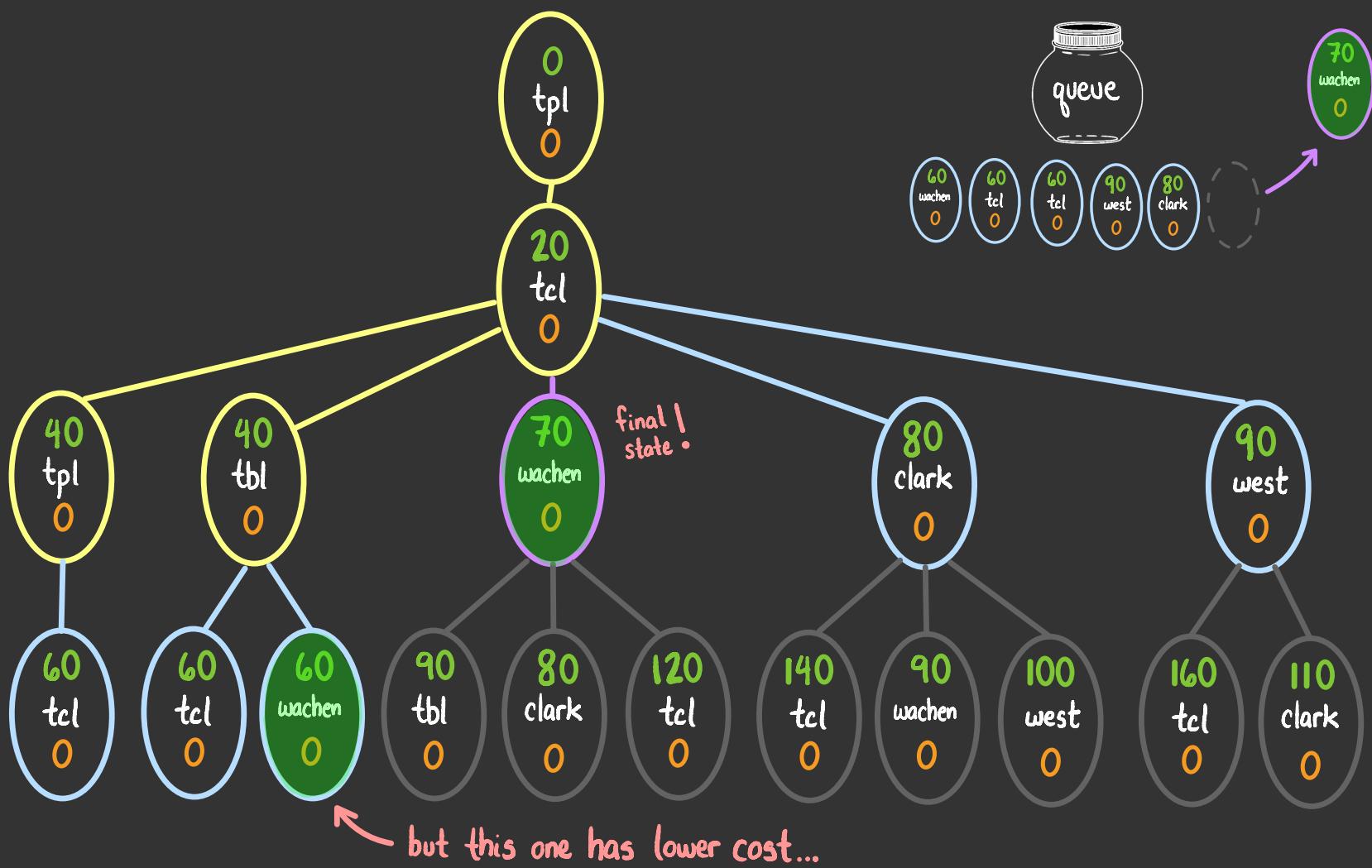
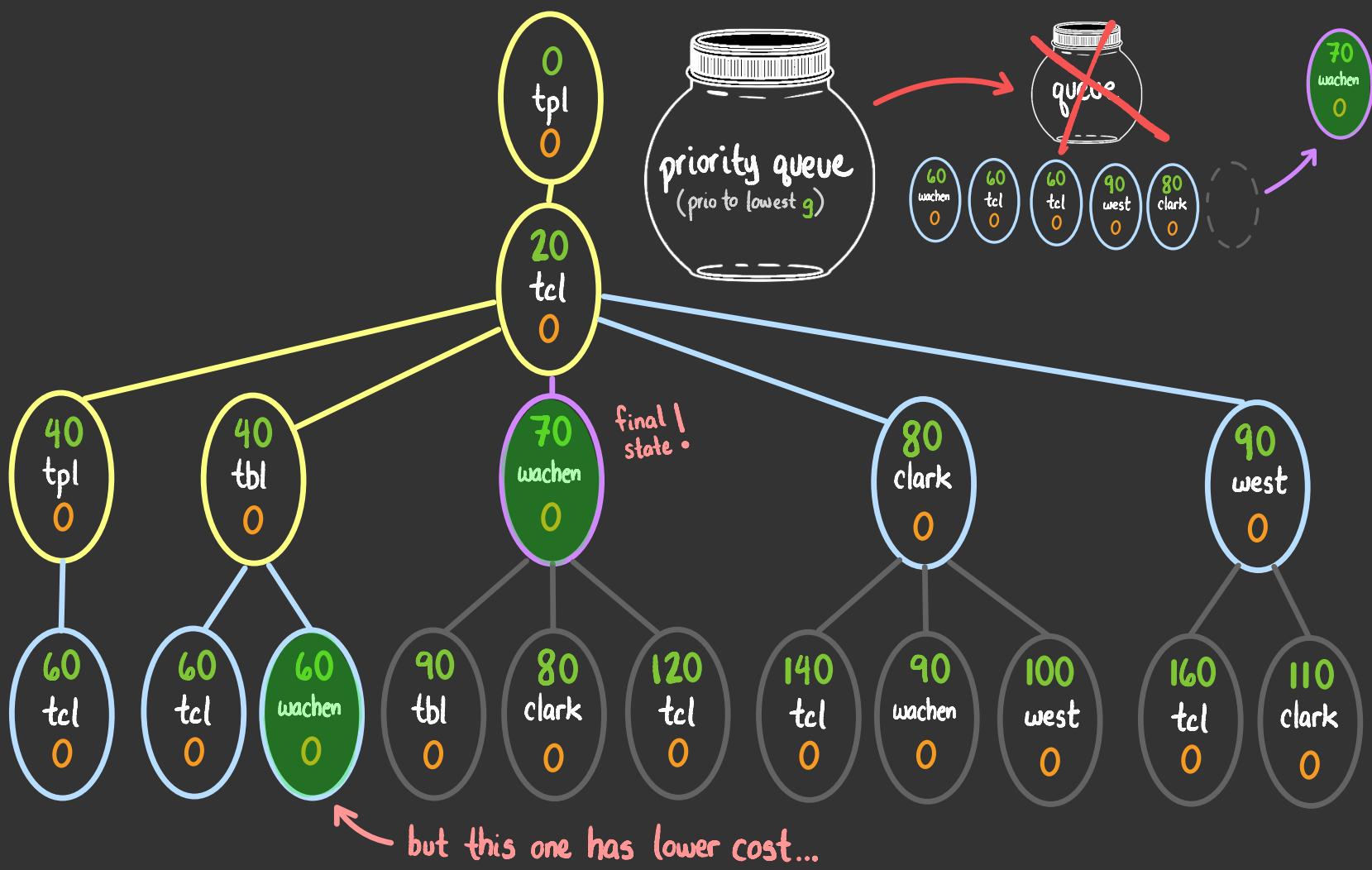


