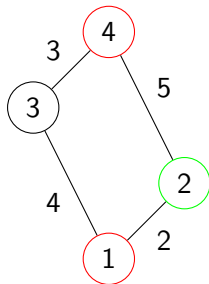
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Let's go from  $v_2$  to  $v_3$



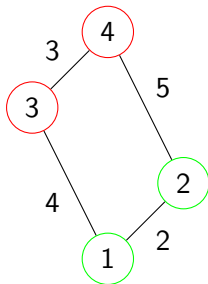
id	dist	settled
2	0	false

Let's go from  $v_2$  to  $v_3$



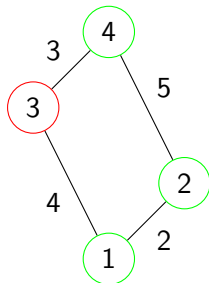
id	dist	settled
2	0	true
1	2	false
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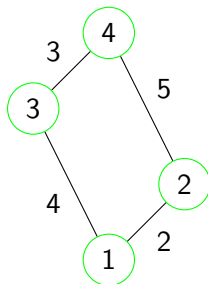
id	dist	settled
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3	6	false

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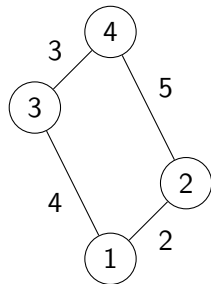


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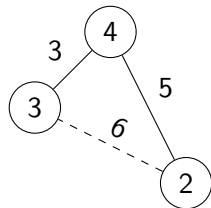
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# Contraction Hierarchies Example

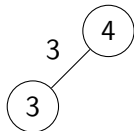




# Contraction Hierarchies Example



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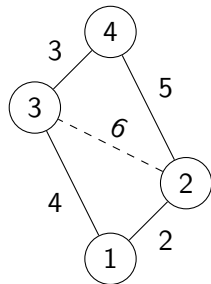


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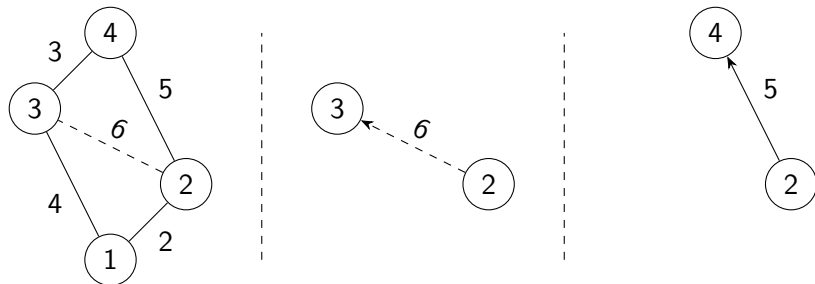
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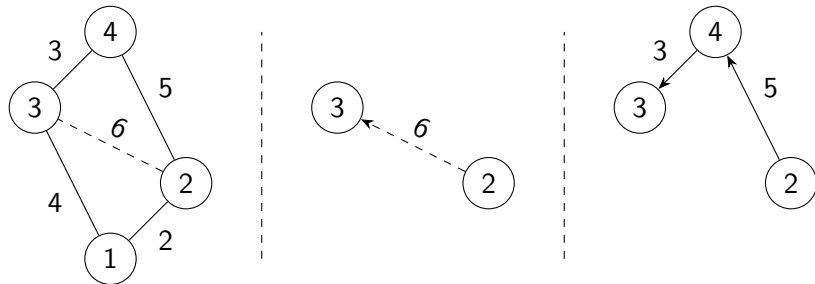
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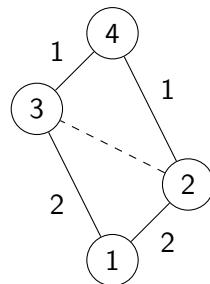
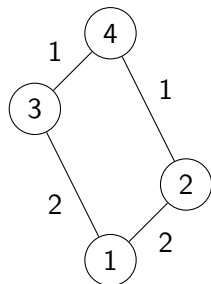
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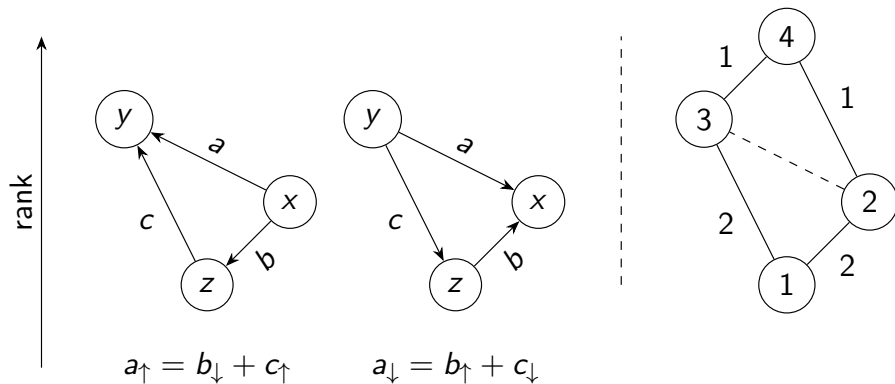
# Customizable Contraction Hierarchies

