

# **CPT\_S 260 Intro to Computer Architecture**

## **Lecture 17**

**qTSPIM Tutorial**  
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# Announcements

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- **Mid term exam 1**
  - Will include everything up to today's lecture
  - I will post practice questions by Monday
  - Homework 3 solutions will also be posted for you to study
- **Quiz 3 is online**
  - Complete by tomorrow

# Procedure Calling

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- **Steps required**
  - 1. Place parameters in registers**
  - 2. Transfer control to procedure**
  - 3. Acquire storage for procedure**
  - 4. Perform procedure's operations**
  - 5. Place result in register for caller**
  - 6. Return to place of call**

# Register Usage

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- **\$a0 – \$a3: arguments (reg's 4 – 7)**
- **\$v0, \$v1: result values (reg's 2 and 3)**
- **\$t0 – \$t9: temporaries**
  - Can be overwritten by callee
- **\$s0 – \$s7: saved**
  - Must be saved/restored by callee
- **\$gp: global pointer for static data (reg 28)**
- **\$sp: stack pointer (reg 29)**
- **\$fp: frame pointer (reg 30)**
- **\$ra: return address (reg 31)**

# Procedure Call Instructions

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- **Procedure call: jump and link**

`jal ProcedureLabel`

- Address of following instruction put in \$ra
- Jumps to target address

- **Procedure return: jump register**

`jr $ra`

- Copies \$ra to program counter
- Can also be used for computed jumps
  - » e.g., for case/switch statements

# Leaf Procedure Example

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- C code:

```
int leaf_example (int g, h, i, j)
{ int f;
  f = (g + h) - (i + j);
  return f;
}
```

- Arguments g, ..., j in \$a0, ..., \$a3
- f in \$s0 (hence, need to save \$s0 on stack)
- Result in \$v0

# Leaf Procedure Example

## ■ MIPS code:

leaf_example:			
addi	\$sp	\$sp	-4
sw	\$s0	0(\$sp)	
add	\$t0	\$a0	\$a1
add	\$t1	\$a2	\$a3
sub	\$s0	\$t0	\$t1
add	\$v0	\$s0	\$zero
lw	\$s0	0(\$sp)	
addi	\$sp	\$sp	4
jr	\$ra		

Save \$s0 on stack

Procedure body

Result

Restore \$s0

Return

```
int leaf_example (int g, h,  
i, j)  
{ int f;  
  f = (g + h) - (i + j);  
  return f;  
}
```

# Non-Leaf Procedures

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- **Procedures that call other procedures**
- **For nested call, caller needs to save on the stack:**
  - Its return address
  - Any arguments and temporaries needed after the call
- **Restore from the stack after the call**



# Non-Leaf Procedure Example

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- C code:

```
int fact (int n)
{
    if (n < 1) return f;
    else return n * fact(n - 1);
}
```

- Argument n in \$a0
- Result in \$v0

# Non-Leaf Procedure Example

- MIPS code:

fact:		
addi	\$sp, \$sp, -8	# adjust stack for 2 items
sw	\$ra, 4(\$sp)	# save return address
sw	\$a0, 0(\$sp)	# save argument
slti	\$t0, \$a0, 1	# test for n < 1
beq	\$t0, \$zero, L1	
addi	\$v0, \$zero, 1	# if so, result is 1
addi	\$sp, \$sp, 8	# pop 2 items from stack
jr	\$ra	# and return
L1:	addi \$a0, \$a0, -1	# else decrement n
	jal fact	# recursive call
lw	\$a0, 0(\$sp)	# restore original n
lw	\$ra, 4(\$sp)	# and return address
addi	\$sp, \$sp, 8	# pop 2 items from stack
mul	\$v0, \$a0, \$v0	# multiply to get result
jr	\$ra	# and return

```
int fact (int n)
{
    if (n < 1) return f;
    else return n *
fact(n - 1);
}
```

# Sequence of Execution with $n = 2$

Main before factorial

```
fact: (n = 2)
    addi $sp, $sp, -8    # adjust stack for 2 items
    sw   $ra, 4($sp)    # save return address
    sw   $a0, 0($sp)    # save argument
    slti $t0, $a0, 1    # test for n < 1
    beq  $t0, $zero, L1
```

```
L1:  addi $a0, $a0, -1    # else decrement n
     jal  fact          # recursive call
```

```
fact: (n = 1)
    addi $sp, $sp, -8    # adjust stack for 2 items
    sw   $ra, 4($sp)    # save return address
    sw   $a0, 0($sp)    # save argument
    slti $t0, $a0, 1    # test for n < 1
    beq  $t0, $zero, L1
```

```
L1:  addi $a0, $a0, -1    # else decrement n
     jal  fact          # recursive call
```

```
fact: (n = 0)
    addi $sp, $sp, -8    # adjust stack for 2 items
    sw   $ra, 4($sp)    # save return address
    sw   $a0, 0($sp)    # save argument
    slti $t0, $a0, 1    # test for n < 1
    beq  $t0, $zero, L1
```