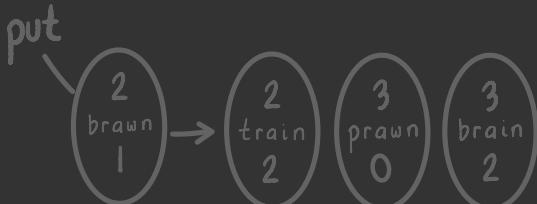
uniform cost  
search  
16 sept  
2022

CSCI  
373

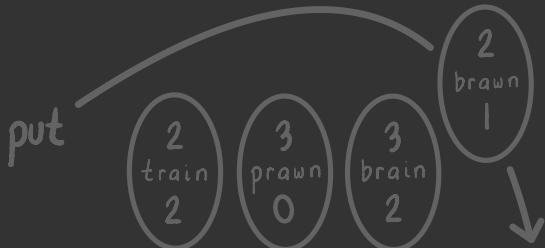




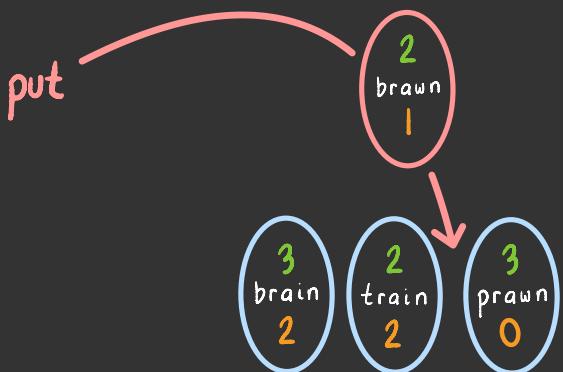
# queue



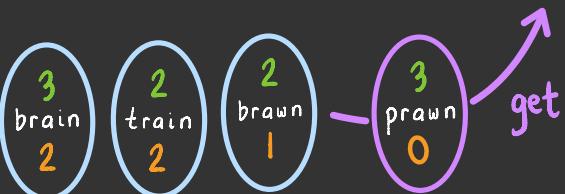
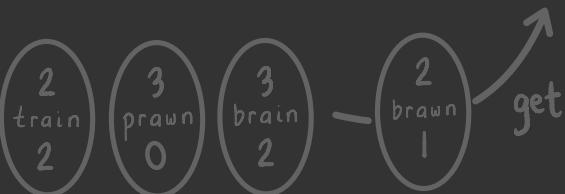
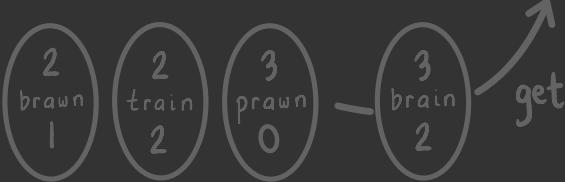
# stack



# priority queue



# example containers

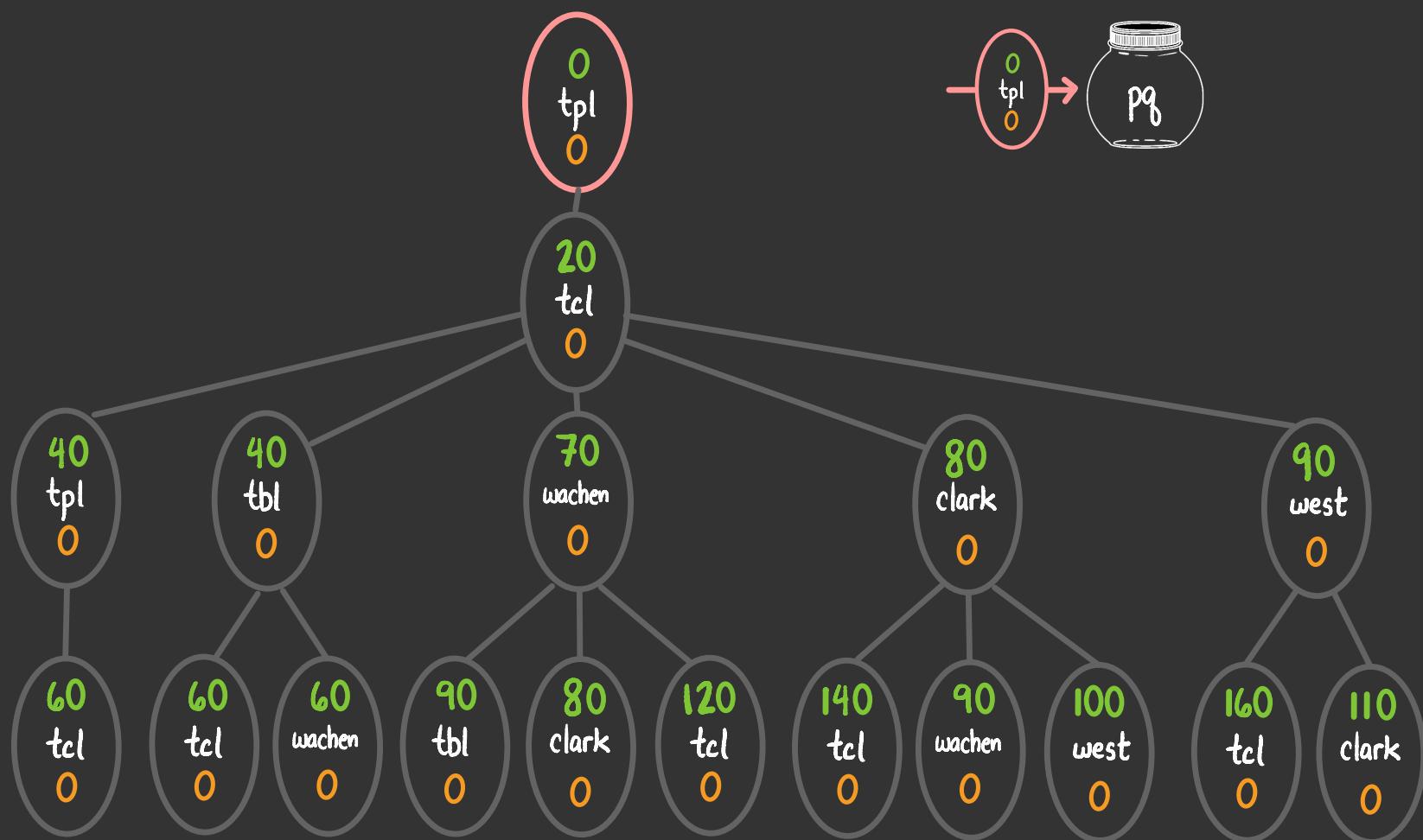


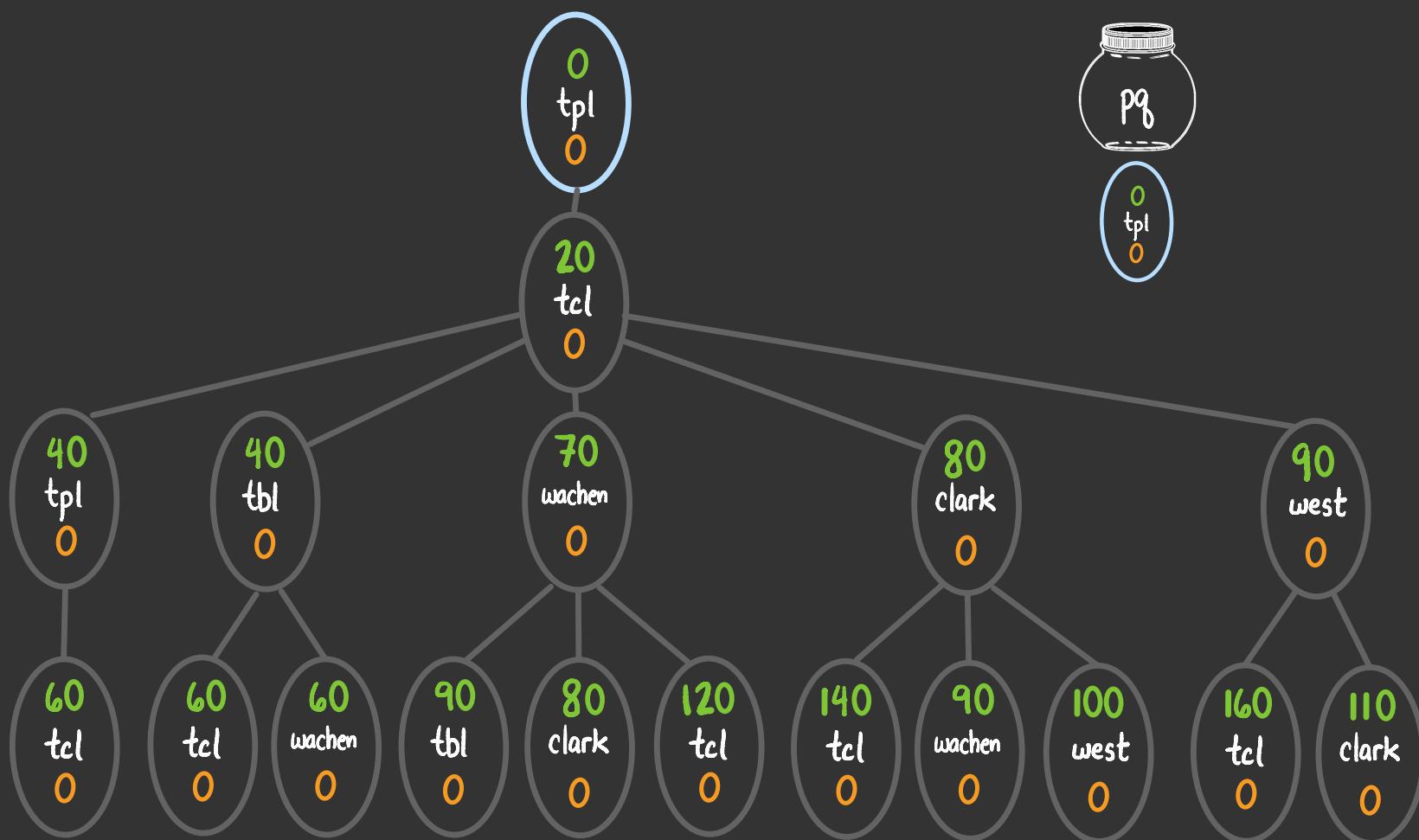
$\text{SEARCH}(M = (Q, \Sigma, \Delta, q_0, F, w), H)$ :