- CH insert shortcut if shortest path property is violated
- CCH insert shortcut if there is no direct connection

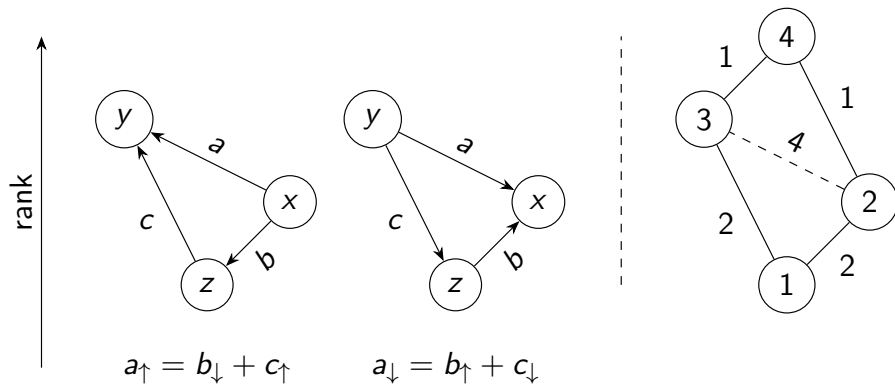




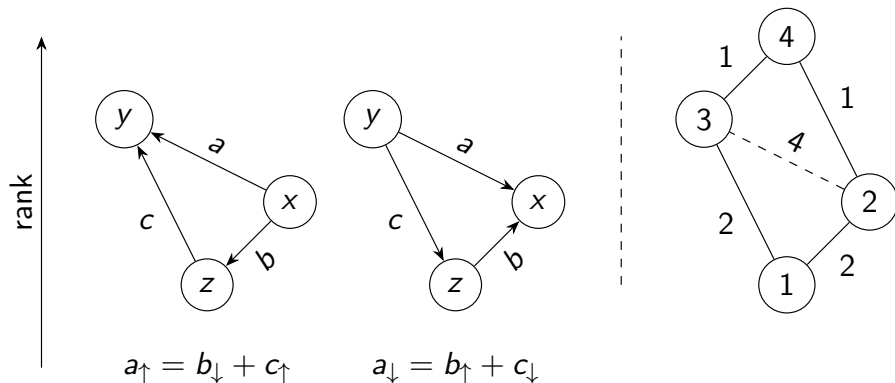
# CCH Weights — Lower Triangles



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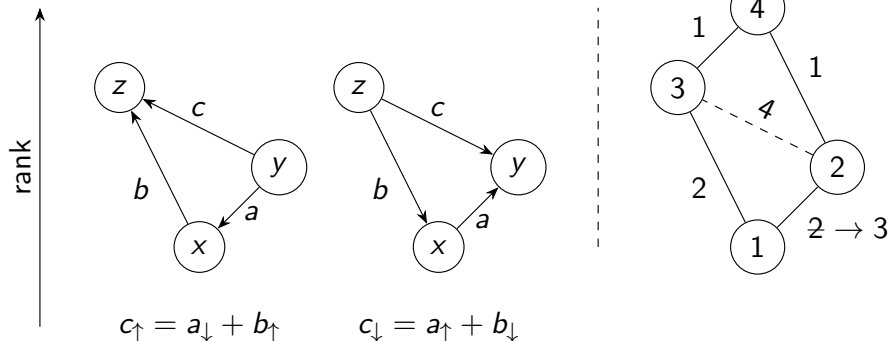


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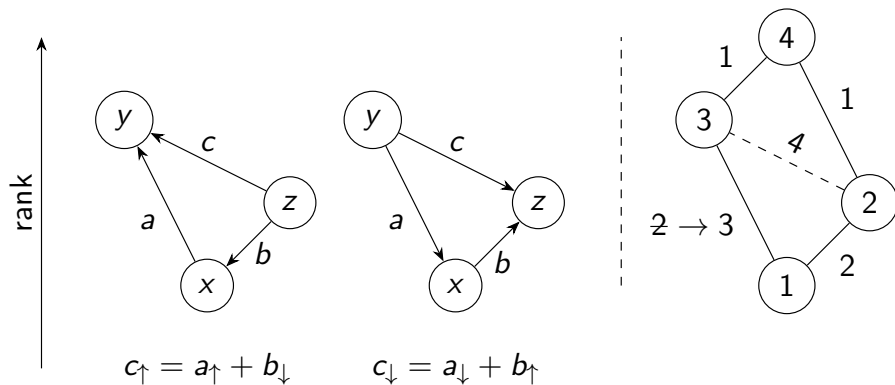


