

$$4. T(n) = 2^n T\left(\frac{n}{2}\right) + n$$

Master method is not applicable since a is not constant

$$5. T(n) = 16T\left(\frac{n}{4}\right) + n$$

$$a = 16, b = 4, f(n) = n$$

$$n^{\log_b a} = n^{\log_2 16} = n^4$$

$$f(n) < n^{\log_b a}$$

$$\therefore \text{rem-root} \quad \therefore f(n) = O(n^4)$$

$$6. T(n) = 2T\left(\frac{n}{2}\right) + n \log n$$

$$a = 2, b = 2, f(n) = n \log n$$

$$n^{\log_b a} = n^{\log_2 2} = n$$

$$f(n) > n^{\log_b a}$$

$$\therefore T(n) = O(n \log n)$$

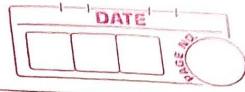
$$7. T(n) = 2T\left(\frac{n}{2}\right) + \frac{n}{\log n}$$

$$a = 2, b = 2$$

$$f(n) = \frac{n}{\log n}$$

master method is not applicable because of non-polynomial difference between $f(n)$ and $n^{\log_b a}$

Master method practice



~~Q1~~

$$1. \quad T(n) = 3T\left(\frac{n}{2}\right) + n^2$$

$$a = 3, b = 2, f(n) = n^2$$

$$n^{\log_b a} = n^{\log_2 3} = n^{1.58}$$

$$n^2 > n^{\log_b a} \quad \therefore f(n) > n^{\log_b a}$$

$$f(n) = \Theta(n^{\log_b a + \epsilon}) \\ \therefore T(n) = \Theta(n^2)$$

$$2. \quad T(n) = 4T\left(\frac{n}{2}\right) + n^2$$

$$a = 4, b = 2, f(n) = n^2$$

$$n^{\log_b a} = n^{\log_2 4} = n^2$$

$$f(n) = n^{\log_b a}$$

$$\therefore T(n) = \Theta(n^2 \log_2 n)$$

$$3. \quad T(n) = T\left(\frac{n}{2}\right) + 2^n$$

$$a = 1, b = 2, f(n) = 2^n$$

$$n^{\log_b a} = n^{\log_2 1} = n^0 = 1$$

$$f(n) > n^{\log_b a}$$

$$f(n) = \Omega(n^{\log_b a} + \epsilon)$$

$$T(n) = O(f(n)) = \Theta(2^n) \quad \Theta(2^n)$$

8. $T(n) = 2T\left(\frac{n}{4}\right) + n^{0.5}$

$$a=2, b=4, f(n)=n^{0.5}$$

$$n^{\log_b a} = 8 n^{\log_4 2} = n^{0.5}$$

$$f(n) \geq n^{\log_b a}$$

$$f(n) = n^{\log_b a} + \epsilon n^{0.5}$$

$$\therefore T(n) = \Theta(n^{0.5})$$

9. $T(n) = \frac{1}{2} T\left(\frac{n}{2}\right) + \frac{1}{n}$

$$a = \frac{1}{2}, b = 2$$

master method is not applicable because
 $a < 1$

10. $T(n) = 16T\left(\frac{n}{4}\right) + n!$

$$a=16, b=4, f(n)=n!$$

$$n^{\log_b a} = n^{\log_4 16} = n^4$$

$$f(n) > n^{\log_b a}$$

$$T(n) = \Theta(n!)$$

$$11. \tau(n) = \sqrt{2} \tau\left(\frac{n}{2}\right) + \log n$$

$$a = \sqrt{2}, b = 2, f(n) = \log n$$

$$n^{\log_b a} = n^{\log_2 \sqrt{2}} = n^{1/2}$$

$$f(n) \leq n^{\log_b a}$$

$$\tau(n) = O(n^{\log_b a}) = O(\sqrt{n})$$

$$12. \tau(n) = 3\tau\left(\frac{n}{2}\right) + n$$

$$a = 3, b = 2, f(n) = n$$

$$n^{\log_b a} = n^{\log_2 3} = n^{1.58}$$

$$f(n) \leq n^{\log_b a}$$

$$\tau(n) = O(n^{\log_b a}) = O(n^{1.58})$$

$$13. \tau(n) = 3\tau\left(\frac{n}{3}\right) + \sqrt{n}$$

$$a = 3, b = 2, f(n) = \sqrt{n} = n^{1/2}$$

$$n^{\log_b a} = n^{\log_3 3} = n \therefore f(n) \leq n^{\log_b a}$$

$$\therefore \tau(n) = O(n^{\log_b a}) = O(n)$$

14. $T(n) = 4T\left(\frac{n}{2}\right) + cn$

$a = 4, b = 2, f(n) = cn$

$$n^{\log_b a} = n^{\log_2 4} = n^2$$
$$f(n) < n^{\log_b a}$$

$$T(n) = O(n^{\log_b a}) = O(n^2)$$

15. $T(n) = 3T\left(\frac{n}{3}\right) + n \log n$

$a = 3, b = 3, f(n) = n \log n$

$$n^{\log_b a} = n^{\log_3 3} = n^{0.17}$$

$$f(n) > n^{\log_b a}$$

$$T(n) = O(n \log n)$$

16. $T(n) = 3T\left(\frac{n}{3}\right) + \frac{n}{2}$

$a = 3, b = 3 \Rightarrow f(n) = \frac{n}{2}$

$$n^{\log_b a} = n^{\log_3 3} = n$$

$$f(n) = n^{\log_b a}$$

$$\therefore T(n) = O(n \log_3 n)$$

$$17. T(n) = 6T\left(\frac{n}{3}\right) + n^2 \log n$$

$$a=6, b=3, f(n)=n^2 \log n$$

$$n^{\log_b a} = n^{\log_3 6} =$$

$$\therefore f(n) > n^{\log_3 6}$$

$$\therefore T(n) = O(n^2 \log n)$$

$$18. T(n) = 4T\left(\frac{n}{2}\right) + \frac{n}{\log n}$$

$$a=4, b=2, f(n) = \frac{n}{\log n}$$

$$n^{\log_b a} = n^{\log_2 4} = n^2$$

$$f(n) = \underline{<} n^{\log_2 4} = O(n^2)$$

$$T(n) = O(n^2)$$

$$19. T(n) = 6T\left(\frac{n}{2}\right) + n^2 \log n$$

$$f(n) = -n^2 \log n$$

it cannot be -ve

\therefore master method not applicable

$$20. T(n) = 7T\left(\frac{n}{3}\right) + n^2$$

$$a=7, b=3, f(n)=n^2$$

$$n^{\log_b a} = n^{\log_3 7} = n^{1.77}$$

$$f(n) > n^{\log_3 7}$$

$$T(n) = O(n^2)$$

a.21 $T(n) = 4T\left(\frac{n}{2}\right) + \log n$

$a = 4, b = 2, f(n) = \log n$

~~PLZ~~ $n^{\log_b a} = n^{\log_2 4} = n^2$

$f(n) < n^{\log_b a}$

$T(n) = O(n^{\log_b a}) = O(n^2)$

a.22 $T(n) = T\left(\frac{n}{2}\right) + n(2 - \cos n)$

~~PLZ~~. $a = 1, b = 2, f(n) = n(2 - \cos n)$

$f(n) =$

$n^{\log_b a} = n^{\log_2 1} = n^0 = 1$

$f(n) > n^{\log_b a}$

$T(n) = O(n(2 - \cos n))$