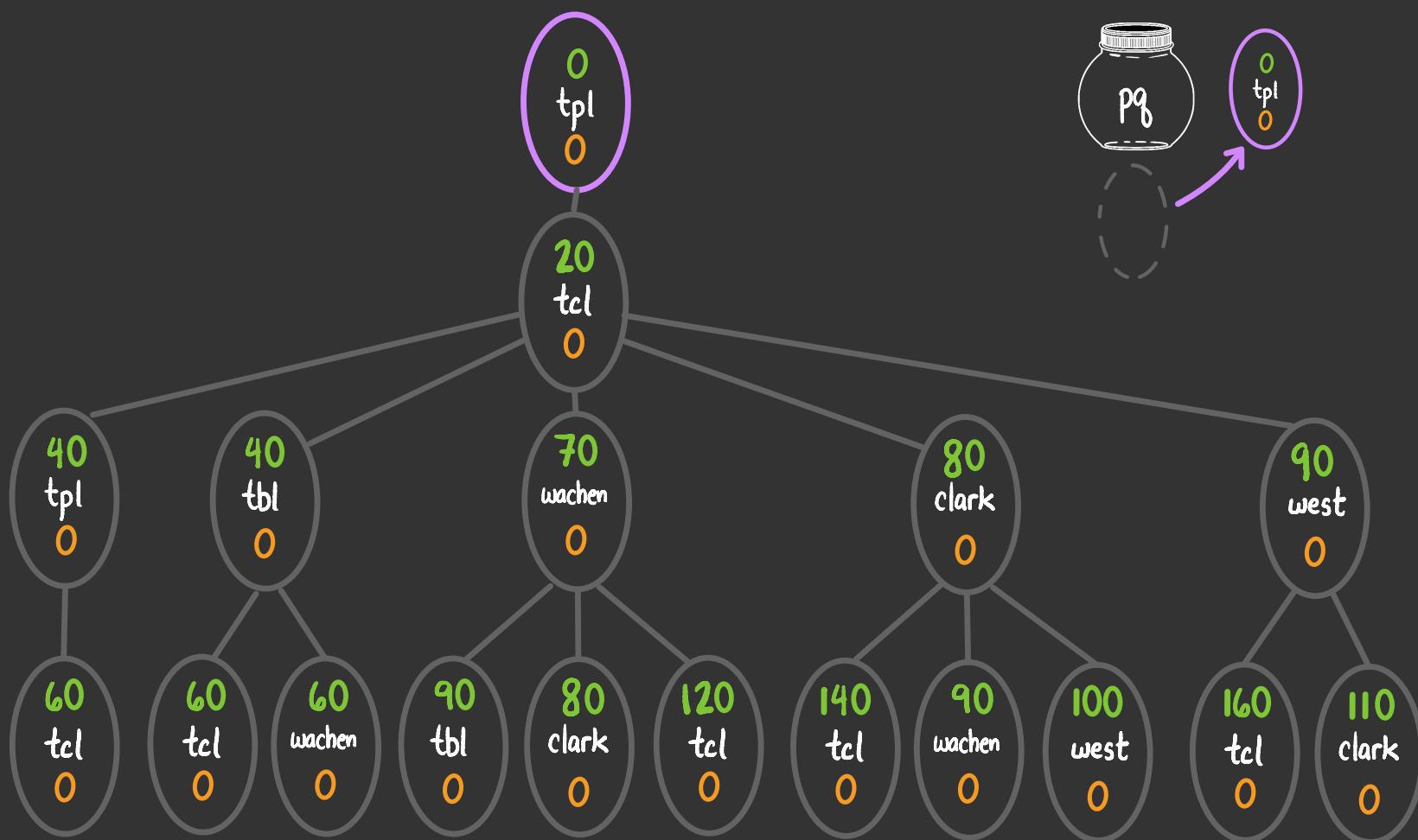
- ▶ container = new Container()
- ▶ container.put(< $q_0, 0, H(q_0)$ >)
- ▶ repeat:
  - ▶ if container.empty() then return  $\infty$
  - ▶ n = container.get()
  - ▶ if  $q(n) \in F$  then return  $g(n)$
  - ▶ let successors <sub>$M, H$</sub> (n) = { < $q', g(n) + w(q(n), \sigma, q'), H(q')$ > |  $(q(n), \sigma, q') \in \Delta$  }
  - ▶ for  $n' \in$  successors <sub>$M, H$</sub> (n):  
    container.put( $n'$ )