Bottom up for every arc! Also original arcs

# CCH Update — Upper Triangle



# CCH Update — Intermediate Triangle

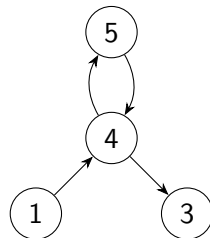
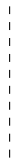
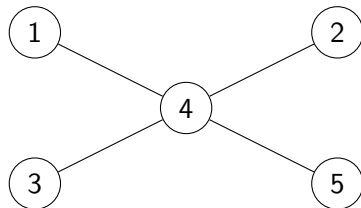


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# Important vertex not contracted Last

- Contracted Using Edge Difference
- Go from  $v(1)$  to  $v(3)$
- Forward and Backward search are deeper than they should be
- Switch contraction order of  $v(4)$  and  $v(5)$



# Linear Contraction

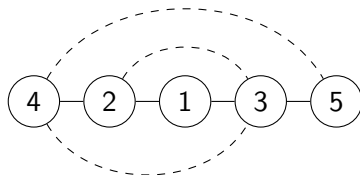
## 1. linear contraction

- No Shortcuts
- Could happen with ED
- four vertices to expand



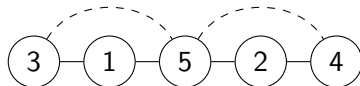
## 2. middle vertex first

- Three Shortcuts
- four vertices to expand



## 3. good contraction order

- Three Shortcuts
- four vertices to expand





# Importance Calculation

- add a level  $l(v) = 0$  to each node
- if contracting  $v$ :  $l(v) = \max\{l(v) + 1, l(w)\} \forall w \in N(v)$
- $A(v)$  set of added arcs
- $D(v)$  set of deleted arcs
- $h(a)$  hops an arc represents if unpacked

$$i(v) = l(v) + \frac{|A(v)|}{|D(v)|} + \frac{\sum_{a \in A(v)} h(a)}{\sum_{a \in D(v)} h(a)}$$

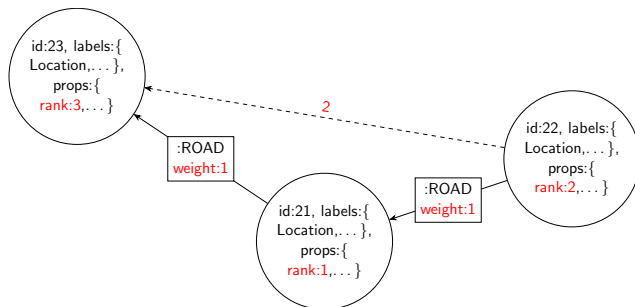
## Important theorem

more shortcuts inserted but improves query time!

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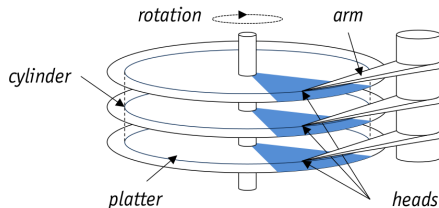
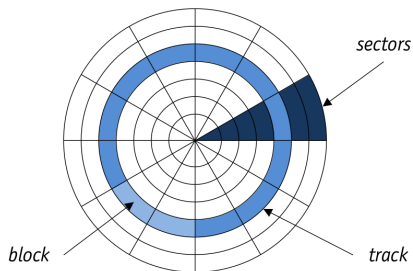
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# Persistence Objectives



- keep only necessary data
  - **rank** → to do the mapping to the input graph
  - arc **weight**
- Store edges that are likely to be request together spacial close
- Use as few space as possible → the less you write the less you read

# Magnetic Disks



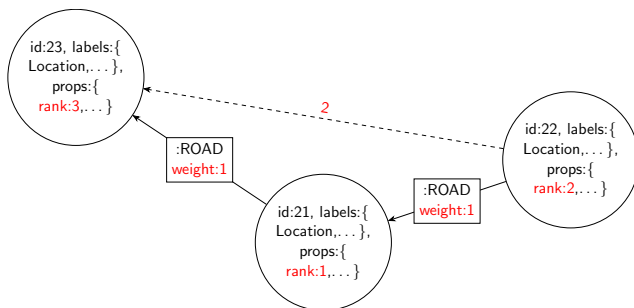
- Data is arranged in concentric rings (tracks) on platters
- Tracks are divided into arc-shaped sectors