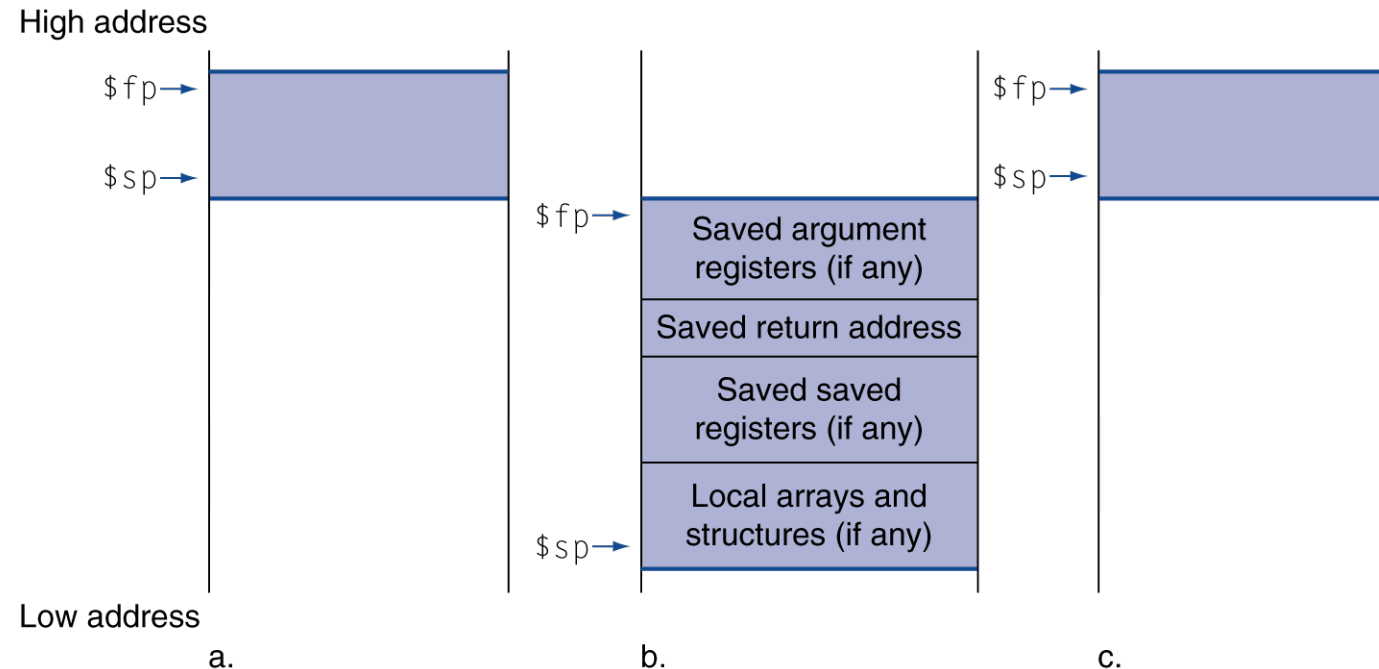
```
    addi $v0, $zero, 1    # if so, result is 1
    addi $sp, $sp, 8      # pop 2 items from stack
    jr   $ra
```

```
lw   $a0, 0($sp)        # restore original n = 1
lw   $ra, 4($sp)        # and return address
addi $sp, $sp, 8        # pop 2 items from stack
mul  $v0, $a0, $v0      # multiply to get result
jr   $ra                # and return
```

```
lw   $a0, 0($sp)        # restore original n = 2
lw   $ra, 4($sp)        # and return address
addi $sp, $sp, 8        # pop 2 items from stack
mul  $v0, $a0, $v0      # multiply to get result
jr   $ra                # and return to main
```

Main after factorial

# Local Data on the Stack



- **Local data allocated by callee**
  - e.g., C automatic variables
- **Procedure frame (activation record)**
  - Used by some compilers to manage stack storage

# Branch Addressing

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- **Branch instructions specify**
  - Opcode, two registers, target address
- **Most branch targets are near branch**
  - Forward or backward

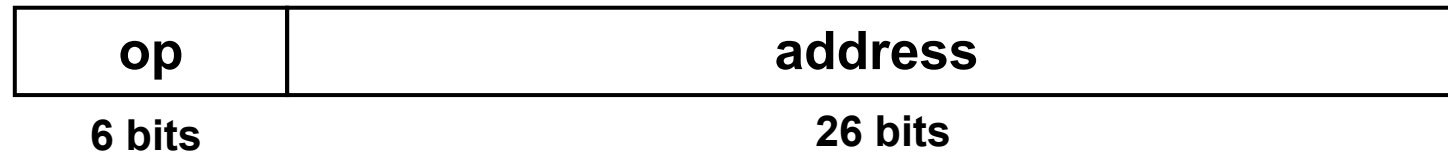


- **PC-relative addressing**
  - **Target address = PC + offset × 4**
  - **PC already incremented by 4 by this time**

