

We have focused on solving congruence equations, e.g.

$$x^n \equiv a \pmod{p} \quad x^n - a \equiv 0 \pmod{p} \text{ for prime modulus.}$$

But what about other moduli? $x^n - a \equiv 0 \pmod{N}$?

What can we say? Plenty!

Thm: If $f(x) = a_n x^n + \dots + a_1 x + a_0$ is a polynomial with integer coefficients, and N, M integers with $(N, M) = 1$, and $f(x) \equiv 0 \pmod{N}$, $f(x) \equiv 0 \pmod{M}$ both have a solution, then so does $f(x) \equiv 0 \pmod{NM}$.

Pf: Suppose $f(a) \equiv 0 \pmod{N}$, $f(b) \equiv 0 \pmod{M}$. We know that if $c \equiv a \pmod{N}$, then $f(c) \equiv f(a) \equiv 0 \pmod{N}$, so c is also a solution. we also know that $f(c) \equiv 0 \pmod{N}$ and $f(c) \equiv 0 \pmod{M} \Rightarrow N | f(c), M | f(c) \Rightarrow NM | f(c) \Rightarrow f(c) \equiv 0 \pmod{NM}$. So it is enough to show that there is a c with $\boxed{c \equiv a \pmod{N} \text{ and } c \equiv b \pmod{M}}$, i.e. there is a simultaneous solution to the system of eqns
$$\begin{aligned} x &\equiv a \pmod{N} \\ x &\equiv b \pmod{M} \end{aligned}$$
 ✓

But this is the Chinese Remainder Thm!

Quick proof of CRT: $(*) \begin{cases} x = a + NK \\ x = b + ML \end{cases}$, so want DD.

$$a + NK = b + ML \text{ some } k, l \text{ i.e. } b - a = ML - NK$$

But $(M, N) = 1 \Rightarrow 1 = Mb - Nk_0$ some b, k_0 so

$$b - a = M(ba)k_0 - N(b-a)k_0, \text{ so set } x = a + N((ba)k_0)$$

Check: any other solution x' of $(*)$ has $x \equiv x' \pmod{MN}$

So for every possible pair of solutions $f(a) \equiv 0 \pmod{N}$, $f(b) \equiv 0 \pmod{M}$ (mod N, M resp.), there is exactly one solution \uparrow to $f(x) \equiv 0 \pmod{MN}$

~~So # of incongruent (mod MN) solutions to $f(x) \equiv 0 \pmod{MN}$~~

So if we call $s(N) = \#$ of incongruent solutions \pmod{N} to $f(x) \equiv 0$, we in fact have

$$\boxed{s(NM) = s(N)s(M)}$$

~~So~~ More generally, if $N = n_1 \cdots n_k$ with

$$(n_i, n_j) = 1 \text{ for } i \neq j \text{ then}$$

$$s(N) = s(n_1) \cdots s(n_k) \quad (\text{by induction!})$$

So to decide if $f(x) \equiv 0$ has any solutions, it is enough, writing $N = p_1^{r_1} \cdots p_k^{r_k}$ $p_1 < \cdots < p_k$ prime, if

$f(x) \equiv 0 \pmod{p_i^{r_i}}$ for each i .

We showed: For p prime, deciding if

$x^p \equiv a \pmod{p}$ has a solution is straightforward:

check if $a^{\frac{p-1}{p}} \equiv 1 \pmod{p}$. Yes \rightarrow ~~many~~ $(p-1)$ solutions
No \rightarrow no solution

But it wasn't really important that p be prime! Only that there is a primitive root of $1 \pmod{p}$ (and its order is $p-1$.) More generally:

If there is a primitive root of $1 \pmod{N}$, then

$x^N \equiv a \pmod{N}$ has a solution (for $(a, N) = 1$) \iff

$a^{\frac{\phi(N)}{(N, N)}} \equiv 1 \pmod{N}$ Our proof goes straight through!

I should have stated that

There is a primitive root of $1 \pmod{n} \iff$

$n = \text{one of } 2, 4, p^k, 2p^k \text{ for } p = \text{odd prime.}$

Combining with the previous result, deciding if

$x^N \equiv a \pmod{N}$ has any solution (and how many!) boils

down to deciding if $x^N \equiv a \pmod{p_i^{r_i}}$ has solutions, which for p_i odd is straightforward; for $p=2$ it is a little trickier!

Why is there a primitive root of 1 mod p^n for p odd and prime?

There is a primitive root mod p , call it a . Then

Claim: $a + tp$ is a primitive root mod p^2 for

all but one value of t .
 $t \neq 0$ and $p \nmid (a + tp)$

$$(a + tp)^{p-1} \equiv 1 \pmod{p^2} \implies (a + tp)^{p-1} \equiv a^{p-1} \pmod{p^2} \text{ so } p-1 \nmid k$$

either $k \equiv p-1$ or $k \equiv 0 \pmod{p-1}$.
 ~~$k \equiv 0 \pmod{p-1}$~~ so either

If $k \equiv p-1$ then $a + tp$ is a primitive root.
 So we want to show that $k \equiv p-1$ for only one value of t .

of t .
 i.e. if $(a + tp)^{p-1} \equiv 1 \pmod{p^2}$ then $(a + tp)^{p-1} \not\equiv 1 \pmod{p^2}$ for $t \neq t$

$$(a + tp)^{p-1} = ((a + tp) + np)^{p-1} \leftarrow \begin{bmatrix} p \\ 0 \\ n \end{bmatrix}$$

$$= (a + tp)^{p-1} + (p-1)(a + tp)^{p-2}(np) + \binom{p-1}{2}(a + tp)^{p-2}(np)^2 + \dots + 0$$

$$\equiv 1 + (p-1)(a + tp)^{p-2}(np) + \dots + 0 \pmod{p^2}$$

$$\equiv 1 + [(p-1)(a + tp)^{p-2}n]p \not\equiv 1 \pmod{p^2}$$

So each primitive root mod p ($\phi(p-1)$ of them) give $\phi(p-1)$ distinct prim roots mod p^2
 \Rightarrow there are $\phi(p-1)\phi(p-1) = \phi(\phi(p^2))$ prim roots mod p^2 .

Then
 If a is a prim root mod p^2 for p odd prime
 then a is a prime root mod p^r for all $r \geq 2$.

By induction: $r \geq 2$ ✓ Assume true for $r-1$

$n = \text{ord}_{p^r}(a)$ then $a^{\frac{n}{p}} \not\equiv 1 \Rightarrow a^{\frac{n}{p^2}} \not\equiv 1 \Rightarrow p(p-1) \mid n$

and $n \mid \phi(p^r) = p^{r-1}(p-1)$, $\therefore n = p^r(p-1)$ for some $1 \leq \beta \leq r-1$.

Claim: $\beta = r-1$. Proof: show $a^{p^{r-2}(p-1)} \not\equiv 1$

We know that $\text{ord}_{p^{r-1}}(a) = \phi(p^{r-1}) = p^{r-2}(p-1)$, and
 $\text{ord}_{p^{r-2}}(a) = p^{r-3}(p-1)$ so $a^{p^{r-2}(p-1)} \not\equiv 1$ but $a^{p^{r-1}} \equiv 1$

so $a^{p^{r-2}(p-1)} = 1 + sp^{r-2}$ with $p \nmid s$. then

$a^{p^{r-2}(p-1)} = (a^{-1})^p = (1 + sp^{r-2})^p = 1^p + psp^{r-2} + \text{high powers of } p$
 $\equiv 1 + sp^{r-1} \not\equiv 1 \pmod{p^r}$ since $p \nmid s$.