## One by One

Data is read from and written to disk one **block** at a time

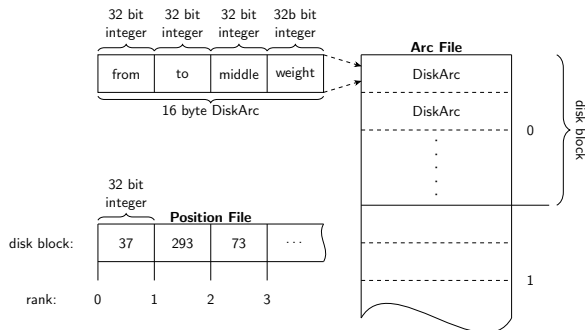
# Transformation to a Table

- Depth-First-Search starting at highest rank
- retrieve only arcs. vertices will be reconstructed from arcs
- remember middle node



start rank	end rank	middle rank	weight
1	3	-1	1
2	3	1	2

# Store Example



- fill all arcs of a vertex into a block
- add block number of rank to position file.  $G_{\uparrow}$  use *from* ;  $G_{\downarrow}$  use *to*
- if next vertex' arcs don't fit anymore  $\rightarrow$  flush block and take next

## Min Block Size

$$\max(d_{\uparrow \max}(v), d_{\downarrow \max}(v)) \leq \frac{\text{diskBlockSize}}{16}$$

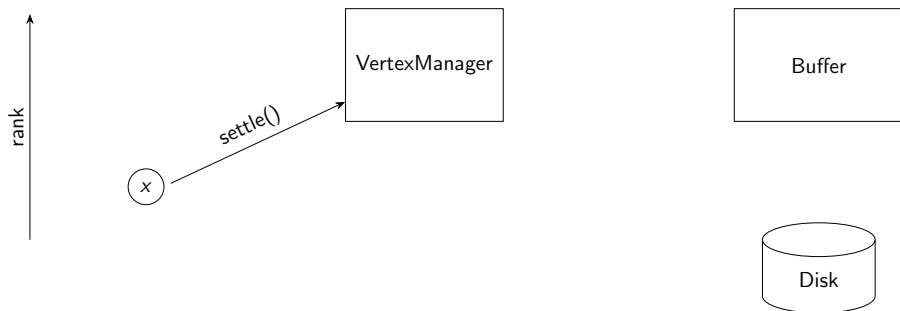
# CCH Disk Search (upwards graph example)

1. lazy load vertices  $\Rightarrow$  only start node is loaded without arcs



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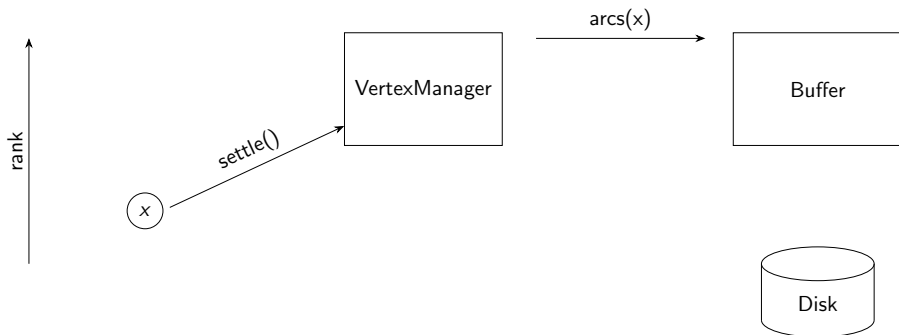
1. lazy load vertices  $\Rightarrow$  only start node is loaded without arcs
2. settle vertex (right before expanding it)





