

# **CPT\_S 260 Intro to Computer Architecture**

## **Lecture 16**

**Intro to MIPS V**  
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**Ganapati Bhat**  
**School of Electrical Engineering and Computer Science**  
**Washington State University**

# Recap: Conditional Operations

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- **Branch to a labeled instruction if a condition is true**
  - Otherwise, continue sequentially
- **beq rs, rt, L1**
  - if (rs == rt) branch to instruction labeled L1;
- **bne rs, rt, L1**
  - if (rs != rt) branch to instruction labeled L1;
- **j L1**
  - unconditional jump to instruction labeled L1

# Recap: Compiling Loop Statements

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- **C code:**

```
while (save[i] == k)
    i += 1;
```

– i in \$s3, k in \$s5, address of save in \$s6

- **Compiled MIPS code:**

```
Loop:  sll $t1,$s3,2           # Temp reg $t1 = i * 4
      add $t1,$t1,$s6         # $t1 = address of save[i]
      lw $t0,0($t1)          # Temp reg $t0 = save[i]
      bne $t0,$s5, Exit      # go to Exit if save[i] ≠ k
      addi $s3,$s3,1         # i = i + 1
      j Loop                 # go to Loop
```

Exit:

# More Conditional Operations

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- **Set result to 1 if a condition is true**
  - Otherwise, set to 0
- **`slt rd, rs, rt`**
  - if ( $rs < rt$ )  $rd = 1$ ; else  $rd = 0$ ;
- **`slti rt, rs, constant`**
  - if ( $rs < \text{constant}$ )  $rt = 1$ ; else  $rt = 0$ ;
- **Use in combination with `beq`, `bne`**
  - `slt $t0, $s1, $s2   # if ($s1 < $s2)`
  - `bne $t0, $zero, L   #    branch to L`

# Recap: Example Switch Statements

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- The simplest way to implement switch is via a sequence of conditional tests, turning the switch statement into a chain of if-then-else statements.

What is the MIPS assembly code assuming f-k correspond to registers \$s0-\$s5 and \$t2 contains 4 and \$t4 contains base address of JumpTable?

```
switch(k) {  
    case 0: f=i+j;break;  
    case 1: f=g+h;break;  
    case 2: f=g-h;break;  
    case 3: f=i-j;break;  
}
```

# MIPS Code for Switch

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```
add $t1,$s5,$s5      #$t1 =2*k
add $t1,$t1,$t1       #$t1 =4*k
add $t1,$t1,$t4       #$t1=address of JumpTable[k]
lw  $t0,0($t1)        #$t0=JumpTable[k]
jr  $t0               #jump based on register $t0
L0:  add $s0,$s3,$s4
      j Exit
L1:  add $s0,$s1,$s2
      j Exit
L2:  sub $s0,$s1,$s2
      j Exit
L3:  sub $s0,$s3,$s4
Exit:

switch(k) {
case 0: f=i+j;break;
case 1: f=g+h;break;
case 2: f=g-h;break;
case 3: f=i-j;break;
}
```

# Procedure Calling

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- **Steps required**
  - 1. Place parameters in registers**
  - 2. Transfer control to procedure**
  - 3. Acquire storage for procedure**
  - 4. Perform procedure's operations**
  - 5. Place result in register for caller**
  - 6. Return to place of call**

# Register Usage

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- **\$a0 – \$a3: arguments (reg's 4 – 7)**
- **\$v0, \$v1: result values (reg's 2 and 3)**
- **\$t0 – \$t9: temporaries**
  - Can be overwritten by callee
- **\$s0 – \$s7: saved**
  - Must be saved/restored by callee
- **\$gp: global pointer for static data (reg 28)**
- **\$sp: stack pointer (reg 29)**
- **\$fp: frame pointer (reg 30)**
- **\$ra: return address (reg 31)**
- **Program counter: A fixed register to hold the address of the instructions being executed**



# Procedure Call Instructions

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- **Procedure call: jump and link**

`jal ProcedureLabel`

- Address of following instruction put in \$ra
- Jumps to target address

- **Procedure return: jump register**

`jr $ra`

- Copies \$ra to program counter
- Can also be used for computed jumps
  - » e.g., for case/switch statements

# Leaf Procedure Example

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- C code:

```
int leaf_example (int g, h, i, j)
{ int f;
  f = (g + h) - (i + j);
  return f;
}
```

- Arguments g, ..., j in \$a0, ..., \$a3
- f in \$s0 (hence, need to save \$s0 on stack)
- Result in \$v0

# Leaf Procedure Example

## ■ MIPS code:

leaf_example:			
addi	\$sp	\$sp	-4
sw	\$s0	0(\$sp)	
add	\$t0	\$a0	\$a1
add	\$t1	\$a2	\$a3
sub	\$s0	\$t0	\$t1
add	\$v0	\$s0	\$zero
lw	\$s0	0(\$sp)	
addi	\$sp	\$sp	4
jr	\$ra		

Save \$s0 on stack

Procedure body

Result

Restore \$s0

Return

