



Definitions, Theorems, and Proofs

Intro to Theory of Computation +
Math Review – Part 3



Definitions, theorems, and proofs

“Theorems and proofs are the heart and soul of mathematics and definitions are its spirit”....Sipser

- **Definitions** describe the objects and notions that we use.
- A **proof** is a *convincing* logical argument that a statement is true
 - Convincing in an absolute sense (“beyond reasonable doubt” is not enough. Mathematics demands proof beyond **any** doubt.)
- A **theorem** is a mathematical statement proved true.



Theorem and its cousins

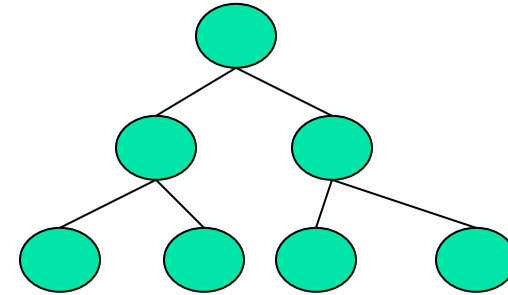
- Generally, we reserve the use of the word *theorem* for statements of special interest.
- Occasionally we prove statements that are interesting only because they assist in the proof of another, more significant statement.
 - Such statements are called **lemmas**
- Occasionally a theorem or its proof may allow us to conclude easily that other, related statements are true.
 - Such statements are called **corollaries** of the theorem

An example

Theorem: *The height of an n -node binary tree is at least $\text{floor}(\log n)$*

Lemma: *Level i of a perfect binary tree has 2^i nodes.*

Corollary: *A perfect binary tree of height h has $2^{h+1}-1$ nodes.*





Quantifiers

“For all” or “For every”

- Universal proofs
- Notation: \forall

“There exists”

- Used in existential proofs
- Notation: \exists

Implication is denoted by \Rightarrow

- E.g., “IF A THEN B” can also be written as “ $A \Rightarrow B$ ”



Finding proofs

- Finding proofs isn't always easy
- Even though no one has a recipe for producing proofs, some helpful general strategies are available
 - Carefully read the statement you want to prove. Rewrite the statement in your own words.
 - Break it down and consider each part separately
 - E.g 1. P if and only if Q statement
 - E.g 2. $A = B$ statement
- Tips for producing a proof:
 - ***Be patient. Come back to it. Be neat. Be concise.***



Proof techniques

- **By construction**

- Many theorems state that a particular type of object exists. One way to prove such a theorem is by demonstrating how to construct the object.

- Example:

Theorem. For each even number n greater than 2, there exists a 3-regular graph with n nodes



Example of proof by construction

THEOREM 0.22

For each even number n greater than 2, there exists a 3-regular graph with n nodes.

PROOF Let n be an even number greater than 2. Construct graph $G = (V, E)$ with n nodes as follows. The set of nodes of G is $V = \{0, 1, \dots, n-1\}$, and the set of edges of G is the set

$$E = \{ \{i, i+1\} \mid \text{for } 0 \leq i \leq n-2 \} \cup \{ \{n-1, 0\} \} \\ \cup \{ \{i, i+n/2\} \mid \text{for } 0 \leq i \leq n/2-1 \}.$$

Picture the nodes of this graph written consecutively around the circumference of a circle. In that case, the edges described in the top line of E go between adjacent pairs around the circle. The edges described in the bottom line of E go between nodes on opposite sides of the circle. This mental picture clearly shows that every node in G has degree 3.

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Proof techniques

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- Example:
Theorem. For each even number n greater than 2, there exists a 3-regular graph with n nodes

- **By contradiction**

- Start with the statement contradictory to the given statement
- Example. Prove that $\sqrt{2}$ is irrational.
 - (Start by claiming that $\sqrt{2}$ is rational, and so it can be written as a ratio of two integers and arrive at a contradiction.)



Details on proof of $\sqrt{2}$ is irrational

- $\sqrt{2} = m/n$
- $n \cdot \sqrt{2} = m$
- $2n^2 = m^2$
- $2n^2 = (2k)^2$
- $2n^2 = 4k^2$
- $n^2 = 2k^2$



Proof techniques

- **By induction**
 - (2 parts) **Basis**, induction hypothesis, **induction step**

The format for writing down a proof by induction is as follows.

Basis: Prove that $\mathcal{P}(1)$ is true.

⋮

Induction step: For each $i \geq 1$, assume that $\mathcal{P}(i)$ is true and use this assumption to show that $\mathcal{P}(i + 1)$ is true.

⋮



Proof techniques

- By induction
 - (2 parts) **Basis**, induction hypothesis, **induction step**
- (By counter-example)
 - Show an example that disproves the claim
- Note: There is no such thing called a “proof by example”!
 - So, when asked to prove a claim, an example that satisfied that claim is *not* a proof



Different ways of saying the same thing

- “*If H then C*”:
 - i. *H implies C*
 - ii. $H \Rightarrow C$
 - iii. *C if H*
 - iv. *H only if C*
 - v. *Whenever H holds, C follows*



“If-and-Only-If” statements

- “A if and only if B” ($A \iff B$)
 - (if part) if B then A (\implies)
 - (only if part) A only if B (\implies)
(same as “if A then B”)
- “If and only if” is abbreviated as “iff”
 - i.e., “A iff B”
- Example:
 - Theorem: *Let x be a real number. Then floor of x = ceiling of x if and only if x is an integer.*
- Proofs for iff have two parts
 - One for the “if part” & another for the “only if part”



Summary (of last three lectures)

- Theory of computation overview
- Mathematical notions and terminology
 - Sets
 - Sequences and tuples
 - Functions and relations
 - Graphs
 - Strings and languages
 - Boolean logic
- Definitions, theorems, and proofs
- Proof techniques
 - By construction
 - By contradiction
 - By induction



HW1 is out

- Due: Fri Jan 28
- Has 6 problems
 - 4 on what we covered so far (look back)
 - 2 on what we will cover next lecture (look forward)
 - 1 involves history of computation
- Submission on Canvas: PDF