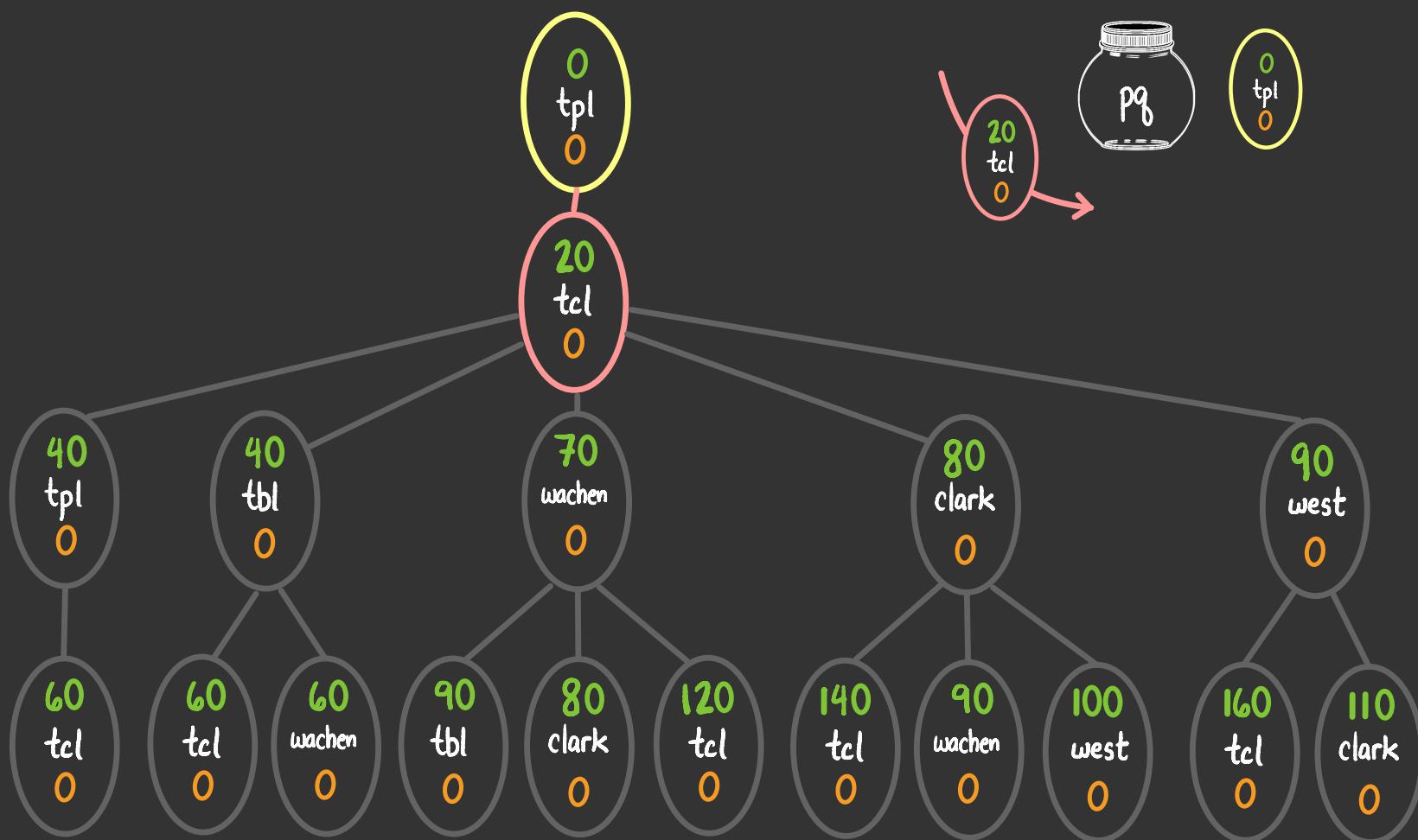


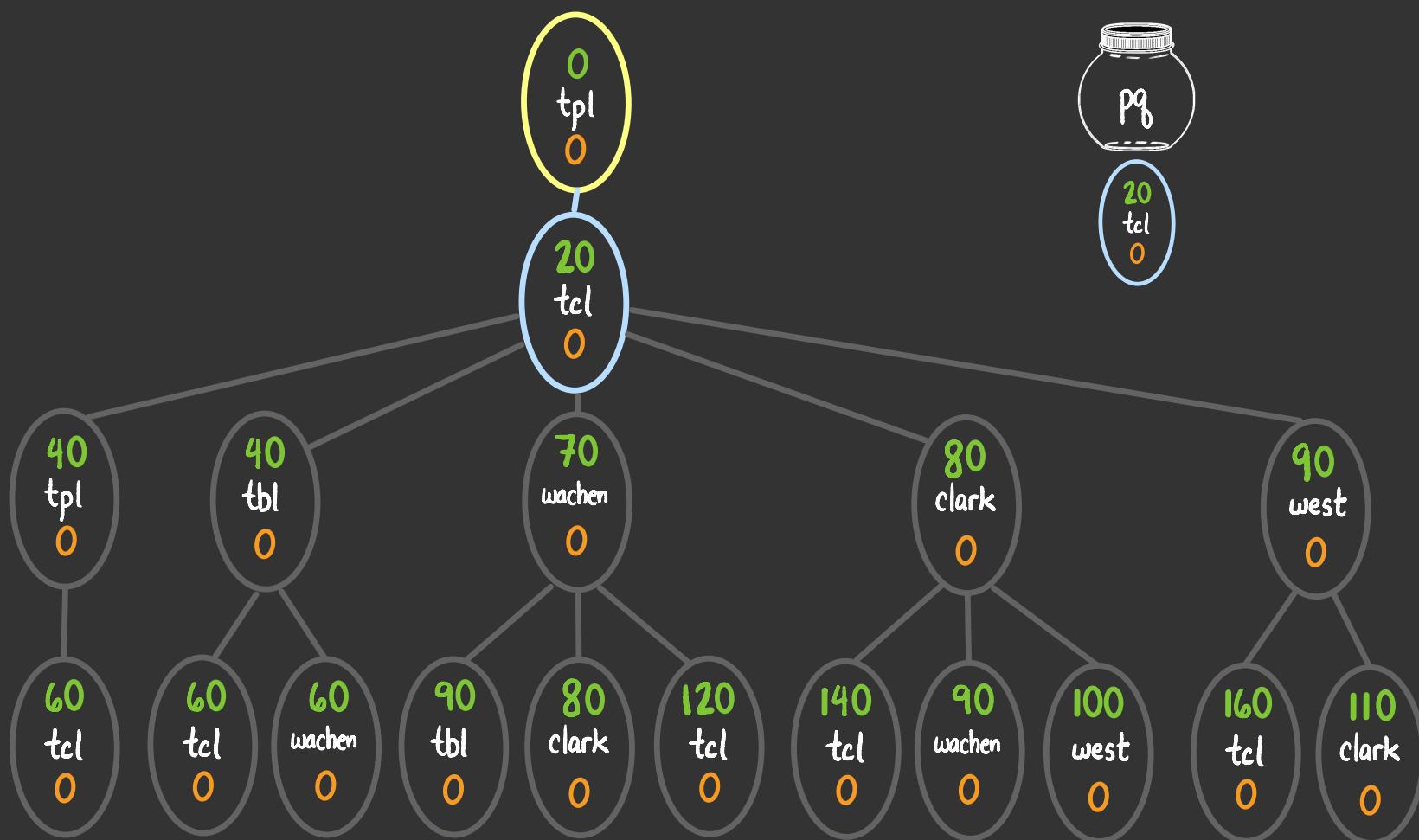
feel free to play along...

