

# CSCI / MATH 2072U - Assignment 3

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## 1(a) Pseudo code

**Input:** integer  $N$

```
1:  $x \leftarrow []$ 
2:  $sum \leftarrow 0$ 
3:  $step \leftarrow 1/N$ 
4:  $V \leftarrow$  empty  $N + 1 \times N + 1$  numpy array
5: for  $i = 0, \dots, N$  do
6:   for  $j = 0, \dots, N$  do
7:      $x \leftarrow sum$  ▷ Append sum to list  $x$ 
8:      $sum \leftarrow sum + step$ 
9:      $V[i, j] \leftarrow x[i]^j$ 
10:   end for
11: end for
12: return  $V$ 
```

**Output:**  $V$ , an  $N + 1 \times N + 1$  array of floats

Counting FLOPs

Line 3: 1 FLOP

Line 4: 2 FLOPs

Line 8:  $1 \times (N + 1) \times (N + 1) = (N + 1)^2$

Line 9: Maximum possible FLOP count =  $N$

**FLOP count:**  $(N + 1)^2 + N + 3$

My algorithm has quadratic time complexity,  $\mathcal{O}(N^2)$ , i.e. its performance is directly proportional to  $N^2$ .

**1(f)** The numerically obtained solution is accurate up until  $N = 18$ . This is because the maximal relative error is larger than 1 for  $N > 18$ . The condition number of  $V$  for  $N = 18$  is  $1.468 \times 10^{15}$  and the relative residual is  $2.250 \times 10^{-16}$ .

## 2(a) Pseudo code

**Input:** arrays of  $n$  floats  $\mathbf{x}$ ,  $\mathbf{a}$ ,  $\mathbf{b}$  and  $\mathbf{c}$

```
1:  $n \leftarrow$  size of vector  $x$ 
2:  $y \leftarrow n \times 1$  empty numpy array
3:  $y[0] \leftarrow b[0] \times x[0] + c[0] \times x[1]$   $\triangleright$  (Setting first entry of  $y$ )
4: for  $i = 1, \dots, N - 2$  do
5:    $y[i] \leftarrow a[i] \times x[i - 1] + b[i] \times x[i] + c[i] \times x[i + 1]$ 
6: end for
7:  $y[n - 1] \leftarrow a[n - 1] \times x[n - 2] + b[n - 1] \times x[n - 1]$   $\triangleright$  (Setting last entry of  $y$ )
8: return  $y$ 
```

**Output:** array of  $n$  floats  $y$  such that  $\mathbf{y} = \mathbf{Ax}$

## 2(b)

Counting FLOPs

Line 3: 3 FLOPs

Line 4: 1 FLOP

Line 5:  $7 \times (n - 2) = 7n - 14$  FLOPs

Line 7: 8 FLOPs

**FLOP count:**  $3 + 1 + 8 + 7n - 14 = 7n - 2$

My algorithm has linear time complexity,  $\mathcal{O}(n)$ . The algorithm's runtime grows directly in proportional to  $n$ .

**2(e)** The plot shows that my algorithm indeed does have linear time complexity, i.e. has a slope of 1.

The plot would show a quadratic time complexity if other matrix-vector product functions such as `np.matmul` and `np.dot` were used. This is because the matrix product of  $\mathbf{A}$  and  $\mathbf{x}$  would be determined using a nested for loop. The plot would still be a straight line but with a slope of 2.