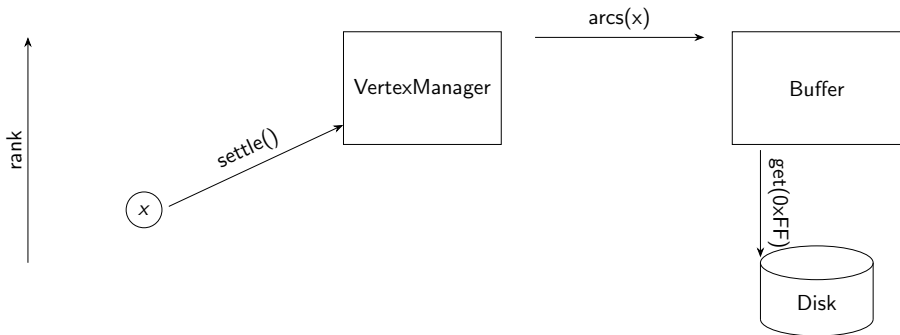
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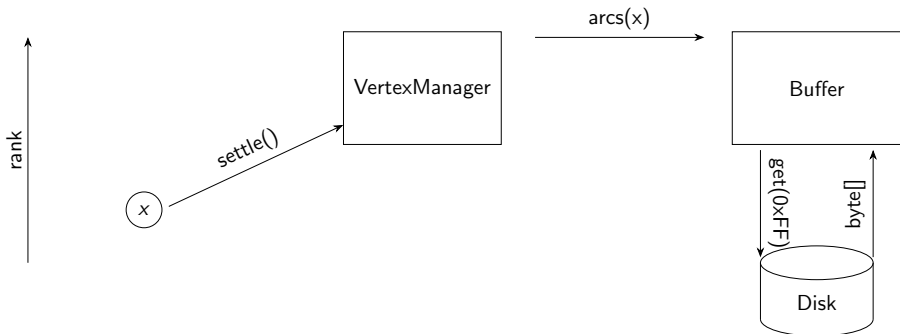
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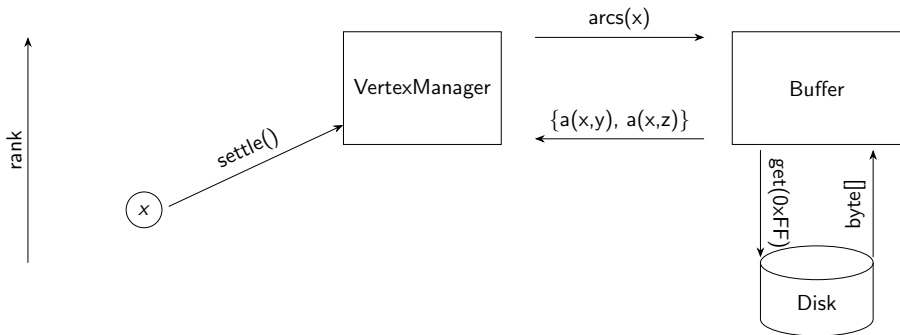
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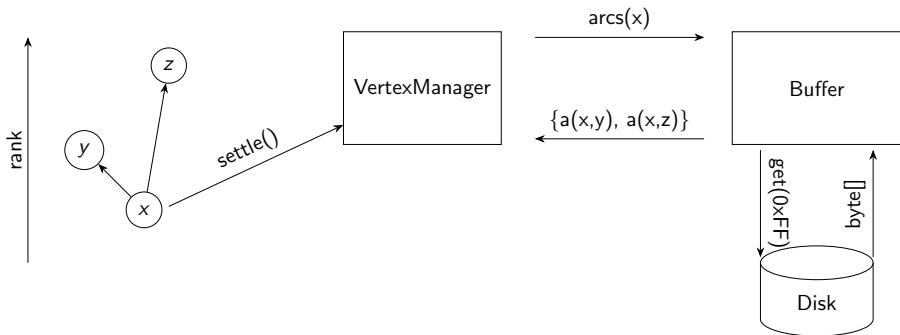
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  - Buffer returns arcs

