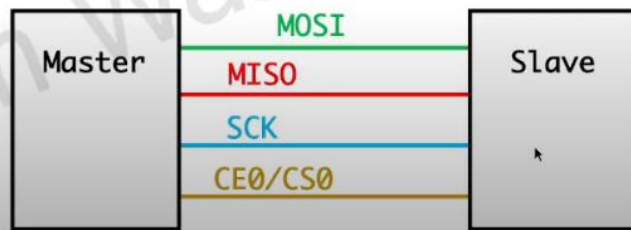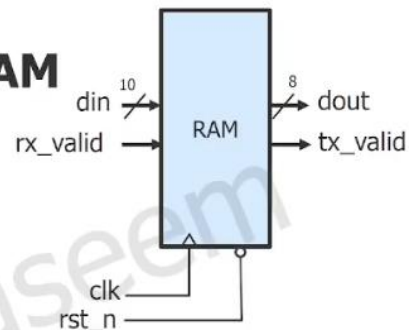# SPI Interface

- One of the most popular Interfaces nowadays
- Stands for Serial-Peripheral Interface
- Four Wires
  - MOSI: Master-Out-Slave-In
  - MISO: Master-In-Slave-Out
  - SCK: Clock
  - SS_n: Slave Select
- High Data Rates

Project: 2- Single-port Sync RAM

- **Parameters**
 - MEM_DEPTH, Default: 256
  - ADDR_SIZE, Default: 8
- **Ports**

Name	Type	Size	Description
din	Input	10 bits	Data Input
clk		1 bit	Clock
rst_n		1 bit	Active low asynchronous reset
rx_valid		1 bit	If HIGH: accept din[7:0] to save the write/read address internally or write a memory word depending on the most significant 2 bits din[9:8]
dout	Output	8 bits	Data Output
tx_valid		1 bit	Whenever the command is memory read the tx_valid should be HIGH

Project: 2- Single-port Sync RAM


Parameters
— MEM_DEPTH, Default: 256
— ADDR_SIZE, Default: 8

- Most significant din bit "din[9]" determines
 if it is a write or read command
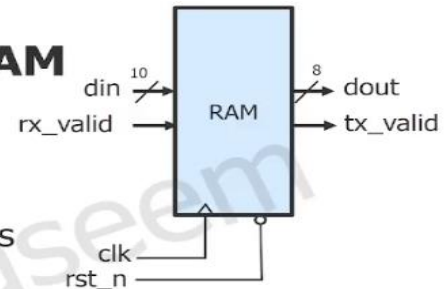
| Port	Din[9:8]	Command	Description
din	00	Write	Hold din[7:0] internally as write address
	01		Write din[7:0] in the memory with write address held previously
	10	Read	Hold din[7:0] internally as read address
	11		Read the memory with read address held previously, tx_valid should be HIGH, dout holds the word read from the memory, ignore din[7:0]