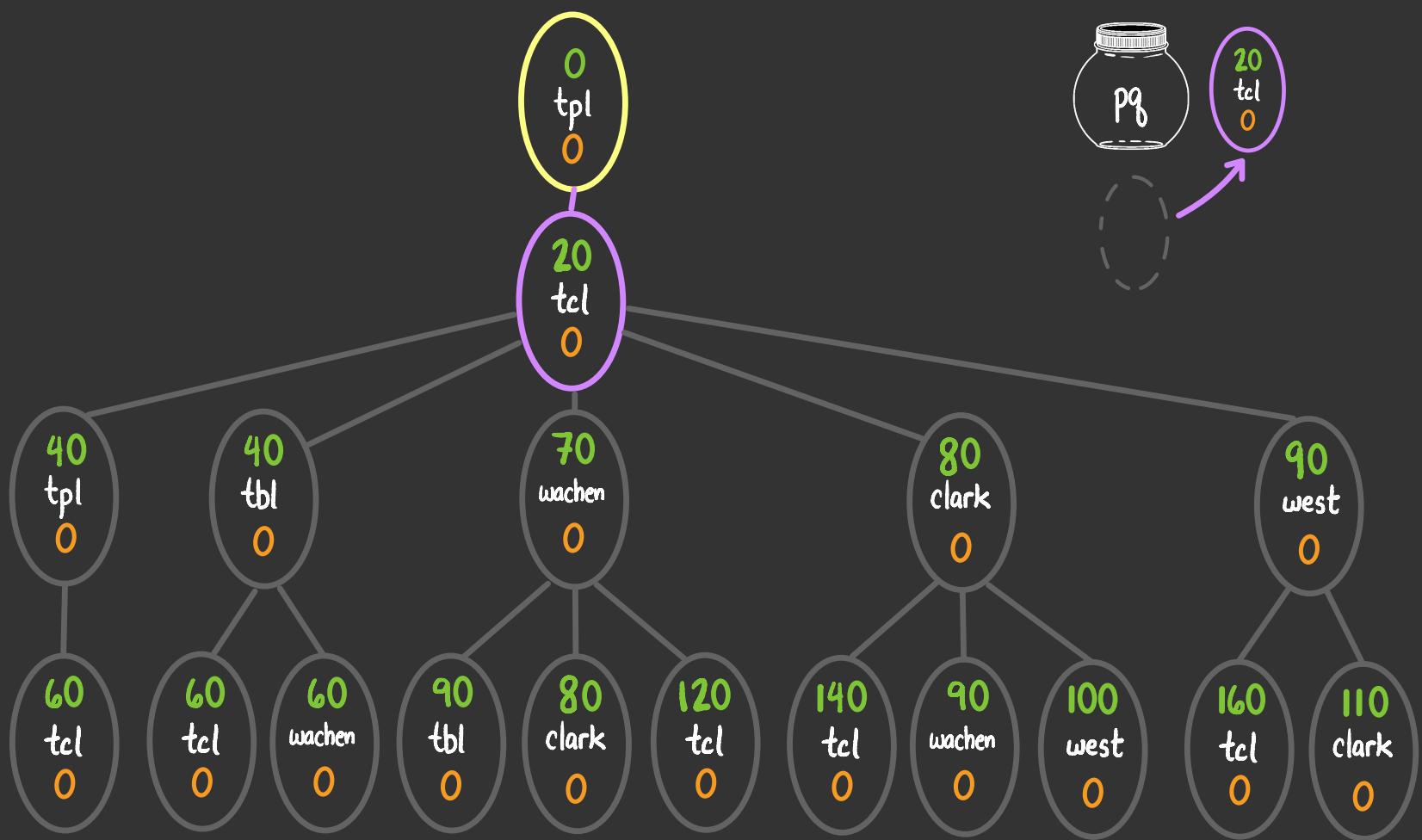


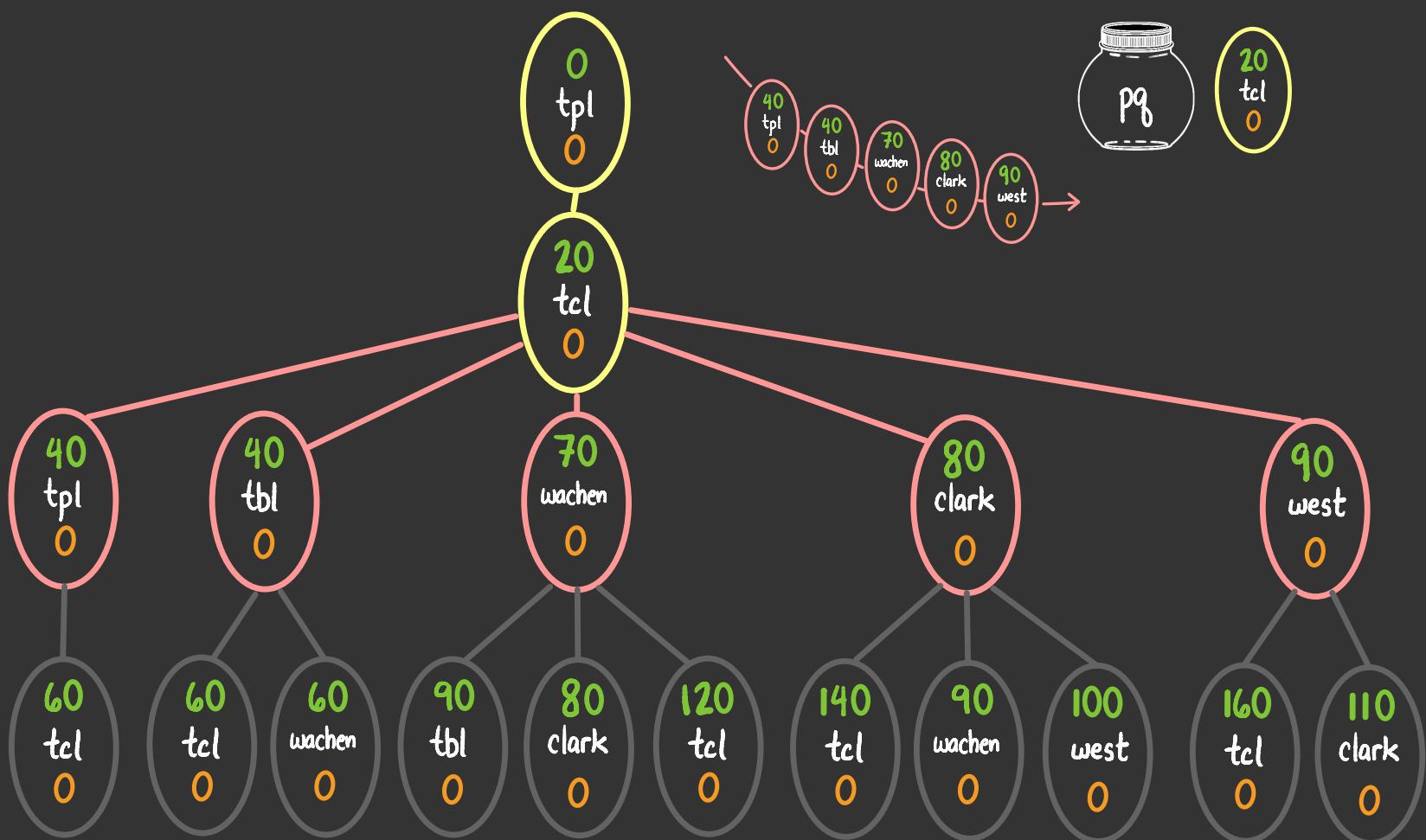


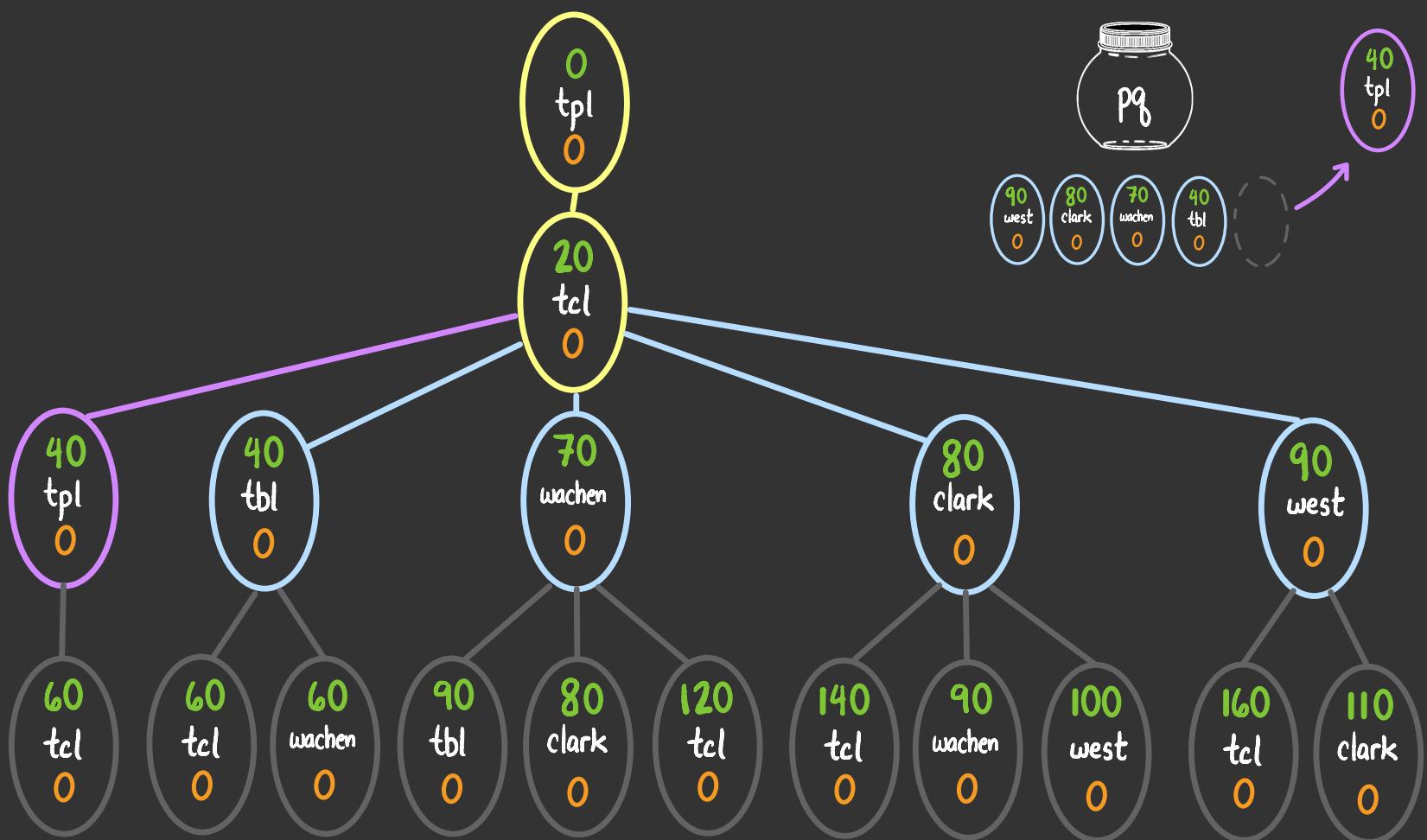