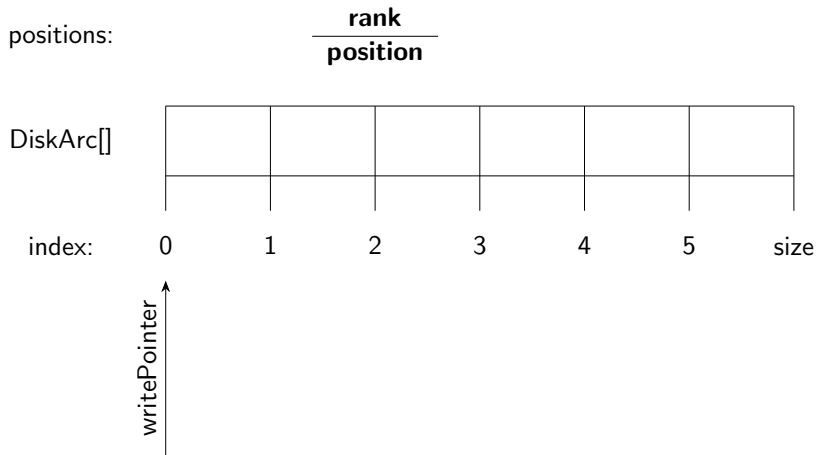


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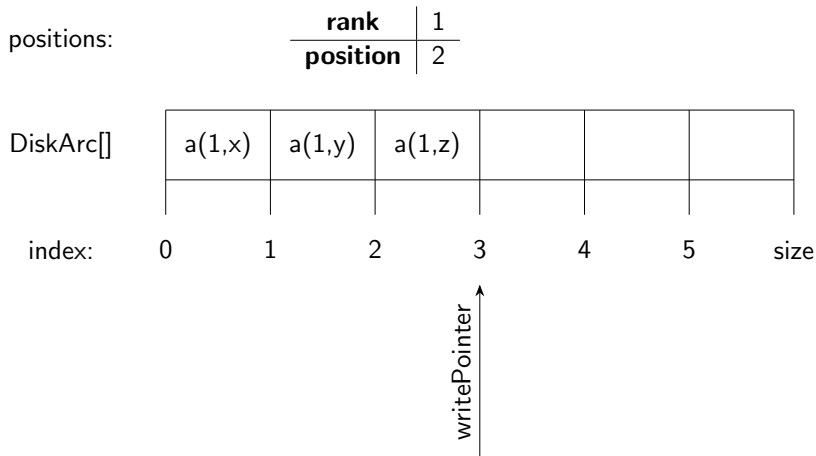
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  - Buffer returns arcs
  - VertexManager attaches arcs to  $v(x)$



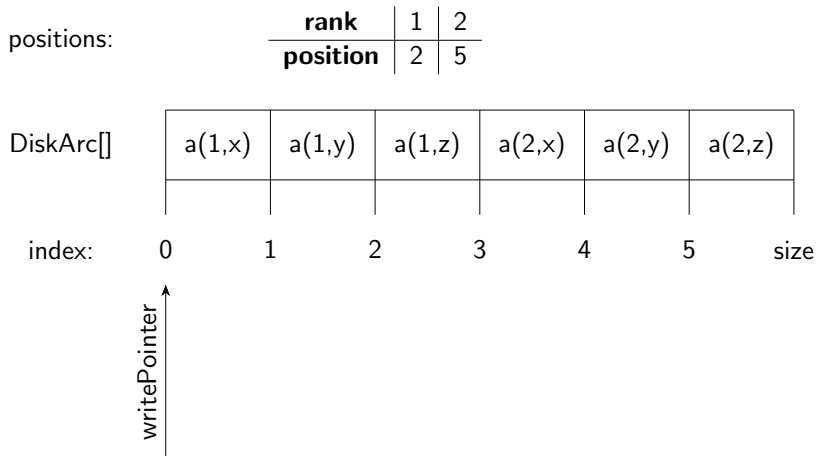
# Circular Buffer



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# Circular Buffer





# Circular Buffer

positions:

rank	1	2	max(rank)
position	2	5	-1

DiskArc[]

a(1,x)	a(1,y)	a(1,z)	a(2,x)	a(2,y)	a(2,z)

index:      0            1            2            3            4            5            size

↑  
writePointer

# Circular Buffer

positions:

rank	1	2	3	max(rank)
position	2	5	3	-1

DiskArc[]

a(3,x)	a(3,y)	a(3,z)	a(3,zx)	a(2,y)	a(2,z)

index:      0            1            2            3            4            5            size

remove incomplete edge set from position

writePointer  
↑

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- Retrieve Arcs  $\Rightarrow$  iterate backwards from position until start vertex differs
- If arc is doesn't start with requested rank  $\Rightarrow$  remove position and refetch

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# General Results

	New York	Colorado	Florida	California +Nevada	Great Lakes	Eastern USA
$ V $	264,346	435,666	1,070,376	1,890,815	2,758,119	3,598,623
$ A $	730,100	1,042,400	2,687,902	4,630,444	6,794,808	8,708,058
$ S $	2,153,002	1,680,290	4,397,804	8,598,552	17,833,050	17,712,722
$\frac{ S }{ A }$	2.93	1.59	1.62	1.85	2.62	2.03
$t_{contraction}$	545 s	233 s	579 s	4,384 s	25.29 h	23.29 h
$\max(d(v))$	1,150	629	785	1,252	2,433	2,391
$ \bigcirc A_{\uparrow} $	23.1 MB	21.9 MB	56.9 MB	107 MB	201 MB	215 MB
pos-file $_{\uparrow}$	1.1 MB	1.8 MB	1.3MB	7.6 MB	11.1 MB	14.4 MB
$t_{dijkstra}$	0.816 s	0.549 s	2.630 s	4.858 s	5.425 s	5.387 s
$t_{cch}^{640kB}$	0.140 s	0.122 s	0.147 s	0.289 s	0.732 s	0.727 s
$I/O^{640kB}$	574	437	500	899	1671	1572
$t_{update}^{640kB}$	90 s	51 s	142 s	444 s	1827s	1557s
$t_{cch-updated}^{640kB}$	0.147 s	0.129 s	0.150 s	0.302 s	0.783 s	0.855 s
$I/O_{cch-upd.}^{640kB}$	569	457	516	924	2779	2716
$t_{cch-updated}^{20\%}$	0.136 s	0.130 s	0.092 s	0.183 s	0.660 s	0.680 s
$I/O_{cch-upd.}^{20\%}$	315	307	226	283	804	680
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$t_{cch-updated}^{640kB}$	0.147 s	0.129 s	0.150 s	0.302 s	0.783 s	0.855 s
$I/O_{cch-upd.}^{640kB}$	569	457	516	924	2779	2716
$t_{cch-updated}^{20\%}$	0.136 s	0.130 s	0.092 s	0.183 s	0.660 s	0.680 s
$I/O_{cch-upd.}^{20\%}$	315	307	226	283	804	680
$t_{cch-updated}^{100\%}$	0.062 s	0.038 s	0.039 s	0.099 s	0.438 s	0.479 s