Project: 3- SPI Wrapper

SPI SLAVE WITH SINGLE PORT RAM
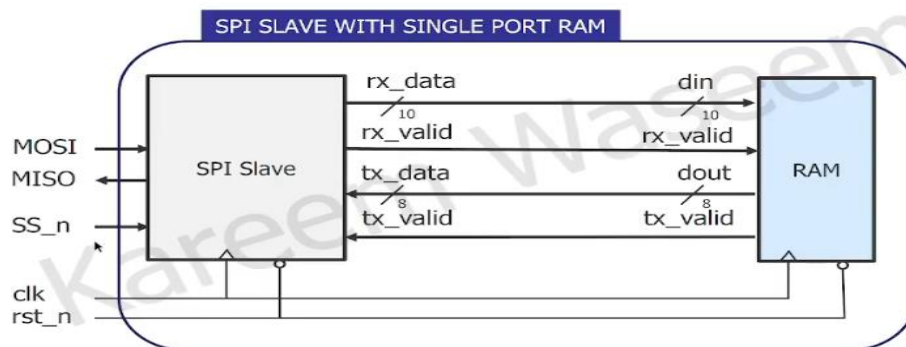
SPI SLAVE STATE TRANSITION DIAGRAM

CHK_CMD transitions will be explained in the class
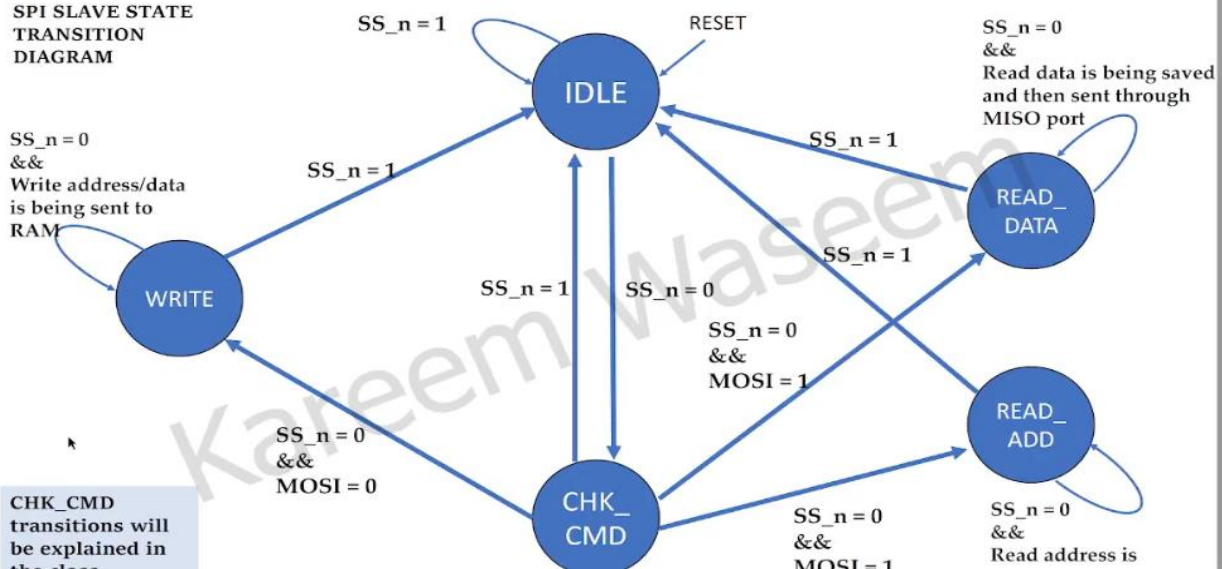
RAM Write Command – Write Address

1. Master will start the write command by sending the write address value, $rx\_data[9:8] = din[9:8] = 2'b00$

2. $SS\_n = 0$ to tell the SPI Slave that the master will begin communication

3. SPI Slave check the first received bit on MOSI port '0' which is a control bit to let the slave determine which operation will take place "write in this case". SPI Slave then expects to receive 10 more bits, the first 2 bits are "00" on two clock cycles and then the wr_address will be sent on 8 more clock cycles

4. Now the data is converted from serial "MOSI" to parallel after writing the $rx\_data[9:0]$ bus

5. rx_valid will be HIGH to inform the RAM that it should expect data on din bus

6. din takes the value of rx_data

7. RAM checks on $din[9:8]$ and find that they hold "00"

8. RAM stores $din[7:0]$ in the internal write address bus

9. $SS\_n = 1$ to end communication from Master side
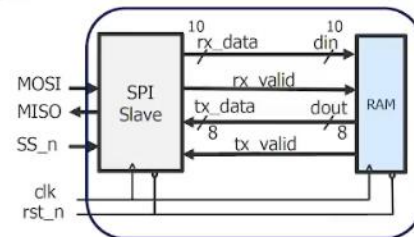
**SPI SLAVE FSM TRANSITIONS IN WRITE ADDRESS**
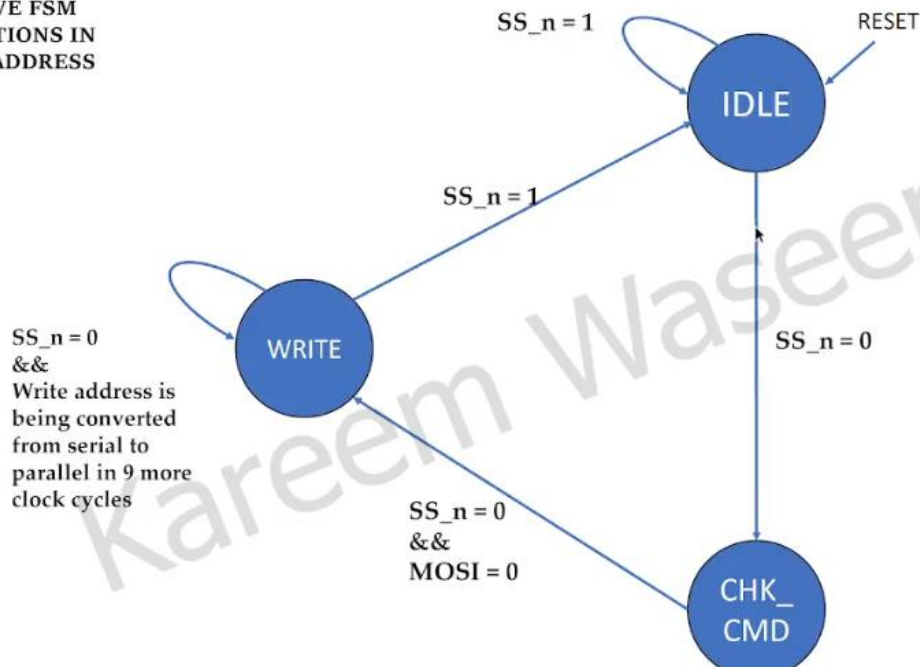
RAM Write Command – Write Data