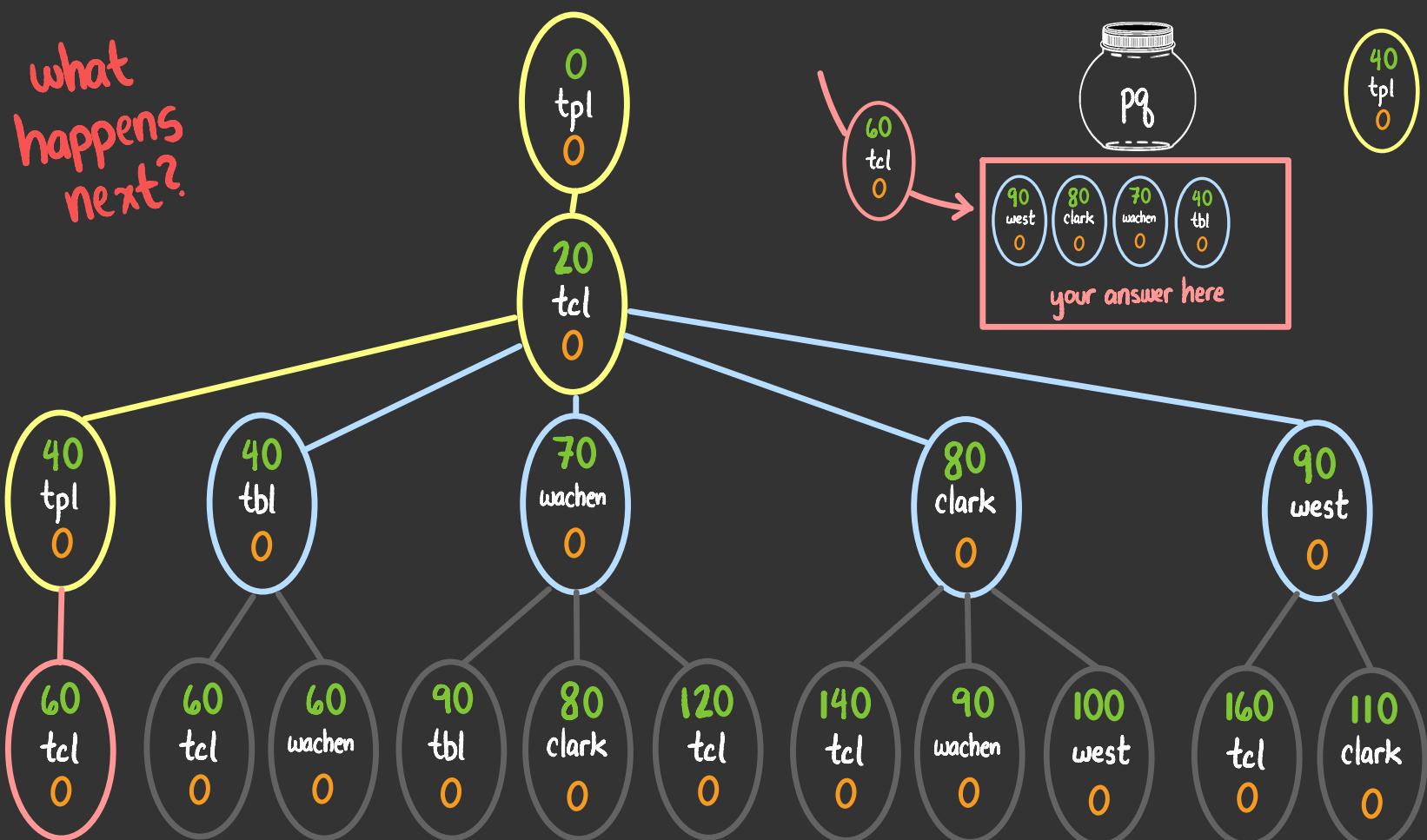


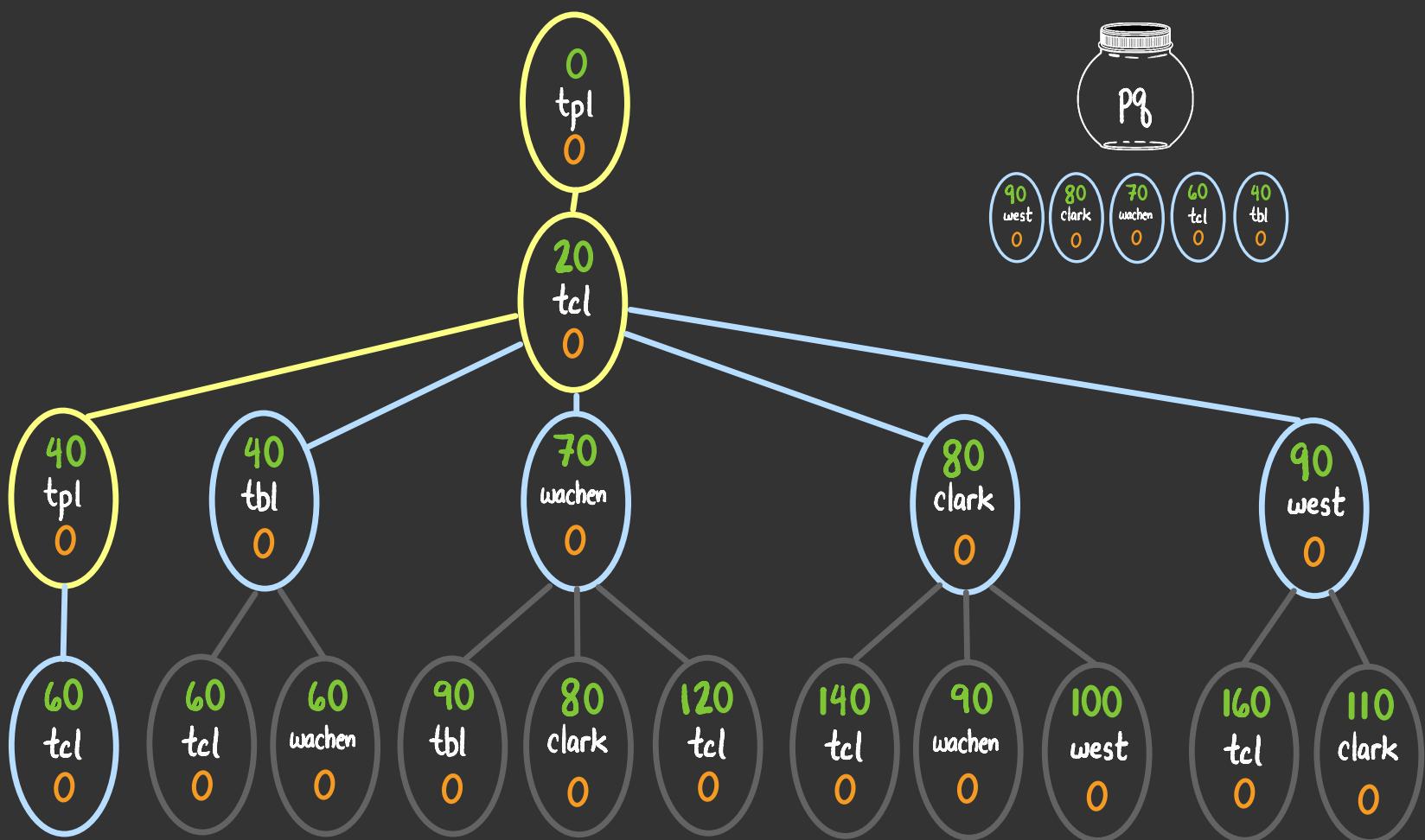


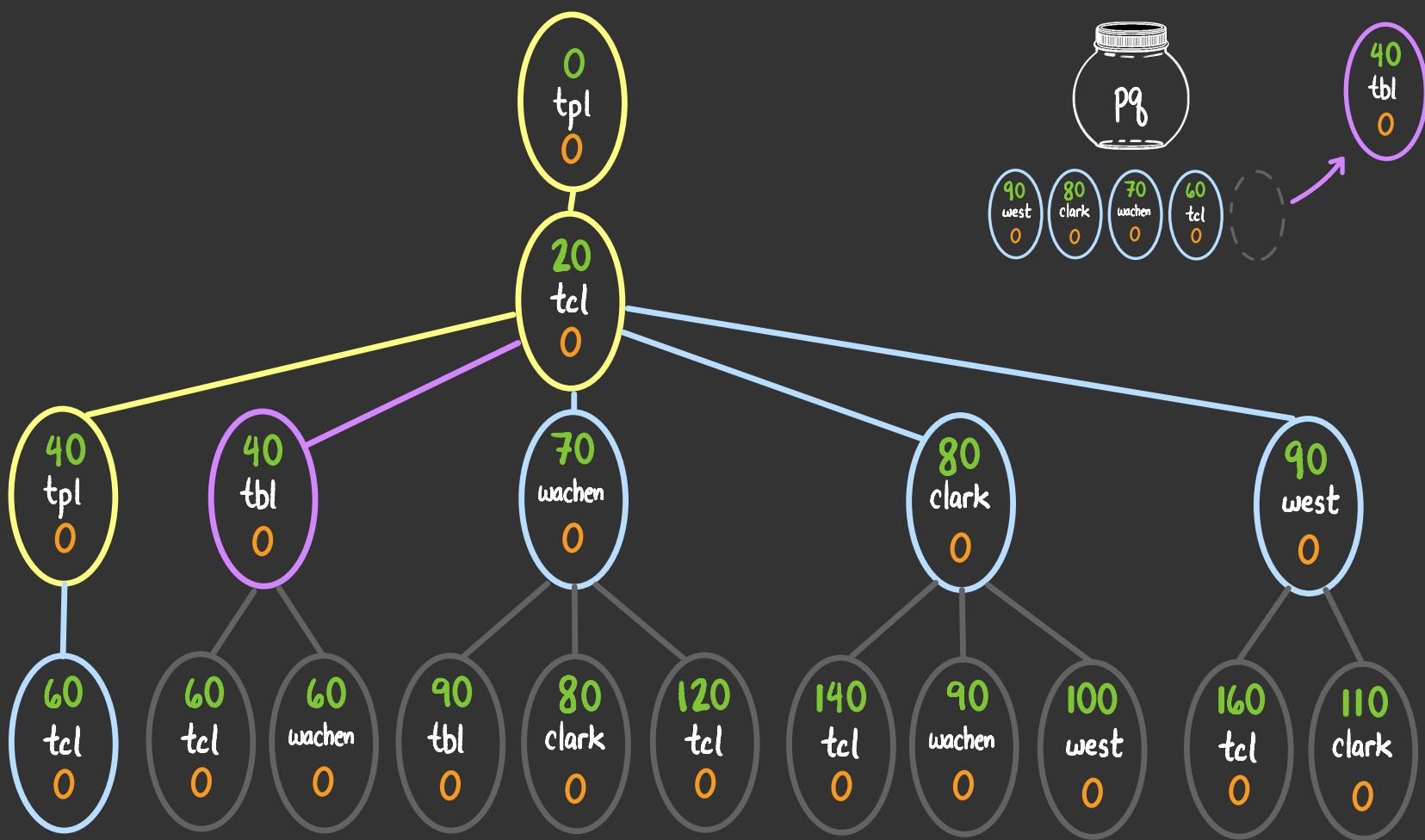


