DRather than treating blackjack as a 2-player game where one player is Fate, we call model it as the following (nondeterministic) state machine:

i.e. 
$$M = (Q, \Sigma, \Delta, g_o, F)$$
 where:

3) In addition, we want to specify a reward Rt for reaching a state after to transitions:

$$\forall t \ge 0$$
  $R_{t}(520) = .58$   $R_{t}(A20) = 0$   $R_{t}(521) = .88$   $R_{t}(Bust) = -1$ 

$$R_{t}(519) = .27$$

3) And binally, we want to know the likelihood of each transition, given we take a particular action in a particular state.

$$W(<19, STAY, S19^{2}) = 1$$
  
 $W(<20, STAY, S20^{2}) = 1$   
 $W(<19, HIT, 20^{2}) = \frac{1}{13}$   
 $W(<19, HIT, S21^{2}) = \frac{1}{13}$   
 $W(<19, HIT, BUST?) = \frac{11}{13}$ 

$$w(19, Hit, 9')$$

$$= P(g'|_{q=19, \sigma=Hit})$$

$$= w(19, Hit, 9') = 1$$

 $w(<20, HiT, 5217) = \frac{1}{13}$  $w(<20, HiT, 8ust>) = \frac{12}{13}$  | Sum w(<20, HiT, q'>) = 1

4) This is called a Markov Decision Process (MDP). Formally ar MDP is a triple (M, R, w), where:

-  $M = (Q, \Sigma, \Delta, q_0, F)$  is a state machine

- R: Q × IN -> IR is a "reward function"

-  $W: \Delta \mapsto \mathbb{R}$  s.t. for all  $q \in \mathbb{Q}$ ,  $\sigma \in \Sigma: \sum_{(q,\sigma,q') \in \Delta} \omega((q,\sigma,q')) = 0$ 

We assume that  $R(q,t) = \mathcal{T} \cdot R(q,t-1)$   $\forall q \in Q, t \ge 1$  for some "discounting factor"  $\mathcal{T}$ . s.t.  $0 < \mathcal{T} \le 1$ .

5) It can be helpful to define the following shorthand for dealing with MDPs:

 $\rightarrow R_{t}(q) \triangleq R(q, t) \quad \forall q \in Q, t > 0$ 

> P(q'|q,0) = w(<q,0,q'>) \(\dagger) \(\dagger) \\ \dagger \q\\ \dagger \(\dagger) \\

The write path ((q, oo, q,), (q, o, q2), ..., (qn-1, on-1, qn) >
q5:
q => q, == , on-1
qn

> reward (q = q, or ... on gn)

 $= R_{o}(q) + R_{i}(q_{i}) + ... + R_{n}(q_{n})$   $P(q = q_{i} = ... = q_{n})$ 

= P(q, 1q, 00) · P(q2 | q1, 01) · ··· · P(qn | qn-1, 0n-1)

6) The main computational challenge, given an MDP, is to determine the best decision to make in each state.

For example:

18 I have 19, should I HIT or STAY? 18 I have 20, should I HIT or STAY?

We can formalize this as a function  $\pi:(Q|F) \to \Sigma$ , which we call a policy.

e.g. T (A19) = HIT T (A20) = HIT Fig. Given a policy or, I can compute my expected reward Ul, starting from various states:

B) These expected rewards (usually called expected utility) can be expressed in terms of each other:

$$\begin{array}{l}
\left(A_{19}\right) = \left(R_{19}\right) + R_{19}\left(A_{20}\right) + R_{2}\left(S_{21}\right) P\left(A_{19}\right) + A_{20}\right) + R_{10}\left(S_{21}\right) P\left(A_{19}\right) + P\left(A_{19}\right) P\left(A_{19}\right) + P\left(A_{19}\right) P\left(A_{1$$

this equals

15 a prob.

1, b/c it

distribution.