1. Master will continue the write command by sending the write data value, rx_data[9:8] = din[9:8] = **2'b01**

2. SS_n = 0 to tell the SPI Slave that the master will begin communication

3. SPI Slave check the first received bit on MOSI port '0' which is a control bit to let the slave determine which operation will take place "write in this case". SPI Slave then expects to receive 10 more bits, the first 2 bits are "01" on two clock cycles and then the wr_**data** will be sent on 8 more clock cycles

4. Now the data is converted from serial "MOSI" to parallel after writing the rx_data[9:0] bus

5. rx_valid will be HIGH to inform the RAM that it should expect data on din bus

6. din takes the value of rx_data

7. RAM checks on din[9:8] and find that they hold **"01"**

8. RAM stores din[7:0] in the RAM with wr_address previously held
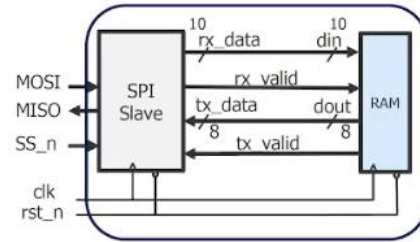
9. SS_n = 1 to end communication from Master side
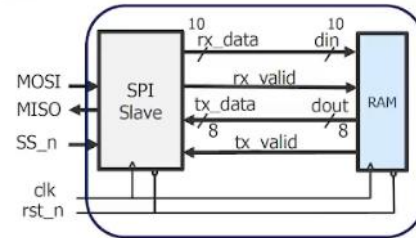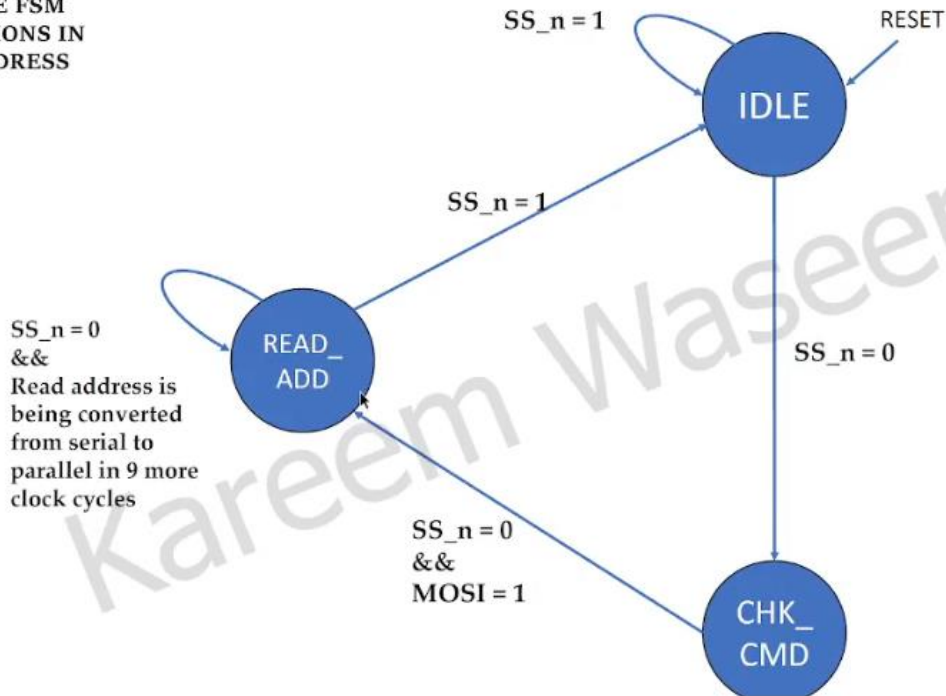