# General Results

	New York	Colorado	Florida	California +Nevada	Great Lakes	Eastern USA
$ V $	264,346	435,666	1,070,376	1,890,815	2,758,119	3,598,623
$ A $	730,100	1,042,400	2,687,902	4,630,444	6,794,808	8,708,058
$ S $	2,153,002	1,680,290	4,397,804	8,598,552	17,833,050	17,712,722
$\frac{ S }{ A }$	2.93	1.59	1.62	1.85	2.62	2.03
$t_{contraction}$	545 s	233 s	579 s	4,384 s	25.29 h	23.29 h
$\max(d(v))$	1,150	629	785	1,252	2,433	2,391
$ \bigcirc A_{\uparrow} $	23.1 MB	21.9 MB	56.9 MB	107 MB	201 MB	215 MB
pos-file $_{\uparrow}$	1.1 MB	1.8 MB	1.3MB	7.6 MB	11.1 MB	14.4 MB
$t_{dijkstra}$	0.816 s	0.549 s	2.630 s	4.858 s	5.425 s	5.387 s
$t_{cch}^{640kB}$	0.140 s	0.122 s	0.147 s	0.289 s	0.732 s	0.727 s
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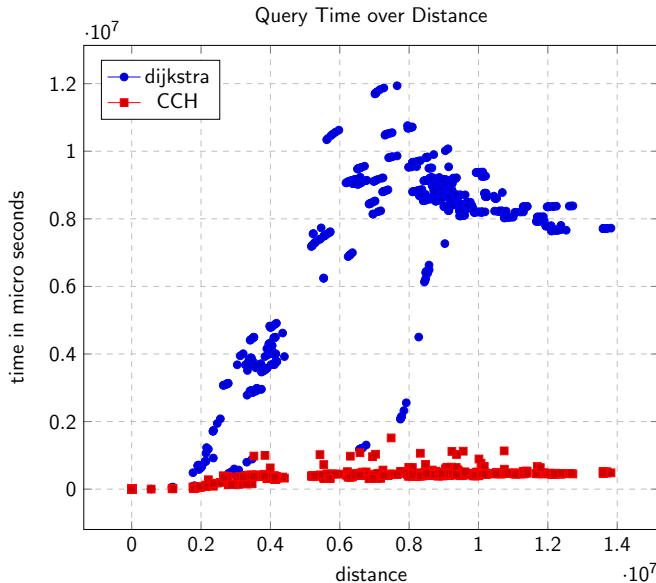
# General Results