# Jump Addressing

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- **Jump (j and jal) targets could be anywhere in text segment**
  - Encode full address in instruction



- **(Pseudo)Direct jump addressing**
  - **Target address =  $PC_{31...28} : (\text{address} \times 4)$**

# Target Addressing Example

- **Loop code from earlier example**

- Assume Loop at location 80000

Loop:	sll	\$t1, \$s3, 2	80000	0	0	19	9	2	0
	add	\$t1, \$t1, \$s6	80004	0	9	22	9	0	32
	lw	\$t0, 0(\$t1)	80008	35	9	8	0		
	bne	\$t0, \$s5, Exit	80012	5	8	21	2		
	addi	\$s3, \$s3, 1	80016	8	19	19	1		
	j	Loop	80020	2	20000				
Exit:	...		80024						

# Branching Far Away

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- If branch target is too far to encode with 16-bit offset, assembler rewrites the code

- Example

```
        beq $s0,$s1, L1
           ↓
        bne $s0,$s1, L2
L2:      j  L1
        ...
```



# Addressing Mode Summary

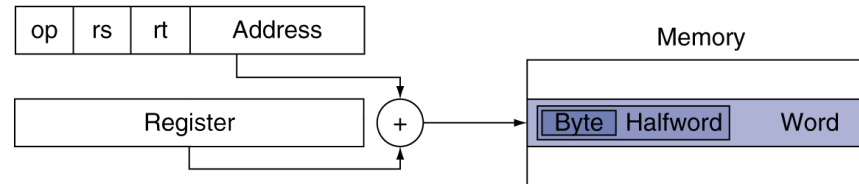
## 1. Immediate addressing



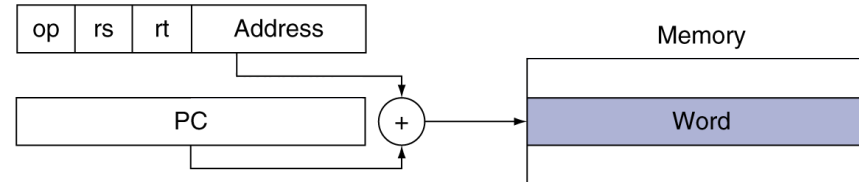
## 2. Register addressing



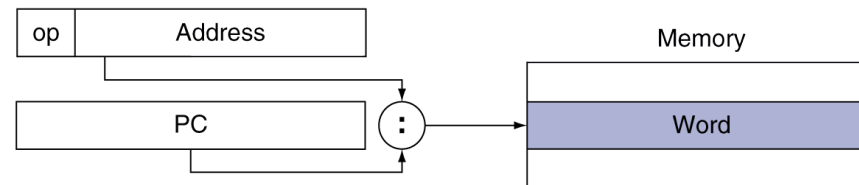
## 3. Base addressing



## 4. PC-relative addressing



## 5. Pseudodirect addressing



# QtSpim Tutorial

# Program Structure

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```
.data      # variable declarations follow this line
           # ...

.text      # instructions follow this line

main:      # indicates start of code (first instruction to execute)
           # ...

# End of program, leave a blank line afterwards to make SPIM happy
```

# System Calls

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- To request a service, a program loads the system call:
- Code into register \$v0 and
- The arguments into registers \$a0, ..., \$a3 (or \$f12 for floating point values).
- System calls that return values put their result in register \$v0 (or \$f0 for floating point results)

# Data Declarations

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- Declares variable names used in program;
- Storage allocated in main memory (RAM)

**format for declarations:**

name:            storage\_type    value(s)

# Code

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- **Placed in section of text identified with assembler directive `.text`**
- **Contains program code (instructions)**
- **Starting point for code execution given label `main`:**
- **Ending point of main code should use `exit` system call (see below under System Calls)**

# System Calls

Service	System Call Code	Arguments	Result
print integer	1	\$a0 = value	(none)
print float	2	\$f12 = float value	(none)
print double	3	\$f12 = double value	(none)
print string	4	\$a0 = address of string	(none)
read integer	5	(none)	\$v0 = value read
read float	6	(none)	\$f0 = value read
read double	7	(none)	\$f0 = value read
read string	8	\$a0 = address where string to be stored \$a1 = number of characters to read + 1	(none)
memory allocation	9	\$a0 = number of bytes of storage desired	\$v0 = address of block
exit (end of program)	10	(none)	(none)
print character	11	\$a0 = integer	(none)
read character	12	(none)	char in \$v0

# Hello World Example

**helloWorld.s or .asm file formats**