SPI SLAVE FSM TRANSITIONS IN WRITE DATA
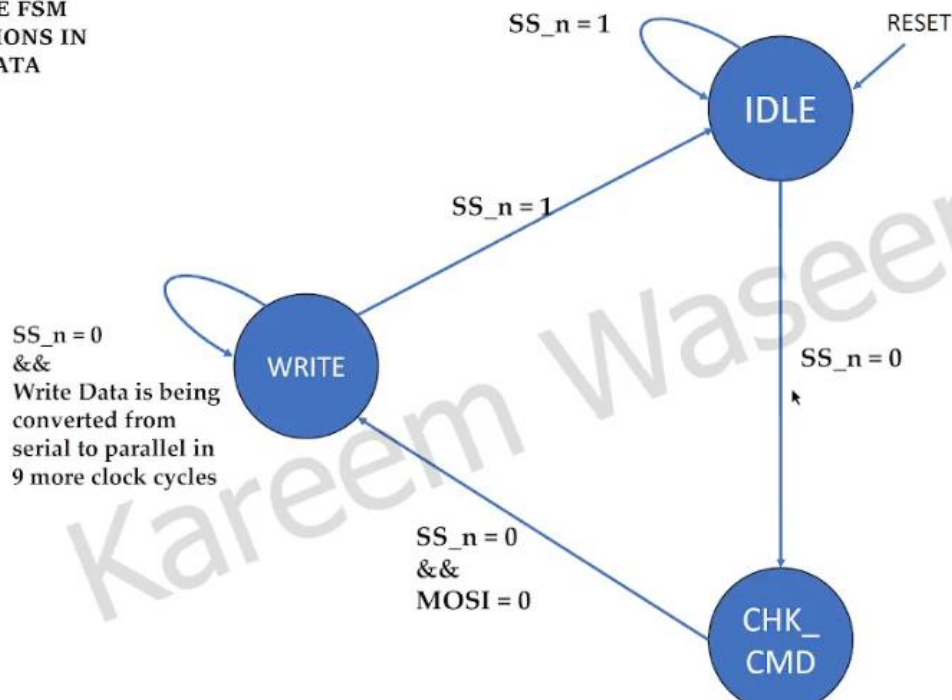
RAM Read Command – Read Address

1. Master will start the write command by sending the read address value, rx_data[9:8] = din[9:8] = 2'b10

2. SS_n = 0 to tell the SPI Slave that the master will begin communication

3. SPI Slave check the first received bit on MOSI port '1' which is a control bit to let the slave determine which operation will take place "read in this case". SPI Slave then expects to receive 10 more bits, the first 2 bits are "10" on two clock cycles and then the rd_address will be sent on 8 more clock cycles

4. Now the data is converted from serial "MOSI" to parallel after writing the rx_data[9:0] bus

5. rx_valid will be HIGH to inform the RAM that it should expect data on din bus

6. din takes the value of rx_data

7. RAM checks on din[9:8] and find that they hold "10"

8. RAM stores din[7:0] in the internal read address bus
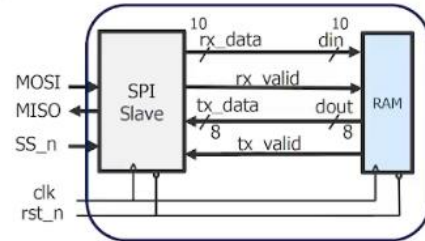
9. SS_n = 1 to end communication from Master side

SPI SLAVE FSM TRANSITIONS IN READ ADDRESS

RAM Read Command – Read Data

1. Master will start the write command by sending the read address value, rx_data[9:8] = din[9:8] = 2'b1**1**

2. SS_n = 0 to tell the SPI Slave that the master will begin communication

3. SPI Slave check the first received bit on MOSI port '**1**' which is a control bit to let the slave determine which operation will take place "read in this case". SPI Slave then expects to receive 10 more bits, the first 2 bits are "11" on two clock cycles and then dummy data will be sent and ignored since the master is waiting for the data to be sent from slave side

4. Now the data is converted from serial "MOSI" to parallel after writing the rx_data[9:0] bus

5. din takes the value of rx_data

6. RAM reads din[9:8] and find that they hold "11"

7. RAM will read from the memory with rd_address previously held

8. RAM will assert tx_valid to inform slave that data out is ready

9. Slave reads tx_data and convert it into serial out data on MISO port

10. SS_n = 1, Master ends communication after receiving data
    "8 clock cycles"

SPI SLAVE FSM TRANSITIONS IN READ ADDRESS