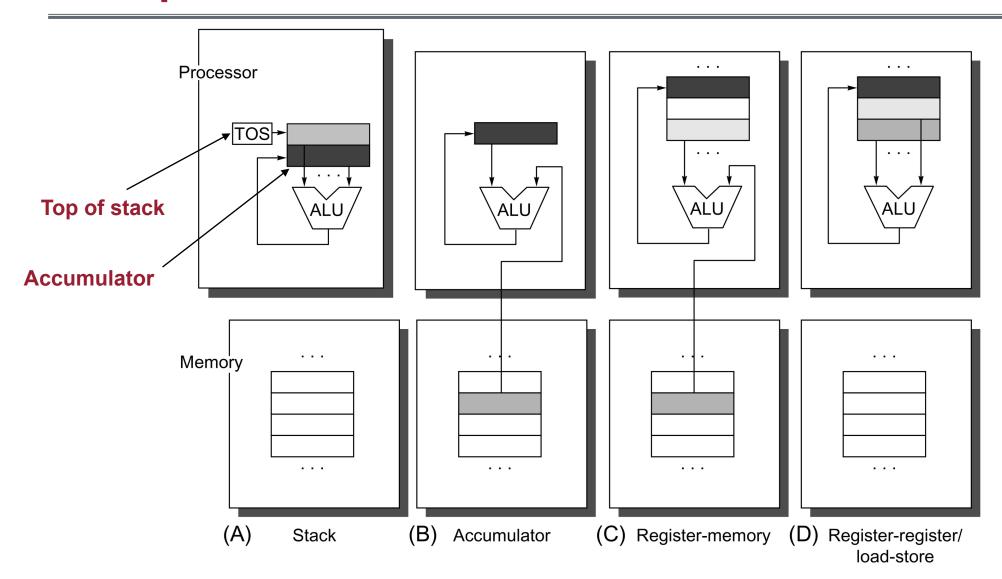
CPT_S 260 Intro to Computer Architecture Lecture 13

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Recap: ISA Classification Visualization



CPT_S 260

Recap: Issues in Instruction Set Design

- Memory addressing
- Type and size of operands
- Operations in the instruction set
- Encoding of ISAs
- Implementation (pipelining, scheduling, etc.)

Recap: MIPS Instruction Goals

- Showing how it is represented in hardware
- The relationship between high-level programming languages and this more primitive one
 - Examples in C and JAVA
- See the impact of programming languages and compiler optimization on performance
- Writing programs in the language of the computer and running them on the simulator

Recap: Arithmetic Operations

- Add and subtract, three operands
 - Two sources and one destination
 - add a, b, c # a gets b + c
- All arithmetic operations have this form
- Design Principle 1: Simplicity favors regularity
 - Regularity makes implementation simpler
 - Simplicity enables higher performance at lower cost

Register Operands

- Arithmetic instructions use register operands
- MIPS has a 32 × 32-bit register file
 - Use for frequently accessed data
 - Numbered 0 to 31
 - 32-bit data called a "word"

Assembler names

- -\$t0, \$t1, ..., \$t9 for temporary values
- -\$s0, \$s1, ..., \$s7 for saved variables

Design Principle 2: Smaller is faster

-c.f. main memory: millions of locations

Register Operand Example

C code:

```
f = (g + h) - (i + j);
For instance, f, ..., j stored in $50, ..., $54
```

Compiled MIPS code:

```
add $t0, $s1, $s2
add $t1, $s3, $s4
sub $s0, $t0, $t1
```

Data Types

- Instructions are all 32 bits
- Byte(8 bits), halfword (2 bytes), word (4 bytes)
- A character requires 1 byte of storage
- An integer requires 1 word (4 bytes) of storage

Memory Operands

Main memory used for composite data

- Arrays, structures, dynamic data

To apply arithmetic operations

- Load values from memory into registers
- -Store result from register to memory

Memory is byte addressed

Each address identifies an 8-bit byte

Words are aligned in memory

Address must be a multiple of 4

MIPS is Big Endian

- Most-significant byte at least address of a word
- c.f. Little Endian: least-significant byte at least address