(3) 
$$U^{\alpha}(A_{1}q) = R_{0}(A_{1}q)$$
+  $P(A_{2}O | A_{1}q, H_{1}r) \left[ (R_{0}(A_{2}O) + R_{0}(S_{2}1)) P(A_{2}O \xrightarrow{1} S_{2}1) + (R_{0}(A_{2}O) + R_{0}(B_{0}S_{1})) P(A_{2}O \xrightarrow{1} B_{0}S_{1}) P(A_{2}O \xrightarrow{1} B_{0}S_{1})$ 

- 9) This is a general result:  $U'''(q) = R_o(q) + Y \sum_{q' \in Q} U'''(q') \cdot P(q'|q, \pi(q))$
- 10) Usually we're not simply interested in computing the expected utility of a state, given some arbitrary policy or. Rather, we'd like to know how much reward we should expect if we execute the best policy not.

 $U(q) = Ro(q) + \delta \max_{q' \in Q} \sum_{q' \in Q} U(q') \cdot P(q'|q, \sigma)$ 

reward from the current State

maximum expected Utility of the next state, given the optimal action

A THE STATE OF THE PROPERTY OF

make a little of the set The real claim of the state of

1) For our blackjack example, we get the following equations.

$$U(A20) = max \left( U(520) \cdot P(520 | A20, STAY), \right)$$
  
 $\left\{ U(521) \cdot P(521 | A20, HIT) + U(BUST) \cdot P(BUST | A20, HIT) \right\}$ 

U(519) = Ro(519) U(520) = Ro(520)

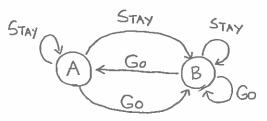
U(521) = R. (521)

U(BUST) = R. (BUST)

We have six equations for six unknowns (U(A19), U(A20), U(S19), U(S20), U(S21), U(BUST)).

<sup>12)</sup> Importantly, these equations are not linear, so we can't use linear algebra techniques.

(13) To motivate our equation-solving technique, let's use a smaller MDP:



$$P(A|A,STAY) = \frac{1}{2}$$

$$P(B|A,STAY) = \frac{1}{2}$$

$$P(B|A,Go) = 1$$

$$P(B|B,STAY) = 1$$

$$P(A|B,G_0) = \frac{1}{5}$$
  
 $P(B|B,G_0) = \frac{4}{5}$ 

$$R_{t}(A) = \lambda^{t} \cdot 1$$

$$R_{t}(B) = \lambda^{t} \cdot (-1)$$
where  $\lambda = \frac{1}{2}$ .

H) We get the following equations:
$$U(A) = R_{o}(A) + Y \cdot \max \left\{ U(B), U(A)P(A|A,STAY) + U(B)P(B|A,STAY) \right\}$$

$$= 1 + \frac{1}{2} \max \left\{ U(B), \frac{1}{2} U(A) + \frac{1}{2} U(B) \right\}$$

$$U(B) = R_{o}(B) + Y \cdot \max \left\{ U(B), U(A)P(A|B,Go) + U(B)P(B|B,Go) \right\}$$

$$= -1 + \frac{1}{2} \max \left\{ U(B), \frac{1}{5} U(A) + \frac{4}{5} U(B) \right\}$$

15) Suppose we guess the values of U(A) and U(B):

Uo is our guess at U(A)

Uo is our guess at U(B)

Consider the following iterative algorithm (called VALUE HERATION)

let  $U_{i}^{A} = 1 + \frac{1}{2} \max \left\{ U_{i-1}^{B}, \frac{1}{2} U_{i-1}^{A} + \frac{1}{2} U_{i-1}^{B} \right\}$ let  $U_{i}^{B} = -1 + \frac{1}{2} \max \left\{ U_{i-1}^{B}, \frac{1}{5} U_{i-1}^{A} + \frac{4}{5} U_{i-1}^{B} \right\}$ 

At each iteration, we assume our guesses for U(A) and U(B) from the previous iteration are correct, and we compute new guesses using our equations.