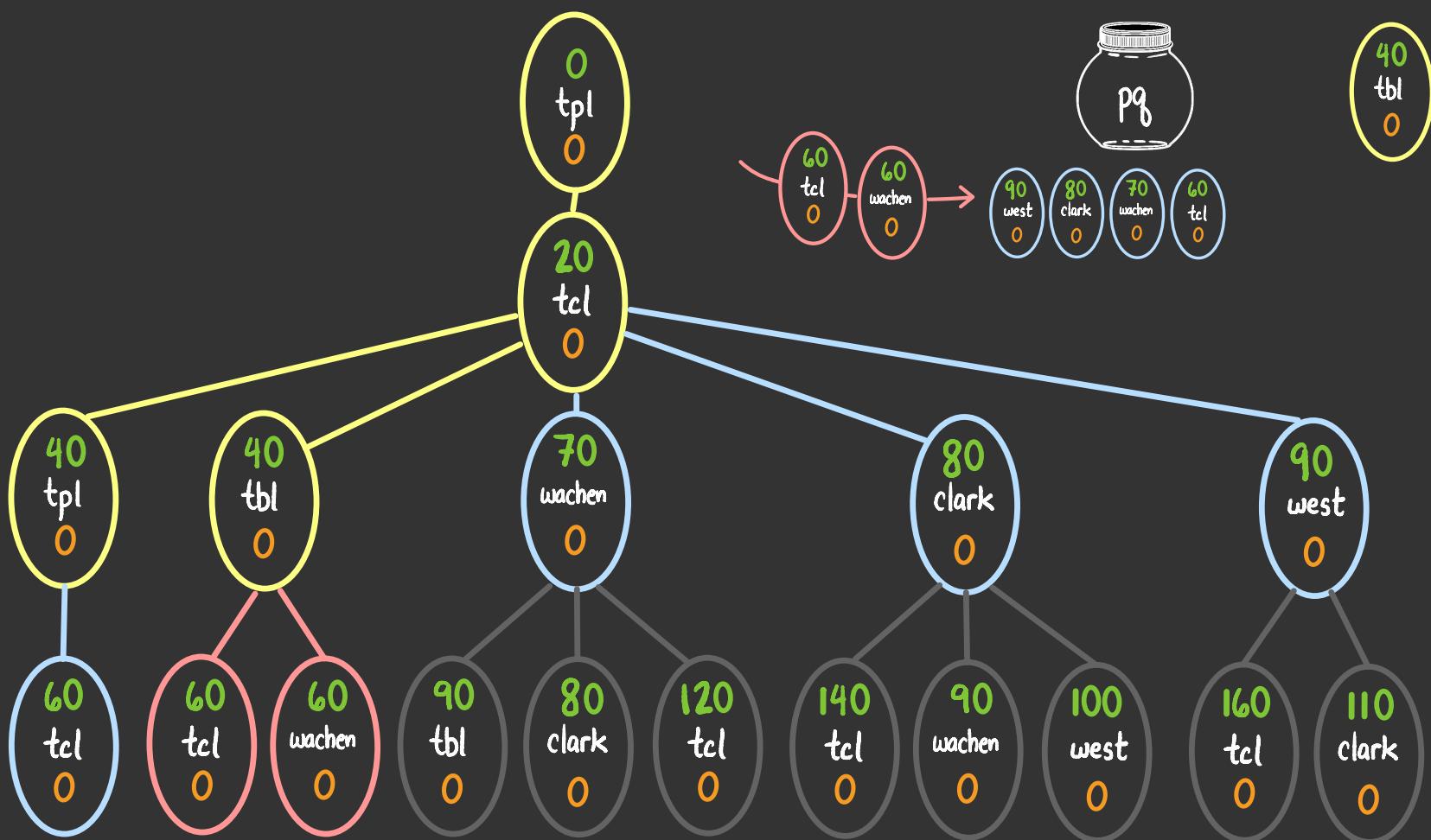


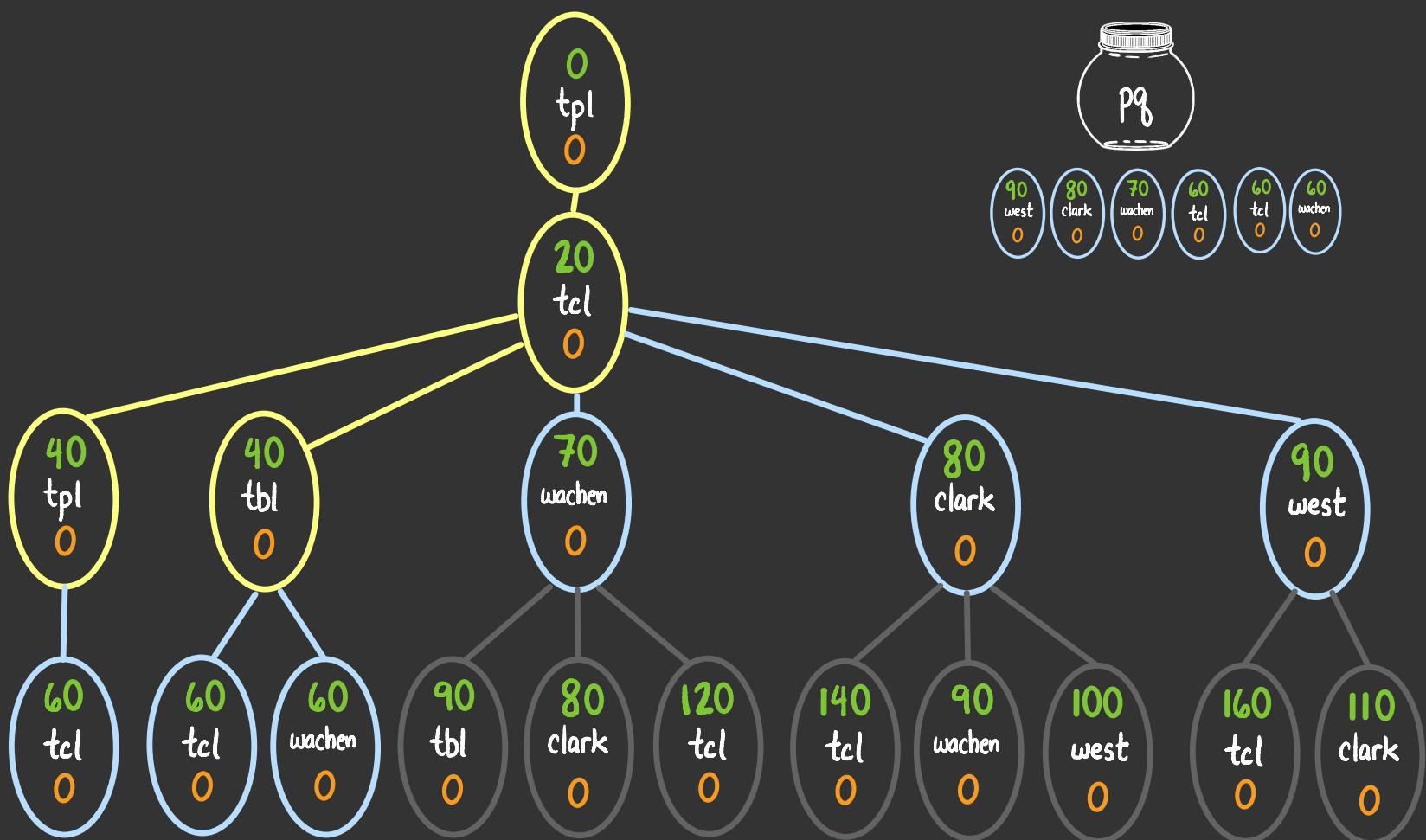
what  
happens  
next?

