Math 445 Number Theory

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Proposition: If f is a polynomial with integer coefficients and (M, N) = 1, then the congruence equation $f(x) \equiv 0 \pmod{MN}$ has a solution \Leftrightarrow the equations $f(x) \equiv 0 \pmod{M}$ and $f(x) \equiv 0 \pmod{N}$ both do.

The direction (\Rightarrow) is immediate; MN|f(x) implies M|f(x) and N|f(x), since M,N|MN. The point to (\Leftarrow) is that the solutions we know of to each of the two equations might be different: $f(x_1) \equiv 0 \pmod M$ and $f(x_2) \equiv 0 \pmod N$. We need that a single number solves both, since then $M|f(x_0)$ and $N|f(x_0)$, and then (M,N)=1 implies that $MN|f(x_0)$. To do this, we use the fact that f is a polynomial, since then if $a \equiv b \pmod n$, then $f(a) \equiv f(b)b \pmod n$. So if we suppose that we have found a and b with $f(a) \equiv 0 \pmod M$ and $f(b) \equiv 0 \pmod N$, then any x satisfying both $x \equiv a \pmod M$ and $x \equiv b \pmod N$ will satisfy both $f(x) \equiv 0 \pmod M$ and $f(x) \equiv 0 \pmod N$ simultaneously, as desired. So it is enough to show that for any a, b, there is an x which simultaneously satisfies

$$x \equiv a \pmod{M}$$
 and $x \equiv b \pmod{N}$

But since (M,N)=1, this is true by the Chinese Remainder Theorem. In fact, finding x is a matter of solving x=a+Mi, x=b+Nj, so we need a+Mi=b+Nj, so b-a=Mi-Nj. But since (M,N)=1, we can use the Euclidean algorithm to write $1=MI_0+NJ_0$, and then $i=(b-a)I_0, j=-(b-a)J_0$ will work, allowing us to solve for x. In fact, since the only other I,J which will work are $I=I_0+kN, J=J_0-kM$, we find that our solution x is unique modulo MN.

For any pair of solutions a,b to $f(a) \equiv 0 \pmod M$ and $f(b) \equiv 0 \pmod N$ there is a unique corresponding $x \mod MN$ (with $x \equiv a \pmod M$) and $x \equiv b \pmod N$) satisfying $f(x) \equiv 0 \pmod MN$. Introducing the notation $S(n) = 0 \pmod N$ the number of solutions, mod n, to the equation $f(x) \equiv 0 \pmod n$, we then have shown that S(MN) = S(M)S(N) whenever (M,N) = 1. So by induction, whenever $N_1, \ldots N_k$ are relatively prime, $S(N_1 \cdots N_k) = S(N_1) \cdots S(N_k)$.

So if $N = p_1^{k_1} \cdots p_r^{k_r}$ is the prime factorization of the odd number N, then if (a, N) = 1 (so $(a, p_i) = 1$ for each i) we have $x^n \equiv a \pmod{N}$ has solutions $\Leftrightarrow x^n \equiv a \pmod{p_i^{k_i}}$ does for every i, and we know how to determine when that occurs.

Quadratic Residues: If $x^2 \equiv a \pmod{n}$ has a solution, a is a quadratic residue modulo n. If it doesn't, a is a quadratic non-residue modulo n. Euler's Criterion gives us a test: if p is a prime, then a is a quadratic residue mod $n \Leftrightarrow a^{\frac{p-1}{2}} \equiv 1 \pmod{p}$. But this may require a lot of calculation if p is large; our next task is to find a quicker way.

To talk about things in a compact manner, we introduce the $Legendre\ symbol$; for p an odd prime,

$$\left(\frac{a}{p}\right) = \begin{cases} 0 & \text{if } p|a\\ 1 & \text{if } a \text{ is a quadratic residue mod } p\\ -1 & \text{if } a \text{ is a quadratic non-residue mod } p \end{cases}$$

By Euler's criterion, this really means $\left(\frac{a}{p}\right) \equiv a^{\frac{p-1}{2}} \pmod{p}$, but our goal is to find a *quicker* way to compute it!