```
int leaf_example (int g, h,  
i, j)  
{ int f;  
  f = (g + h) - (i + j);  
  return f;  
}
```

# Non-Leaf Procedures

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- **Procedures that call other procedures**
- **For nested call, caller needs to save on the stack:**
  - Its return address
  - Any arguments and temporaries needed after the call
- **Restore from the stack after the call**

# Non-Leaf Procedure Example

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- C code:

```
int fact (int n)
{
    if (n < 1) return f;
    else return n * fact(n - 1);
}
```

- Argument n in \$a0
- Result in \$v0

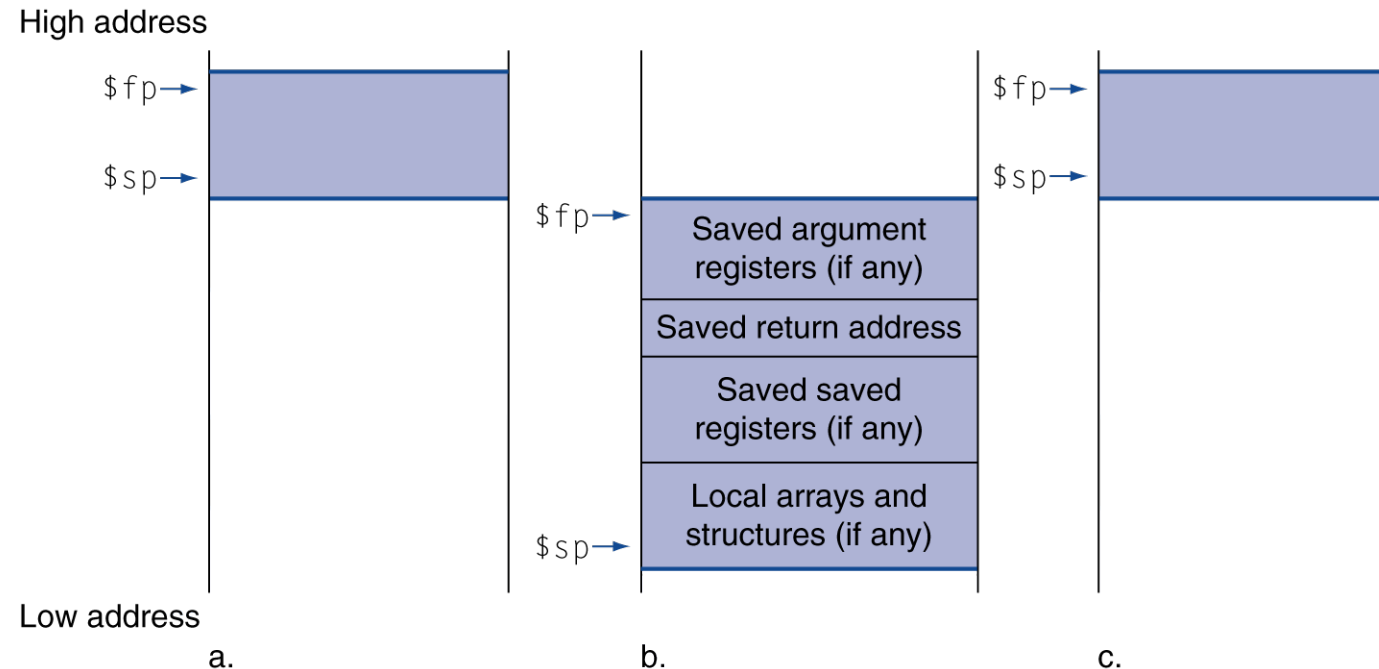
# Non-Leaf Procedure Example

- MIPS code:

fact:		
addi	\$sp, \$sp, -8	# adjust stack for 2 items
sw	\$ra, 4(\$sp)	# save return address
sw	\$a0, 0(\$sp)	# save argument
slti	\$t0, \$a0, 1	# test for n < 1
beq	\$t0, \$zero, L1	
addi	\$v0, \$zero, 1	# if so, result is 1
addi	\$sp, \$sp, 8	# pop 2 items from stack
jr	\$ra	# and return
L1:	addi \$a0, \$a0, -1	# else decrement n
	jal fact	# recursive call
lw	\$a0, 0(\$sp)	# restore original n
lw	\$ra, 4(\$sp)	# and return address
addi	\$sp, \$sp, 8	# pop 2 items from stack
mul	\$v0, \$a0, \$v0	# multiply to get result
jr	\$ra	# and return