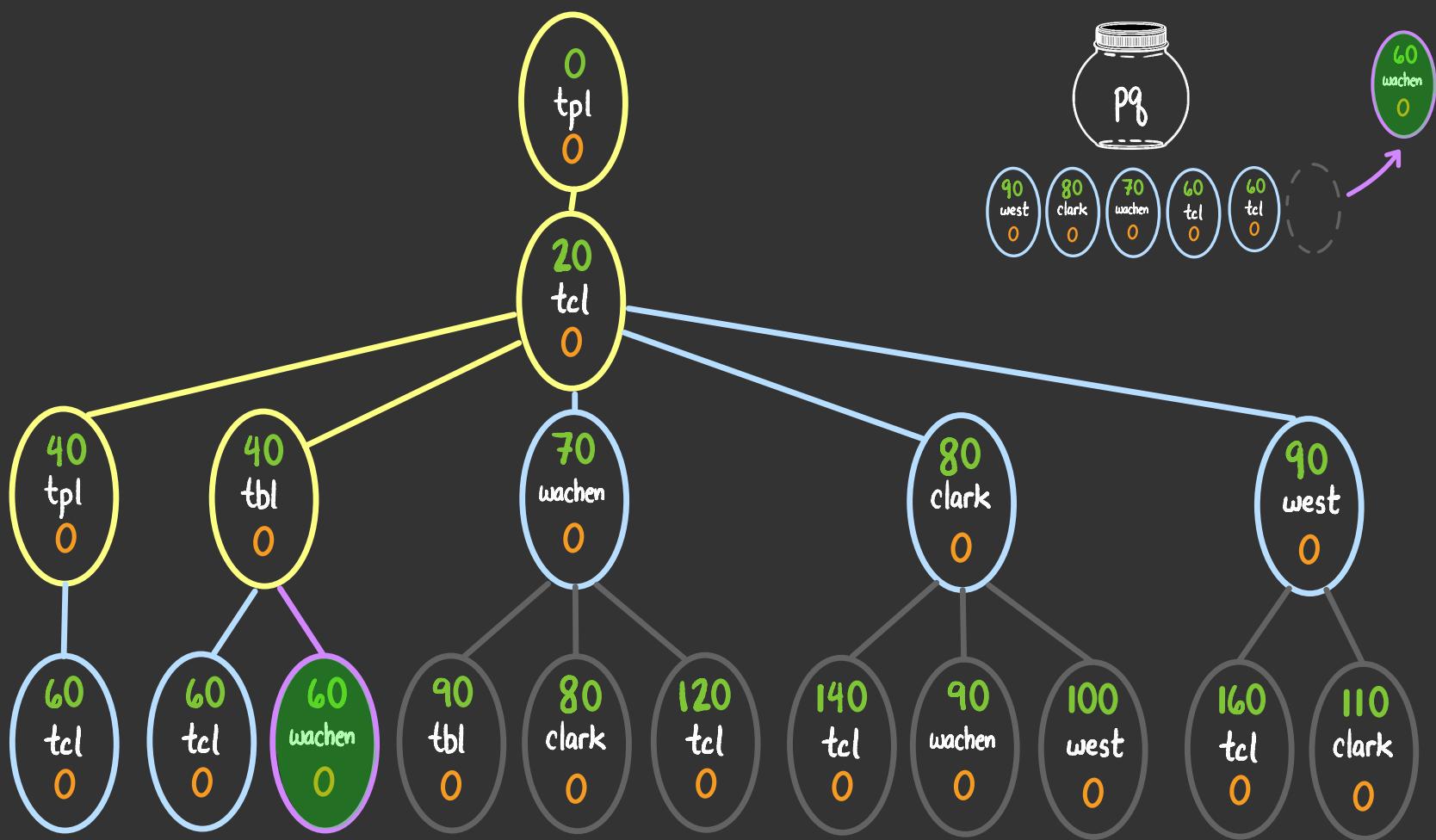


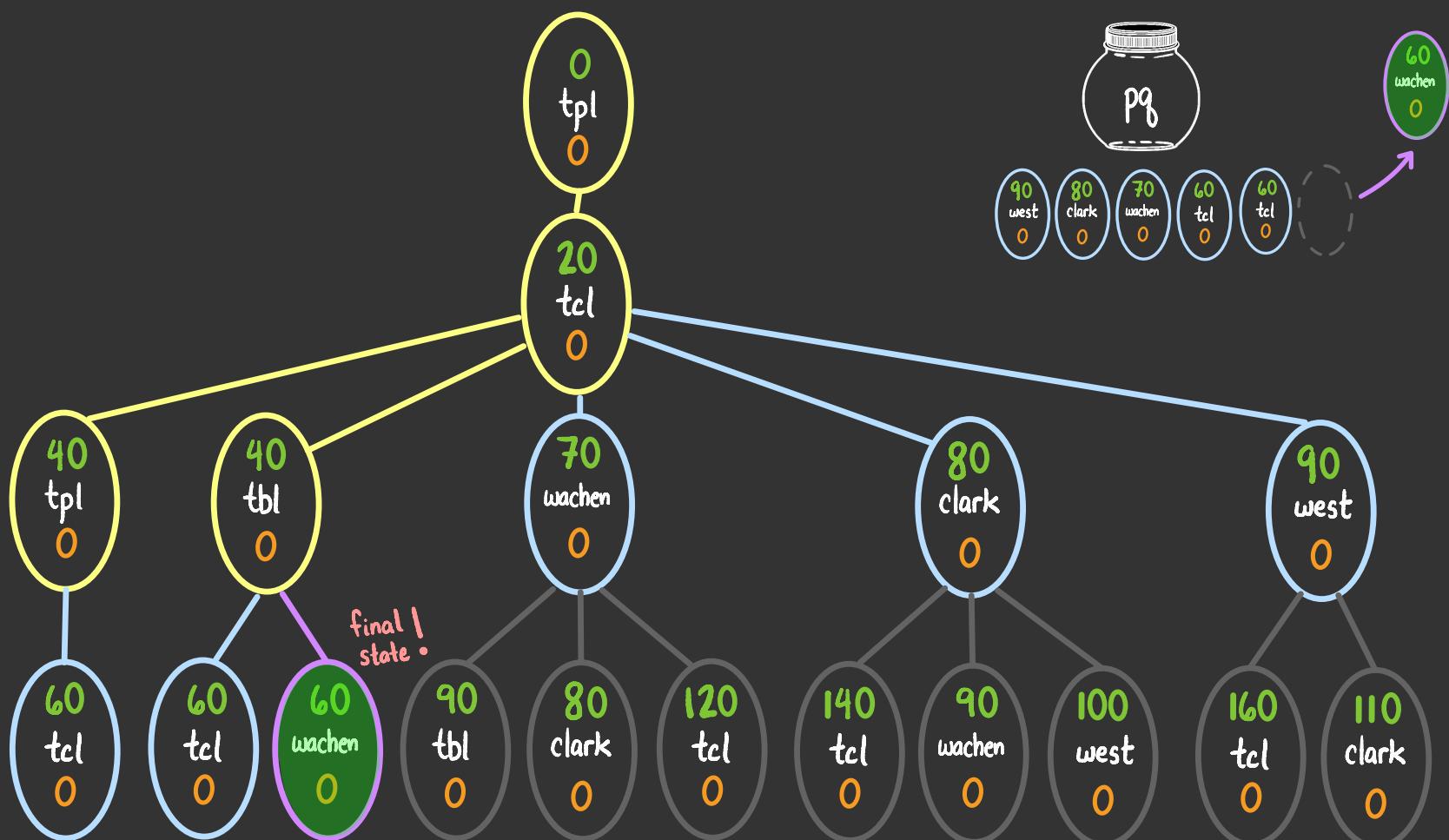












and it has the **lowest cost...**

uniform cost search is guaranteed to find the optimal solution

why? if search node  $n_1$  is visited before search node  $n_2$ , then  $g(n_1) \leq g(n_2)$

(can be proven by contradiction)