Memory Operand Example #1

C code:

$$g = h + A[8];$$

- Assume that g is in \$s1, h in \$s2, base address of A in \$s3
- Compiled MIPS code:
 - Index 8 requires offset of 32
 - » 4 bytes per word

offset

base register

Memory Operand Example #2

C code:

$$A[12] = h + A[8];$$

– h in \$s2, base address of A in \$s3

Compiled MIPS code:

Index 8 requires offset of 32

```
lw $t0, 32($s3)  # load word
add $t0, $s2, $t0
sw $t0, 48($s3)  # store word
```

Registers vs. Memory

- Registers are faster to access than memory
- Operating on memory data requires loads and stores
 - More instructions to be executed
- Compiler must use registers for variables as much as possible
 - Only spill to memory for less frequently used variables
 - Register optimization is important!

Immediate Operands

Constant data specified in an instruction

```
addi $s3, $s3, 4
```

- No subtract immediate instruction
 - Just use a negative constantaddi \$\$2\$, \$\$\$1\$, \$\$-1\$
- Design Principle 3: Make the common case fast
 - Small constants are common
 - Immediate operand avoids a load instruction

The Constant Zero

- MIPS register 0 (\$zero) is the constant 0
 - Cannot be overwritten
- Useful for common operations
 - E.g., move between registersadd \$t2, \$s1, \$zero

Register Number	Alternative Name	Description		
0	zero	the value 0		
1	\$at	(assembler temporary) reserved by the assembler		
2-3	\$v0 - \$v1	(values) from expression evaluation and function results		
4-7	\$a0 - \$a3	(arguments) First four parameters for subroutine. Not preserved across procedure calls		
8-15	\$t0 - \$t7	(temporaries) Caller saved if needed. Subroutines can use w/out saving Not preserved across procedure calls		
16-23	\$s0 - \$s7	(saved values) - Callee saved. A subroutine using one of these must save original and restore it before exiting. Preserved across procedure calls		
24-25	\$t8 - \$t9	(temporaries) Caller saved if needed. Subroutines can use w/out saving These are in addition to \$t0 - \$t7 above. Not preserved across procedure calls.		
26-27	\$k0 - \$k1	reserved for use by the interrupt/trap handler		
28	\$gp	global pointer. Points to the middle of the 64K block of memory in the static data segment.		
29	\$sp	stack pointer Points to last location on the stack.		
30	\$s8/\$fp	saved value / frame pointer Preserved across procedure calls		
31	\$ra	return address		

Sign Extension

Representing a number using more bits

Preserve the numeric value

In MIPS instruction set

- addi: extend immediate value
- 1b, 1h: extend loaded byte/halfword
- beq, bne: extend the displacement

Replicate the sign bit to the left

- Unsigned values: extend with 0s
- Examples: 8-bit to 16-bit
 - -+2: 0000 0010 => 0000 0000 0000 0010
 - --2: **1**111 1110 => **1**111 **1**111 1111 1110

Logical Operations

Instructions for bitwise manipulation

Operation	С	Java	MIPS
Shift left	<<	<<	sll
Shift right	>>	>>>	srl
Bitwise AND	&	&	and, andi
Bitwise OR			or, ori
Bitwise NOT	~	~	nor

Useful for extracting and inserting groups of bits in a word

AND Operations

- Useful to mask bits in a word
 - Select some bits, clear others to 0

and \$t0, \$t1, \$t2

\$t2 | 0000 0000 0000 0000 00<mark>00 11</mark>01 1100 0000

\$t0 | 0000 0000 0000 00<mark>00 11</mark>00 0000 0000

OR Operations

- Useful to include bits in a word
 - Set some bits to 1, leave others unchanged

or \$t0, \$t1, \$t2

\$t2 | 0000 0000 0000 000<mark>00 11</mark>01 1100 0000

\$t1 | 0000 0000 0000 000<mark>11 11</mark>00 0000 0000

\$t0 | 0000 0000 0000 000<mark>11 11</mark>01 1100 0000

NOT Operations

- Useful to invert bits in a word
 - Change 0 to 1, and 1 to 0
- MIPS has NOR 3-operand instruction
 - a NOR b == NOT (a OR b)

nor \$t0, \$t1, \$zero ← Register 0: always read as zero