```
int fact (int n)
{
    if (n < 1) return f;
    else return n *
fact(n - 1);
}
```

# Local Data on the Stack



- **Local data allocated by callee**
  - e.g., C automatic variables
- **Procedure frame (activation record)**
  - Used by some compilers to manage stack storage

# Branch Addressing

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- **Branch instructions specify**
  - Opcode, two registers, target address
- **Most branch targets are near branch**
  - Forward or backward



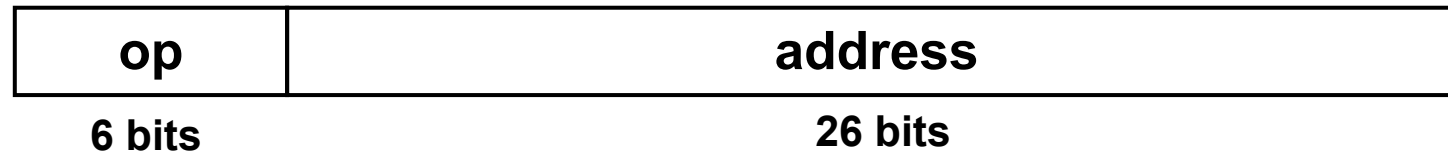
- **PC-relative addressing**
  - **Target address = PC + offset × 4**
  - **PC already incremented by 4 by this time**



# Jump Addressing

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- **Jump (j and jal) targets could be anywhere in text segment**
  - Encode full address in instruction



- **(Pseudo)Direct jump addressing**
  - **Target address =  $PC_{31...28} : (\text{address} \times 4)$**

# Target Addressing Example

- **Loop code from earlier example**

- Assume Loop at location 80000

Loop: sll	\$t1, \$s3, 2	80000	0	0	19	9	2	0
add	\$t1, \$t1, \$s6	80004	0	9	22	9	0	32
lw	\$t0, 0(\$t1)	80008	35	9	8	0		
bne	\$t0, \$s5, Exit	80012	5	8	21	2		
addi	\$s3, \$s3, 1	80016	8	19	19	1		
j	Loop	80020	2	20000				
Exit: ...		80024						

# Branching Far Away

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- If branch target is too far to encode with 16-bit offset, assembler rewrites the code

- Example

```
        beq $s0,$s1, L1
           ↓
        bne $s0,$s1, L2
L2:      j  L1
        ...
```

# Addressing Mode Summary

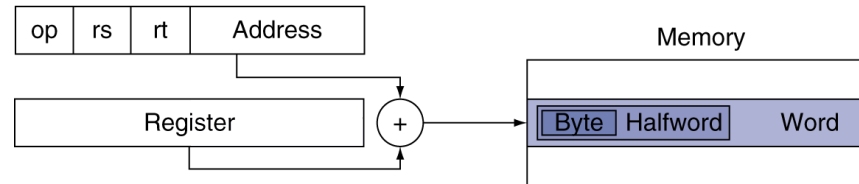
## 1. Immediate addressing



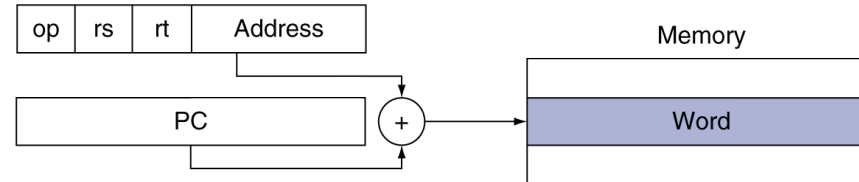
## 2. Register addressing



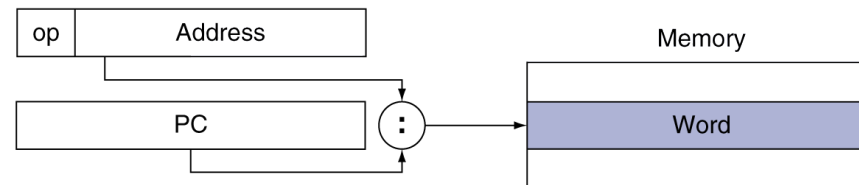
## 3. Base addressing



## 4. PC-relative addressing



## 5. Pseudodirect addressing



# QtSpim Tutorial

# Program Structure

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```
.data      # variable declarations follow this line
           # ...

.text      # instructions follow this line

main:      # indicates start of code (first instruction to execute)
           # ...

# End of program, leave a blank line afterwards to make SPIM happy
```

# System Calls

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- To request a service, a program loads the system call:
- Code into register \$v0 and
- The arguments into registers \$a0, ..., \$a3 (or \$f12 for floating point values).
- System calls that return values put their result in register \$v0 (or \$f0 for floating point results)

# Data Declarations

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- Declares variable names used in program;
- Storage allocated in main memory (RAM)

**format for declarations:**

name:            storage\_type    value(s)



# Code

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- **Placed in section of text identified with assembler directive `.text`**
- **Contains program code (instructions)**
- **Starting point for code execution given label `main`:**
- **Ending point of main code should use `exit` system call (see below under System Calls)**

# System Calls

Service	System Call Code	Arguments	Result
print integer	1	\$a0 = value	(none)
print float	2	\$f12 = float value	(none)
print double	3	\$f12 = double value	(none)
print string	4	\$a0 = address of string	(none)
read integer	5	(none)	\$v0 = value read
read float	6	(none)	\$f0 = value read
read double	7	(none)	\$f0 = value read
read string	8	\$a0 = address where string to be stored \$a1 = number of characters to read + 1	(none)
memory allocation	9	\$a0 = number of bytes of storage desired	\$v0 = address of block
exit (end of program)	10	(none)	(none)
print character	11	\$a0 = integer	(none)
read character	12	(none)	char in \$v0

# Hello World Example

**helloWorld.s or .asm file formats**