16) Why would this ever work? Well, it soems to converge...

(17) But can we prove it? Assume for the moment that a solution exists, i.e. there's a vector

$$U^* = \begin{bmatrix} U^A \\ U^B \end{bmatrix}$$

such that:

$$U^{A} = 1 + \frac{1}{2} \max \{ U^{B}, \frac{1}{2} U^{A} + \frac{1}{2} U^{B} \}$$

$$U^{B} = -1 + \frac{1}{2} \max \{ U^{B}, \frac{1}{5} U^{A} + \frac{4}{5} U^{B} \}$$

(18) Let's also measure how bad our guesses are.

Suppose the solution is [2] and our guess is [5].

We'll measure the distance of our guesses to the solution as the absolute difference between our wast guess and the solution, i.e.

$$dist([5],[2]) = max {[5-2], |3-4|} = 3$$

19) Now what if we could show that our guesses get better every iteration? i.e. that

$$dist\left(\begin{bmatrix} U_{i+1}^{A} \\ U_{i+1}^{B} \end{bmatrix}, \begin{bmatrix} U_{i}^{A} \\ U_{i}^{B} \end{bmatrix}\right) \leqslant K \cdot dist\left(\begin{bmatrix} U_{i}^{A} \\ U_{i}^{B} \end{bmatrix}, \begin{bmatrix} U_{i}^{A} \\ U_{i}^{B} \end{bmatrix}\right)$$

for 0 < K < 1.

That would mean:

Since dist 
$$\left(\begin{bmatrix} U_i^A \end{bmatrix}, \begin{bmatrix} U_i^A \end{bmatrix}\right) \ge 0$$
, thus  $\lim_{i \to \infty} \operatorname{dist} \left(\begin{bmatrix} U_i^A \end{bmatrix}, \begin{bmatrix} U_i^A \end{bmatrix}\right) = 0$ 

So our guesses would converge te the solution.

20 So then let's show it.

dist 
$$\left(\begin{bmatrix} U_{i+1}^A \\ U_{i+1}^B \end{bmatrix}, \begin{bmatrix} U_B^A \end{bmatrix}\right) = \max_{q \in \{A,B\}} \left| U_{i+1}^8 - U_b^8 \right|$$

| We simplify 
$$| U_{i+1}^{g} - U_{i}^{g} |$$
 we get:  
 $| U_{i+1}^{g} - U_{i}^{g} | = | Y \cdot \left[ \max_{\sigma \in Q'} \sum_{q'} P(q'|q,\sigma) - \max_{\sigma \in Q'} \sum_{q'} P(q'|q,\sigma) \right]$ 

because 
$$\forall f,g$$

max  $f(\sigma)$ -max  $g(\sigma)$ 
 $= x \cdot \max_{g'} |x' \cdot P(g'|g,\sigma) - x \cdot U_{g'}P(g'|g,\sigma)|$ 
 $= x \cdot \max_{g'} |x' \cdot P(g'|g,\sigma)|$ 
 $= x \cdot \max_{g'} |x' \cdot P(g'|g,\sigma)|$ 

(because the weighted (average of a set of numbers is at most the max)

Thus:

$$dist\left(\begin{bmatrix} U_{i+1}^{A} \\ U_{i+1}^{B} \end{bmatrix}, \begin{bmatrix} U_{i}^{A} \end{bmatrix}\right) \leq \max_{g} \left[ X \cdot \max_{g} \left| U_{i}^{g'} - U_{i}^{g'} \right| \right]$$

$$= \sum_{g} \max_{g} \left[ U_{i}^{g'} - U_{i}^{g'} \right]$$

$$= \sum_{g} dist\left(\begin{bmatrix} U_{i}^{A} \\ U_{i}^{B} \end{bmatrix}, \begin{bmatrix} U_{i}^{A} \end{bmatrix}\right)$$

(not-including 1), then our iterative technique (called "value iteration") will converge to the correct solution.