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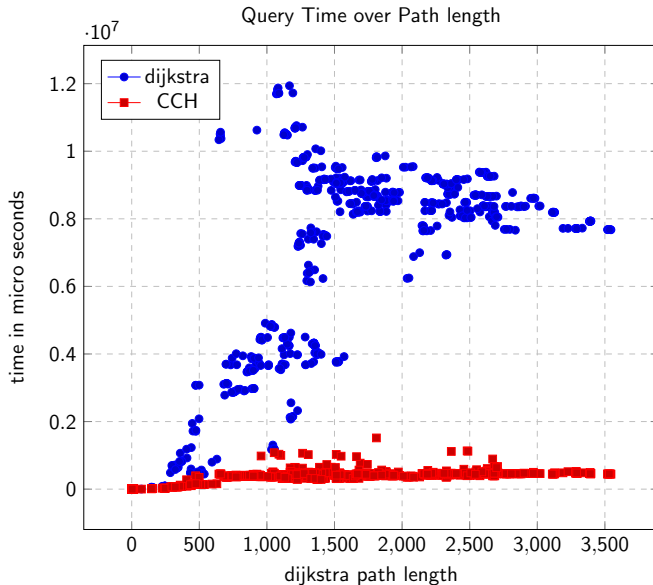
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	New York	Colorado	Florida	California +Nevada	Great Lakes	Eastern USA
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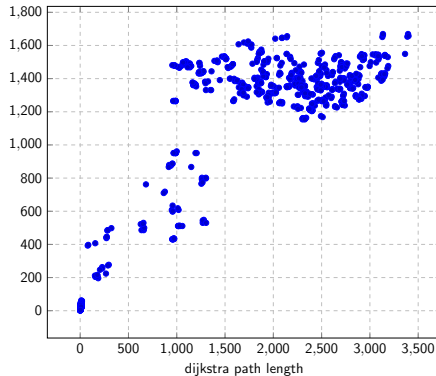
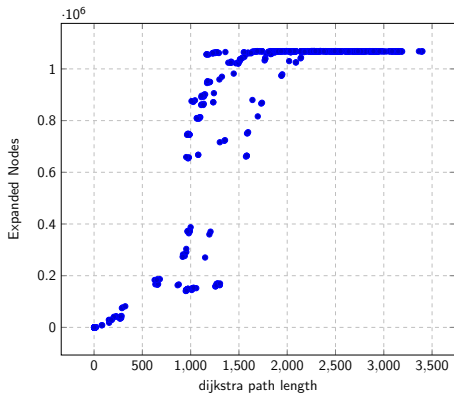
# Dijkstra vs. CCH — Query Time



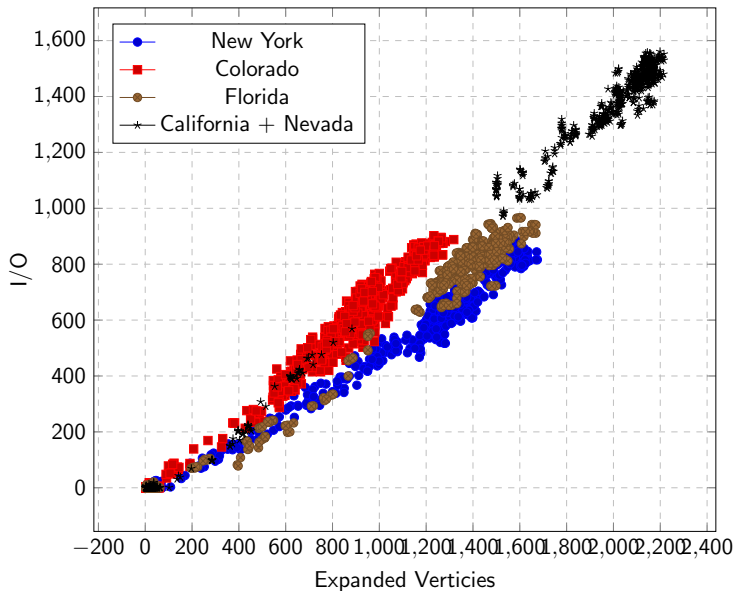
# Dijkstra vs. CCH — Query Time



# Expanded Vertices

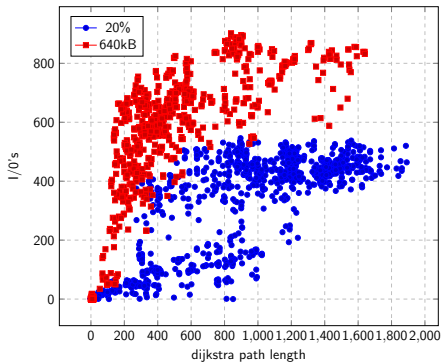


# Buffer 640 kB

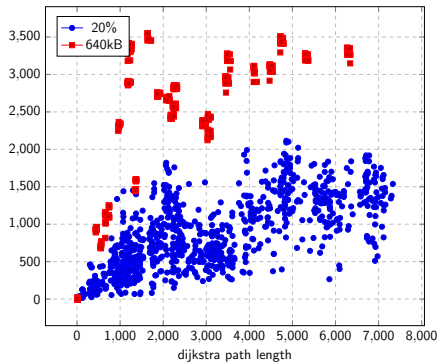


# Buffer 20% Arc Count

Colorado I/O Circular Buffer



Great Lakes I/O Circular Buffer





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- 1 Introduction
- 2 Dijkstra
- 3 Contraction Hierarchies
- 4 Vertex Order
- 5 External Memory
- 6 Experiments
- 7 Conclusion**

- We can accelerate Graph Databases with CCH
- The major problem is to flatten the graph
- Try it with a Relational Database